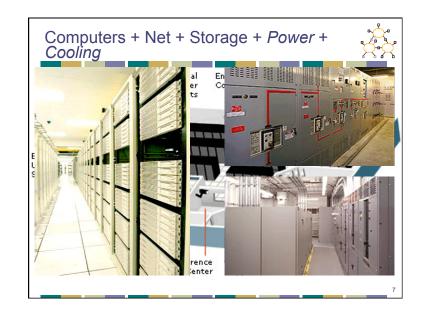


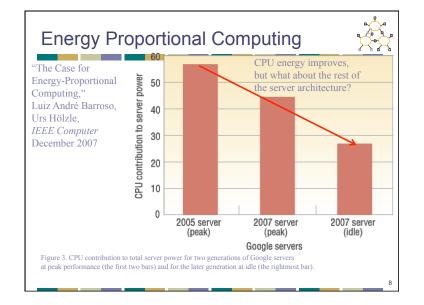
Datacenter Arms Race

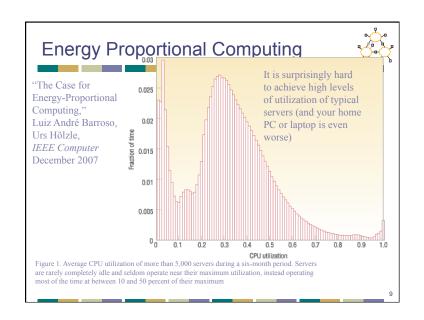


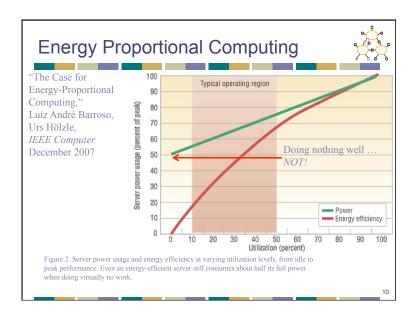
- Amazon, Google, Microsoft, Yahoo!, ... race to build next-gen mega-datacenters
 - Industrial-scale Information Technology
 - 100,000+ servers
 - Located where land, water, fiber-optic connectivity, and cheap power are available
- E.g., Microsoft Quincy
 - 43600 sq. ft. (10 football fields), sized for 48 MW
 - · Also Chicago, San Antonio, Dublin @\$500M each
- · E.g., Google:
 - The Dalles OR, Pryor OK, Council Bluffs, IW, Lenoir 5 NC, Goose Creek, SC

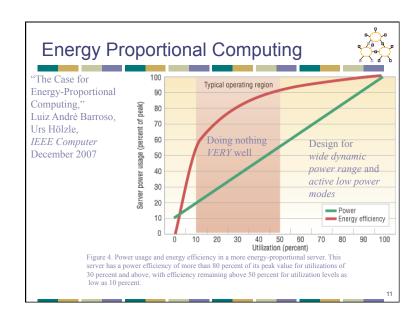


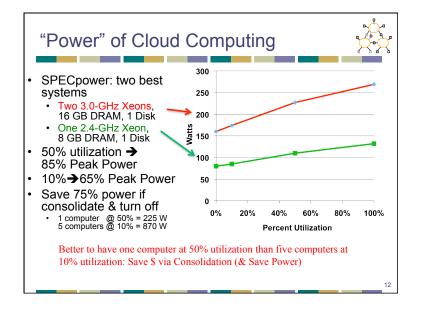












Bringing Resources On-/Off-line



- Save power by taking DC "slices" off-line
 - · Resource footprint of applications hard to model
 - Dynamic environment, complex cost functions require measurement-driven decisions -- opportunity for statistical machine learning
 - Must maintain Service Level Agreements, no negative impacts on hardware reliability
 - Pervasive use of virtualization (VMs, VLANs, VStor) makes feasible rapid shutdown/migration/restart
- Recent results suggest that conserving energy may actually improve reliability
 - MTTF: stress of on/off cycle vs. benefits of off-hours

Aside: Disk Power

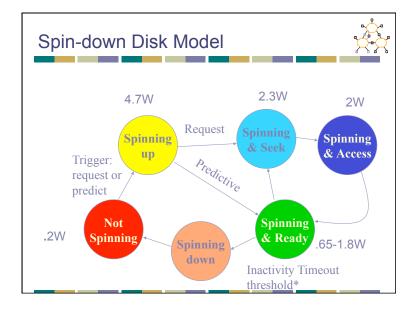


IBM Microdrive (1inch)

- writing 300mA (3.3V) 1W
- standby 65mA (3.3V) .2W

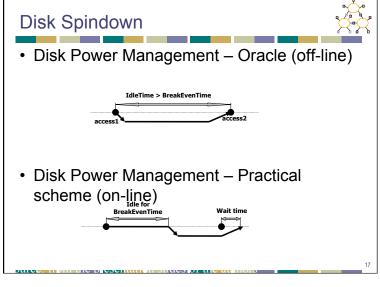
IBM TravelStar (2.5inch)

- · read/write 2W
- spinning 1.8W
- low power idle .65W
- standby .25W
- sleep .1W
- startup 4.7 W
- seek 2.3W



1

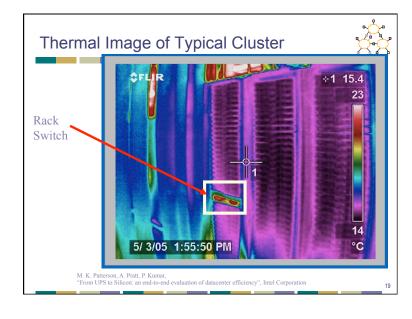
Disk Spindown • Disk Power Management – Oracle (off-line) IdleTime > BreakEvenTime • Disk Power Management – Practical scheme (on-line)



Spin-Down Policies Fixed Thresholds



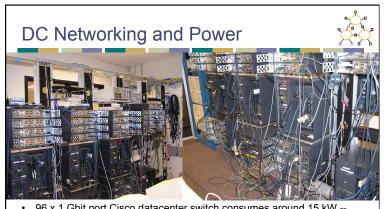
- T_{out} = spin-down cost s.t. 2*E_{transition} = P_{spin}*T_{out}
- Adaptive Thresholds: T_{out} = f (recent accesses)
 - Exploit burstiness in Tidle
- Minimizing Bumps (user annoyance/latency)
 - · Predictive spin-ups
- Changing access patterns (making burstiness)
 - Caching
 - · Prefetching



DC Networking and Power

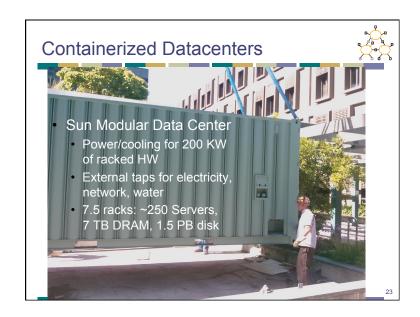


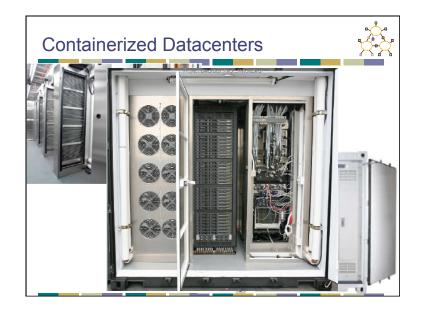
- · Within DC racks, network equipment often the "hottest" components in the hot spot
- Network opportunities for power reduction
 - Transition to higher speed interconnects (10 Gbs) at DC scales and densities
 - High function/high power assists embedded in network element (e.g., TCAMs)



- 96 x 1 Gbit port Cisco datacenter switch consumes around 15 kW -approximately 100x a typical dual processor Google server @ 145 W
- High port density drives network element design, but such high power density makes it difficult to tightly pack them with servers
- Alternative distributed processing/communications topology under investigation by various research groups





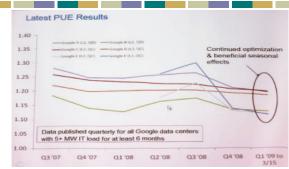


Google

- Since 2005, its data centers have been composed of standard shipping containers-each with 1,160 servers and a power consumption that can reach 250 kilowatts
- Google server was 3.5 inches thick--2U, or 2 rack units, in data center parlance. It had two processors, two hard drives, and eight memory slots mounted on a motherboard built by Gigabyte

Google's PUE





- In the third quarter of 2008, Google's PUE was 1.21, but it dropped to 1.20 for the fourth quarter and to 1.19 for the first quarter of 2009 through March 15
- · Newest facilities have 1.12

Summary

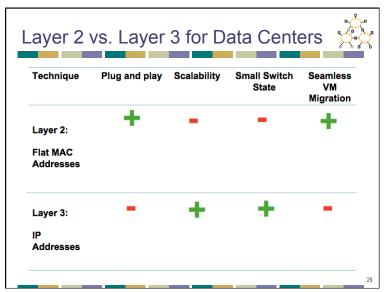


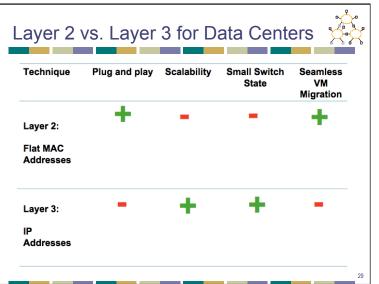
- Energy Consumption in IT Equipment
 - Energy Proportional Computing
 - Inherent inefficiencies in electrical energy distribution
- Energy Consumption in Internet Datacenters
 - · Backend to billions of network capable devices
 - Enormous processing, storage, and bandwidth supporting applications for huge user communities
 - Resource Management: Processor, Memory, I/O, Network to maximize performance subject to power constraints: "Do Nothing Well"
 - New packaging opportunities for better optimization of computing + communicating + power + mechanical

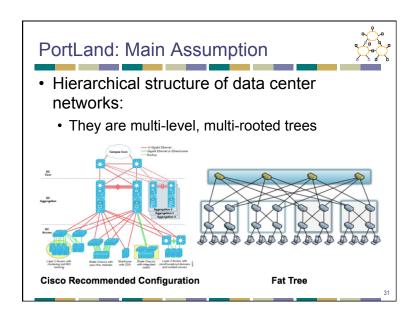
Overview



- Data Center Overview
- Networking in the DC

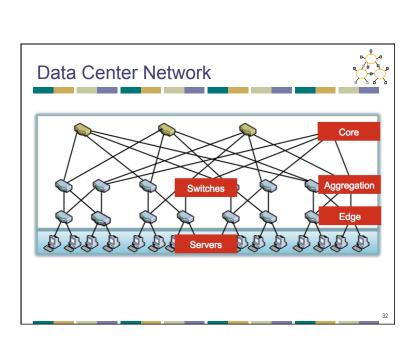




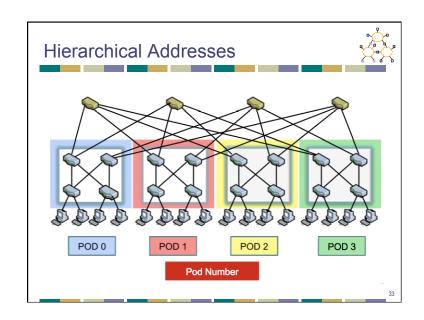


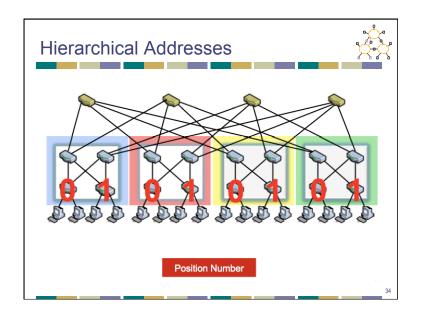
Flat vs. Location Based Addresses Commodity switches today have ~640 KB of low latency, power hungry, expensive on chip memory • Stores 32 – 64 K flow entries Assume 10 million virtual endpoints in 500,000 servers in datacenter Flat addresses → 10 million address mappings → ~100 MB on chip memory > ~150 times the memory size that can be put on chip today

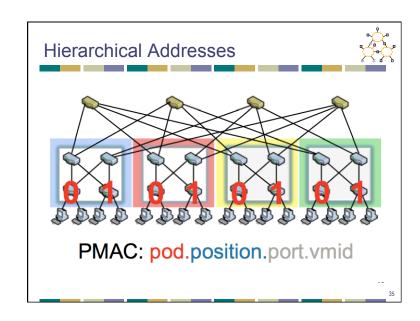
accommodated in switches today

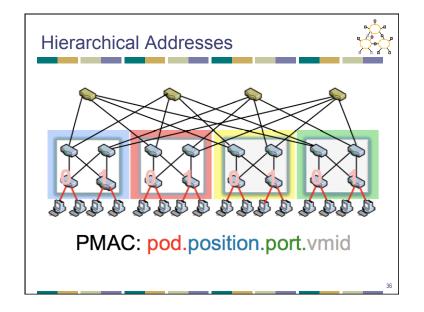


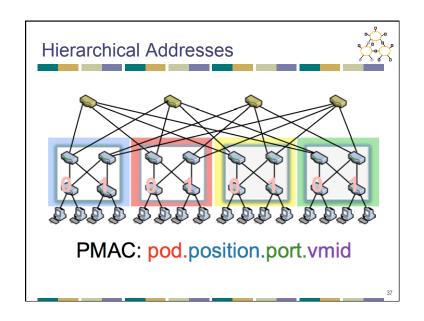
• Location based addresses → 100 – 1000 address mappings → ~10 KB of memory → easily

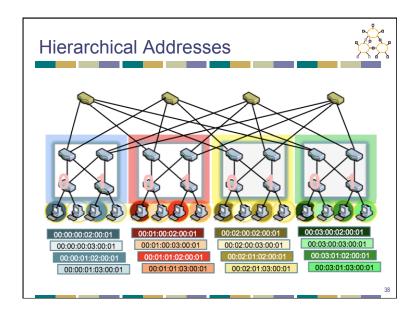


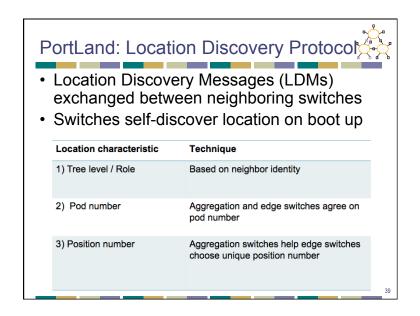


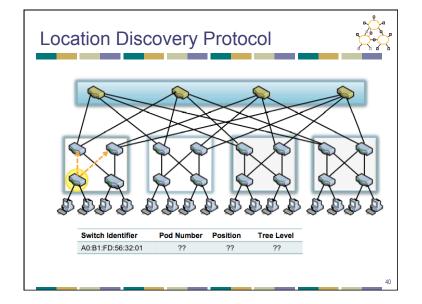


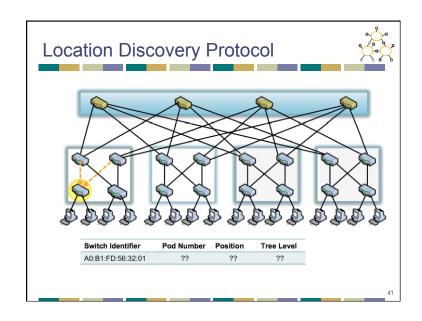


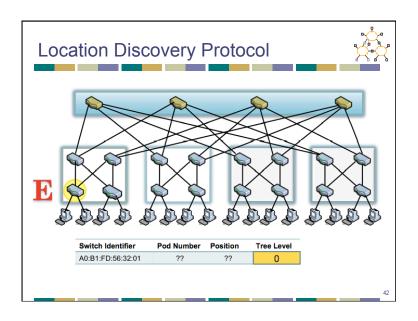


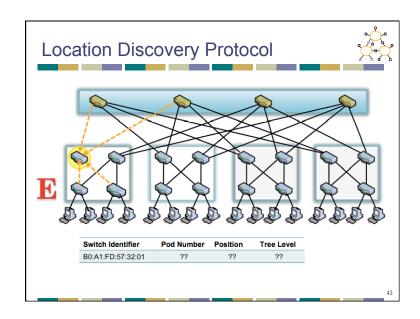


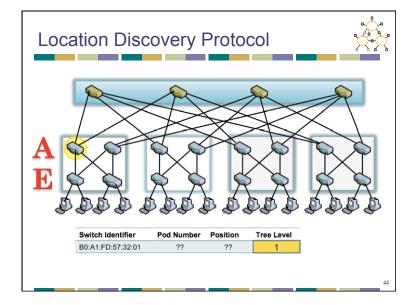


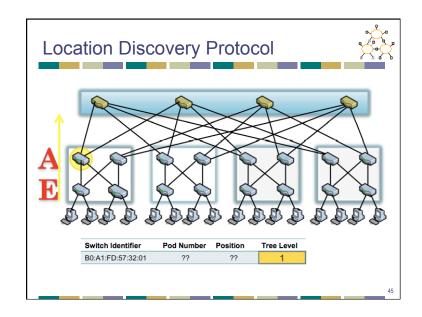


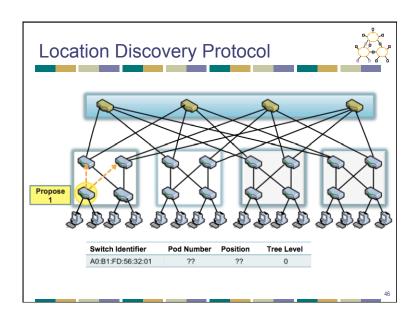


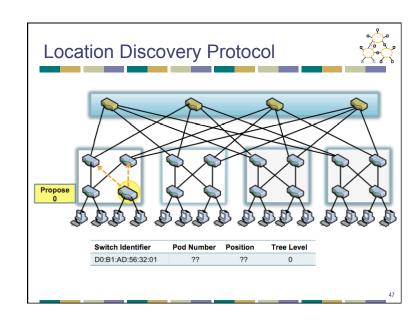


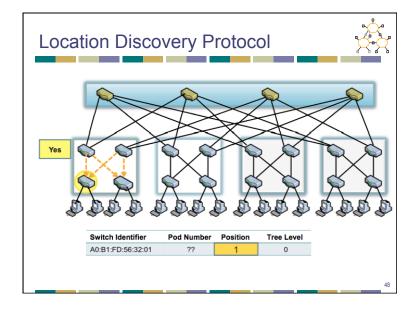


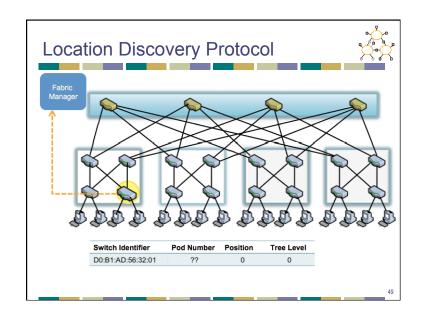


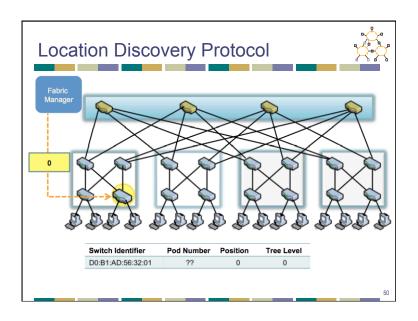


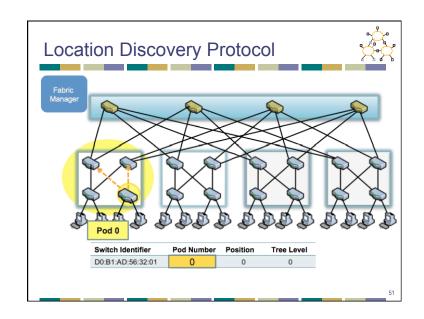


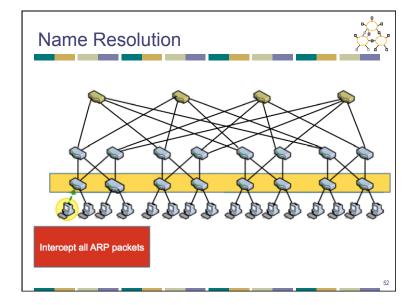


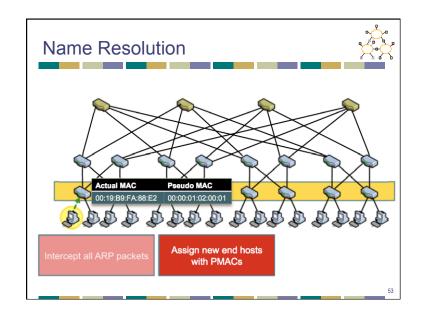


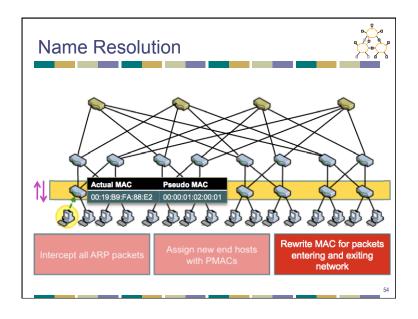


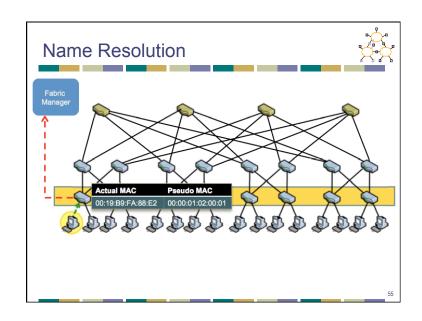


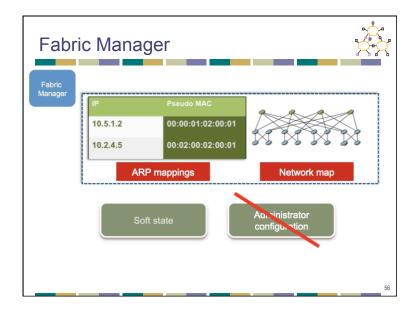


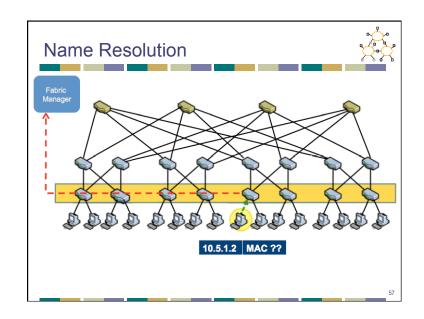


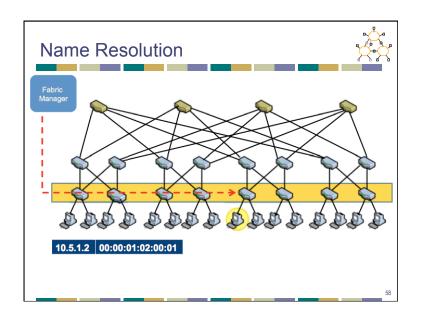


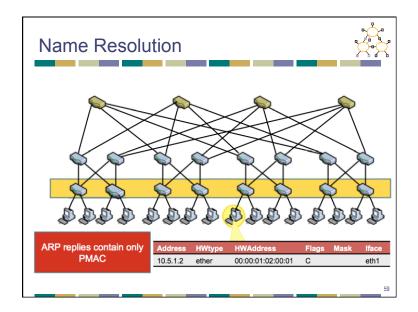












Other Schemes



- SEATTLE [SIGCOMM '08]:
 - Layer 2 network fabric that works at enterprise scale
 - Eliminates ARP broadcast, proposes one-hop DHT
 - · Eliminates flooding, uses broadcast based LSR
 - Scalability limited by
 - · Broadcast based routing protocol
 - · Large switch state
- VL2 [SIGCOMM '09]
 - Network architecture that scales to support huge data centers
 - Layer 3 routing fabric used to implement a virtual layer 2
 - Scale Layer 2 via end host modifications
 - Unmodified switch hardware and software
 - End hosts modified to perform enhanced resolution to assist routing and forwarding

