

Constraint Handling Rules with Multiset Comprehension Patterns

Edmund S. L. Lam Iliano Cervesato
sllam@qatar.cmu.edu iliano@cmu.edu

Carnegie Mellon University

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Outline

- 1 Introduction
- 2 Comprehensions in CHR^{cp}
- 3 Monotonicity
- 4 Semantics
- 5 Status

Constraint Handling Rules (CHR)

Constraint Logic Programming Languages

- Rule-based multiset rewriting
 - Declarative
 - Forward chaining
 - Committed choice
 - Concurrent
- CHR is a specific instance

Constraint Handling Rules (CHR)

$$r @ \bar{H}_p \setminus \bar{H}_s \iff g \mid \bar{B}$$

- \bar{H}_p , \bar{H}_s , \bar{B} are multisets of *atomic constraint patterns*: $p(\vec{t})$
 - r : Rule name
 - \bar{H}_p : Propagated head constraints
 - \bar{H}_s : Simplified head constraints
 - g : Guard conditions
 - \bar{B} : Body constraints
- $\left. \begin{array}{l} \text{LHS} \\ \text{RHS} \end{array} \right\}$
- A *program* \mathcal{P} is a set of rules

$$\mathcal{P} \triangleright St \mapsto_{\alpha} St'$$

- The *stores* St and St' are multiset of constraints

CHR with Multiset Comprehension Patterns (CHR^{cp})

- CHR rule patterns are strictly *atomic* predicates
- Pure forward chaining + atomic patterns are good
 - declarative and concise
 - high-level
- but not ideal when
 - performing aggregated operation
 - matching dynamic numbers of facts
- This work:
 - Extend CHR with *multiset comprehension patterns* (CHR^{cp})
 - Define an operational semantics for CHR^{cp}

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Example: Pivoted Swapping (in “Vanilla” CHR)

- LHS cannot match a dynamic numbers of constraints
- Two agents X and Y want to swap data D w.r.t. pivot P
- Standard CHR implementation:

$$\begin{array}{ll}
 \textit{init} @ \textit{swap}(X, Y, P) & \iff \textit{grabGE}(X, P, Y, []), \textit{grabLT}(Y, P, X, []) \\
 \textit{ge1} @ \textit{grabGE}(X, P, Y, Ds), \textit{data}(X, D) & \iff D \geq P \mid \textit{grabGE}(X, P, Y, [D \mid Ds]) \\
 \textit{ge2} @ \textit{grabGE}(X, P, Y, Ds) & \iff \textit{unrollData}(Y, Ds) \\
 \textit{lt1} @ \textit{grabLT}(Y, P, X, Ds), \textit{data}(Y, D) & \iff D < P \mid \textit{grabLT}(Y, P, X, [D \mid Ds]) \\
 \textit{lt2} @ \textit{grabLT}(Y, P, X, Ds) & \iff \textit{unrollData}(X, Ds) \\
 \textit{unroll1} @ \textit{unrollData}(L, [D \mid Ds]) & \iff \textit{unrollData}(L, Ds), \textit{data}(L, D) \\
 \textit{unroll2} @ \textit{unrollData}(L, []) & \iff \textit{true}
 \end{array}$$

- Verbose: 7 rules and 3 auxiliary constraints

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$ge1 @ grabGE(X, P, Y, Ds), data(X, D)$	\iff	$D \geq P \mid grabGE(X, P, Y, [D \mid Ds])$
$ge2 @ grabGE(X, P, Y, Ds)$	\iff	$unrollData(Y, Ds)$
$lt1 @ grabLT(Y, P, X, Ds), data(Y, D)$	\iff	$D < P \mid grabLT(Y, P, X, [D \mid Ds])$
$lt2 @ grabLT(Y, P, X, Ds)$	\iff	$unrollData(X, Ds)$
$unroll1 @ unrollData(L, [D \mid Ds])$	\iff	$unrollData(L, Ds), data(L, D)$
$unroll2 @ unrollData(L, [])$	\iff	$true$

- Verbose: 7 rules and 3 auxiliary constraints
- Accumulators incrementally store matches

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$ge2 @ grabGE(X, P, Y, Ds)$	\iff	$unrollData(Y, Ds)$
$lt1 @ grabLT(Y, P, X, Ds), data(Y, D)$	\iff	$D < P \mid grabLT(Y, P, X, [D \mid Ds])$
$lt2 @ grabLT(Y, P, X, Ds)$	\iff	$unrollData(X, Ds)$
$unroll1 @ unrollData(L, [D \mid Ds])$	\iff	$unrollData(L, Ds), data(L, D)$
$unroll2 @ unrollData(L, [])$	\iff	$true$

- Verbose: 7 rules and 3 auxiliary constraints
- Accumulators incrementally store matches
- Relies on rule priorities: apply $ge1$ before $ge2$, $lt1$ before $lt2$

CHR with Comprehension Patterns

- Additional form of constraint patterns:
 - *Comprehension patterns* M : $\{p(\vec{t}) \mid g\}_{\vec{x} \in t}$
 - Collects *all* $p(\vec{t})$ in the store that satisfy g
 - t is the multiset of all bindings to \vec{x}
- Pivoted Swapping in CHR^{cp} :

$$\begin{array}{l}
 \text{pivotSwap @ } \begin{array}{l} \text{swap}(X, Y, P) \\ \{data(X, D) \mid D \geq P\}_{D \in Xs} \\ \{data(Y, D) \mid D < P\}_{D \in Ys} \end{array} \iff \begin{array}{l} \{data(Y, D)\}_{D \in Xs} \\ \{data(X, D)\}_{D \in Ys} \end{array}
 \end{array}$$

- Xs and Ys built from the store — *output*
- Xs and Ys used to unfold the comprehensions — *input*

CHR with Comprehension Patterns

$$\text{pivotSwap} @ \begin{array}{l} \text{swap}(X, Y, P) \\ \{ \text{data}(X, D) \mid D \geq P \}_{D \in X_s} \\ \{ \text{data}(Y, D) \mid D < P \}_{D \in Y_s} \end{array} \iff \begin{array}{l} \{ \text{data}(Y, D) \}_{D \in X_s} \\ \{ \text{data}(X, D) \}_{D \in Y_s} \end{array}$$

- An example of CHR^{CP} rule application:

$$\mapsto_{\alpha} \begin{array}{l} \{ \text{swap}(a, b, 5), \text{data}(a, 1), \text{data}(a, 6), \text{data}(a, 7), \text{data}(b, 2), \text{data}(b, 8) \} \\ \{ \text{data}(a, 1), \text{data}(b, 6), \text{data}(b, 7), \text{data}(a, 2), \text{data}(b, 8) \} \end{array}$$

- Corresponds to the rule instance:

$$\text{pivotSwap} @ \begin{array}{l} \text{swap}(a, b, 5) \\ \{ \text{data}(a, D) \mid D \geq 5 \}_{D \in \{6, 7\}} \\ \{ \text{data}(b, D) \mid D < 5 \}_{D \in \{2\}} \end{array} \iff \begin{array}{l} \{ \text{data}(b, D) \}_{D \in \{6, 7\}} \\ \{ \text{data}(a, D) \}_{D \in \{2\}} \end{array}$$

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Maximality of Comprehension vs. Monotonicity

- Semantics of CHR^{cp} enforces *maximality* of comprehensions
- Example

$$\text{pivotSwap} @ \begin{array}{l} \text{swap}(X, Y, P) \\ \{ \text{data}(X, D) \mid D \geq P \}_{D \in X_s} \\ \{ \text{data}(Y, D) \mid D < P \}_{D \in Y_s} \end{array} \iff \begin{array}{l} \{ \text{data}(Y, D) \}_{D \in X_s} \\ \{ \text{data}(X, D) \}_{D \in Y_s} \end{array}$$

$$\begin{array}{l} \triangleright \{ \text{swap}(a, b, 5), \text{data}(a, 1), \boxed{\text{data}(a, 6)}, \text{data}(a, 7), \text{data}(b, 2), \text{data}(b, 8) \} \\ \not\vdash_{\alpha} \{ \text{data}(a, 1), \boxed{\text{data}(a, 6)}, \text{data}(b, 7), \text{data}(a, 2), \text{data}(b, 8) \} \end{array}$$

- Invalid — $\text{data}(a, 6)$ is omitted
- CHR^{cp} is not monotonic!

$\mathcal{P} \triangleright St \vdash_{\alpha} St'$ does not imply $\mathcal{P} \triangleright \{St, St''\} \vdash_{\alpha} \{St', St''\}$

Non-monotonicity

$\mathcal{P} \triangleright St \mapsto_{\alpha} St'$ does not imply $\mathcal{P} \triangleright \{St, St''\} \mapsto_{\alpha} \{St', St''\}$

- So what?
 - We cannot incrementally process constraints
 - At least not in the traditional manner ([Duck et al., 2004])
- Refined Operational Semantics

$$\mathcal{P} \triangleright \langle Gs ; St \rangle \mapsto_{\omega} \langle Gs' ; St' \rangle$$

- Gs : list of constraints to process (*goals*)
- St : constraint store

Coping with Non-monotonicity

- Let's try to process incrementally:

$$\text{pivotSwap} @ \begin{array}{l} \text{swap}(X, Y, P) \\ \textcolor{brown}{\lambda data(X, D) \mid D \geq P} \int_{D \in X_s} \\ \textcolor{blue}{\lambda data(Y, D) \mid D < P} \int_{D \in Y_s} \end{array} \iff \begin{array}{l} \textcolor{brown}{\lambda data(Y, D) \int_{D \in X_s}} \\ \textcolor{blue}{\lambda data(X, D) \int_{D \in Y_s}} \end{array}$$

\mapsto_ω $\langle [d(a, 1), d(a, 6), s(a, b, 5), d(a, 7), d(b, 2), d(b, 8)] ; \int \rangle$
 \mapsto_ω $\langle [\text{act } d(a, 1), d(a, 6), s(a, b, 5), d(a, 7), d(b, 2), d(b, 8)] ; \int d(a, 1) \int \rangle$
 \mapsto_ω $\langle [d(a, 6), s(a, b, 5), d(a, 7), d(b, 2), d(b, 8)] ; \int d(a, 1) \int \rangle$
 \mapsto_ω $\langle [\text{act } d(a, 6), s(a, b, 5), d(a, 7), d(b, 2), d(b, 8)] ; \int d(a, 1), d(a, 6) \int \rangle$
 \mapsto_ω $\langle [s(a, b, 5), d(a, 7), d(b, 2), d(b, 8)] ; \int d(a, 1), d(a, 6) \int \rangle$
 \mapsto_ω $\langle [\text{act } s(a, b, 5), d(a, 7), d(b, 2), d(b, 8)] ; \int d(a, 1), \textcolor{brown}{d(a, 6)}, \textcolor{green}{s(a, b, 5)} \int \rangle$
 \mapsto_ω $\langle [\textcolor{red}{d(b, 6)}, s(a, b, 5), d(a, 7), \boxed{d(b, 2)}, d(b, 8)] ; \int d(a, 1) \int \rangle$
 ...

* Abbreviating *swap* with *s* and *data* with *d*.

Coping with Non-monotonicity

- Let's try to process incrementally:

$$\text{pivotSwap} @ \begin{array}{l} \text{swap}(X, Y, P) \\ \textcolor{red}{\lambda \text{data}(X, D) \mid D \geq P} \int_{D \in X_s} \\ \textcolor{blue}{\lambda \text{data}(Y, D) \mid D < P} \int_{D \in Y_s} \end{array} \iff \begin{array}{l} \textcolor{red}{\lambda \text{data}(Y, D)} \int_{D \in X_s} \\ \textcolor{blue}{\lambda \text{data}(X, D)} \int_{D \in Y_s} \end{array}$$

$$\begin{array}{l} \langle [d(a, 1), d(a, 6), s(a, b, 5), d(a, 7), d(b, 2), d(b, 8)] ; \int \rangle \\ \mapsto_{\omega} \langle [\text{act } d(a, 1), d(a, 6), s(a, b, 5), d(a, 7), d(b, 2), d(b, 8)] ; \int d(a, 1) \rangle \\ \mapsto_{\omega} \langle [d(a, 6), s(a, b, 5), d(a, 7), d(b, 2), d(b, 8)] ; \int d(a, 1) \rangle \\ \mapsto_{\omega} \langle [\text{act } d(a, 6), s(a, b, 5), d(a, 7), d(b, 2), d(b, 8)] ; \int d(a, 1), d(a, 6) \rangle \\ \mapsto_{\omega} \langle [s(a, b, 5), d(a, 7), d(b, 2), d(b, 8)] ; \int d(a, 1), d(a, 6) \rangle \\ \mapsto_{\omega} \langle [\text{act } s(a, b, 5), d(a, 7), d(b, 2), d(b, 8)] ; \int d(a, 1), \textcolor{red}{d(a, 6)}, \textcolor{green}{s(a, b, 5)} \rangle \\ \mapsto_{\omega} \langle [\textcolor{red}{d(b, 6)}, s(a, b, 5), d(a, 7), \boxed{d(b, 2)}, d(b, 8)] ; \int d(a, 1) \rangle \\ \dots \end{array}$$

- Valid derivation w.r.t. operational semantics
($\int \textcolor{blue}{\text{data}(Y, D) \mid D < P} \int_{D \in Y_s}$ such that $Y_s = \emptyset$)
- But what about $\boxed{d(b, 2)}$? Its not yet visible in the store!

* Abbreviating *swap* with *s* and *data* with *d*.

Coping with Non-monotonicity

- Let's try to process incrementally:

$$\text{pivotSwap} @ \begin{array}{l} \text{swap}(X, Y, P) \\ \textcolor{brown}{\lambda \text{data}(X, D) \mid D \geq P} \textcolor{blue}{\lambda \text{data}(Y, D) \mid D < P} \end{array} \Leftrightarrow \begin{array}{l} \textcolor{brown}{\lambda \text{data}(Y, D)} \\ \textcolor{blue}{\lambda \text{data}(X, D)} \end{array}$$

\mapsto_{ω} $\langle [d(a, 1), d(a, 6), s(a, b, 5), d(a, 7), d(b, 2), d(b, 8)] ; \lambda \rangle$
 \mapsto_{ω} $\langle [\text{act } d(a, 1), d(a, 6), s(a, b, 5), d(a, 7), d(b, 2), d(b, 8)] ; \lambda d(a, 1) \rangle$
 \mapsto_{ω} $\langle [d(a, 6), s(a, b, 5), d(a, 7), d(b, 2), d(b, 8)] ; \lambda d(a, 1) \rangle$
 \mapsto_{ω} $\langle [\text{act } d(a, 6), s(a, b, 5), d(a, 7), d(b, 2), d(b, 8)] ; \lambda d(a, 1), d(a, 6) \rangle$
 \mapsto_{ω} $\langle [s(a, b, 5), d(a, 7), d(b, 2), d(b, 8)] ; \lambda d(a, 1), d(a, 6) \rangle$
 \mapsto_{ω} $\langle [\text{act } s(a, b, 5), d(a, 7), d(b, 2), d(b, 8)] ; \lambda d(a, 1), \textcolor{brown}{d(a, 6)}, \textcolor{green}{s(a, b, 5)} \rangle$
 \mapsto_{ω} $\langle [\textcolor{red}{d(b, 6)}, s(a, b, 5), d(a, 7), \boxed{d(b, 2)}, d(b, 8)] ; \lambda d(a, 1) \rangle$
 ...

- This is not sound w.r.t. abstract semantics!

$$\begin{array}{l} \lambda d(a, 1), \textcolor{brown}{d(a, 6)}, \textcolor{green}{s(a, b, 5)}, d(a, 7), \boxed{d(b, 2)}, d(b, 8) \\ \not\vdash_{\alpha} \lambda d(a, 1), \textcolor{red}{d(b, 6)}, d(a, 7), \boxed{d(b, 2)}, d(b, 8) \end{array}$$

- $\lambda \text{data}(Y, D) \mid D < P$ is non-maximal! $\boxed{d(b, 2)}$ was omitted!

Coping with Non-monotonicity

$\mathcal{P} \triangleright St \mapsto_{\alpha} St'$ does not imply $\mathcal{P} \triangleright \{St, St''\} \mapsto_{\alpha} \{St', St''\}$

Conditional monotonicity

If St'' contains only constraints *monotone w.r.t.* \mathcal{P}
 then $\mathcal{P} \triangleright St \mapsto_{\alpha} St'$ implies $\mathcal{P} \triangleright \{St, St''\} \mapsto_{\alpha} \{St', St''\}$

- A constraint $p(\vec{t})$ is *monotone w.r.t.* \mathcal{P} (denoted $\mathcal{P} \triangleq_{\text{unf}}^{\neg} p(\vec{t})$), iff $p(\vec{t})$ cannot be unified with any head comprehension patterns in \mathcal{P}

Monotone Constraints: $\mathcal{P} \triangleq_{\text{unf}}^{\neg} C$

- $g \triangleright C \sqsubseteq_{\text{unf}} M$: a constraint that matches C may match M under guard g :

$g \triangleright p(\vec{t}) \sqsubseteq_{\text{unf}} \{p'(\vec{t}') \mid g'\}$ iff $\theta p(\vec{t}) \equiv \theta p'(\vec{t}'), \models \theta(g' \wedge g)$ for some θ

$g'' \triangleright \{p(\vec{t}) \mid g\} \sqsubseteq_{\text{unf}} \{p'(\vec{t}') \mid g'\}$ iff $\theta p(\vec{t}) \equiv \theta p'(\vec{t}'), \models \theta(g'' \wedge g' \wedge g)$ for some θ

- $\mathcal{P} \triangleq_{\text{unf}}^{\neg} C$: for each rule $r @ \bar{H}_p \setminus \bar{H}_s \iff g \mid \bar{B} \in \mathcal{P}$, there is no comprehension pattern M in $\{\bar{H}_p, \bar{H}_s\}$ s.t. $g \triangleright C \sqsubseteq_{\text{unf}} M$

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Processing constraints incrementally in CHR^{cp}

- Operational semantics of CHR^{cp} extends [Duck et al., 2004] by
 - Identify (non-)monotone constraints (s.t., $\mathcal{P} \triangleq_{\text{unf}}^{\neg} C$)
 - Incrementally store *monotone* constraints
 - Eagerly store *non-monotone* constraints
- Hence we exploit conditional monotonicity where possible (the glass is half-full)

Processing constraints incrementally in CHR^{cp}

$$\begin{array}{l}
 \text{pivotSwap } @ \quad \textcolor{green}{swap}(X, Y, P) \\
 \quad \textcolor{brown}{\lambda data(X, D) \mid D \geq P} \int_{D \in X_s} \iff \textcolor{red}{\lambda data(Y, D)} \int_{D \in X_s} \\
 \quad \textcolor{blue}{\lambda data(Y, D) \mid D < P} \int_{D \in Y_s} \quad \textcolor{blue}{\lambda data(X, D)} \int_{D \in Y_s} \\
 \langle [\text{init } \lambda d(a, 6), s(a, b, 5), d(a, 7), d(b, 2)] ; \lambda \rangle
 \end{array}$$

\implies Initial store is contained in an “init” goal

Processing constraints incrementally in CHR^{cp}

$$\text{pivotSwap } @ \quad \textcolor{green}{swap}(X, Y, P) \quad \textcolor{brown}{\lambda data(X, D) \mid D \geq P} \int_{D \in X_s} \iff \textcolor{red}{\lambda data(Y, D)} \int_{D \in X_s} \\ \textcolor{blue}{\lambda data(Y, D) \mid D < P} \int_{D \in Y_s} \quad \textcolor{blue}{\lambda data(X, D)} \int_{D \in Y_s}$$

$$\begin{aligned} & \langle [\text{init } \lambda d(a, 6), s(a, b, 5), d(a, 7), d(b, 2)] ; \lambda \rangle \\ \mapsto_\omega & \langle [\text{eager } d(a, 6), \text{lazy } s(a, b, 5), \text{eager } d(a, 7), \text{eager } d(b, 2)] \\ & ; \lambda d(a, 6), d(a, 7), d(b, 2) \rangle \end{aligned}$$

\implies *Non-monotone constraints (e.g., $d(-, -)$) are eagerly stored*

Processing constraints incrementally in CHR^{cp}

$$\text{pivotSwap } @ \text{ } \overset{\text{swap}(X, Y, P)}{\underbrace{\underbrace{\lambda \text{data}(X, D) \mid D \geq P}_{D \in X_s} \iff \underbrace{\lambda \text{data}(Y, D)}_{D \in X_s}}_{\underbrace{\lambda \text{data}(Y, D) \mid D < P}_{D \in Y_s}} \underbrace{\lambda \text{data}(X, D)}_{D \in Y_s}}$$

$\langle [\text{init } \lambda d(a, 6), s(a, b, 5), d(a, 7), d(b, 2)] ; \lambda \rangle$
 $\mapsto_\omega \langle [\text{eager } d(a, 6), \text{lazy } s(a, b, 5), \text{eager } d(a, 7), \text{eager } d(b, 2)] ; \lambda d(a, 6), d(a, 7), d(b, 2) \rangle$
 $\mapsto_\omega \langle [\text{act } d(a, 6), \text{lazy } s(a, b, 5), \text{eager } d(a, 7), \text{eager } d(b, 2)] ; \lambda d(a, 6), d(a, 7), d(b, 2) \rangle$

\implies Eager goal is activated. Already added: No changes to store

Processing constraints incrementally in CHR^{cp}

$$\text{pivotSwap} @ \text{swap}(X, Y, P) \quad \begin{array}{l} \textcolor{brown}{\lambda \text{data}(X, D) \mid D \geq P} \int_{D \in X_s} \\ \textcolor{blue}{\lambda \text{data}(Y, D) \mid D < P} \int_{D \in Y_s} \end{array} \iff \begin{array}{l} \textcolor{brown}{\lambda \text{data}(Y, D)} \int_{D \in X_s} \\ \textcolor{blue}{\lambda \text{data}(X, D)} \int_{D \in Y_s} \end{array}$$

$$\begin{aligned} & \langle [\text{init } \lambda d(a, 6), s(a, b, 5), d(a, 7), d(b, 2)] ; \lambda \rangle \\ \mapsto_{\omega} & \langle [\text{eager } d(a, 6), \text{lazy } s(a, b, 5), \text{eager } d(a, 7), \text{eager } d(b, 2)] \\ & ; \lambda d(a, 6), d(a, 7), d(b, 2) \rangle \\ \mapsto_{\omega} & \langle [\text{act } d(a, 6), \text{lazy } s(a, b, 5), \text{eager } d(a, 7), \text{eager } d(b, 2)] \\ & ; \lambda d(a, 6), d(a, 7), d(b, 2) \rangle \\ \mapsto_{\omega} & \langle [\text{lazy } s(a, b, 5), \text{eager } d(a, 7), \text{eager } d(b, 2)] \\ & ; \lambda d(a, 6), d(a, 7), d(b, 2) \rangle \end{aligned}$$

\implies No match to *pivotSwap*, hence we drop *act* $d(a, 6)$

Processing constraints incrementally in CHR^{cp}

$$\text{pivotSwap} @ \text{ } \overset{\text{swap}(X, Y, P)}{\text{ } } \quad \begin{array}{l} \textcolor{brown}{\lambda \text{data}(X, D) \mid D \geq P} \int_{D \in X_s} \\ \textcolor{blue}{\lambda \text{data}(Y, D) \mid D < P} \int_{D \in Y_s} \end{array} \iff \begin{array}{l} \textcolor{brown}{\lambda \text{data}(Y, D)} \int_{D \in X_s} \\ \textcolor{blue}{\lambda \text{data}(X, D)} \int_{D \in Y_s} \end{array}$$

$$\begin{aligned} & \langle [\text{init } \lambda d(a, 6), s(a, b, 5), d(a, 7), d(b, 2)] ; \lambda \rangle \\ \mapsto_{\omega} & \langle [\text{eager } d(a, 6), \text{lazy } s(a, b, 5), \text{eager } d(a, 7), \text{eager } d(b, 2)] \\ & ; \lambda d(a, 6), d(a, 7), d(b, 2) \rangle \\ \mapsto_{\omega} & \langle [\text{act } d(a, 6), \text{lazy } s(a, b, 5), \text{eager } d(a, 7), \text{eager } d(b, 2)] \\ & ; \lambda d(a, 6), d(a, 7), d(b, 2) \rangle \\ \mapsto_{\omega} & \langle [\text{lazy } s(a, b, 5), \text{eager } d(a, 7), \text{eager } d(b, 2)] \\ & ; \lambda d(a, 6), d(a, 7), d(b, 2) \rangle \\ \mapsto_{\omega} & \langle [\text{act } s(a, b, 5), \text{eager } d(a, 7), \text{eager } d(b, 2)] \\ & ; \lambda \textcolor{green}{s(a, b, 5)}, \textcolor{brown}{d(a, 6)}, \textcolor{brown}{d(a, 7)}, \textcolor{blue}{d(b, 2)} \rangle \end{aligned}$$

\implies *Lazy goal is activated. We add $s(a, b, 5)$ to the store*

Processing constraints incrementally in CHR^{cp}

$$\text{pivotSwap} @ \begin{array}{l} \text{swap}(X, Y, P) \\ \textcolor{brown}{\lambda \text{data}(X, D) \mid D \geq P} \textcolor{brown}{\int_{D \in X_s}} \\ \textcolor{blue}{\lambda \text{data}(Y, D) \mid D < P} \textcolor{blue}{\int_{D \in Y_s}} \end{array} \iff \begin{array}{l} \textcolor{red}{\lambda \text{data}(Y, D) \int_{D \in X_s}} \\ \textcolor{blue}{\lambda \text{data}(X, D) \int_{D \in Y_s}} \end{array}$$

$$\begin{aligned} & \langle [\text{init } \lambda d(a, 6), s(a, b, 5), d(a, 7), d(b, 2)] ; \lambda \int \rangle \\ \mapsto_{\omega} & \langle [\text{eager } d(a, 6), \text{lazy } s(a, b, 5), \text{eager } d(a, 7), \text{eager } d(b, 2)] \\ & ; \lambda d(a, 6), d(a, 7), d(b, 2) \int \rangle \\ \mapsto_{\omega} & \langle [\text{act } d(a, 6), \text{lazy } s(a, b, 5), \text{eager } d(a, 7), \text{eager } d(b, 2)] \\ & ; \lambda d(a, 6), d(a, 7), d(b, 2) \int \rangle \\ \mapsto_{\omega} & \langle [\text{lazy } s(a, b, 5), \text{eager } d(a, 7), \text{eager } d(b, 2)] \\ & ; \lambda d(a, 6), d(a, 7), d(b, 2) \int \rangle \\ \mapsto_{\omega} & \langle [\text{act } s(a, b, 5), \text{eager } d(a, 7), \text{eager } d(b, 2)] \\ & ; \lambda \textcolor{green}{s(a, b, 5)}, \textcolor{brown}{d(a, 6)}, \textcolor{brown}{d(a, 7)}, \textcolor{blue}{d(b, 2)} \int \rangle \\ \mapsto_{\omega} & \langle [\text{init } \lambda \textcolor{red}{d(b, 6)}, \textcolor{red}{d(b, 7)}, \textcolor{blue}{d(b, 2)}], \text{eager } d(a, 7), \text{eager } d(b, 2)] ; \lambda \int \rangle \\ & \dots \end{aligned}$$

\implies Apply *pivotSwap*: rule body added as an “*init*” goal. Notice that $d(b, 2)$ is visible, hence maximality is guaranteed

Processing constraints incrementally in CHR^{cp}

$$\text{pivotSwap} @ \begin{array}{l} \text{swap}(X, Y, P) \\ \{ \text{data}(X, D) \mid D \geq P \}_{D \in X_s} \\ \{ \text{data}(Y, D) \mid D < P \}_{D \in Y_s} \end{array} \iff \begin{array}{l} \{ \text{data}(Y, D) \}_{D \in X_s} \\ \{ \text{data}(X, D) \}_{D \in Y_s} \end{array}$$

$$\mapsto_{\omega} \quad \dots \quad \langle [\text{init } \wr d(b, 6), d(b, 7), d(b, 2)], \text{eager } d(a, 7), \text{eager } d(b, 2)] ; \wr \rangle$$

$$\implies \text{Process "init" goal}$$

Processing constraints incrementally in CHR^{cp}

$$\text{pivotSwap} @ \begin{array}{l} \text{swap}(X, Y, P) \\ \textcolor{brown}{\lambda \text{data}(X, D) \mid D \geq P} \textcolor{brown}{\int_{D \in X_s}} \\ \textcolor{blue}{\lambda \text{data}(Y, D) \mid D < P} \textcolor{blue}{\int_{D \in Y_s}} \end{array} \iff \begin{array}{l} \textcolor{red}{\lambda \text{data}(Y, D) \int_{D \in X_s}} \\ \textcolor{blue}{\lambda \text{data}(X, D) \int_{D \in Y_s}} \end{array}$$

...

$$\begin{array}{l} \mapsto_{\omega} \langle [\text{init } \lambda d(b, 6), d(b, 7), d(b, 2) \int, \text{eager } d(a, 7), \text{eager } d(b, 2)] ; \int \rangle \\ \mapsto_{\omega} \langle [\text{eager } d(b, 6), \text{eager } d(b, 7), \text{eager } d(b, 2), \text{eager } d(a, 7), \text{eager } d(b, 2)] \\ ; \lambda d(b, 6), d(b, 7), d(b, 2) \int \rangle \end{array}$$

\implies All constraints in rule body are non-monotone

Processing constraints incrementally in CHR^{cp}

$$\text{pivotSwap} @ \begin{array}{l} \text{swap}(X, Y, P) \\ \textcolor{brown}{\lambda \text{data}(X, D) \mid D \geq P} \textcolor{brown}{\int_{D \in X_s}} \\ \textcolor{blue}{\lambda \text{data}(Y, D) \mid D < P} \textcolor{blue}{\int_{D \in Y_s}} \end{array} \iff \begin{array}{l} \textcolor{red}{\lambda \text{data}(Y, D) \int_{D \in X_s}} \\ \textcolor{blue}{\lambda \text{data}(X, D) \int_{D \in Y_s}} \end{array}$$

...

$$\begin{aligned} &\mapsto_{\omega} \langle [\text{init } \lambda d(b, 6), d(b, 7), d(b, 2) \int, \text{eager } d(a, 7), \text{eager } d(b, 2)] ; \lambda \int \rangle \\ &\mapsto_{\omega} \langle [\text{eager } d(b, 6), \text{eager } d(b, 7), \text{eager } d(b, 2), \text{eager } d(a, 7), \text{eager } d(b, 2)] \\ &\quad ; \lambda d(b, 6), d(b, 7), d(b, 2) \int \rangle \\ &\mapsto_{\omega}^* \langle [\text{eager } d(a, 7), \text{eager } d(b, 2)] ; \lambda d(b, 6), d(b, 7), d(b, 2) \int \rangle \end{aligned}$$

\implies No match for $d(b, 6)$, $d(b, 7)$ and $d(b, 2)$: we drop the goals

Processing constraints incrementally in CHR^{cp}

$$\text{pivotSwap} @ \begin{array}{l} \text{swap}(X, Y, P) \\ \textcolor{brown}{\lambda \text{data}(X, D) \mid D \geq P} \textcolor{brown}{\int_{D \in X_s}} \\ \textcolor{blue}{\lambda \text{data}(Y, D) \mid D < P} \textcolor{blue}{\int_{D \in Y_s}} \end{array} \iff \begin{array}{l} \textcolor{red}{\lambda \text{data}(Y, D) \int_{D \in X_s}} \\ \textcolor{blue}{\lambda \text{data}(X, D) \int_{D \in Y_s}} \end{array}$$

...

$$\begin{aligned} &\mapsto_{\omega} \langle [\text{init } \lambda d(b, 6), d(b, 7), d(b, 2) \int, \text{eager } d(a, 7), \text{eager } d(b, 2)] ; \lambda \int \rangle \\ &\mapsto_{\omega} \langle [\text{eager } d(b, 6), \text{eager } d(b, 7), \text{eager } d(b, 2), \text{eager } d(a, 7), \text{eager } d(b, 2)] \\ &\quad ; \lambda d(b, 6), d(b, 7), d(b, 2) \int \rangle \\ &\mapsto_{\omega}^* \langle [\text{eager } d(a, 7), \text{eager } d(b, 2)] ; \lambda d(b, 6), d(b, 7), d(b, 2) \int \rangle \\ &\mapsto_{\omega}^* \langle [] ; \lambda d(b, 6), d(b, 7), d(b, 2) \int \rangle \end{aligned}$$

$\implies d(a, 7)$ and $d(b, 2)$ no longer exists: we drop them as well

Refined Operational Semantics of CHR^{cp}

- Extends the refined operational semantics of CHR [Duck et al., 2004]
- Supports incremental processing of *monotone* constraints
- For each rule in \mathcal{P} , we identify which $p(\vec{t}) \in \bar{B}$ is monotone
 - Non-monotone constraints are made visible immediately
 - Monotone constraints can be incrementally stored as usual
- Sound w.r.t. abstract semantics of CHR^{cp}
- See paper for details!
 - See TR for even more details!!

Outline

- 1 Introduction
- 2 Comprehensions in CHR^{cp}
- 3 Monotonicity
- 4 Semantics
- 5 Status

Current Status of CHR^{cp}

- Prototype implementation
 - Available for download at:
<https://github.com/sllam/chrcp>
 - Compiler implemented in Python
 - Generates C++ codes that implements CHR^{cp} run-time
 - Utilizes Z3 SMT solver for type checking and monotonicity testing [Lam and Cervesato, 2014b] (i.e., $\mathcal{P} \triangleq_{\text{unf}}^{\neg} C$).
- See technical report [Lam and Cervesato, 2014a] for
 - more examples
 - preliminary experimental results
 - optimized compilation details (e.g., join ordering, etc)

Related Works

- Aggregates in CHR [Sneyers et al., 2007]
- CHR with Negation as Absence [Weert et al., 2006]
- Meld [Cruz et al., 2012]

Conclusion

- CHR^{cp} : Constraint Handling Rules with Multiset Comprehension Patterns
- Abstract semantics: \mapsto_{α}
- Operational semantics: \mapsto_{ω}
 - supports partial incremental processing of constraints
 - sound w.r.t. abstract semantics

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