



APPLICATION: Variable Elimination

Matt Gormley
Lecture 11
Nov. 28, 2018

Reminders

- Homework C: Data Structures
 - Out: Mon, Nov. 26
 - Due: Mon, Dec. 3 at 11:59pm
- Quiz B: Computation; Programming & Efficiency
 - Wed, Dec. 5, in-class
 - Covers Lectures 7 – 12

APPLICATION: EXACT INFERENCE IN GRAPHICAL MODELS

EXACT INFERENCE

Exact Inference

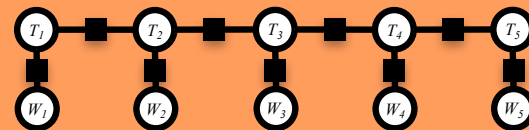
1. Data

$$\mathcal{D} = \{\mathbf{x}^{(n)}\}_{n=1}^N$$

Sample 1:	n time	v flies	p like	d an	n irrov
Sample 2:	n time	n flies	v like	d an	n irrov
Sample 3:	n flies	v fly	p with	n their	n rings
Sample 4:	p with	n time	n you	v will	v see

2. Model

$$p(\mathbf{x} \mid \boldsymbol{\theta}) = \frac{1}{Z(\boldsymbol{\theta})} \prod_{C \in \mathcal{C}} \psi_C(\mathbf{x}_C)$$



3. Objective

$$\ell(\boldsymbol{\theta}; \mathcal{D}) = \sum_{n=1}^N \log p(\mathbf{x}^{(n)} \mid \boldsymbol{\theta})$$

5. Inference

1. Marginal Inference

$$p(\mathbf{x}_C) = \sum_{\mathbf{x}': \mathbf{x}'_C = \mathbf{x}_C} p(\mathbf{x}' \mid \boldsymbol{\theta})$$

2. Partition Function

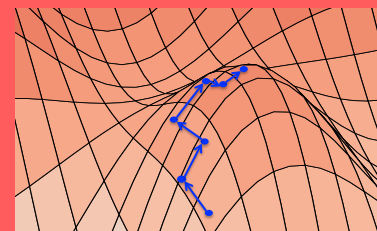
$$Z(\boldsymbol{\theta}) = \sum_{\mathbf{x}} \prod_{C \in \mathcal{C}} \psi_C(\mathbf{x}_C)$$

3. MAP Inference

$$\hat{\mathbf{x}} = \operatorname{argmax}_{\mathbf{x}} p(\mathbf{x} \mid \boldsymbol{\theta})$$

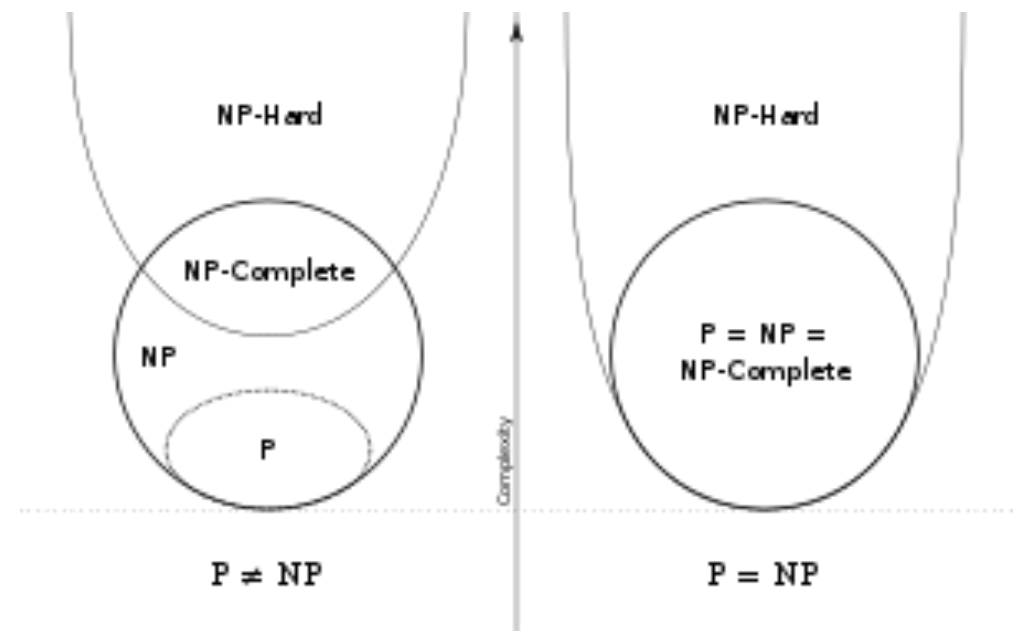
4. Learning

$$\boldsymbol{\theta}^* = \operatorname{argmax}_{\boldsymbol{\theta}} \ell(\boldsymbol{\theta}; \mathcal{D})$$



Complexity Classes

- An algorithm runs in **polynomial time** if its runtime is a polynomial function of the input size (e.g. $O(n^k)$ for some fixed constant k)
- The **class P** consists of all problems that can be solved in polynomial time
- A problem for which the answer is binary (e.g. yes/no) is called a **decision problem**
- The **class NP** contains all decision problems where 'yes' answers can be verified (proved) in polynomial time
- A problem is **NP-Hard** if given an $O(1)$ oracle to solve it, every problem in NP can be solved in polynomial time (e.g. by reduction)
- A problem is **NP-Complete** if it belongs to both the classes NP and NP-Hard



5. Inference

Three Tasks:

1. Marginal Inference (#P-Hard)

Compute marginals of variables and cliques

$$p(x_i) = \sum_{\mathbf{x}' : x'_i = x_i} p(\mathbf{x}' \mid \boldsymbol{\theta}) \quad \Bigg| \quad p(\mathbf{x}_C) = \sum_{\mathbf{x}' : \mathbf{x}'_C = \mathbf{x}_C} p(\mathbf{x}' \mid \boldsymbol{\theta})$$

2. Partition Function (#P-Hard)

Compute the normalization constant

$$Z(\boldsymbol{\theta}) = \sum_{\mathbf{x}} \prod_{C \in \mathcal{C}} \psi_C(\mathbf{x}_C)$$

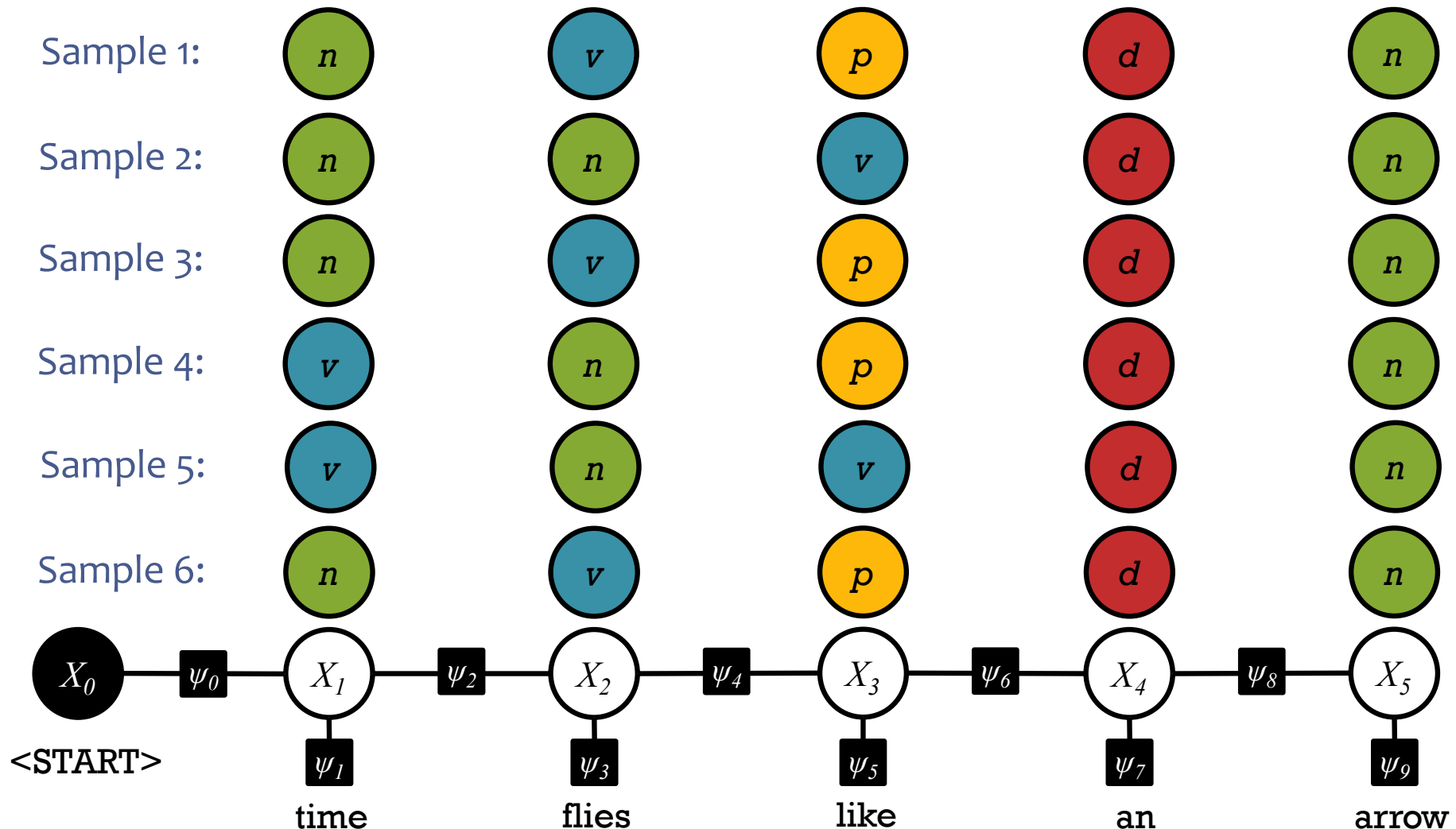
3. MAP Inference (NP-Hard)

Compute variable assignment with highest probability

$$\hat{\mathbf{x}} = \underset{\mathbf{x}}{\operatorname{argmax}} p(\mathbf{x} \mid \boldsymbol{\theta})$$

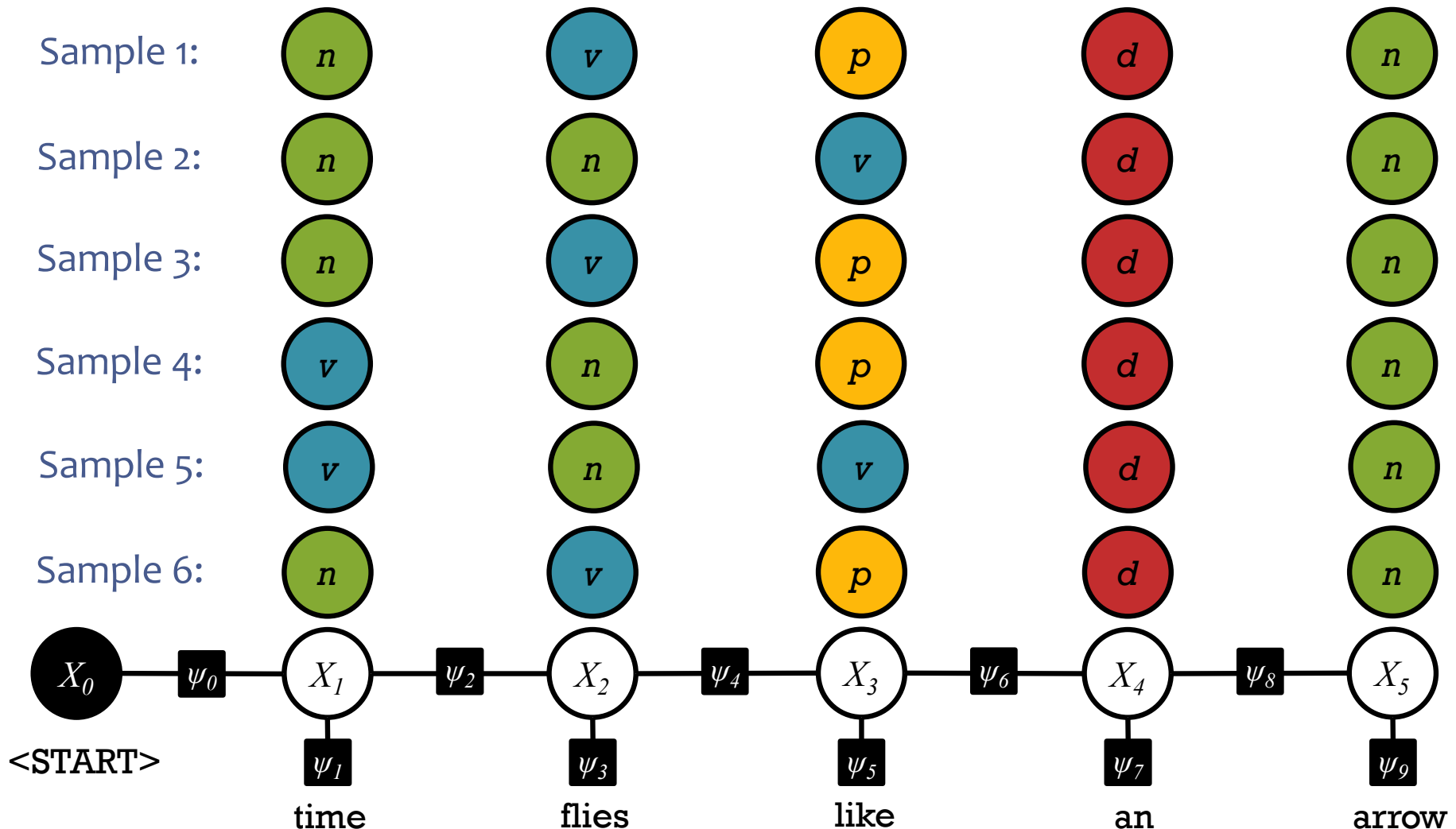
Marginals by Sampling on Factor Graph

Suppose we took many samples from the distribution over taggings: $p(\mathbf{x}) = \frac{1}{Z} \prod_{\alpha} \psi_{\alpha}(\mathbf{x}_{\alpha})$



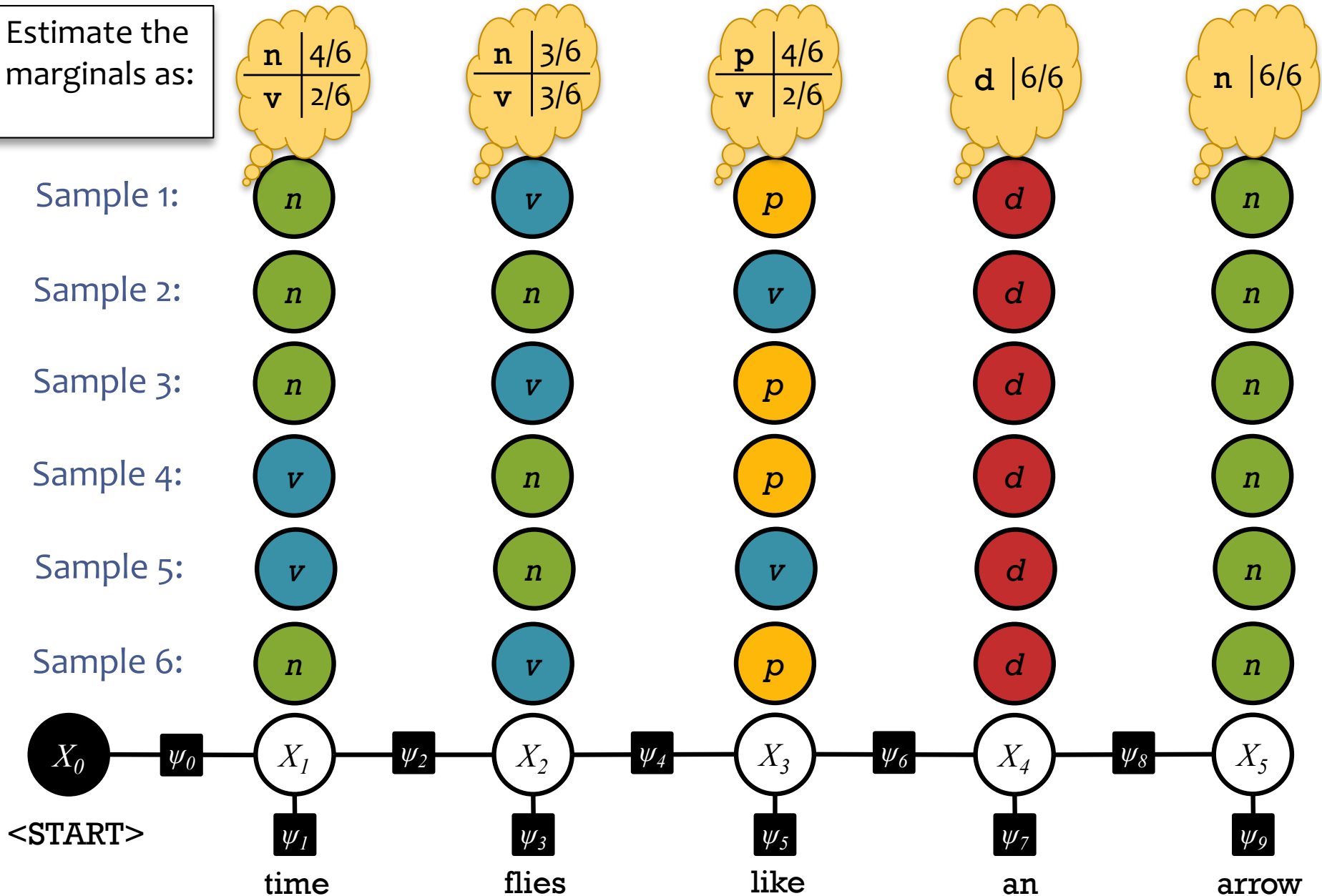
Marginals by Sampling on Factor Graph

The marginal $p(X_i = x_i)$ gives the probability that variable X_i takes value x_i in a random sample



Marginals by Sampling on Factor Graph

Estimate the
marginals as:



Simple and general exact inference for graphical models

VARIABLE ELIMINATION

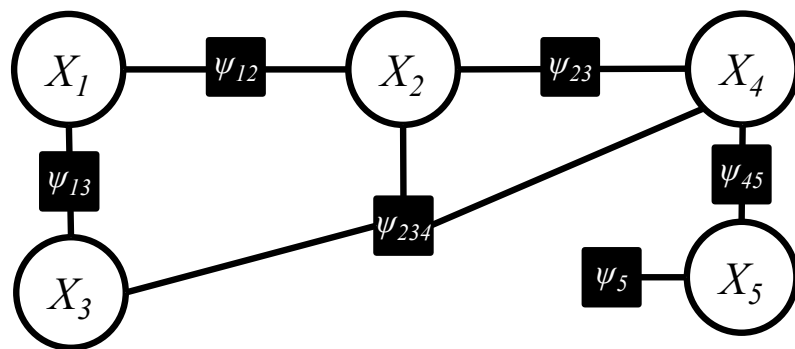
Brute Force (Naïve) Inference

For all i , suppose the **range** of X_i is $\{0, 1, 2\}$.

Let $k=3$ denote the **size of the range**.

The distribution **factorizes** as:

$$p(\mathbf{x}) = \psi_{12}(x_1, x_2) \psi_{13}(x_1, x_3) \psi_{24}(x_2, x_4) \\ \psi_{234}(x_2, x_3, x_4) \psi_{45}(x_4, x_5) \psi_5(x_5)$$



Naively, we compute the **partition function** as:

$$Z = \sum_{x_1} \sum_{x_2} \sum_{x_3} \sum_{x_4} \sum_{x_5} p(\mathbf{x})$$

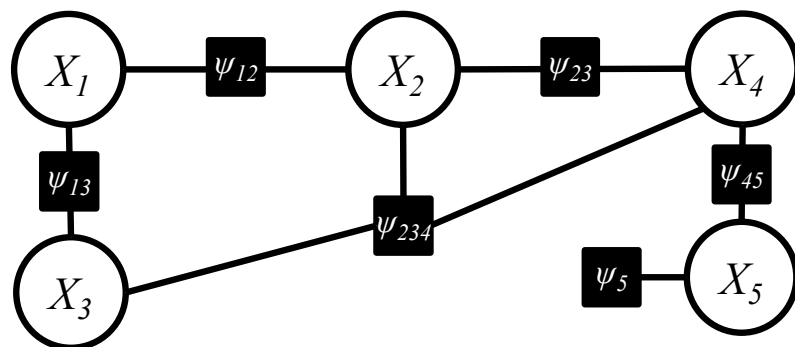
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$p(\mathbf{x})$ can be represented as a joint probability table with 3^5 entries:

x_1	x_2	x_3	x_4	x_5	$p(\mathbf{x})$
0	0	0	0	0	0.019517693
0	0	0	0	1	0.017090249
0	0	0	0	2	0.014885825
0	0	0	1	0	0.024117638
0	0	0	1	1	0.000925849
0	0	0	1	2	0.028112576
0	0	0	2	0	0.028050205
0	0	0	2	1	0.004812689
0	0	0	2	2	0.007987737
0	0	1	0	0	0.028433687
0	0	1	0	1	0.037073469
0	0	1	0	2	0.013558227
0	0	1	1	0	0.019479016
0	0	1	1	1	0.012312901
0	0	1	1	2	0.023439775
0	0	1	2	0	0.038206131
0	0	1	2	1	0.038996005
0	0	1	2	2	0.041458783
0	0	2	0	0	0.044616806
0	0	2	0	1	0.020846989
0	0	2	0	2	0.03006475
0	0	2	1	0	0.048436964
0	0	2	1	1	0.02854376
0	0	2	1	2	0.029191506
0	0	2	2	0	0.031531118
0	0	2	2	1	0.005132392
0	0	2	2	2	0.032027091
...

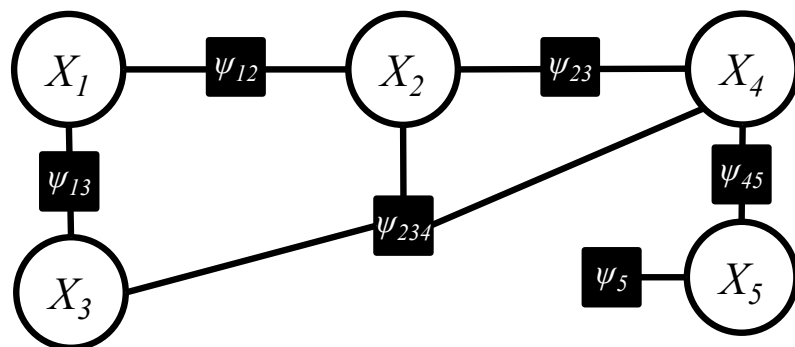
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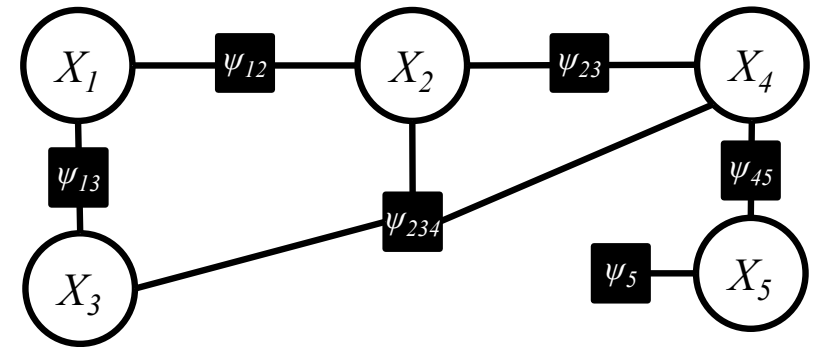
x_1	x_2	x_3	x_4	x_5	$p(\mathbf{x})$
0	0	0	0	0	0.019517693
0	0	0	0	1	0.017090249
0	0	0	0	2	0.014885825
0	0	0	1	0	0.024117638
0	0	0	1	1	0.000925849
0	0	0	1	2	0.028112576
0	0	0	2	0	0.028050205
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0	0	0	2	2	0.007987737
0	0	1	0	0	0.028433687
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0	0	1	1	0	0.019479016
0	0	1	1	1	0.012312901
0	0	1	1	2	0.023439775
0	0	1	2	0	0.038206131
0	0	1	2	1	0.038996005
0	0	1	2	2	0.041458783
0	0	2	0	0	0.044616806
0	0	2	0	1	0.020846989
0	0	2	0	2	0.03006475
0	0	2	1	0	0.048436964
0	0	2	1	1	0.02854376

Naïve computation of Z requires 3^5 additions.

Can we do better?

The Variable Elimination Algorithm

Instead, capitalize on the factorization of $p(\mathbf{x})$.



$$\begin{aligned}
 Z &= \sum_{x_1} \sum_{x_2} \sum_{x_3} \sum_{x_4} \sum_{x_5} \psi_{12}(x_1, x_2) \psi_{13}(x_1, x_3) \psi_{24}(x_2, x_4) \psi_{234}(x_2, x_3, x_4) \psi_{45}(x_4, x_5) \psi_5(x_5) \\
 &= \sum_{x_1} \sum_{x_2} \sum_{x_3} \sum_{x_4} \psi_{12}(x_1, x_2) \psi_{13}(x_1, x_3) \psi_{24}(x_2, x_4) \psi_{234}(x_2, x_3, x_4) \underbrace{\sum_{x_5} \psi_{45}(x_4, x_5) \psi_5(x_5)}_{\text{factor}}
 \end{aligned}$$

Only 3^2 additions are needed to marginalize out x_5 .

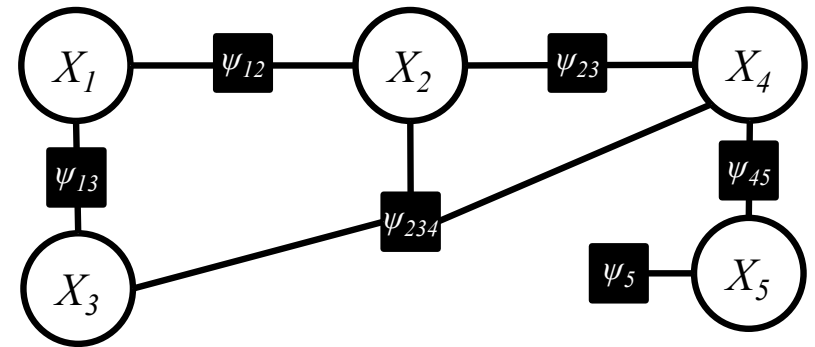
We denote the **marginal's table** by $m_5(x_4)$.

This “factor” is a much smaller table with 3^2 entries:

x_4	x_5	$p(\mathbf{x})$
0	0	0.019517693
0	1	0.017090249
0	2	0.014885825
1	0	0.024117638
1	1	0.000925849
1	2	0.028112576
2	0	0.028050205
2	1	0.004812689
2	2	0.007987737

The Variable Elimination Algorithm

Instead, capitalize on the factorization of $p(\mathbf{x})$.

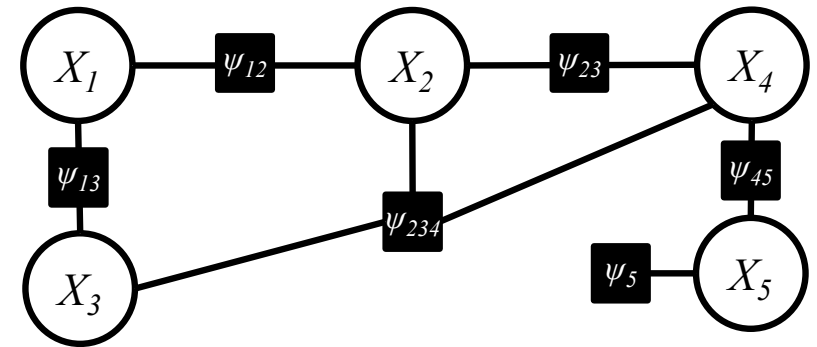


$$\begin{aligned}
 Z &= \sum_{x_1} \sum_{x_2} \sum_{x_3} \sum_{x_4} \sum_{x_5} \psi_{12}(x_1, x_2) \psi_{13}(x_1, x_3) \psi_{24}(x_2, x_4) \psi_{234}(x_2, x_3, x_4) \psi_{45}(x_4, x_5) \psi_5(x_5) \\
 &= \sum_{x_1} \sum_{x_2} \sum_{x_3} \sum_{x_4} \psi_{12}(x_1, x_2) \psi_{13}(x_1, x_3) \psi_{24}(x_2, x_4) \psi_{234}(x_2, x_3, x_4) \sum_{x_5} \psi_{45}(x_4, x_5) \psi_5(x_5) \\
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 \end{aligned}$$

$$m_5(x_4) \triangleq \sum_{x_5} \psi_{45}(x_4, x_5) \psi_5(x_5)$$

The Variable Elimination Algorithm

Instead, capitalize on the factorization of $p(\mathbf{x})$.



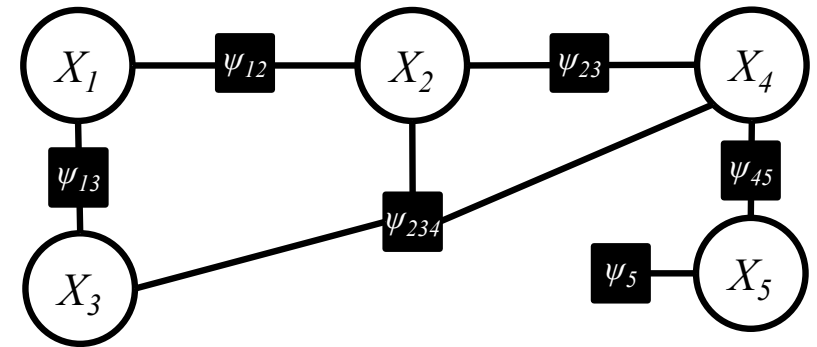
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 &= \sum_{x_1} \sum_{x_2} \sum_{x_3} \sum_{x_4} \psi_{12}(x_1, x_2) \psi_{13}(x_1, x_3) \psi_{24}(x_2, x_4) \psi_{234}(x_2, x_3, x_4) \sum_{x_5} \psi_{45}(x_4, x_5) \psi_5(x_5) \\
 &= \sum_{x_1} \sum_{x_2} \sum_{x_3} \sum_{x_4} \underbrace{\psi_{12}(x_1, x_2) \psi_{13}(x_1, x_3) \psi_{24}(x_2, x_4) \psi_{234}(x_2, x_3, x_4)}_{m_5(x_4)} m_5(x_4)
 \end{aligned}$$

This “factor” is still a 3^4 table so apply the same trick again.

$$m_5(x_4) \triangleq \sum_{x_5} \psi_{45}(x_4, x_5) \psi_5(x_5)$$

The Variable Elimination Algorithm

Instead, capitalize on the factorization of $p(\mathbf{x})$.



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 Z &= \sum_{x_1} \sum_{x_2} \psi_{12}(x_1, x_2) \sum_{x_3} \psi_{13}(x_1, x_3) \sum_{x_4} \psi_{24}(x_2, x_4) \psi_{234}(x_2, x_3, x_4) \sum_{x_5} \psi_{45}(x_4, x_5) \psi_5(x_5) \\
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 &= \sum_{x_1} \sum_{x_2} \psi_{12}(x_1, x_2) \sum_{x_3} \psi_{13}(x_1, x_3) m_4(x_2, x_3) \\
 &= \sum_{x_1} \sum_{x_2} \psi_{12}(x_1, x_2) m_3(x_1, x_2) \\
 &= \sum_{x_1} m_2(x_1)
 \end{aligned}$$

3^2 additions

3^3 additions

3^3 additions

3^2 additions

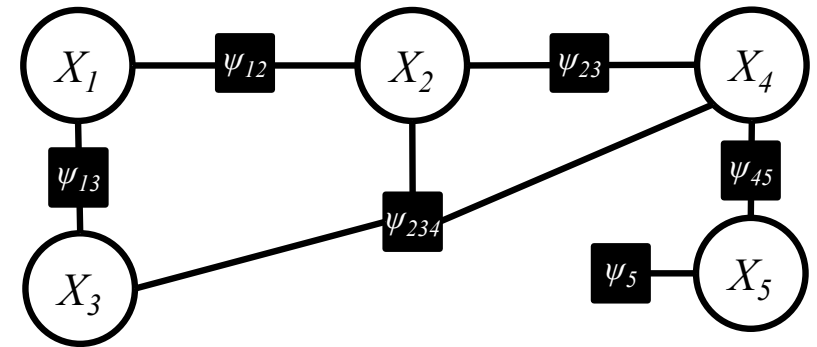
3 additions

Naïve solution requires $3^5 = 243$ additions.

Variable elimination only requires $3 + 3^2 + 3^3 + 3^3 + 3^2 = 75$ additions.

The Variable Elimination Algorithm

The same trick can be used to compute **marginal probabilities**. Just choose the variable elimination order such that the query variables are last.



$$\begin{aligned}
 p(x_1) &= \frac{1}{Z} \sum_{x_2} \psi_{12}(x_1, x_2) \sum_{x_3} \psi_{13}(x_1, x_3) \sum_{x_4} \psi_{24}(x_2, x_4) \psi_{234}(x_2, x_3, x_4) \sum_{x_5} \psi_{45}(x_4, x_5) \psi_5(x_5) \\
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 &= \frac{1}{Z} \sum_{x_2} \psi_{12}(x_1, x_2) \sum_{x_3} \psi_{13}(x_1, x_3) m_4(x_2, x_3) \\
 &= \frac{1}{Z} \sum_{x_2} \psi_{12}(x_1, x_2) m_3(x_1, x_2) \\
 &= \frac{1}{Z} m_2(x_1)
 \end{aligned}$$

3 different values on LHS

3^2 additions

3^3 additions

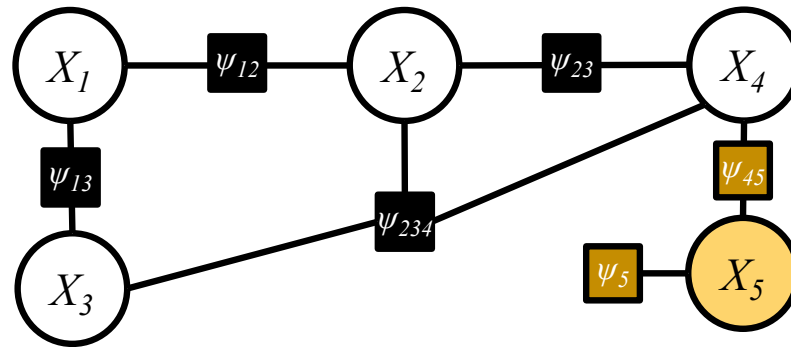
3^3 additions

3^2 additions

For directed graphs, $Z = 1$.

For undirected graphs, if we compute each (unnormalized) value on the LHS, we can sum them to get Z .

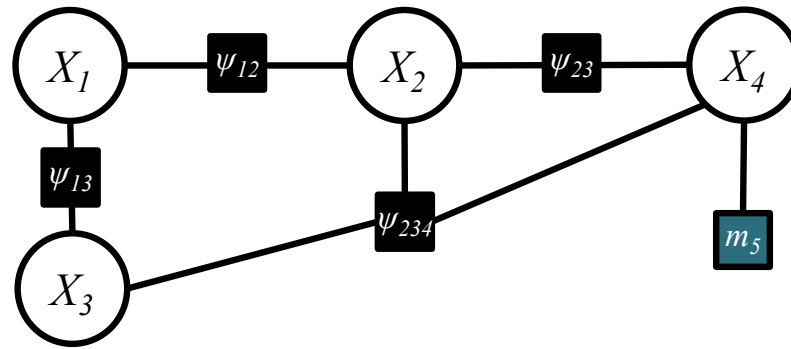
The Variable Elimination Algorithm



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 &= \sum_{x_1} \sum_{x_2} \psi_{12}(x_1, x_2) \sum_{x_3} \psi_{13}(x_1, x_3) m_4(x_2, x_3)
 \end{aligned}$$

In a factor graph, variable **elimination** corresponds to replacement of a subgraph with a factor.

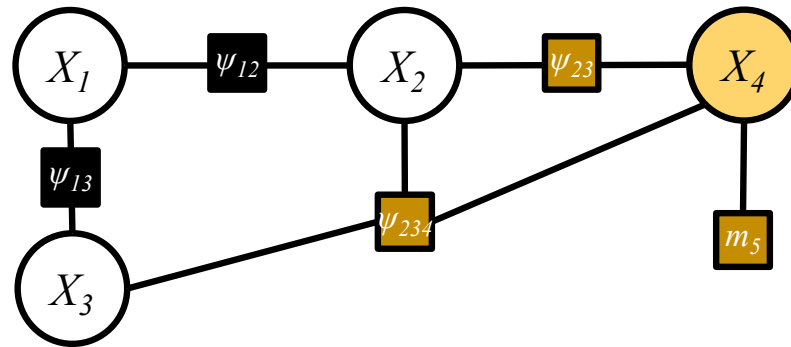
The Variable Elimination Algorithm



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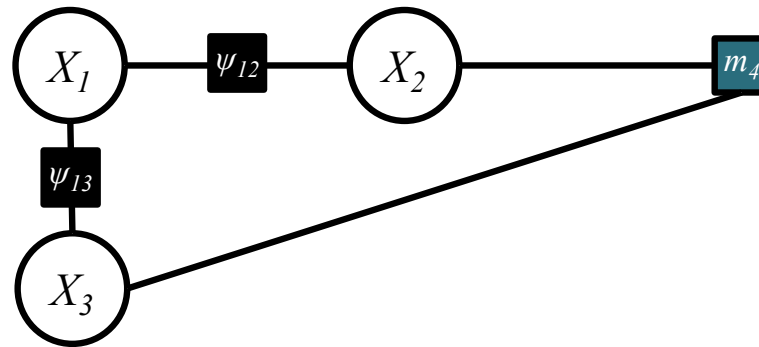
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The Variable Elimination Algorithm



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 \end{aligned}$$

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Variable Elimination for Marginal Inference

Algorithm 1: Variable Elimination for Marginal Inference

Input: the factor graph and the query variable

Output: the marginal distribution for the query variable

- a. Run a breadth-first-search starting at the query variable to obtain an ordering of the variable nodes
- b. Reverse that ordering
- c. Eliminate each variable in the reversed ordering using Algorithm 2

Algorithm 2: Eliminate One Variable

Input: the variable to be eliminated

Output: new factor graph with the variable marginalized out

- a. Find the input variable and its neighboring factors -- call this set the eliminated set
- b. Replace the eliminated set with a new factor
 - a. The neighbors of the new factor should be all the neighbors of all the factors in the eliminated set
 - b. The new factor should assign a score to each possible assignment of its neighboring variables
 - c. Said score should be identical to the product of the factors it is replacing, summing over the eliminated variable

Variable Elimination for Marginal Inference

Algorithm 3: Variable Elimination for the Partition Function

Input: the factor graph

Output: the partition function

- a. Run a breadth-first-search starting at an arbitrary variable to obtain an ordering of the variable nodes
- b. Eliminate each variable in the ordering using Algorithm 2

Algorithm 2: Eliminate One Variable

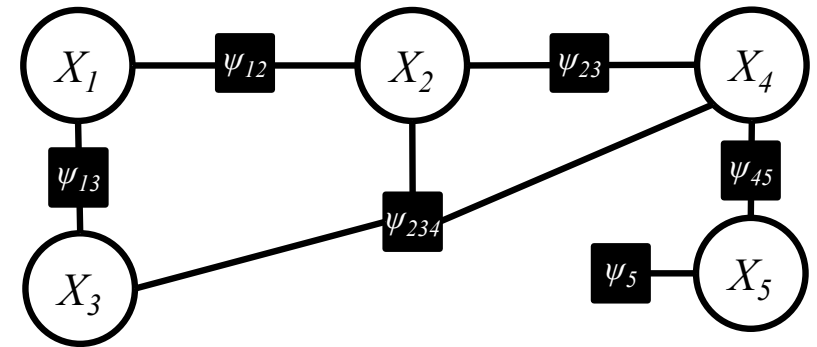
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 - a. The neighbors of the new factor should be all the neighbors of all the factors in the eliminated set
 - b. The new factor should assign a score to each possible assignment of its neighboring variables
 - c. Said score should be identical to the product of the factors it is replacing, summing over the eliminated variable

Variable Elimination Complexity

Instead, capitalize on the factorization of $p(\mathbf{x})$.



In-Class Exercise: *Fill in the blank*

Brute force, naïve,
inference is $O(\underline{\hspace{1cm}})$

Variable elimination
is $O(\underline{\hspace{1cm}})$

where $n = \#$ of variables
 $k = \max \#$ values a variable can take
 $r = \#$ variables participating in
largest “intermediate” table

PROFILING FOR EFFICIENCY

Software Profiling

CPU Profiler:

- **Intermediate Goal:** Analyze the CPU usage of a program at a fine-grained level (e.g. time spent within each function)
- **End Goal:** To make the program more CPU efficient by optimizing most time consuming parts of program

Memory Profiler:

- **Intermediate Goal:** Analyze the memory consumption of a program (e.g. how much space does a particular type of object use on the heap)
- **End Goal:** To make the program more memory by utilizing different data structures or data storage techniques to reduce memory load

Software Profiling

Deterministic CPU Profiler

- Augments the code with **additional bookkeeping** calls
- Provides **exact number** of times each function is called, and **exact amount** of time spent in each function
- Comes at the cost of much **slower runtime**

Statistical CPU Profiler

- Leaves the code nearly unchanged, and instead **takes samples** (hundreds or more) of the stacktrace
- Provides the **proportion of samples** that landed in each function and **estimates** the total time spent in each function
- Typically yields **little to no slowdown** of the code

Line Profiler

- Same as above for each type, but counts the **number of times each line is executed** and provides the amount of **time spent on each line**
- Increases complexity of the profiler, but provides **much more detailed analysis**

Python Profilers

Name	Type	Level of Detail	Output	Notes
cProfile	deterministic	function-level	console	built into Python standard library; C-based implementation
profile	deterministic	function-level	console	same as cProfile, but implemented in pure Python
line_profiler	deterministic + statistical	line-level	console	C-based implementation
pprofile	deterministic + statistical	line-level	console	pure Python implementation (few users)
PyFlame by Uber	deterministic + statistical	line-level	flame graph	Linux only as of 2018
Plop by Dropbox	deterministic + statistical	line-level	bubble plot	(few users)

cProfile Output

```
$ python -m cProfile -s cumtime lwn2pocket.py
72270 function calls (70640 primitive calls) in 4.481 seconds

Ordered by: cumulative time
```

ncalls	tottime	percall	cumtime	percall	filename:lineno(function)
1	0.004	0.004	4.481	4.481	lwn2pocket.py:2(<module>)
1	0.001	0.001	4.296	4.296	lwn2pocket.py:51(main)
3	0.000	0.000	4.286	1.429	api.py:17(request)
3	0.000	0.000	4.268	1.423	sessions.py:386(request)
4/3	0.000	0.000	3.816	1.272	sessions.py:539(send)
4	0.000	0.000	2.965	0.741	adapters.py:323(send)
4	0.000	0.000	2.962	0.740	connectionpool.py:421(urlopen)
4	0.000	0.000	2.961	0.740	connectionpool.py:317(_make_request)
2	0.000	0.000	2.675	1.338	api.py:98(post)
30	0.000	0.000	1.621	0.054	ssl.py:727(recv)
30	0.000	0.000	1.621	0.054	ssl.py:610(read)
30	1.621	0.054	1.621	0.054	{method 'read' of '_ssl._SSLSocket' objects}
1	0.000	0.000	1.611	1.611	api.py:58(get)
4	0.000	0.000	1.572	0.393	httplib.py:1095(getresponse)
4	0.000	0.000	1.572	0.393	httplib.py:446(begin)
60	0.000	0.000	1.571	0.026	socket.py:410(readline)
4	0.000	0.000	1.571	0.393	httplib.py:407(_read_status)
1	0.000	0.000	1.462	1.462	pocket.py:44(wrapped)
1	0.000	0.000	1.462	1.462	pocket.py:152(make_request)
1	0.000	0.000	1.462	1.462	pocket.py:139(_make_request)
1	0.000	0.000	1.459	1.459	pocket.py:134(_post_request)

[...]

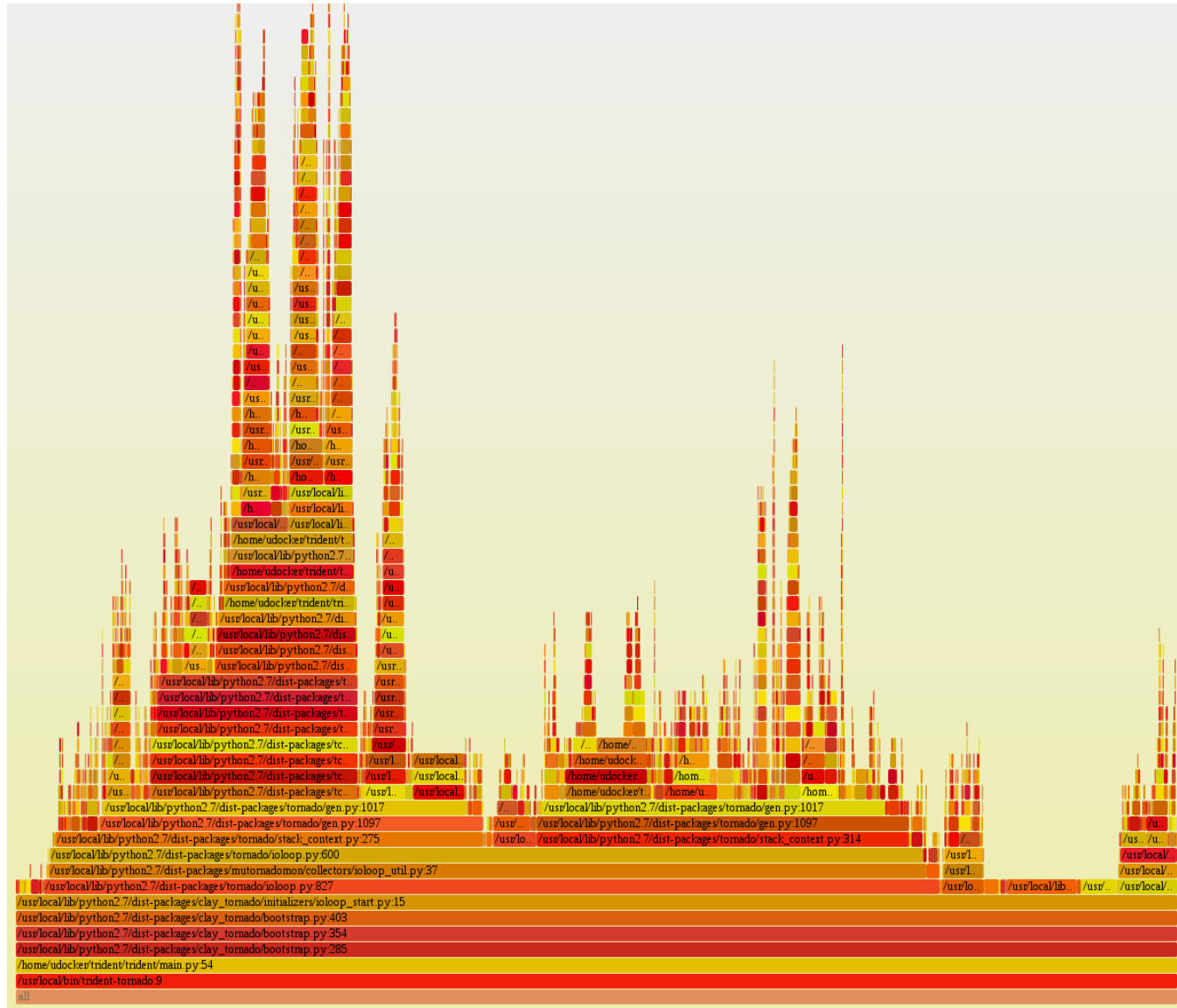
line_profiler Output

```
Pystone(1.1) time for 50000 passes = 2.48
This machine benchmarks at 20161.3 pystones/second
Wrote profile results to pystone.py.lprof
Timer unit: 1e-06 s
```

```
File: pystone.py
Function: Proc2 at line 149
Total time: 0.606656 s
```

Line #	Hits	Time	Per Hit	% Time	Line Contents
149					@profile
150					def Proc2(IntParIO):
151	50000	82003	1.6	13.5	IntLoc = IntParIO + 10
152	50000	63162	1.3	10.4	while 1:
153	50000	69065	1.4	11.4	if Char1Glob == 'A':
154	50000	66354	1.3	10.9	IntLoc = IntLoc - 1
155	50000	67263	1.3	11.1	IntParIO = IntLoc - IntGlob
156	50000	65494	1.3	10.8	EnumLoc = Ident1
157	50000	68001	1.4	11.2	if EnumLoc == Ident1:
158	50000	63739	1.3	10.5	break
159	50000	61575	1.2	10.1	return IntParIO

PyFlame Output



Plop Output

