

Graduate Course on Computer Security

Lecture 6: Case Study II - WEP

Iliano Cervesato iliano@itd.nrl.navy.mil

ITT Industries, Inc @ NRL - Washington DC

http://www.cs.stanford.edu/~iliano/



Outline

- The 802.11 wireless communication standard
- WEP: Wired Equivalent Privacy
 - > Architecture
 - > Security goals
 - > Attacks
 - Confidentiality
 - Authentication
 - Integrity
 - > Lessons Learned



Integrity





WEP Secrecy Access

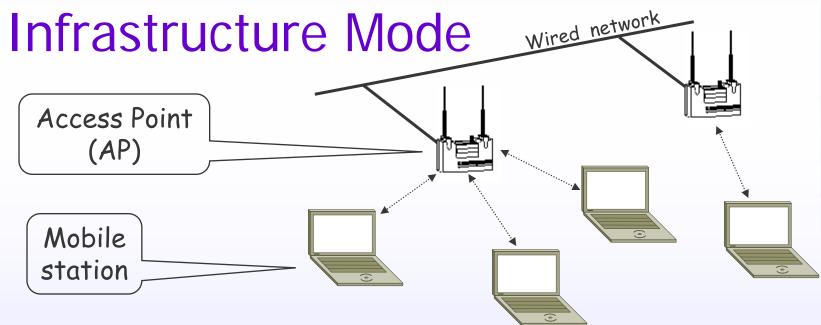
Integrity Lessons

The IEEE 802.11 Standard

Specifies standard networking functions over radio waves

- > Transparent layer for upper network protocols (IP, TCP, Novell NetWare, ...)
 - Implements wireless networks (WLAN)
 - Integrates seamlessly into a LAN
 - Works on any platform, given drivers
- > Fast: up to 11Mbit/s
 - Ethernet is 10Mbit/s, fast Ethernet 100Mbit/s
 - Range about 30m/100feet
- > Widely deployed
 - PCMCIA cards, ISA bus cards, embedded solutions, ...
 - Offered by major vendors





- 802.11
- WEP
 Secrecy
 Access
 Integrity

- Access points connect to wired network
- Multiple mobile stations per AP
 - > Full internet connection for mobile users
 - University campus
 - Coffee shops
 - airport lounges, ...



Ad Hoc Mode

- On the fly networking
 - Impromptu meeting
- LAN set up is difficult
 - Monitoring volcanoes
 - Study of jungle canopy
- LAN set up is dangerous
 - War zones

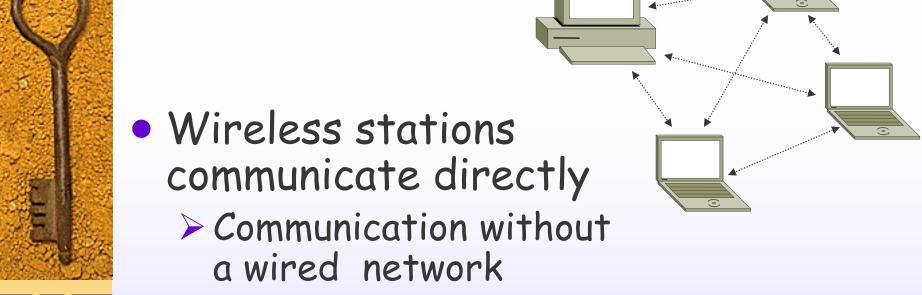


WEP

Secrecy

Access

Integrity





Data Transmission

For both LANs and WLANs

- Communication broken into frames
 - Variable length (up to ~ 1,500 byte)
- Header associated with frame
 - > Source address
 - > Destination address
 - > Frame length, ...
- Packet = header + frame

WEP Secrecy Access

Integrity Lessons



Subverting Communication

WLAN

- Eavesdropping
 - > Hardware widely sold
 - > Proximity of source
 - Parking lot attack
- Injecting traffic
 - > Just send to network
 - May need to modify driver setup
- Removing traffic
 - > Scramble radio signal

LAN

- Eavesdropping
 - > Plug in laptop
 - > Need access to wire
 - Hardly unnoticeable
- Injecting traffic
 - > Just send to network
 - May need to modify driver setup
- Removing traffic
 - > Feasible



WEP

Secrecy Access Integrity Lessons



WEP – Wired Equivalent Privacy

Security mechanism for WLANs

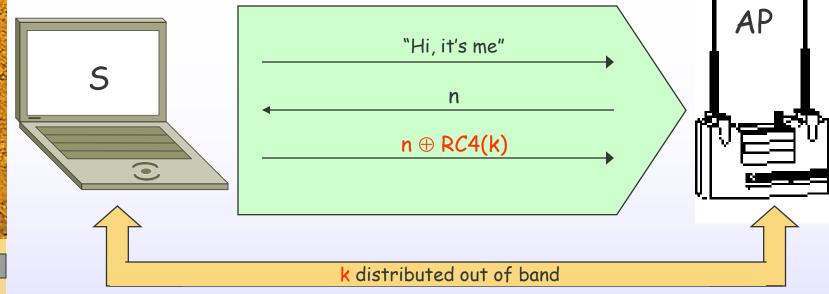
- 2 subsystems
 - > Station authentication
 - Simulate wired access control
 - > Data encapsulation
 - Create privacy of wired network
- Part of 802.11 standard

WEP

Secrecy Access Integrity

9

WEP Authentication



802.11

WEP

Secrecy Access Integrity

- S and AP share key k
 - > 802.11 standard: 40 bit
 - Most vendors now offer 104 bits (advertised as 128 bit!)
- n is randomly generated nonce
- S is accepted only if last message decrypts to n



Data Encapsulation

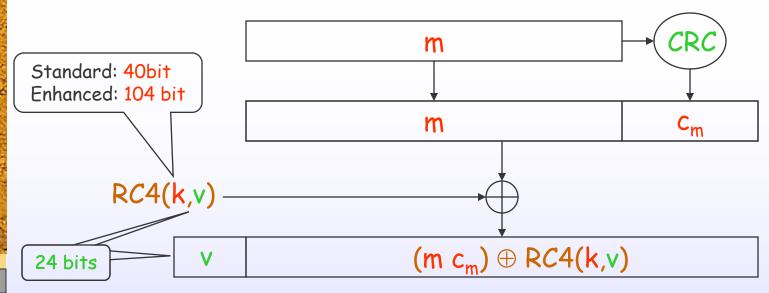
A wants to send frame m to B

- Encapsulation (A)
 - > Compute CRC-32 integrity checksum cm of m
 - Public algorithm, does not depend on k
 - Compute keystream RC4(k,v)
 - RC4 is secure keystream function (proprietary RSA)
 - v is 24 bit initialization vector (IV)
 - \triangleright Broadcast $v,x = v, ((m c_m) \oplus RC4(k,v))$
- Decapsulation (B)
 - $\rightarrow x \oplus RC4(k,v)) = m c_m$

WEP

Secrecy Access Integrity Lessons

... Pictorially



- Checksum guarantees data integrity
- IV
 - Prevents reuse of keystream
 - WEP does not prescribe modification of IVs
 - > Sent with each packet



WEP

Secrecy Access Integrity Lessons



WEP Security Goals

- Confidentiality
 - > Prevent eavesdropping
- Access control
 - > Prevent unauthorized access
- Integrity
 - > Prevent tempering with messages

802.11

WEP

Secrecy Access

Integrity

Lessons

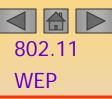
WEP does not achieve any of them!



Keystream Reuse

WEP collision

- Figure 1. If $x_1 = ((m_1 c_{m_1}) \oplus RC4(k,v))$ and $x_2 = ((m_2 c_{m_2}) \oplus RC4(k,v))$
- > Then $x_1 \oplus x_2 = (m_1 c_{m_1}) \oplus (m_2 c_{m_2})$
- Independent from key length!
- Recognizing collisions
 - k changes very seldom, if ever
 - Generally, all stations use same k
 - v sent in clear with every packet
 - > Look for packets with the same IV



Secrecy



Likelihood of Keystream Reuse

Given r_1 , ... $r_n \in [0, 1, ..., B]$ If $n \ge 1.2\sqrt{B}$, then $Prob[\exists i \ne j : r_i = r_j] > \frac{1}{2}$

- Ideal case
 - > By birthday paradox
 - 50% chances of collision after ~5000 packets
 - 4 minutes at 5Mbit/s (packets of 1500 bytes)
 - All 2^{24} keystreams recovered in $\frac{1}{2}$ day
- In practice, IVs are poorly generated
 - > Many PCMCIA cards
 - IV=0 when inserted
 - incremented by 1 at each packet
 - > Few thousand IVs determine most traffic
- 802.11 does not require changing IV



Secrecy



Attacks

```
If x_1 = ((m_1 c_{m_1}) \oplus RC4(k,v))
and x_2 = ((m_2 c_{m_2}) \oplus RC4(k,v))
then x_1 \oplus x_2 = (m_1 c_{m_1}) \oplus (m_2 c_{m_2})
```

- Passive attacks
 - > Exploit message redundancy
 - Many fields of IP header are predictable
 - Login sequences (e.g. Password:)
 - Transfer of shared libraries, ...
- Active attacks
 - > Send spam to mobile host
 - > Have mobile host send you email, ...
- Dumb attacks
 - > Some APs send frames unencrypted also

802.11 WEP

Secrecy

Access

Integrity Lessons

Computer Security: 6 - Case Study II, WEP



Decryption Dictionaries

- Once packet is revealed, keystream is known
- Build table of intercepted keystreams
 - \triangleright Maps every v to RC4(k,v))
 - > Requires ~24Gb for 2²⁴ for 1,500 byte frames
 - > Less than 1Gb with PCMCIA IV generation
- Then, can decrypt all traffic

Secrecy



Key Management

- 802.11 does not specify how to
 - > Generate
 - > Distribute
 - > Update shared key (and how often)
- In practice
 - > Key is loaded in device by hand when set up
 - Often keep manufacturer's default
 - Never updated again
 - > Attacker has years to compromise key
 - A few hours are enough for 40 bit version



Secrecy



Restoring Confidentiality

- IV is too short
 - > Collisions frequency reduced with longer IVs
 - > Relatively small decryption dictionary
- IV update unspecified (and non required)
 - > Force collision resistant IV generation
 - > From keyed random number generator
- Key management inexistent
 - > Introduce mandatory key update protocol
 - > Force different key for each host

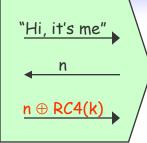


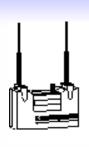
Secrecy



Gaining Access







Trivial!

- Record one authentication exchange
 - \succ from (n, n \oplus RC4(k)), recover RC4(k)
 - Use it to encrypt all future authentication challenges
- Remedy
 - > Use different cipher for authentication
 - A block cipher would do

802.11
WEP
Secrecy

Access

Integrity



Perturbing Traffic

Integrity protected by CRC-32 checksum

Checksums are linear w.r.t. ⊕

$$c_{m \oplus m'} = c_m \oplus c_{m'}$$

• Then for any Δ , xor'ing any ciphertext \times with (Δc_{Δ}) will go undetected

- Remedy
 - > ... exercise

Integrity



Targeted Traffic Alteration

- Linearity of CRC limited to flipping bits
- Use format of frames to force bit values
 - E.g. IP header
- Build decryption dictionary



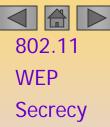
Integrity



Analysis of a Débacle

Why is WEP so bad??

- > International standard
- Backed by big vendors (IBM, 3COM, Apple, ...)
- Written by communication engineers
 - "Keep packet length small"
 - "Be conservative in what you send, liberal in what you accept"
 - Not security people involved
 - > Opaque design (no public review before standardization)
 - > Could have profited from IPSec experience
- Should operate with limited resource
 - > Cell phones, PDAs, ...



Access

Integrity



The Future of WEP

Proposal for a new standard 802.1X

- Use stream cipher based on AES
- Sequence number to avoid replays
- Replace CRC with MAC
- Authentication based on Kerberos





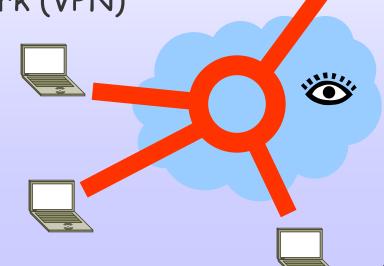
Should You Go Wireless?

YES!

• 802.11 is a fine communication suite

Handle security at higher levels

- Virtual Private Network (VPN)
- > IPSec
- > ... or just what you normally use!



802.11
WEP
Secrecy
Access
Integrity
Lessons



Readings

- N. Borisov, I. Goldberg and D. Wagner, *Intercepting Mobile Communications: the Insecurity of 802.11*, 2001
- W. Arbaugh, N. Shankar, and Y. Wan, Your 802.11
 Wireless Network has no Clothes, 2001
- IEEE 802.11 Working Group web page, http://grouper.ieee.org/groups/802/11
- Jesse Walker, "Overview of 802.11 Security", 2001



Exercises for Lecture 6

- Prove that
 - \rightarrow if $x = ((m c_m) \oplus RC4(k,v)),$
 - Then $x \oplus (\Delta c_{\Delta})$ has a correct checksum for every Δ
- Suggest a remedy for traffic perturbation

802.11
WEP
Secrecy
Access
Integrity
Lessons



Next ...

Specification Languages

