

Proof Systems and Proof Complexity

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Automated Reasoning and Satisfiability

September 29, 2021

Proofs of Unsatisfiability

Beyond Resolution

Propagation Redundancy

Satisfaction-Driven Clause Learning

Challenges

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Certifying Satisfiability and Unsatisfiability

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■ Certifying **unsatisfiability** is not so easy:

- If a formula has n variables, there are 2^n possible assignments.

➡ Checking whether **every** assignment falsifies the formula is **costly**.

- More compact certificates of unsatisfiability are desirable.

➡ Proofs

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- In general, a **proof** is a **string** that **certifies the unsatisfiability** of a formula.
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... but can be of exponential size with respect to a formula.
- **Example:** Resolution proofs
 - A **resolution proof** is a sequence C_1, \dots, C_m of clauses.
 - Every clause is either contained in the formula or derived from two earlier clauses via the **resolution rule**:

$$\frac{C \vee x \quad \bar{x} \vee D}{C \vee D}$$

- C_m is the **empty clause** (containing no literals), denoted by \perp .
- There exists a resolution proof for every unsatisfiable formula.

Resolution Proofs

- **Example:** $F = (\bar{x} \vee \bar{y} \vee z) \wedge (\bar{z}) \wedge (x \vee \bar{y}) \wedge (\bar{u} \vee y) \wedge (u)$
- **Resolution proof:**
 $(\bar{x} \vee \bar{y} \vee z), (\bar{z}), (\bar{x} \vee \bar{y}), (x \vee \bar{y}), (\bar{y}), (\bar{u} \vee y), (\bar{u}), (u), \perp$

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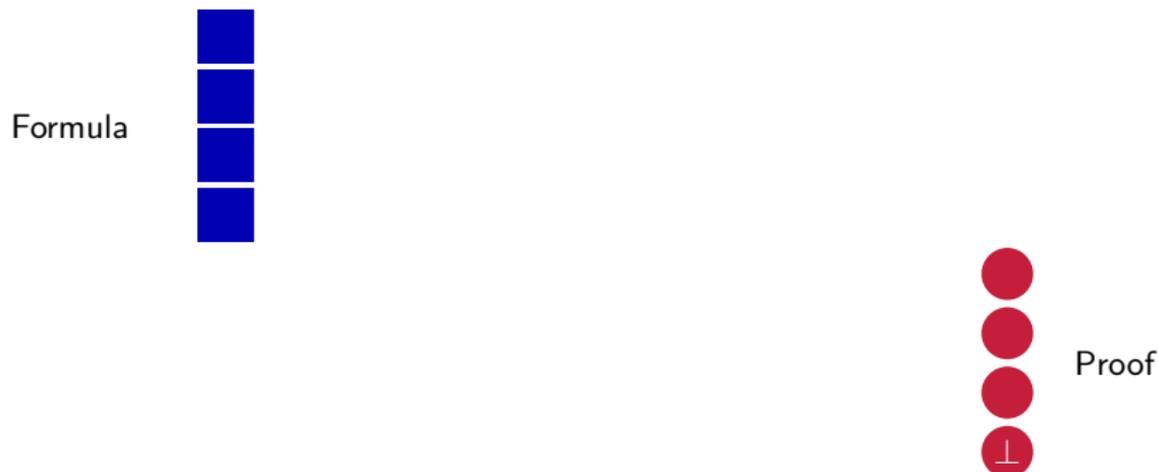
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■ **Drawbacks** of resolution:

- For **many** seemingly simple formulas, there are **only** resolution proofs of **exponential size**.
- **State-of-the-art solving techniques** are **not succinctly expressible**.

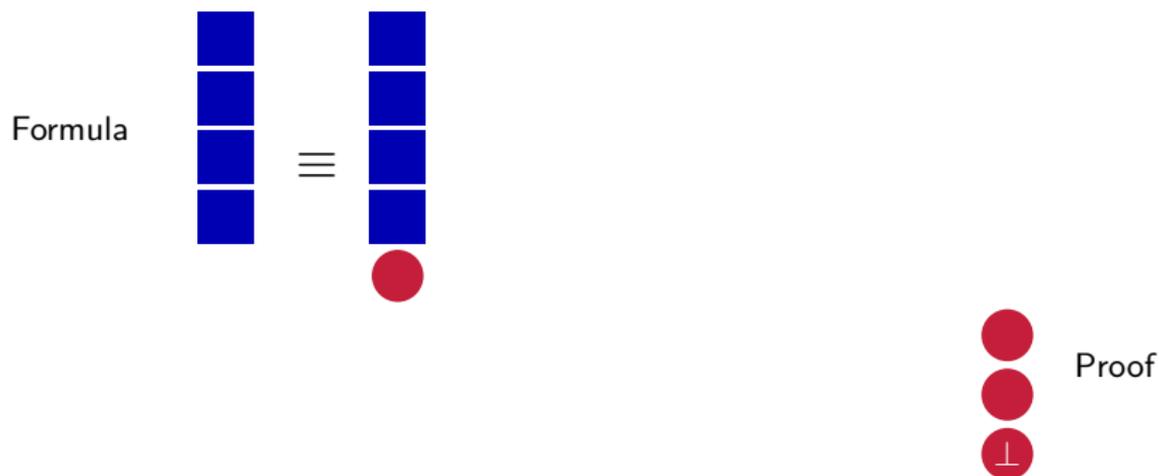
Clausal Proofs

Reduce the size of the proof by only storing added clauses



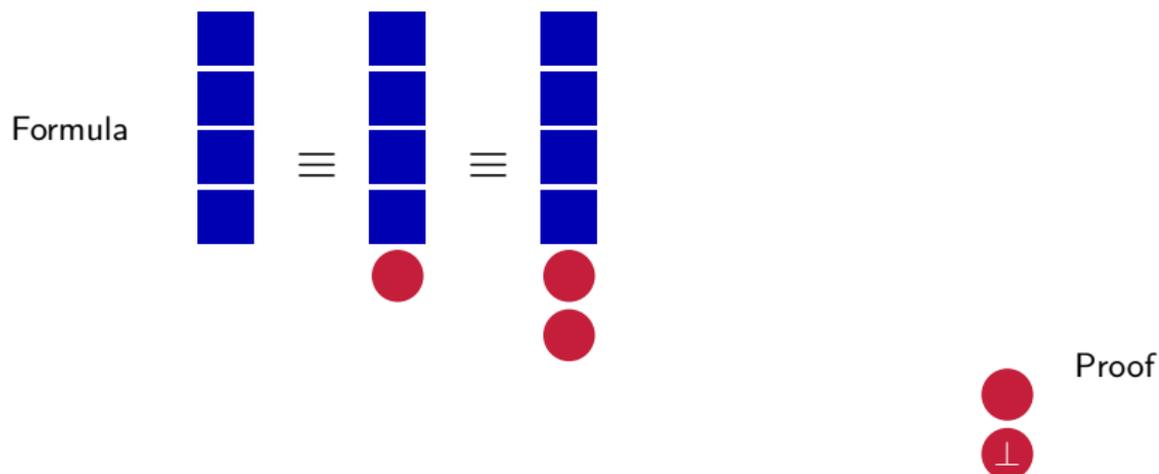
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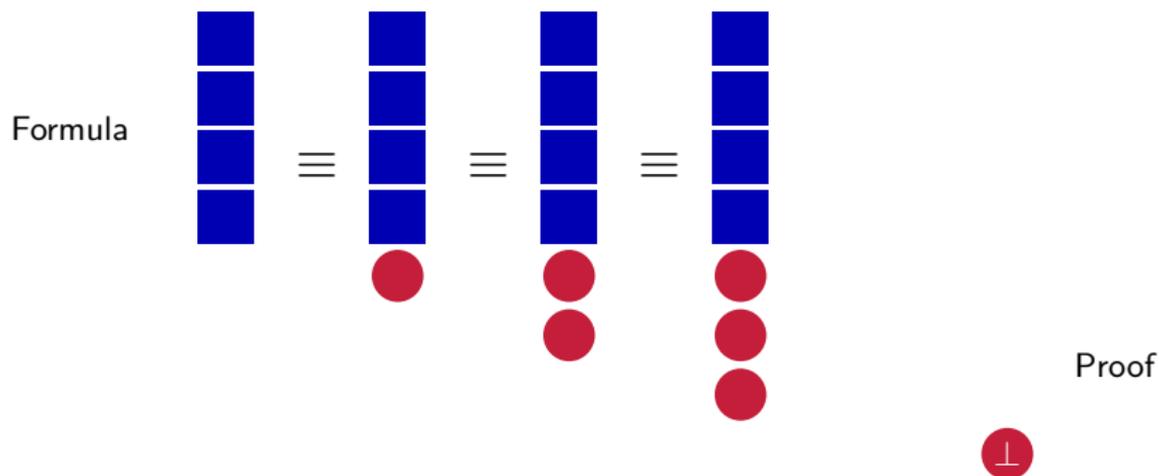
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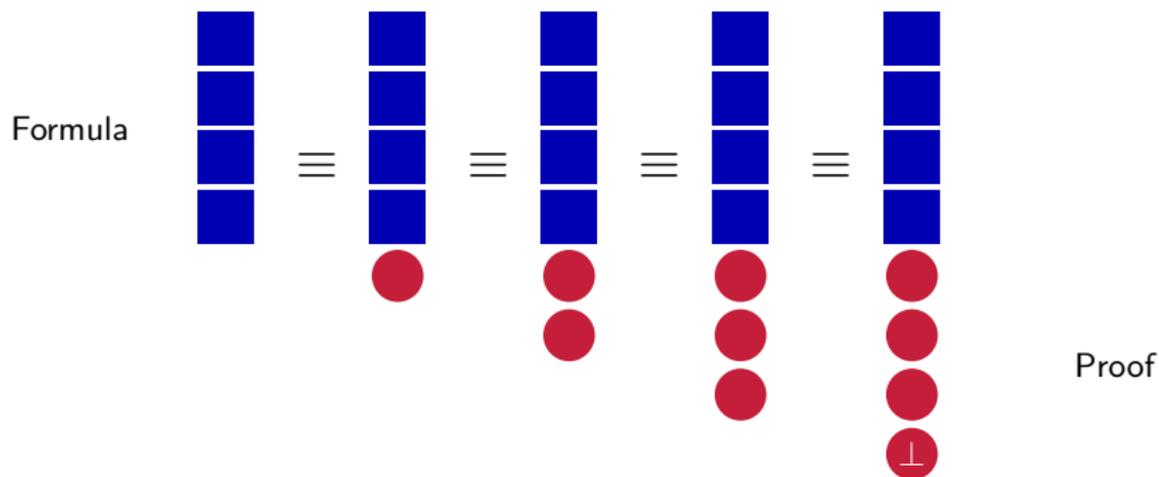
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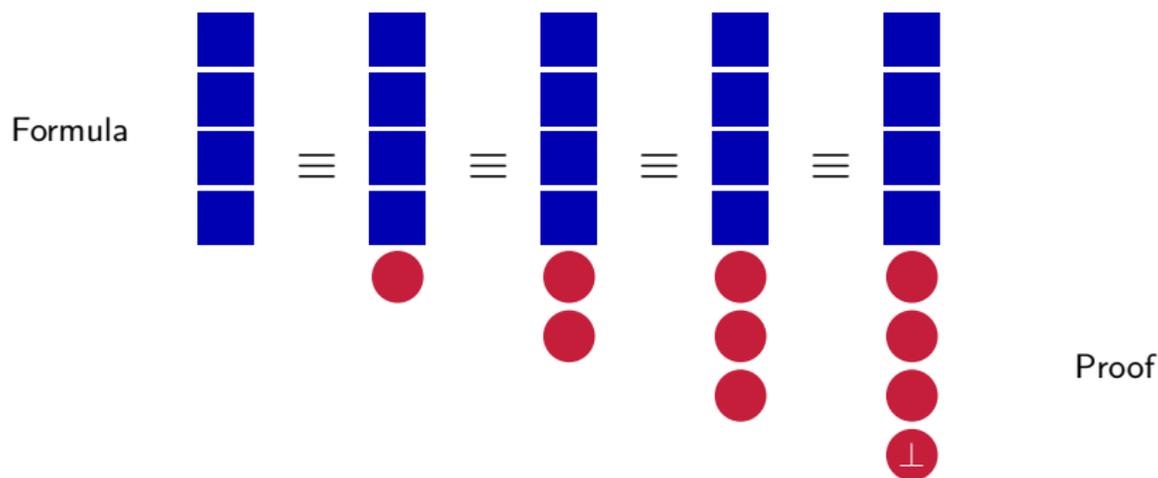
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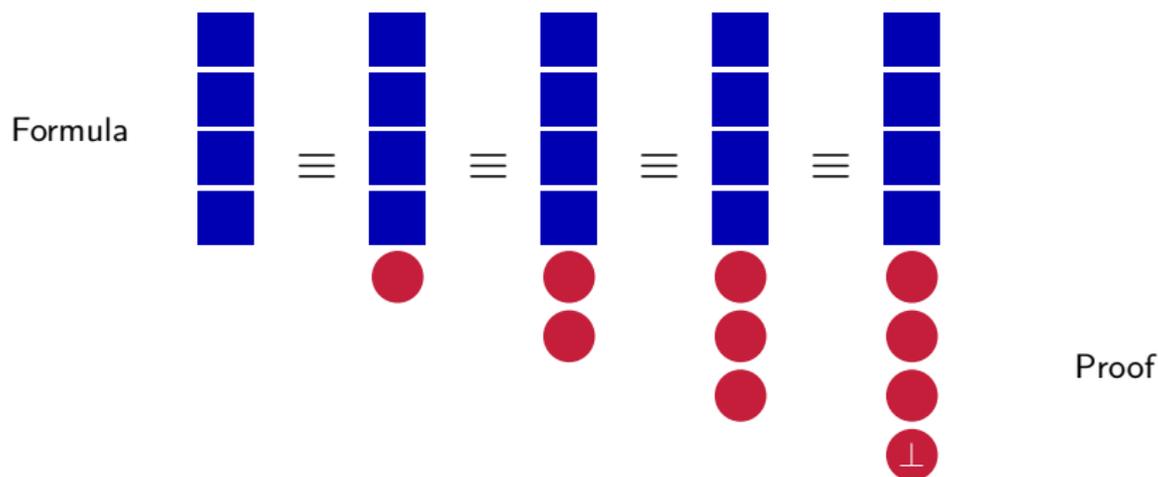
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- Checking redundancy should be **efficient**.

Clausal Proofs

Reduce the size of the proof by only storing added clauses



- Clauses whose addition preserves satisfiability are **redundant**.
- Checking redundancy should be **efficient**.
- ➔ **Idea**: Only add clauses that fulfill an **efficiently checkable redundancy criterion**.

Reverse Unit Propagation

- **Unit propagation** (UP) satisfies unit clauses by assigning their literal to true (until fixpoint or a conflict).
- Let F be a formula, C a clause, and α the smallest assignment that falsifies C . C is **implied by F via UP** (denoted by $F \vdash_1 C$) if UP on $F|\alpha$ results in a conflict.

Example

$$F = (a \vee b \vee \bar{c}) \wedge (\bar{a} \vee \bar{b} \vee c) \wedge (b \vee c \vee \bar{d}) \wedge (\bar{b} \vee \bar{c} \vee d) \wedge (a \vee c \vee d) \wedge (\bar{a} \vee \bar{c} \vee \bar{d}) \wedge (\bar{a} \vee b \vee d) \wedge (a \vee \bar{b} \vee \bar{d})$$

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$$\frac{\frac{(a \vee c \vee d) \quad (b \vee c \vee \bar{d})}{(a \vee b \vee c)} \quad (a \vee b \vee \bar{c})}{(a \vee b)}$$

Proofs of Unsatisfiability

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Satisfaction-Driven Clause Learning

Challenges

Traditional Proofs vs. Interference-Based Proofs

- In **traditional** proof systems, everything that is **inferred**, is **logically implied** by the premises.

$$\frac{C \vee x \quad \bar{x} \vee D}{C \vee D} \text{ (RES)} \qquad \frac{A \quad A \rightarrow B}{B} \text{ (MP)}$$

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- If certain premises are present, infer the conclusion.

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- ➔ Inference rules reason about the **presence** of facts.
 - If certain premises are present, infer the conclusion.
- **Different approach**: Allow **not only implied conclusions**.
 - **Require only** that the addition of facts preserves **satisfiability**.
 - Reason also about the **absence** of facts.
- ➔ This leads to **interference-based proof systems**.

Early work on reasoning beyond resolution

The early SAT decision procedures used the **Pure Literal rule** [Davis and Putnam 1960; Davis, Logemann and Loveland 1962]:

$$\frac{\bar{x} \notin F}{(x)} \text{ (pure)}$$

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Extended Resolution (ER) [Tseitin 1966]

- Combines resolution with the **Extension rule**:

$$\frac{x \notin F \quad \bar{x} \notin F}{(x \vee \bar{a} \vee \bar{b}) \wedge (\bar{x} \vee a) \wedge (\bar{x} \vee b)} \text{ (ER)}$$

- Equivalently, adds the definition $x := \text{AND}(a, b)$
- Can be considered the **first interference-based proof system**
- Is very powerful: **No known lower bounds**

Short Proofs of Pigeon Hole Formulas [Cook 1967]

Can $n+1$ pigeons be in n holes (at-most-one pigeon per hole)?

$$PHP_n := \bigwedge_{1 \leq p \leq n+1} (x_{1,p} \vee \dots \vee x_{n,p}) \wedge \bigwedge_{1 \leq h \leq n} \bigwedge_{1 \leq p < q \leq n+1} (\bar{x}_{h,p} \vee \bar{x}_{h,q})$$

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However, these proofs require introducing new variables:

- Hard to find such proofs automatically
- Existing ER approaches produce exponentially large proofs
- How to get rid of this hurdle? First approach: blocked clauses...

Blocked Clauses [Kullmann 1999]

Definition (Block Clause)

A clause $(C \vee x)$ is a **blocked** on x w.r.t. a CNF formula F if for every clause $(D \vee \bar{x}) \in F$, resolvent $C \vee D$ is a **tautology**.

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Example

Consider the formula $(a \vee b) \wedge (a \vee \bar{b} \vee \bar{c}) \wedge (\bar{a} \vee c)$.

First clause is not blocked.

Second clause is blocked by both a and \bar{c} .

Third clause is blocked by c

Theorem

Adding or removing a blocked clause preserves (un)satisfiability.

Blocked Clause Addition and Blocked Clause Elimination

The Blocked Clause proof system (BC) combines the resolution rule with the addition of blocked clauses.

- BC generalizes ER [Kullmann 1999]

- Recall

$$\frac{x \notin F \quad \bar{x} \notin F}{(x \vee \bar{a} \vee \bar{b}) \wedge (\bar{x} \vee a) \wedge (\bar{x} \vee b)} \text{ (ER)}$$

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Blocked clause elimination used in preprocessing and inprocessing

- Simulates many circuit optimization techniques
- Removes redundant Pythagorean Triples

DRAT: An Interference-Based Proof System

- DRAT is a popular interference-based proof system
- DRAT allows adding RATs (defined below) to a formula.
 - It can be efficiently checked if a clause is a RAT.
 - RATs are not necessarily implied by the formula.
 - But RATs are redundant: their addition preserves satisfiability.
- DRAT also allows clause deletion
 - Initially introduced to check proofs more efficiently
 - Clause deletion may introduce clause addition options (interference)

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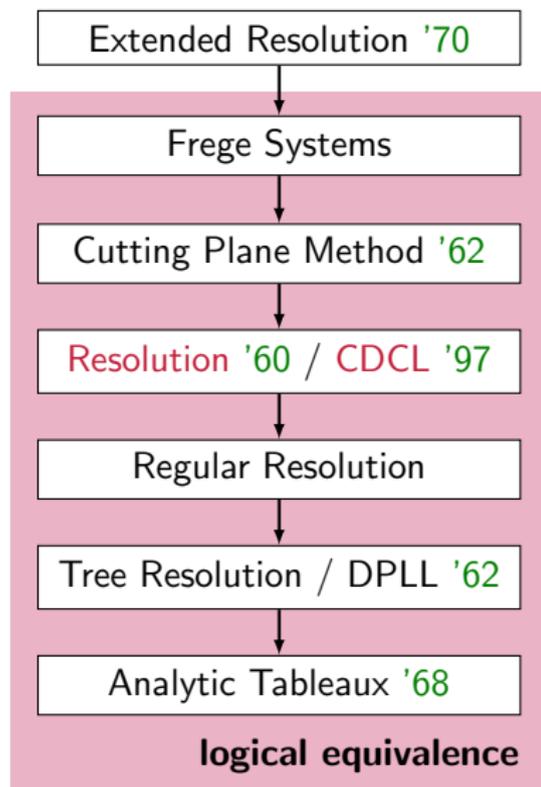
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Definition (Resolution Asymmetric Tautology)

A clause $(C \vee x)$ is a resolution asymmetric tautology (RAT) on x w.r.t. a CNF formula F if for every clause $(D \vee \bar{x}) \in F$, $C \vee D$ is implied by F via unit-propagation, i.e., $F \vdash_1 C \vee D$.

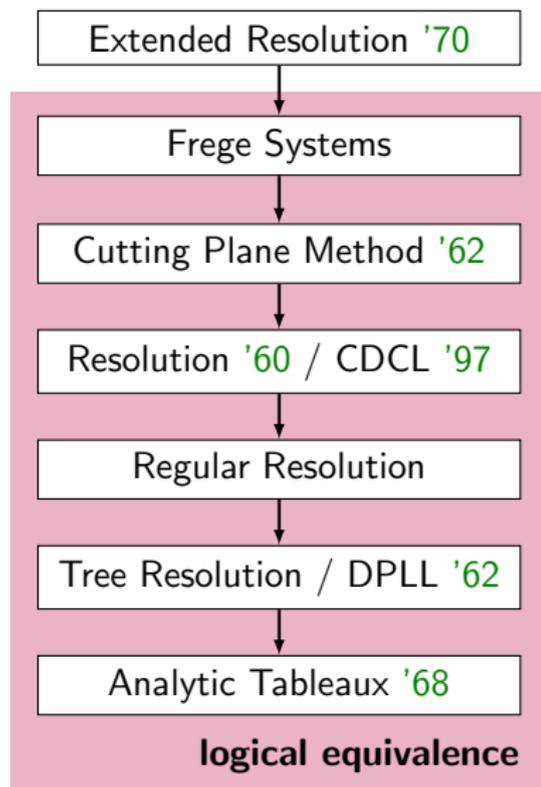
Proof Search in Strong Proof Systems

Existence of Short Proofs

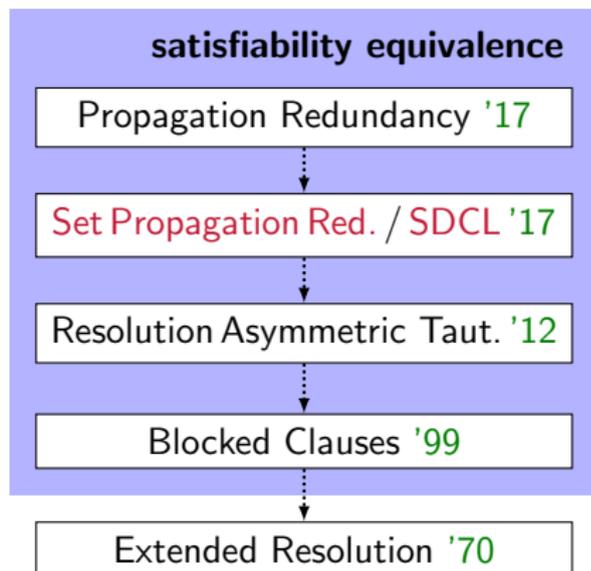


Proof Search in Strong Proof Systems

Existence of Short Proofs



Finding Short Proofs



Express solving techniques compactly
[Järvisalo, Heule, and Biere '12]
Short proofs without new variables
[Heule, Kiesel, and Biere '17A]

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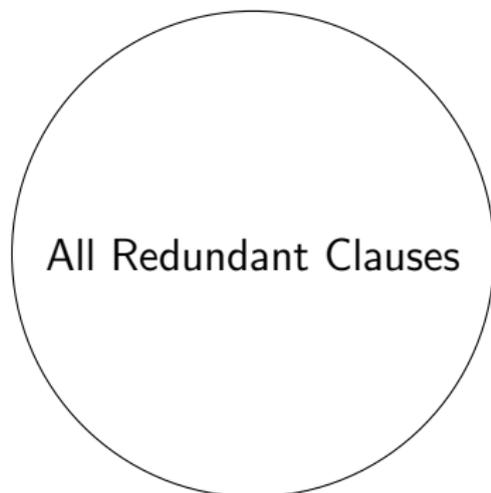
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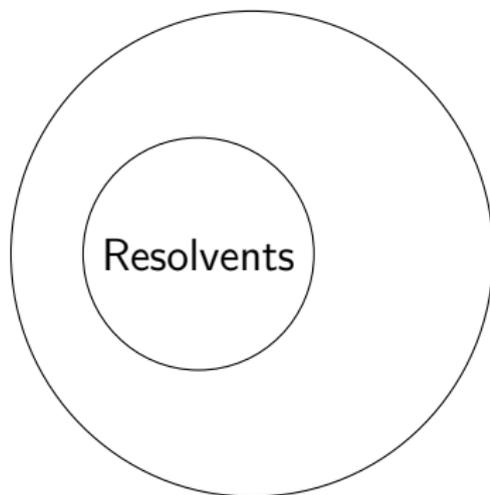
Redundant Clauses

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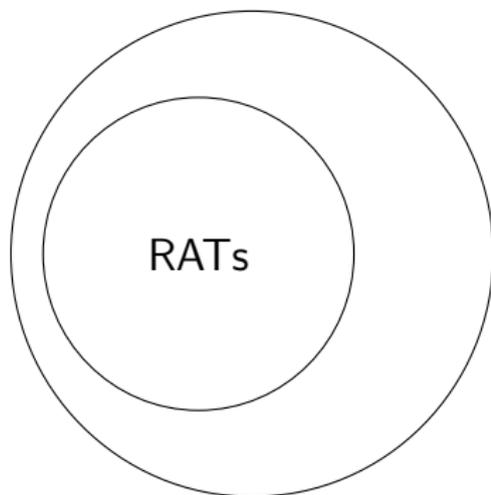
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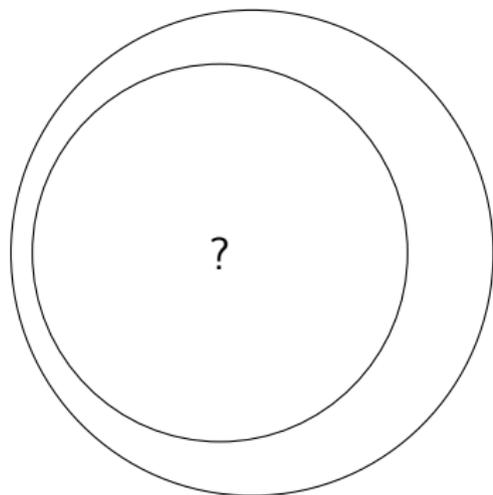
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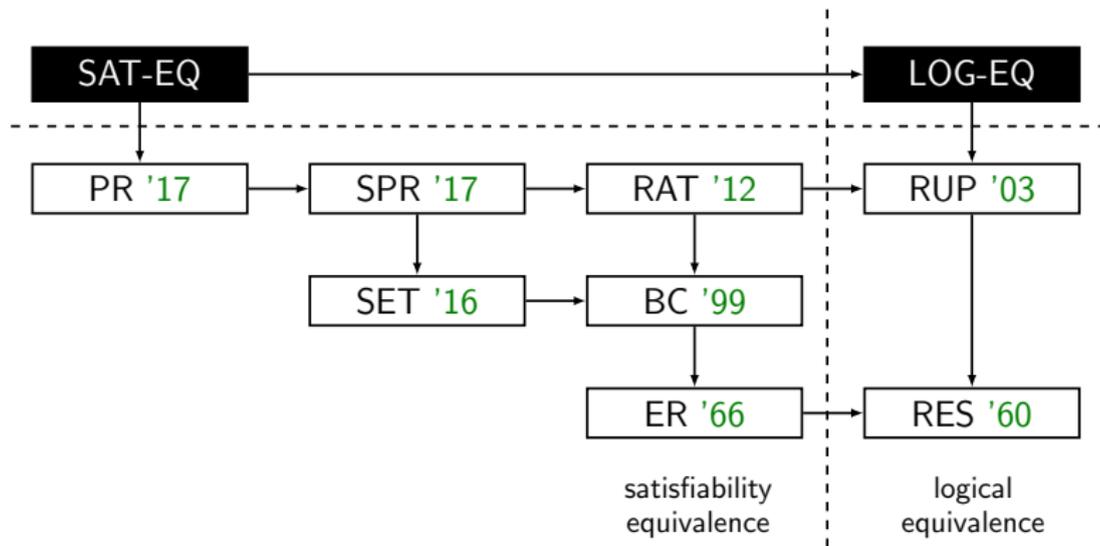


- Are **stronger** redundancy notions still **efficiently checkable**?

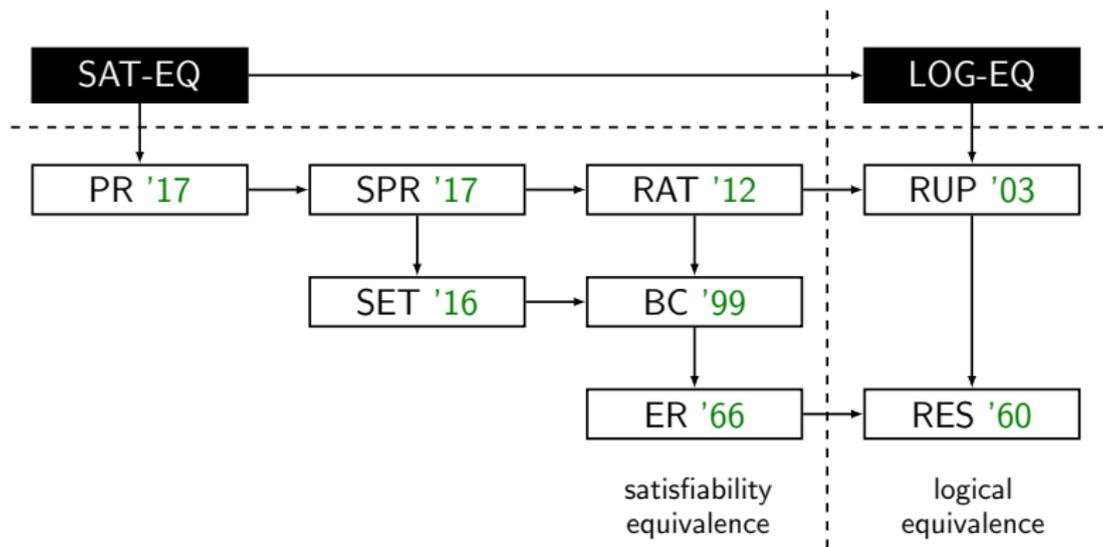
New Propositional Proof Systems

- We introduced **new clause-redundancy notions**:
 - Propagation-redundant (PR) clauses
 - Set-propagation-redundant (SPR) clauses
 - Literal-propagation-redundant (LPR) clauses
- LPR clauses coincide with RAT.
- SPR clauses strictly generalize RATs.
- PR clauses strictly generalize SPR clauses.
- The redundancy notions provide the basis for **new proof systems**.

New Proof Systems for Propositional Logic



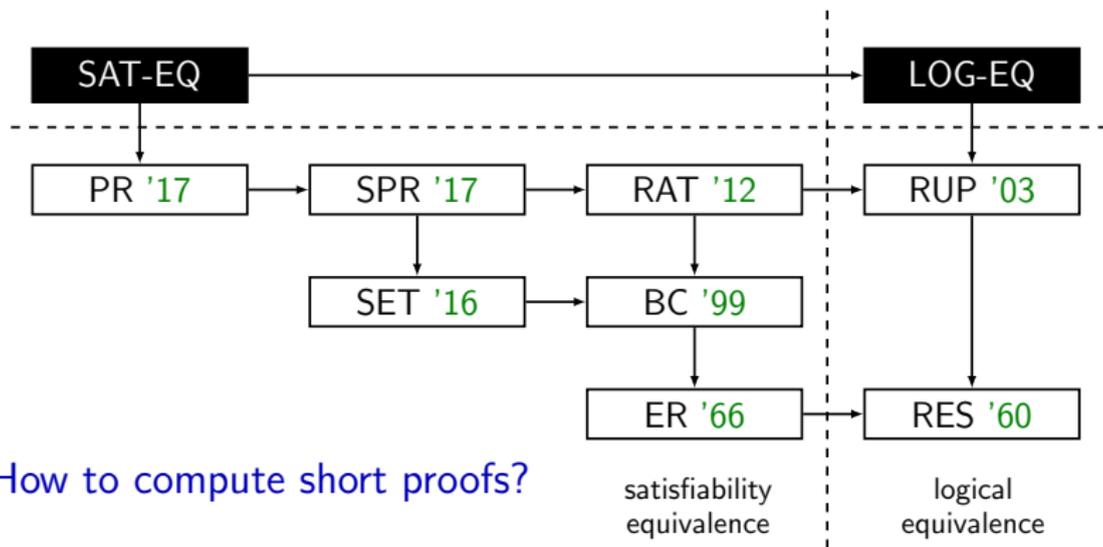
New Proof Systems for Propositional Logic



RAT simulates PR [Heule and Biere 2018]

ER simulates RAT [Kiesl, Rebola-Pardo, Heule 2018]

New Proof Systems for Propositional Logic



How to compute short proofs?

RAT simulates PR [Heule and Biere 2018]

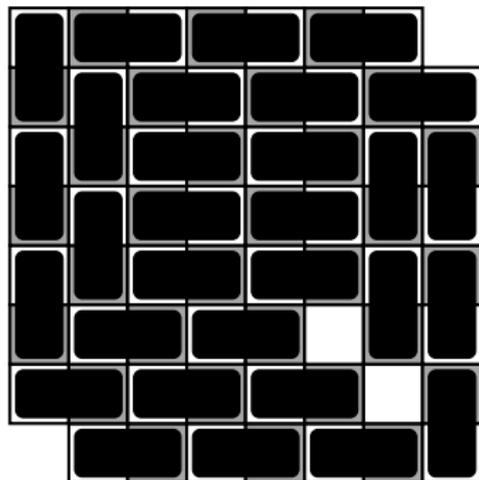
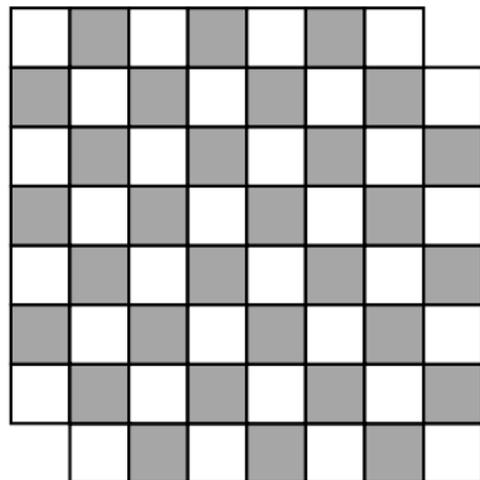
ER simulates RAT [Kiesl, Rebola-Pardo, Heule 2018]

Stronger Proof Systems: What Are They Good For?

- The new proof systems can give **short proofs** of formulas that are considered **hard**.
- We have **short SPR and PR proofs** for the well-known **pigeon hole formulas** (linear in the size of the input).
 - Pigeon hole formulas have **only exponential-size resolution proofs**.
 - If the **addition of new variables via definitions** is allowed, there are polynomial-size proofs.
- Strong proof systems do **not** require new variables.
 - ➔ Search space of possible clauses is **finite**.
 - ➔ Makes **search** for such clauses **easier**.

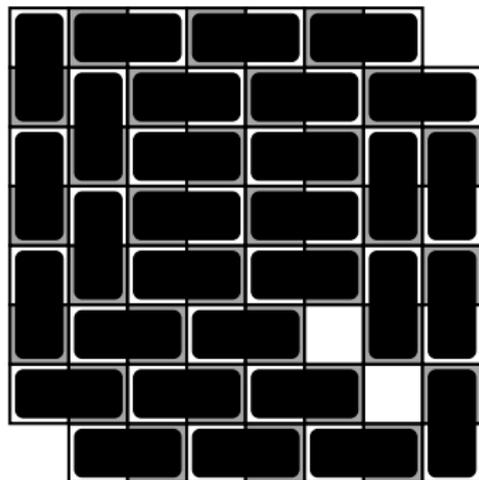
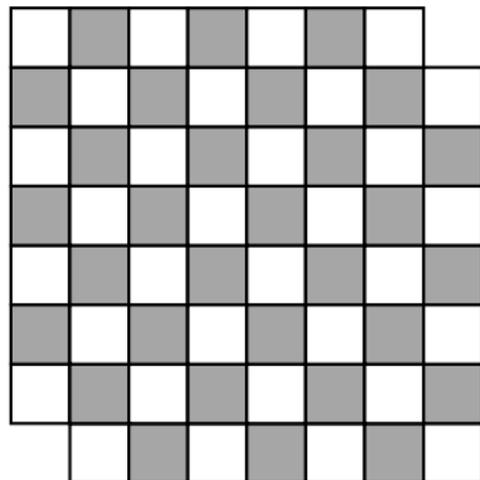
Mutilated Chessboards: “A Tough Nut to Crack” [McCarthy]

Can a chessboard be fully covered with dominos after removing two diagonally opposite corner squares?



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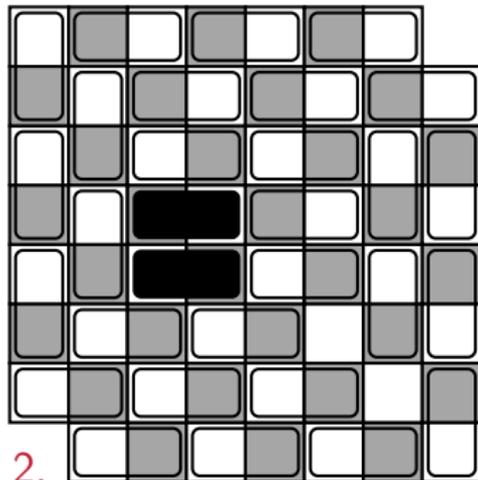
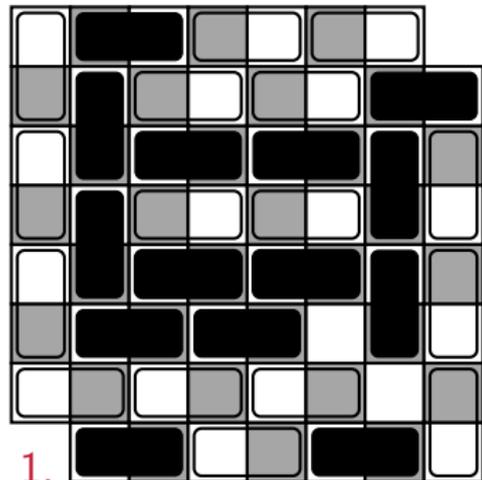


Easy to refute based on the following two observations:

- There are more white squares than black squares; and
- A domino covers exactly one white and one black square.

Without Loss of Satisfaction

One of the crucial techniques in SAT solvers is to **generalize a conflicting state** and use it to constrain the problem.

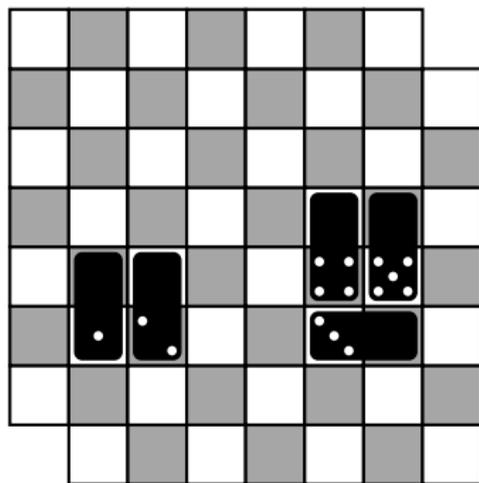
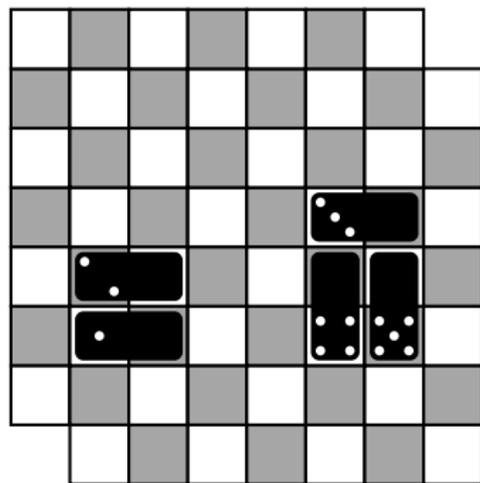


The used proof system can have a big impact on the size:

1. Resolution can only reduce the 30 dominos to 14 (left); and
2. “Without loss of satisfaction” can reduce them to 2 (right).

Mutilated Chessboards: An alternative proof

Satisfaction-Driven Clause Learning (SDCL) is a new solving paradigm that finds proofs in the PR proof system [HKB '17]



SDCL can detect that the above two patterns can be **blocked**

- This reduces the number of explored states **exponentially**
- We produced **SPR** proofs that are linear in the formula size

Redundancy as an Implication

A formula G is **at least as satisfiable** as a formula F if $F \models G$.

Given a formula F and assignment α , we denote with $F|\alpha$ the **reduced formula** after removing from F all clauses satisfied by α and all literals falsified by α .

Theorem

*Let F be a formula, C a clause, and α the smallest assignment that falsifies C . Then, C is **redundant** w.r.t. F iff there exists an assignment ω such that 1) ω satisfies C ; and 2) $F|\alpha \models F|\omega$.*

This is the **strongest notion** of redundancy. However, entailment (\models) cannot be checked in polynomial time (assuming $P \neq NP$), unless **bounded**.

Checking Redundancy Using Unit Propagation

- **Unit propagation** (UP) satisfies unit clauses by assigning their literal to true (until fixpoint or a conflict).
- Let F be a formula, C a clause, and α the smallest assignment that falsifies C . C is **implied by F via UP** (denoted by $F \vdash_1 C$) if UP on $F|\alpha$ results in a conflict.
- Implied by UP is used in SAT solvers to determine redundancy of learned clauses and therefore \vdash_1 is a **natural restriction** of \models .
- We bound $F|\alpha \models F|\omega$ by $F|\alpha \vdash_1 F|\omega$.
- **Example:**
 $F = (x \vee y \vee z) \wedge (\bar{x} \vee y \vee z) \wedge (x \vee \bar{y} \vee z) \wedge (\bar{x} \vee \bar{y} \vee z)$
and $G = (z)$. Observe that $F \models G$, but that $F \not\vdash_1 G$.

Hand-crafted PR Proofs of Pigeon Hole Formulas

We manually constructed PR proofs of the famous pigeon hole formulas and the two-pigeons-per-hole family.

- The proofs consist only of **binary and unit** clauses.
- Only **original variables** appear in the proof.
- All proofs are **linear** in the size of the formula.
- ➔ The PR proofs are smaller than Cook's **ER proofs**.
- All resolution proofs of these formulas are **exponential** in size.

Proofs of Unsatisfiability

Beyond Resolution

Propagation Redundancy

Satisfaction-Driven Clause Learning

Challenges

Finding Redundant Clauses: The Positive Reduct

Determining whether a clause C is SET or PR w.r.t. a formula F is an NP-complete problem.

How to find SET and PR clauses? Encode it in **SAT**!

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Given a formula F and a clause C . Let α denote the smallest assignment that falsifies C . The **positive reduct** of F and α is a formula which is satisfiable if and only if C is SET w.r.t. F .

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Key Idea: While solving a formula F , check whether the positive reduct of F and the current assignment α is **satisfiable**. In that case, **prune** the branch α .

The Positive Reduct: An Example

Given a formula F and a clause C . Let α denote the smallest assignment that falsifies C . The **positive reduct** of F and α , denoted by $p(F, \alpha)$, is the formula that contains C and all assigned(D, α) with $D \in F$ and D is satisfied by α .

Example

Consider the formula $F := (x \vee y) \wedge (x \vee \bar{y}) \wedge (\bar{y} \vee \bar{z})$.

Let $C_1 = (\bar{x})$, so $\alpha_1 = x$.

The positive reduct $p(F, \alpha_1) = (\bar{x}) \wedge (x) \wedge (x)$ is **unsatisfiable**.

Let $C_2 = (\bar{x} \vee \bar{y})$, so $\alpha_2 = x y$.

The positive reduct $p(F, \alpha_2) = (\bar{x} \vee \bar{y}) \wedge (x \vee y) \wedge (x \vee \bar{y})$ is **satisfiable**.

Autarkies

A non-empty assignment α is an **autarky** for formula F if every clause $C \in F$ that is **touched** by α is also **satisfied** by α .

A **pure literal** and a **satisfying assignment** are autarkies.

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Assignment $\alpha_1 = \bar{z}$ is an autarky:

$(x \vee y) \wedge (x \vee \bar{y}) \wedge (\bar{y} \vee \bar{z})$. Assignment $\alpha_2 = x \bar{y} z$ is an autarky: $(x \vee y) \wedge (x \vee \bar{y}) \wedge (\bar{y} \vee \bar{z})$.

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Given an assignment α , $F|_{\alpha}$ denotes a formula F without the clauses satisfied by α and without the literals falsified by α .

Theorem ([Monien and Speckenmeyer 1985])

Let α be an autarky for formula F .

Then, F and $F|_{\alpha}$ are satisfiability equivalent.

Conditional Autarkies

An assignment $\alpha = \alpha_{\text{con}} \cup \alpha_{\text{aut}}$ is a **conditional autarky** for formula F if α_{aut} is an autarky for $F|\alpha_{\text{con}}$.

Example

Consider the formula $F := (x \vee y) \wedge (x \vee \bar{y}) \wedge (\bar{y} \vee \bar{z})$.

Let $\alpha_{\text{con}} = x$ and $\alpha_{\text{aut}} = \bar{y}$, then $\alpha = \alpha_{\text{con}} \cup \alpha_{\text{aut}} = x\bar{y}$ is a conditional autarky for F :

$$\alpha_{\text{aut}} = \bar{y} \text{ is an autarky for } F|\alpha_{\text{con}} = (\bar{y} \vee \bar{z}).$$

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Let $\alpha = \alpha_{\text{con}} \cup \alpha_{\text{aut}}$ be a **conditional autarky** for formula F . Then F and $F \wedge (\alpha_{\text{con}} \rightarrow \alpha_{\text{aut}})$ are satisfiability-equivalent.

In the above example, we could therefore learn $(\bar{x} \vee \bar{y})$.

Learning PR clauses

Theorem

Given a formula F and an assignment α . Every *satisfying assignment* ω of $p(F, \alpha)$ is a *conditional autarky* of F .

Recall: Given a formula F and a clause C . Let α denote the smallest assignment that falsifies C . C is SET w.r.t. F if and only if $p(F, \alpha)$ is *satisfiable*.

Let assignment ω satisfy $p(F, \alpha)$. Removing all but one of the literals in C that are satisfied by ω results in a *PR clause* w.r.t. F .

Pseudo-Code of CDCL (formula F)

```
1    $\alpha := \emptyset$ 
2   forever do
3      $\alpha := \text{Simplify} (F, \alpha)$ 
4     if  $F|_{\alpha}$  contains a falsified clause then
5        $C := \text{AnalyzeConflict} ()$ 
6       if C is the empty clause then return unsatisfiable
7        $F := F \cup \{C\}$ 
8        $\alpha := \text{BackJump} (C, \alpha)$ 
13  else
14     $l := \text{Decide} ()$ 
15    if  $l$  is undefined then return satisfiable
16     $\alpha := \alpha \cup \{l\}$ 
```

Pseudo-Code of SDCL (formula F)

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2  forever do
3     $\alpha := \text{Simplify}(F, \alpha)$ 
4    if  $F|\alpha$  contains a falsified clause then
5       $C := \text{AnalyzeConflict}()$ 
6      if C is the empty clause then return unsatisfiable
7       $F := F \cup \{C\}$ 
8       $\alpha := \text{BackJump}(C, \alpha)$ 
9    else if  $p(F, \alpha)$  is satisfiable then
10      $C := \text{AnalyzeWitness}()$ 
11      $F := F \cup \{C\}$ 
12      $\alpha := \text{BackJump}(C, \alpha)$ 
13   else
14      $l := \text{Decide}()$ 
15     if  $l$  is undefined then return satisfiable
16      $\alpha := \alpha \cup \{l\}$ 
```

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Theoretical Challenges

Lower bounds for interference-based proof systems with new variables will be hard, but what about **without new variables**?

- Lower bound for BC w/o new variables? Pigeon-hole formulas?
- Lower bound for SET w/o new variables? Tseitin formulas?
- Lower bound for PR w/o new variables?!

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Can the new proof systems without new variables **simulate** old ones, in particular **Frege systems** (or the other way around)?

What about **cutting planes**?

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Can the new proof systems without new variables **simulate** old ones, in particular **Frege systems** (or the other way around)?

What about **cutting planes**?

Can we design stronger proof systems that make it even **easier to compute** short proofs?

Practical Challenges

The current version of SDCL is just the beginning:

- Which heuristics allow learning short PR clauses?
- How to construct an AnalyzeWitness procedure?
- Can the positive reduct be improved?

Can **local search** be used to find short proofs of unsatisfiability?

Constructing positive reducts (or similar formulas) efficiently:

- Generating a positive reduct is **more costly** than solving them
- Can we design **data-structures** to cheaply compute them?