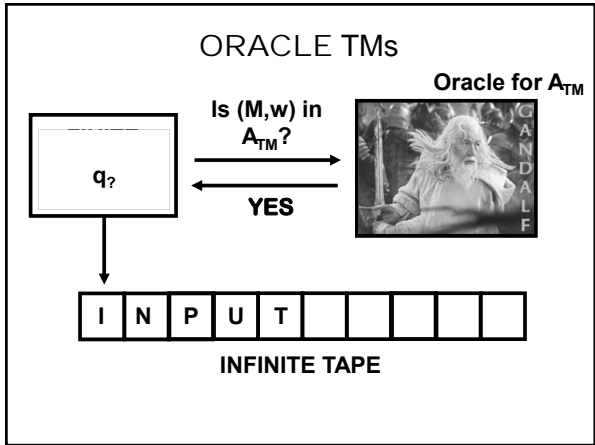
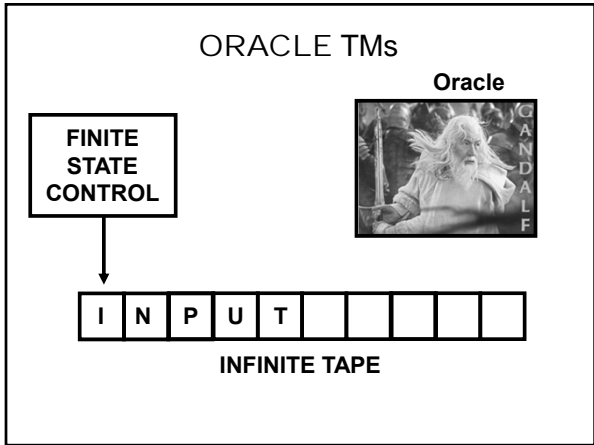


15-453

FORMAL LANGUAGES, AUTOMATA AND COMPUTABILITY

ORACLE TURING MACHINES AND TURING REDUCIBILITY

TUESDAY, MAR 4



ORACLE MACHINES

An **ORACLE** is a set B to which the TM may pose membership questions “Is w in B ?”
(formally: TM enters state $q_?$)
 and the TM always receives a correct answer in one step
(formally: if the string on the tape is in B , state $q_?$ is changed to q_{YES} , otherwise q_{NO})

This makes sense even if B is not decidable!
(We do not assume that the oracle B is a computable set!)

We say A is semi-decidable in B
if there is an oracle TM M with oracle B that
***semi-decides* A**

We say A is decidable in B
if there is an oracle TM M with oracle B that
***decides* A**

HALT_{TM} is DECIDABLE in A_{TM}

On input (M,w), decide if M halts on w as follows:

1. Ask the oracle for A_{TM} whether M accepts w. If yes, then ACCEPT
2. Switch the accept and reject states of M to get M'. Ask the oracle for A_{TM} whether M' accepts w. If yes, then ACCEPT
3. REJECT

A_{TM} is DECIDABLE in HALT_{TM}

On input (M,w), decide if M accepts w as follows:

Ask the oracle for HALT_{TM} whether M halts on w. If yes, then run M(w) and output its answer. If no, then REJECT.

Language A “Turing Reduces” to Language B

if A is decidable in B, ie if there is an oracle TM M with oracle B that decides A

A ≤_T B

≤_T VERSUS ≤_m

Theorem: If A ≤_m B then A ≤_T B

Proof:

If A ≤_m B then there is a computable function $f : \Sigma^* \rightarrow \Sigma^*$, where for every w,

$w \in A \Leftrightarrow f(w) \in B$

We can thus use an oracle for B to decide A

Theorem: $\neg \text{HALT}_{\text{TM}} \leq_T \text{HALT}_{\text{TM}}$

Theorem: $\neg \text{HALT}_{\text{TM}} \not\leq_m \text{HALT}_{\text{TM}}$

THE ARITHMETIC HIERARCHY

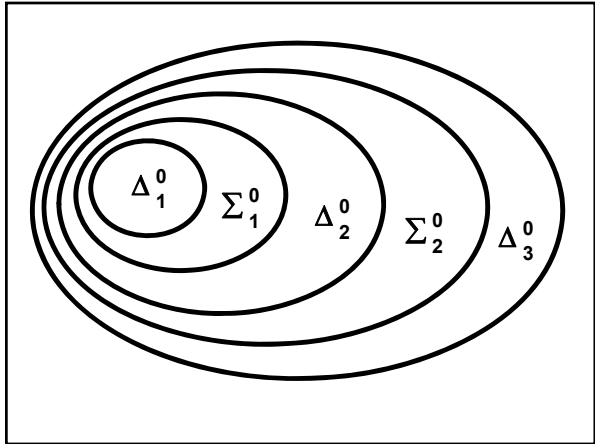
$\Delta_1^0 = \{ \text{decidable sets} \}$ (sets = languages)

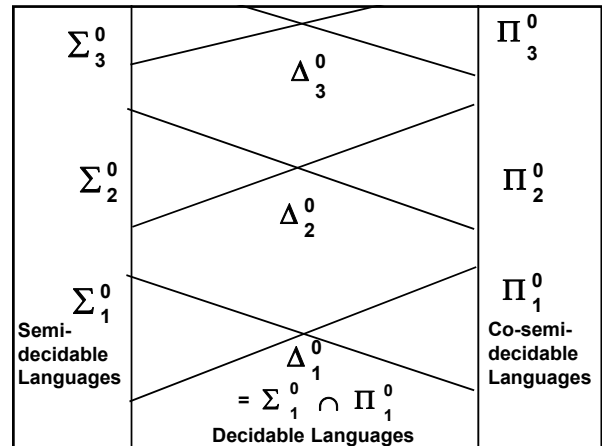
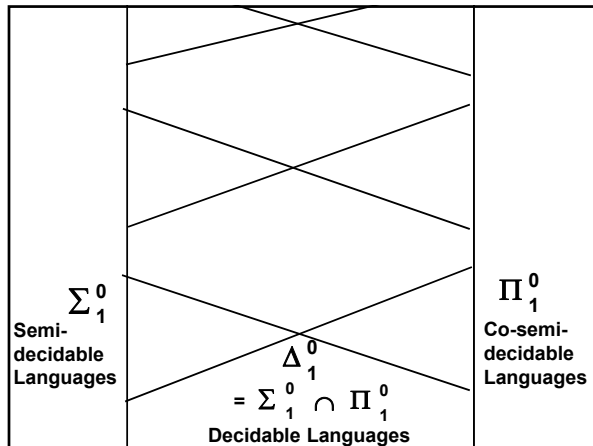
$\Sigma_1^0 = \{ \text{semi-decidable sets} \}$

$\Sigma_{n+1}^0 = \{ \text{sets semi-decidable in some } B \in \Sigma_n^0 \}$

$\Delta_{n+1}^0 = \{ \text{sets decidable in some } B \in \Sigma_n^0 \}$

$\Pi_n^0 = \{ \text{complements of sets in } \Sigma_n^0 \}$





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 Read Chapter 6.4 for next time