

18-452/18-750
Wireless Networks and Applications
Lecture 21: Sensor Networks

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<http://www.cs.cmu.edu/~prs/wirelessS17/>

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1

Wireless Sensors

Low-power microscopic sensors with wireless communication capability

- Miniaturization of computer hardware
→ Intelligence
- Micro Electro-Mechanical Structures (MEMS)
→ Sensing
- Low-cost CMOS-based RF Radios
→ Wireless Communication

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3

Outline

- Example applications?
- Characteristics and design issues:
 - » Power management
 - » Reliable data collection
 - » Hybrid architectures
- Are there size limitations?
- Conclusion
- Based on slides by Prof JP Hubaux, EPFL, and Dr. Lama Nachman, Intel

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2

Wireless Sensor Networks(WSN)

- Even though wireless sensors has limited resources in memory, computation power, bandwidth, and energy.
- With small physical size→Can be embedded in the physical environment.
- Support powerful service in aggregated form (interacting/collaborating among nodes)
- ***Self-organizing multi-hop ad-hoc networks***
- Pervasive computing/sensoring

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4

WSN Applications

- **Commercial Applications**
 - » Light/temperature control
 - » Precision agriculture (optimize watering schedule)
 - » Asset management (tracking freight movement/storage)
- **Wide area monitoring tools supporting Scientific Research**
 - » Wild life Habitat monitoring projects Great Duck Island (UCB), James Reserve (UCLA), ZebraNet (Princeton).
 - » Building/Infrastructure structure (Earthquake impact)
- **Military Applications**
 - » Shooter Localization
 - » Perimeter Defense (Oil pipeline protection)
 - » Insurgent Activity Monitoring (MicroRadar)

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Cold Chain Management

- Supermarket chains need to track the storage temperature of perishable goods in their warehouses and stores.
- Tens if not hundreds of fridges should be monitored in real-time
- Whenever the temperature of a monitored item goes above a threshold
 - » An alarm is raised and an attendant is warned (pager, sms)
 - » The refrigeration system is turned on
- History of data is kept in the system for legal purpose
- Similar concept can be applied to pressure and temperature monitoring in
 - » Production chains, containers, pipelines

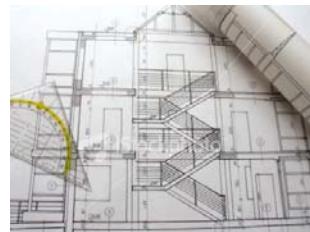
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Home automation

- **Temperature management**
 - » Monitor heating and cooling of a building in an integrated way
 - » Temperature in different rooms is monitored centrally
 - » A power consumption profile is to be drawn in order to save energy in the future
- **Lighting management:**
 - » Detect human presence in a room to automatically switch lights on and off
 - » Responds to manual activation/deactivation of switches
 - » Tracks movement to anticipate the activation of light-switches on the path of a person
- **Similar concept can be applied to**
 - » Security cameras, controlling access, ...



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Precision Agriculture Management

- Farming decisions depend on environmental data (typically photosynthesis):
 - » Solar radiation
 - » Temperature
 - » Humidity
 - » Soil moisture
- Data evolve continuously over time and space
- A farmer's means of action to influence crop yield :
 - » Irrigation
 - » Fertilization
 - » Pest treatment
- To be optimal, these actions should be highly localized (homogenous parcels can be as small as one hectare or less)
- Environmental impact is also to be taken into account
 - » Salinization of soils, groundwater depletion, well contamination, etc.



8

WaterSense

- Goal: Help define and implement farming strategies for farmers in a situation of water scarcity.

- » Crop assessment
- » Water conservation measures
- » Time of farming operations
- » Real-time monitoring of the field conditions



- Desired Outcome: farming decision support system based on environmental data

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Earthquake detection

- The occurrence of an earthquake can be detected automatically by accelerometers
- Earthquake speed: around 5-10km/s
- If the epicenter of an earthquake is in an unpopulated area 200km from a city center, instantaneous detection can give a warning up to 30 sec before the shockwave hits the city
- If a proper municipal actuation network is in place:
 - » Sirens go off
 - » Traffic lights go to red
 - » Elevators open at the nearest floor
 - » Pipeline valves are shut
- Even a warning of a few seconds, can reduce the effects of the earthquake
- Similar concept can be applied to
 - » Forest fire, landslides, etc.

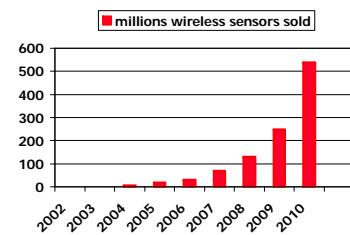


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Economic Forecast

- Industrial Monitoring (35% – 45%)
 - Monitor and control production chain
 - Storage management
 - Monitor and control distribution
- Building Monitoring and Control (20 – 30%)
 - Alarms (fire, intrusion etc.)
 - Access control
- Home Automation (15 – 25%)
 - Energy management (light, heating, AC etc.)
 - Remote control of appliances
- Automated Meter Reading (10-20%)
 - Water meter, electricity meter, etc.
- Environmental Monitoring (5%)
 - Agriculture
 - Wildlife monitoring
- One recent forecast: 1.8 Billion \$ by 2024



Other areas:

- Performance monitoring in sports
- Patient monitoring in health and medicine
- Sensor going wireless in vehicular networks

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11

WSN Characteristics and Design Issues

- Characteristics
 - » Distributed data collection
 - » Many-to-one (rarely peer-to-peer)
 - » Limited mobility
 - » Data collection (time and space resolution)
 - » Event detection
 - » Minimal intrusiveness
- Design issues
 - » Low-cost (hardware and communication)
 - » Extended life-time
 - » Reliable communication
 - » Efficient integrated data processing
 - » Hybrid network infrastructure
 - » Security

Wireless helps but may not be required!

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12

Outline

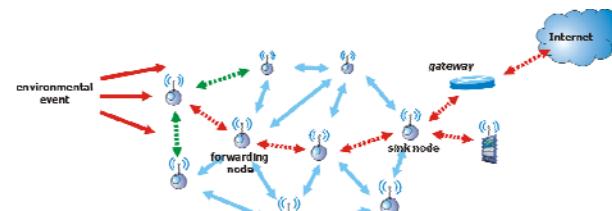
- Example applications?
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13

Wireless Sensor Network architecture

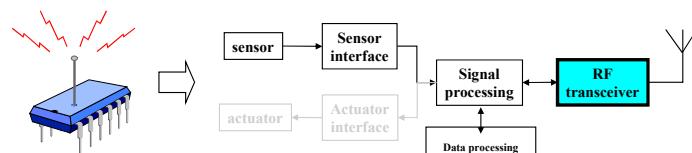
- Numerous sensor devices
 - » Modest wireless communication, processing, memory capabilities
 - » Form Ad Hoc Network (self-organized)
 - » Report the measured data to the user



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14

Sensor Node architecture

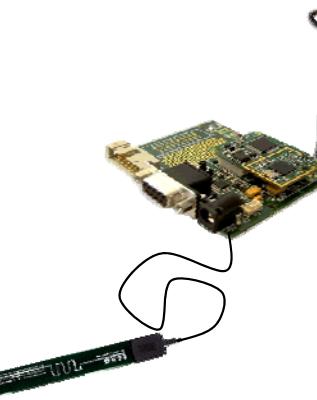


- A sensor node can be an *information source*, a *sink* and a *router*
- Autonomous ⇒ *low-power*
- Combine *sensing, signal conditioning, signal processing, control* and *communication* capabilities

(courtesy of Swiss Center for Electronics and Microelectronics, Neuchâtel)

15

Example of a Low Power Transceiver: Tinynode™



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- 868 MHz multi-channel transceiver
- 8 MHz µ-Controller
- 10KB RAM
- 48 kB Program space
- 512 External Flash
- 115 kbps data rate
- 3 V supply voltage
- Current consumption
 - » Transmit 33 mA
 - » Receive 14 mA
 - » Sleep < µA
- -121 dBm sensitivity
- Radio range 200m (outdoor)
- 39 MHz quartz reference

16

Design Issue: Low-cost

- **Hardware**
 - » Low-cost radio
 - » Low cost internal clock
 - » Limited storage and processing capabilities
 - » Not tamper-proof
 - » May have to withstand tough environmental conditions
- **Communication**
 - » Cannot rely on existing pay-per-use cellular infrastructure
 - » Use unlicensed spectrum to reach a “gateway”, which has internet connectivity
 - Wired, WiFi, drive-by, cellular, ...

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17

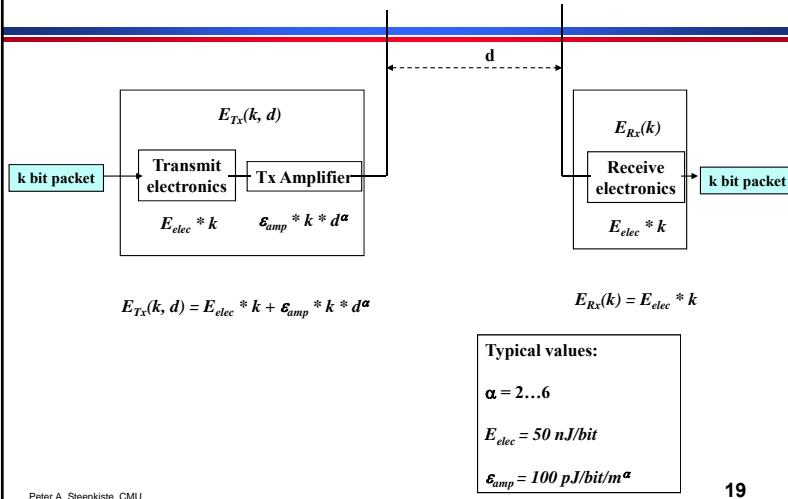
Design Issue: Power Management

- **Energy-efficient routing**
 - » Minimum-cost spanning tree
- **Load-balancing**
 - » Mobility
 - » In-network aggregation
- **Medium-access control**

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18

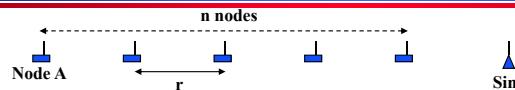
Simple Model for Energy Consumption



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19

Energy-efficient Routing : Example



Transmitting a single k -bit message from node A (located at distance r from Sink) to Sink:

Direct transmission: $E_{direct} = E_{Tx}(k, d = n \cdot r) = E_{elec} * k + \epsilon_{amp} * k * (nr)^\alpha = k(E_{elec} + \epsilon_{amp} n^\alpha r^\alpha)$

Multi-Hop Transmission: $E_{multi-hop} = n * E_{Tx}(k, d = r) + (n-1) * E_{Rx}(k)$
 $= n(E_{elec} * k + \epsilon_{amp} * k * r^\alpha) + (n-1) * E_{elec} * k = k((2n-1)E_{elec} + \epsilon_{amp} nr^\alpha)$

MultiHop routing requires less energy than direct communication if: $\frac{E_{elec}}{\epsilon_{amp}} < \frac{r^\alpha (n^{\alpha-1} - 1)}{2}$

Assuming $\alpha = 3, r = 10m$, we get $E_{multi-hop} < E_{direct}$ as soon as $n \geq 2$

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Minimum Energy in a Wireless Network

- Problem: for an arbitrary set of nodes, find (in a fully distributed way) the minimum cost spanning tree to and from a given *sink* node
- Assumptions
 - Each node knows its own exact location (e.g., using GPS)
 - The power decreases with distance according to a power law with a known and uniform exponent α
 - Each node can communicate with another node located at an arbitrary distance
 - Nodes do not move
 - Slightly different power model

sending: td^α
receiving: c

- Example:



Power to send from A to C via B:

$$td_{AB}^\alpha + td_{BC}^\alpha + c$$

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21

Relay region

Relay region:

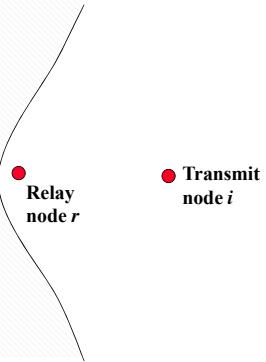
$$R_{i \rightarrow r} \equiv \{(x, y) \mid P_{i \rightarrow r \rightarrow (x, y)} < P_{i \rightarrow (x, y)}\}$$

We can expand this to:

$$td_{i,r}^\alpha + td_{r,(x,y)}^\alpha + c < td_{i,(x,y)}^\alpha$$

$$t((i_x - x)^2 + (i_y - y)^2)^{\alpha/2} - t((r_x - x)^2 + (r_y - y)^2)^{\alpha/2} >$$

$$t((i_x - r_x)^2 + (i_y - r_y)^2)^{\alpha/2} + c$$

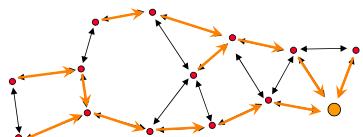


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22

Distributed Network Protocol

- Finds the minimum power topology for a stationary set of nodes with a single sink
- Assumption: each node is equipped with a GPS-receiver and transmits its position to its neighbors
- The protocol proceeds in 2 phases:
 - Each node computes its own enclosure
 - Each node computes its optimal cost distribution



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23

Load-balancing

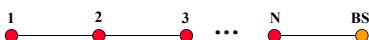
- Assumption: in a multi-hop many-to-one sensor network, the data collection follows a spanning tree.
- Power consumption due to transmission/reception grows exponentially from the leaves to the root of the tree
- Consequence: the power sources of the nodes close to the sink deplete faster. Since they relay all the network's traffic, they pull the network lifetime down.

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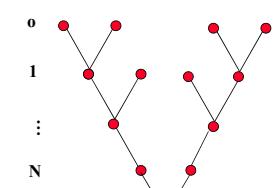
24

Load-balancing

Line topology



Tree topology



P_a : Average transmission power consumption

P_{ra} : Average reception power consumption

P_{pr} : Average processing power consumption

$P_t(k)$: Total power consumption of node k

$P = P_{pr} + P_a + (k-1)(P_a + P_{ra})$

P grows linearly with the distance from the leaf node

d : distance from leaf

F : number of messages forwarded

P : Power consumption

Assumptions:

1) all nodes have either 0 or $n_c > 2$ children

2) all leaves are at the same distance from the sink

$F(d) \geq 2^d$

$P(d) \geq P_a + 2^d (P_a + P_{ra})$

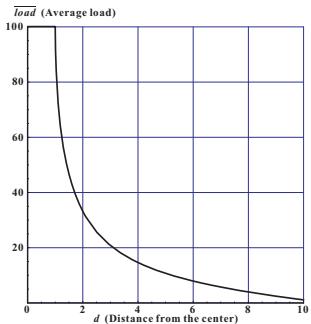
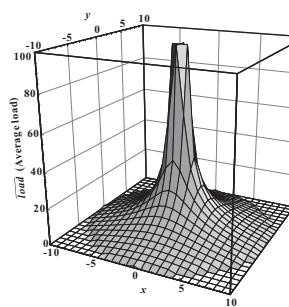
P grows exponentially with distance from leaf node

25

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Load balancing

- Power consumption increases at least linearly when nodes are closer to the sink
- Typical case is much worse

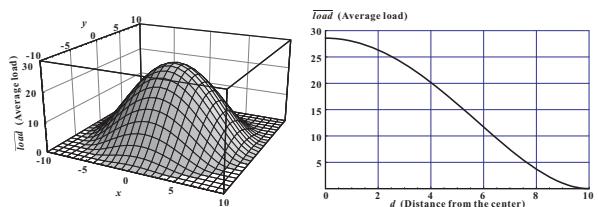


26

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Use Mobility for Load-balancing

- Move the base station to distribute the role of “hot spots” (i.e., nodes around the base station) over time
- The data collection continues through multi-hop routing wherever the base station is, so the solution does not sacrifice latency



27

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In-network Data Aggregation

- To mitigate cost of forwarding, compute relevant statistics along the way: *mean*, *max*, *min*, *median* etc.
- Forwarding nodes aggregate the data they receive with their own and send one message instead of relaying an exponentially growing number of messages
- Issues
 - » Location-based information (which nodes sent what) is lost
 - » Distributed computation of statistics
 - *mean*: node needs to know both the mean values and the sizes of samples to aggregate correctly
 - *median*: only an approximated computation is possible
- Especially useful in a query-based data collection system
 - » Queries regard a known subset of nodes
 - » Aggregation function can be specified

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28

Medium-Access Control

- MAC attributes:
 - » Collision avoidance
 - » Energy efficiency
 - » Scalability and adaptivity
- Nodes transmit very intermittently, but once a transmission is taking place, we must ensure that the intended receiver gets it.
- Current-consumption in receive state or in radio-on idle state are comparable
- Idle state (idle listening) is a dominant factor in power consumption

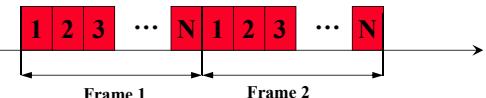
Goal is to put nodes to sleep most of the time, and wake them up only to receive a packet

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29

<http://research.cens.ucla.edu/people/estrin/resources/conferences/2002jun-Ye-Estrin-Energy.pdf>

Synchronous MACs

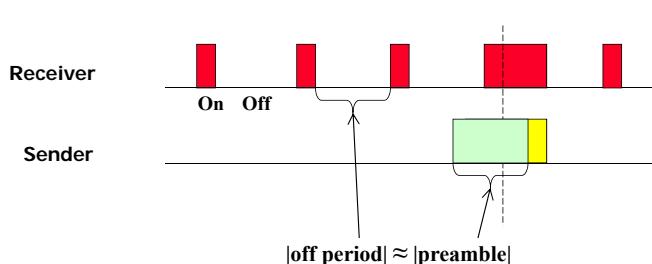
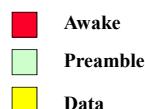
- TDMA (similar to cellular networks)
- Shortcomings
 - » Necessity to organize nodes in clusters and cluster hierarchies
 - » High control traffic cost
- Possible solution
 - » Each node maintains two schedules
 - Its parent schedule
 - The schedule it sets for its children
 - » Beacons are used to compensate for clock drifts

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30

Asynchronous: B-MAC

- Asynchronous
- Low Power listening



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31

Shortcomings

- Transmitting a packet is very expensive
- Overhearing is expensive
- Relaying packets is expensive (multihop)

Simple Improvement:

- Aggregating packets before sending them
 - » In low duty cycle data collection network, gain may be substantial
 - » Price to pay : real-time

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32

Design Issue: Reliable data collection

- Many-to-one communication paradigm
- Multi-hop communication
- Nodes select one parent to send their data packets (tree topology)

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33

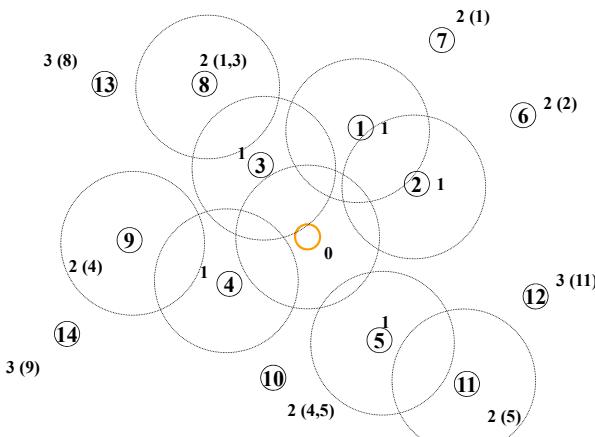
MintRoute: A Data Collection Tree Routing Protocol

- Distance-vector routing protocol: one routing path per node
- A shortest path: Minimum number of transmissions
- The base station send periodic beacons that are broadcast by each node after incrementing a hop count
- Node select beacons with lowest hop count from the ones it received, and adds its sender among a list of potential parents
- Neighboring nodes exchange periodic beacons for link quality evaluation (gaps within the sequence # of packets → packet losses)
- Nodes select their parent based on hop count, link quality and load.
- Volatile routing topology → load balancing
- Cycle avoidance : Link quality must not vary too rapidly

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34

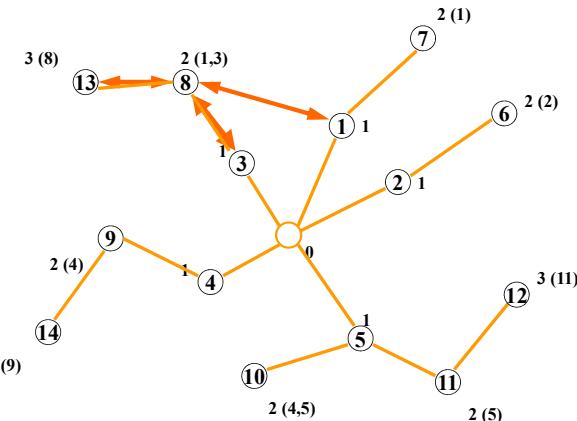
MintRoute : Root Beacons



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35

MintRoute : Link Estimation



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36

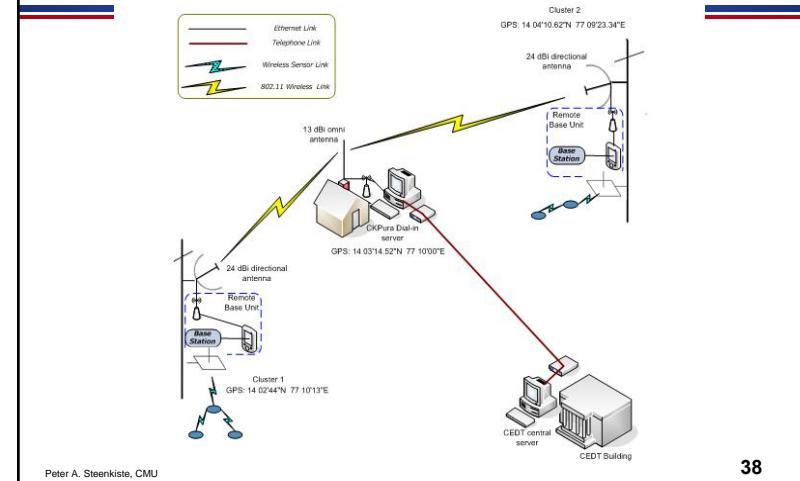
How Large can Sensor Networks Feasibly Grow?

- Multi-tier is the typical approach to scaling
 - » Multi-tier architecture divides the scaling problem into manageable segments
 - » Segmented networks (2nd tier aggregation, high throughput apps)
 - Static allocation of networks simplifies the scalability problem but complicates the management problem
 - Dynamic allocation of networks is more appealing
- Clusters on the orders of 10-50 nodes are adequate for most applications
 - » In most industrial applications, having access to power for 2nd tier nodes is realistic (or at least intermittent power sufficient to charge a battery)
 - » Outdoor deployments can also leverage solar power for recharging 2nd tier batteries
 - » Pipeline monitoring applications look different, will probably need higher scaling

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37

2-tier Architecture with 802.11 Bridge



Delay Tolerant Network with Data Mules

- Clusters are not directly connected to the server
- Cluster heads store data from the cluster nodes
- “Data mules” collect the data periodically
 - » Cars, robots, plane, etc.
- When a cluster-head detects a mule, it uploads to it the data it had in store

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39

Conclusion

- WSNs are an emerging technology which is poised to grow exponentially in the coming years
- This new communication paradigm introduces a new set of design constraints
 - » They must be extremely low-cost
 - Both to purchase and to operate
 - » They must be extremely energy efficient since their lifetime is potentially years
 - Hardware design
 - Routing and topology mechanisms
 - Specialized Medium Access Control mechanisms
 - » Despite their low-cost and power management features, they must implement reliable communication protocols
 - » They must integrate versatile middle-ware and provide data processing
 - » They will often rely on a hybrid network infrastructure

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40