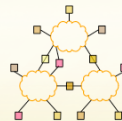


## 15-446 Distributed Systems Spring 2009



L-22 Sensor Networks

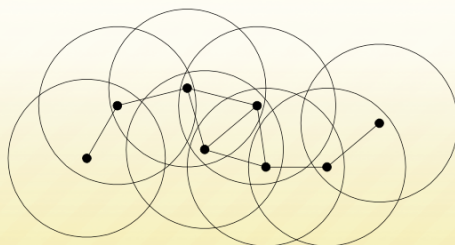
## Overview

- Ad hoc routing
- Sensor Networks
- Directed Diffusion
- Aggregation
  - TAG
  - Synopsis Diffusion

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## Ad Hoc Routing

- Goal: Communication between wireless nodes
  - No external setup (self-configuring)
  - Often need multiple hops to reach dst



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## Ad Hoc Routing

- Create multi-hop connectivity among set of wireless, possibly moving, nodes
- Mobile, wireless hosts act as forwarding nodes as well as end systems
- Need routing protocol to find multi-hop paths
  - Needs to be dynamic to adapt to new routes, movement
  - Interesting challenges related to interference and power limitations
  - Low consumption of memory, bandwidth, power
  - Scalable with numbers of nodes
  - Localized effects of link failure

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## Challenges and Variants

- Poorly-defined "links"
  - Probabilistic delivery, etc. Kind of  $n^2$  links
- Time-varying link characteristics
- No oracle for configuration (no ground truth configuration file of connectivity)
- Low bandwidth (relative to wired)
- Possibly mobile
- Possibly power-constrained

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## Problems Using DV or LS

- DV protocols may form loops
  - Very wasteful in wireless: bandwidth, power
  - Loop avoidance sometimes complex
- LS protocols: high storage and communication overhead
- More links in wireless (e.g., clusters) - may be redundant → higher protocol overhead

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## Problems Using DV or LS

- Periodic updates waste power
  - Tx sends portion of battery power into air
  - Reception requires less power, but periodic updates prevent mobile from "sleeping"
- Convergence may be slower in conventional networks but must be fast in ad-hoc networks and be done without frequent updates

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## Proposed Protocols

- Destination-Sequenced Distance Vector (DSDV)
  - DV protocol, destinations advertise sequence number to avoid loops, not on demand
- Temporally-Ordered Routing Algorithm (TORA)
  - On demand creation of hbh routes based on link-reversal
- **Dynamic Source Routing (DSR)**
  - On demand source route discovery
- Ad Hoc On-Demand Distance Vector (AODV)
  - Combination of DSR and DSDV: on demand route discovery with hbh routing

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## DSR Concepts

- Source routing
  - No need to maintain up-to-date info at intermediate nodes
- On-demand route discovery
  - No need for periodic route advertisements

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## DSR Components

- Route discovery
  - The mechanism by which a sending node obtains a route to destination
- Route maintenance
  - The mechanism by which a sending node detects that the network topology has changed and its route to destination is no longer valid

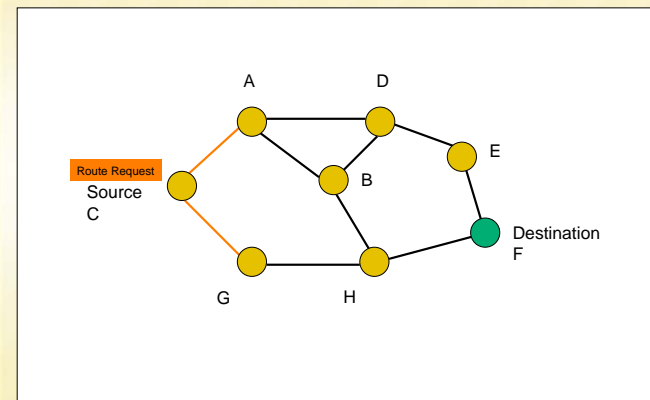
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## DSR Route Discovery

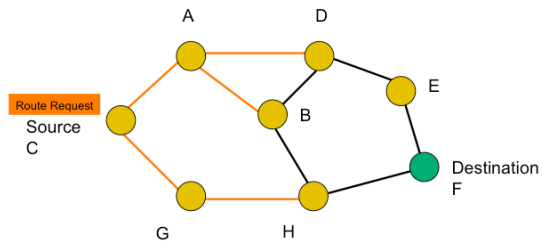
- Route discovery - basic idea
  - **Source** broadcasts route-request to **Destination**
  - Each node forwards request by adding own address and re-broadcasting
  - Requests propagate outward until:
    - Target is found, or
    - A node that has a route to Destination is found

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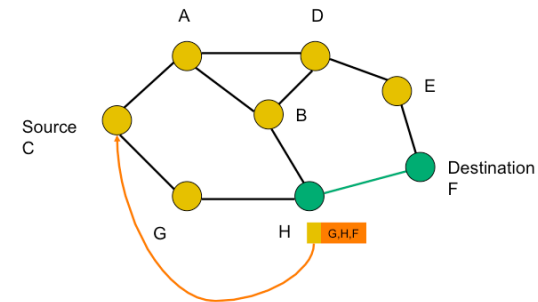
## C Broadcasts Route Request to F



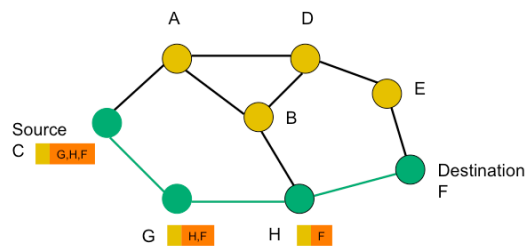
### C Broadcasts Route Request to F



### H Responds to Route Request



### C Transmits a Packet to F



### Forwarding Route Requests

- A request is forwarded if:
  - Node is not the destination
  - Node not already listed in recorded source route
  - Node has not seen request with same sequence number
  - IP TTL field may be used to limit scope
- Destination copies route into a Route-reply packet and sends it back to **Source**



## Route Cache

- All source routes learned by a node are kept in Route Cache
  - Reduces cost of route discovery
- If intermediate node receives RR for destination and has entry for destination in route cache, it responds to RR and does not propagate RR further
- Nodes overhearing RR/RP may insert routes in cache

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## Sending Data

- Check cache for route to destination
- If route exists then
  - If reachable in one hop
    - Send packet
  - Else insert routing header to destination and send
- If route does not exist, buffer packet and initiate route discovery

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## Discussion

- Source routing is good for on demand routes instead of a priori distribution
- Route discovery protocol used to obtain routes on demand
  - Caching used to minimize use of discovery
- Periodic messages avoided

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## Overview

- Ad Hoc Routing
- **Sensor Networks**
- Directed Diffusion
- Aggregation
  - TAG
  - Synopsis Diffusion

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## Sensor Net Sample Apps

Habitat Monitoring: Storm petrels on great duck island, microclimates on James Reserve.



Earthquake monitoring in shake-test sites.



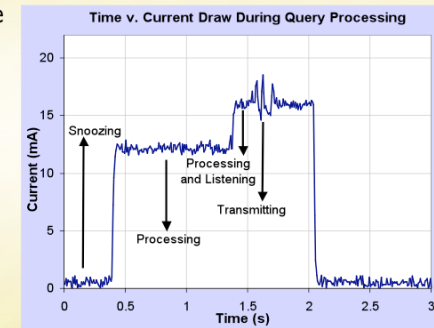
Traditional monitoring apparatus.

Vehicle detection: sensors along a road, collect data about passing vehicles.



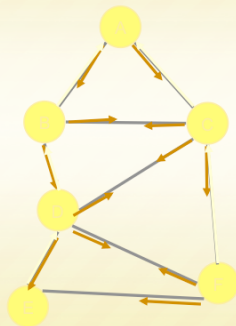
## Metric: Communication

- Lifetime from one pair of AA batteries
  - 2-3 days at full power
  - 6 months at 2% duty cycle
- Communication dominates cost
  - < few mS to compute
  - 30mS to send message



## Communication In Sensor Nets

- Radio communication has high link-level losses
  - typically about 20% @ 5m
- Ad-hoc neighbor discovery
- Tree-based routing

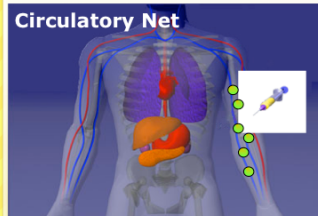


## Overview

- Ad Hoc Routing
- Sensor Networks
- Directed Diffusion
- Aggregation
  - TAG
  - Synopsis Diffusion

## The long term goal

Embed numerous distributed devices to monitor and interact with physical world: in work-spaces, hospitals, homes, vehicles, and "the environment" (water, soil, air...)



### Disaster Response

Network these devices so that they can coordinate to perform higher-level tasks.

Requires robust distributed systems of tens of thousands of devices.

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## Motivation

- Properties of Sensor Networks
  - Data centric, but not node centric
  - Have no notion of central authority
  - Are often resource constrained
- Nodes are tied to physical locations, but:
  - They may not know the topology
  - They may fail or move arbitrarily
- Problem: How can we get data from the sensors?

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## Directed Diffusion

- Data centric – nodes are unimportant
- Request driven:
  - Sinks place requests as interests
  - Sources are eventually found and satisfy interests
  - Intermediate nodes route data toward sinks
- Localized repair and reinforcement
- Multi-path delivery for multiple sources, sinks, and queries

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## Motivating Example

- Sensor nodes are monitoring a flat space for animals
- We are interested in receiving data for all 4-legged creatures seen in a rectangle
- We want to specify the data rate

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## Interest and Event Naming

- Query/interest:
  1. Type=four-legged animal
  2. Interval=20ms (event data rate)
  3. Duration=10 seconds (time to cache)
  4. Rect=[-100, 100, 200, 400]
- Reply:
  1. Type=four-legged animal
  2. Instance = elephant
  3. Location = [125, 220]
  4. Intensity = 0.6
  5. Confidence = 0.85
  6. Timestamp = 01:20:40
- Attribute-Value pairs, no advanced naming scheme

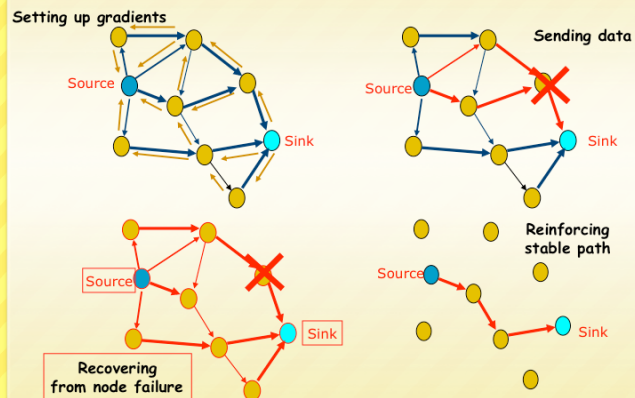
33

## Diffusion (High Level)

- Sinks broadcast interest to neighbors
- Interests are cached by neighbors
- Gradients are set up pointing back to where interests came from at low data rate
- Once a sensor receives an interest, it routes measurements along gradients

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## Illustrating Directed Diffusion



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## Summary

- Data Centric
  - Sensors net is queried for specific data
  - Source of data is irrelevant
  - No sensor-specific query
- Application Specific
  - In-sensor processing to reduce data transmitted
  - In-sensor caching
- Localized Algorithms
  - Maintain minimum local connectivity – save energy
  - Achieve global objective through local coordination
- Its gains due to aggregation and duplicate suppression may make it more viable than ad-hoc routing in sensor networks

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## Overview

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## TAG Introduction

- Programming sensor nets is hard!
- Declarative queries are easy
  - Tiny Aggregation (TAG): In-network processing via declarative queries
- In-network processing of aggregates
  - Common data analysis operation
  - Communication reducing
    - Operator dependent benefit
  - Across nodes during same epoch
- Exploit semantics improve efficiency!
- Example:
  - Vehicle tracking application: 2 weeks for 2 students
  - Vehicle tracking query: took 2 minutes to write, worked just as well!

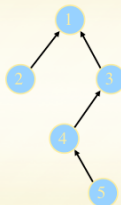


```
SELECT MAX(mag)
FROM sensors
WHERE mag > thresh
EPOCH DURATION 64ms
```

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## Basic Aggregation

- In each epoch:
  - Each node samples local sensors once
  - Generates **partial state record (PSR)**
    - local readings
    - readings from children
  - Outputs PSR during its comm. slot.
- At end of epoch, PSR for whole network output at root
- (In paper: pipelining, grouping)

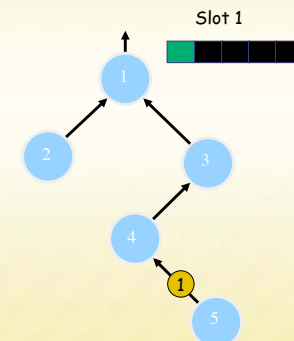


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## Illustration: Aggregation

```
SELECT COUNT(*) FROM
sensors
```

		Sensor #				
		1	2	3	4	5
Slot #	1				1	
	2					
	3					
	4					
	5					

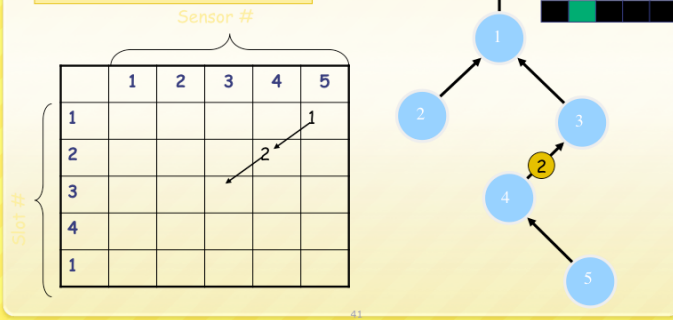


40



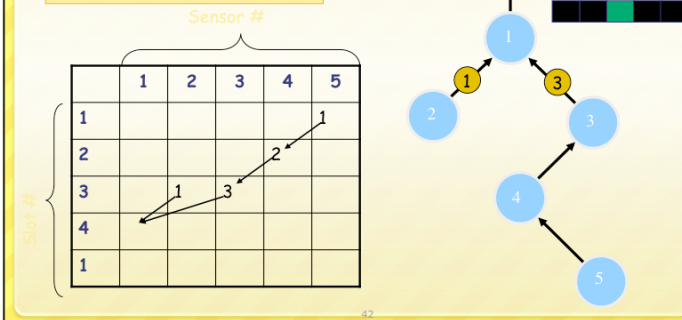
## Illustration: Aggregation

SELECT COUNT(\*) FROM  
sensors



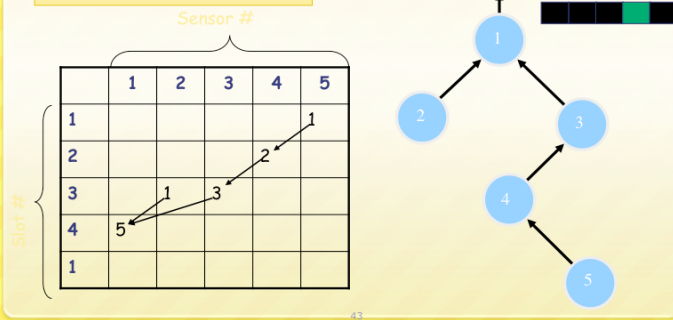
## Illustration: Aggregation

SELECT COUNT(\*) FROM  
sensors



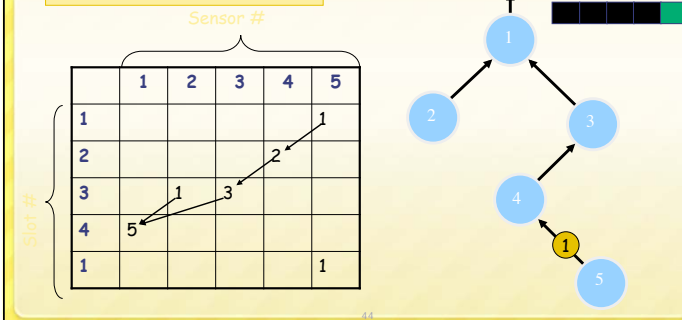
## Illustration: Aggregation

SELECT COUNT(\*) FROM  
sensors



## Illustration: Aggregation

SELECT COUNT(\*) FROM  
sensors



## Types of Aggregates

- SQL supports MIN, MAX, SUM, COUNT, AVERAGE
- Any function *can* be computed via TAG
- In network benefit for many operations
  - E.g. Standard deviation, top/bottom N, spatial union/intersection, histograms, etc.
  - Compactness of PSR

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## Taxonomy of Aggregates

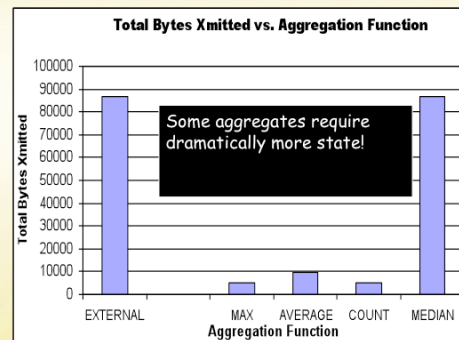
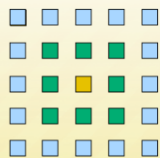
- TAG insight: classify aggregates according to various functional properties
  - Yields a general set of optimizations that can automatically be applied

Property	Examples	Affects
Partial State	MEDIAN : unbounded, MAX : 1 record	Effectiveness of TAG
Duplicate Sensitivity	MIN : dup. insensitive, AVG : dup. sensitive	Routing Redundancy
Exemplary vs. Summary	MAX : exemplary COUNT : summary	Applicability of Sampling, Effect of Loss
Monotonic	COUNT : monotonic AVG : non-monotonic	Hypothesis Testing, Snooping

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## Benefit of In-Network Processing

Simulation Results  
2500 Nodes  
50x50 Grid  
Depth = ~10  
Neighbors = ~20



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## Optimization: Channel Sharing (“Snooping”)

- Insight: Shared channel enables optimizations
- Suppress messages that won't affect aggregate
  - E.g., MAX
  - Applies to all exemplary, monotonic aggregates

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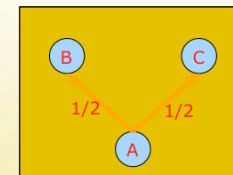
## Optimization: Hypothesis Testing

- Insight: Guess from root can be used for suppression
  - E.g. 'MIN < 50'
  - Works for monotonic & exemplary aggregates
    - Also summary, if imprecision allowed
- How is hypothesis computed?
  - Blind or statistically informed guess
  - Observation over network subset

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## Optimization: Use Multiple Parents

- For duplicate insensitive aggregates
- Or aggregates that can be expressed as a linear combination of parts
  - Send (part of) aggregate to all parents
    - In just one message, via broadcast
  - Decreases variance



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## Overview

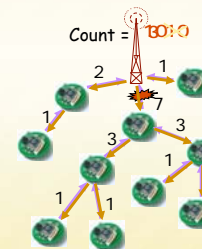
- Ad Hoc Routing
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## Aggregation in Wireless Sensors

Aggregate data is often more important

**In-network aggregation**  
over tree with unreliable communication



Used by current systems,  
 TinyDB [Madden et al. OSDI'02]  
 Cougar [Bonnet et al. MDM'01]

Not robust against  
 node- or link-failures

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## Traditional Approach

- Reliable communication
  - E.g., RMST over Directed Diffusion [Stann'03]
- High resource overhead
  - 3x more energy consumption
  - 3x more latency
  - 25% less channel capacity
- Not suitable for resource constrained sensors

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## Exploiting Broadcast Medium



- ✓ Robust multi-path
- ✓ Energy-efficient

- ✗ Double-counting
- ✗ Different ordering

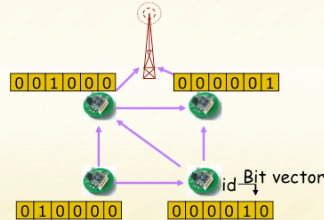


➤ Challenge: order and duplicate insensitivity (ODI)

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## A Naïve ODI Algorithm

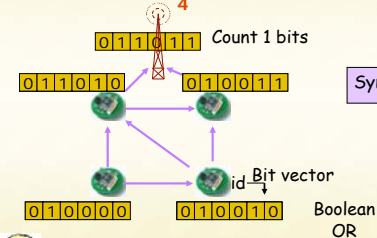
- Goal: count the live sensors in the network



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## Synopsis Diffusion (SenSys'04)

- Goal: count the live sensors in the network



Synopsis should be small

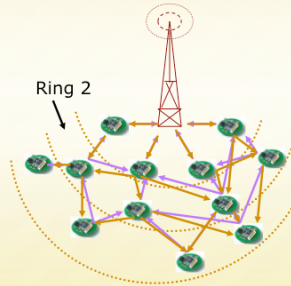


Approximate COUNT algorithm: logarithmic size bit vector

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## Synopsis Diffusion over Rings

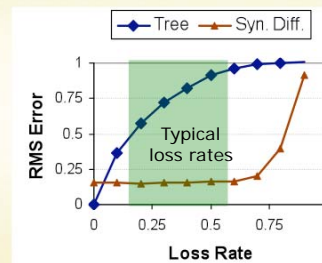
- A node is in ring  $i$  if it is  $i$  hops away from the base-station
- Broadcasts by nodes in ring  $i$  are received by neighbors in ring  $i-1$
- Each node transmits once = optimal energy cost (same as Tree)



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## Evaluation

Approximate COUNT with Synopsis Diffusion



More robust than Tree

Scheme	Energy
Tree	41.8 mJ
Syn. Diff.	42.1 mJ

Per node energy

Almost as energy efficient as Tree

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