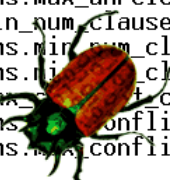


# Lecture 1: Model Checking

Edmund Clarke  
School of Computer Science  
Carnegie Mellon University

```
++CDatabase::_stats.mem_used_u
_params.max_unrelevance = (int
if (_params.max_unrelevance <
_params.max_unrelevance =
_params.min_num_clause_lits_fo
if (_params.min_num_clause_lit
_params.min_num_clause_lit
_params.max_num_clause_le
if (_params.min_num_conflict_claus
_params.max_num_conflict_claus
CHECK(
cout << "Forced to reduce unre
cout << "MaxUnrel: " << _params
<< " MinLenDel: " << _pa
<< " MaxLenCL : " << _pa
);
```





June 2002

“Software bugs, or errors, are so prevalent and so detrimental that they cost the U.S. economy an estimated \$59.5 billion annually, or about 0.6 percent of the gross domestic product

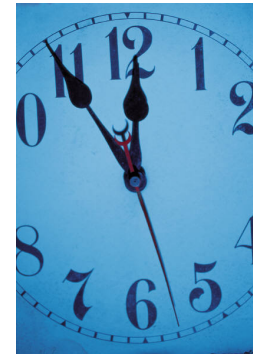
...

At the national level, over half of the costs are borne by software users and the remainder by software developers/vendors.”



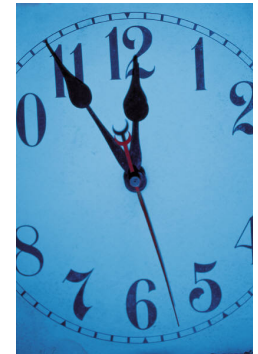
“The study also found that, although all errors cannot be removed, more than a third of these costs, or an estimated **\$22.2 billion, could be eliminated** by an improved testing infrastructure that enables earlier and more effective identification and removal of software defects.”

# Model Checking



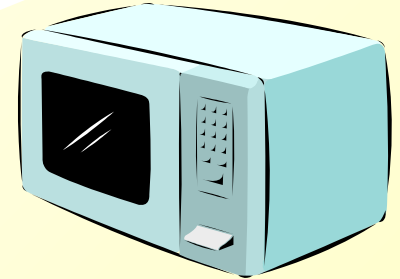
- Developed independently by **Clarke and Emerson** and by **Queille and Sifakis** in early 1980' s.
- **Properties** are written in **propositional temporal logic**.
- Systems are modeled by **finite state machines**.
- Verification procedure is an **exhaustive search of the state space** of the design.
- Model checking **complements** testing/simulation.

# Advantages of Model Checking



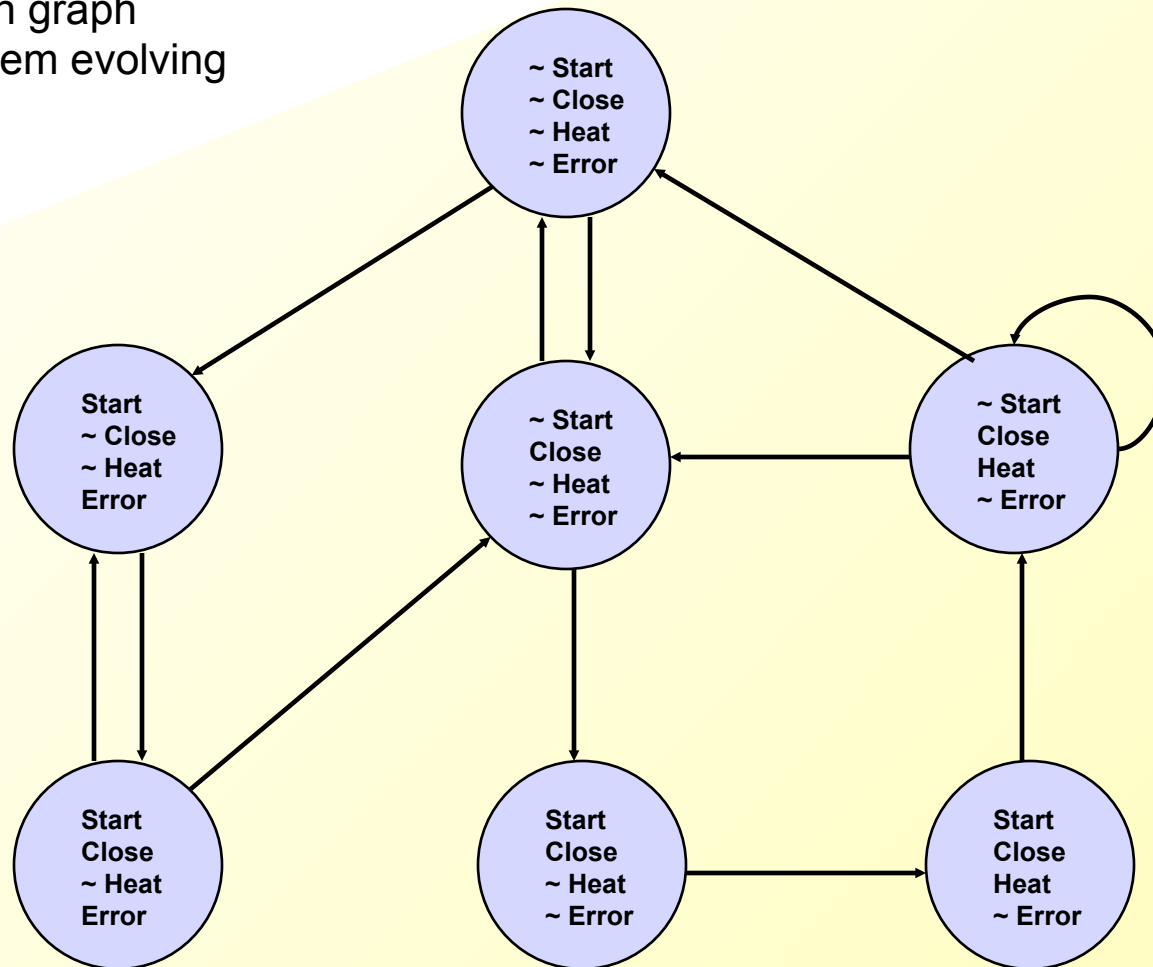
- **No proofs!!!**
- **Fast (compared to other rigorous methods)**
- **Diagnostic counterexamples**
- **No problem with partial specifications / properties**
- **Logics can easily express many concurrency properties**

# Model of computation



## Microwave Oven Example

State-transition graph describes system evolving over time.



# Temporal Logic



- The oven doesn't **heat up** until the **door is closed**.
- **Not heat\_up** holds **until door\_closed**
- $(\sim \text{heat\_up}) \text{ U } \text{door\_closed}$

# Basic Temporal Operators

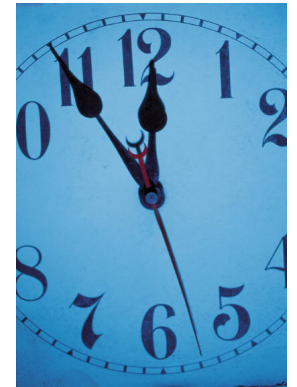


The symbol “**p**” is an atomic proposition, e.g. “**heat\_up**” or “**door\_closed**”.

- **F**p      - p holds sometime in the *future*.
- **G**p      - p holds *globally* in the future.
- **X**p      - p holds *next* time.
- p**U**q     - p holds *until* q holds.



# Model Checking Problem



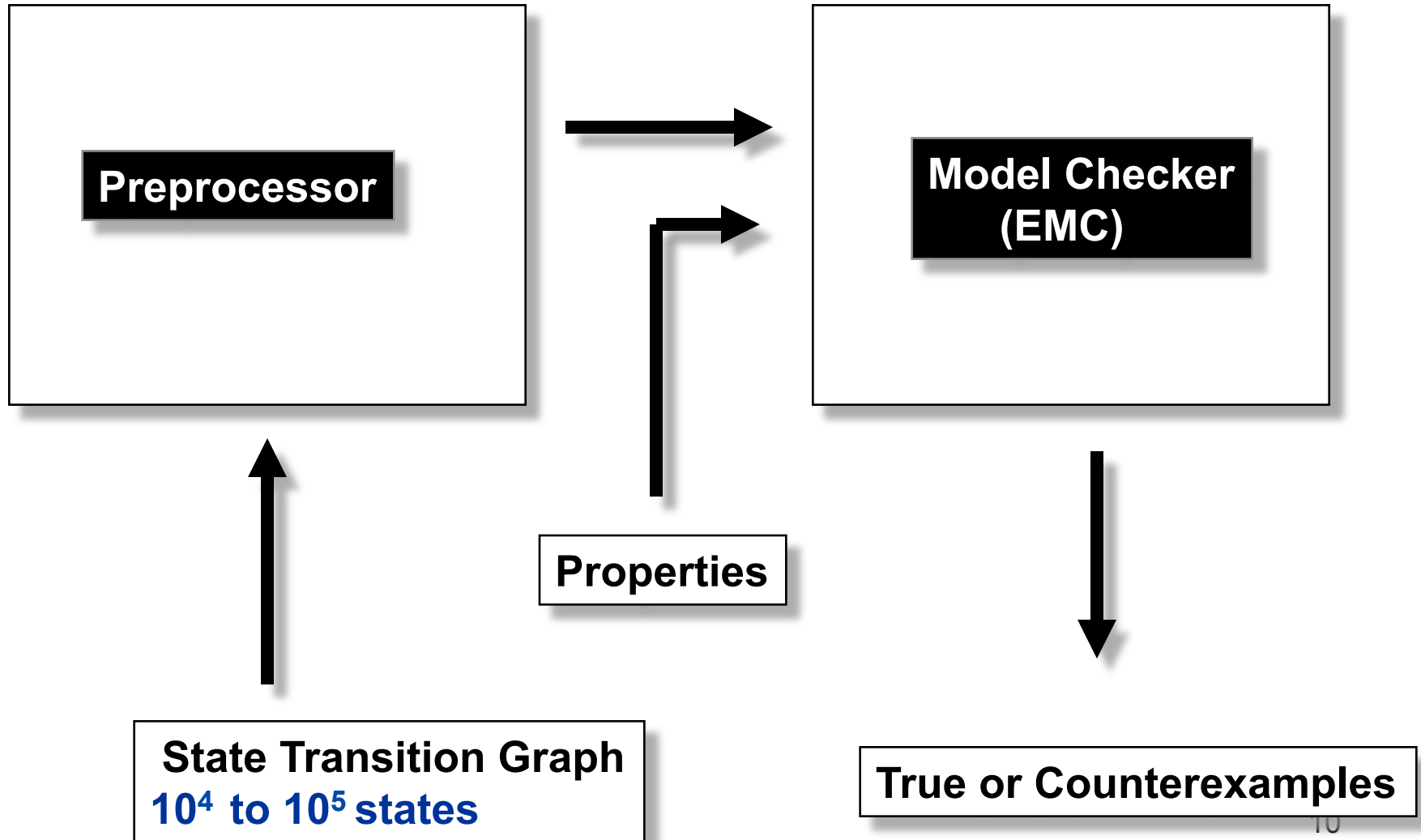
Let  $M$  be a model, i.e., a **state-transition graph**.

Let  $f$  be the **property** in temporal logic.

Find all states  $s$  such that  $M$  has property  $f$  at state  $s$ .

Efficient Algorithms: CE81, CES83

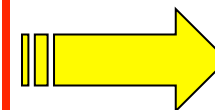
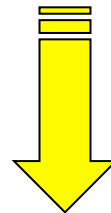
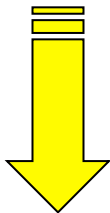
# The EMC System 1982/83



# Model Checker Architecture

System Description

Formal Specification

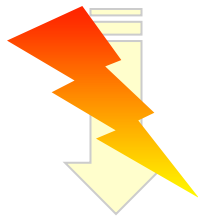


Validation  
or  
Counterexample

Model Checker

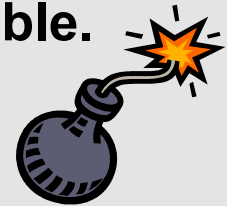
# The State Explosion Problem

System Description



State Transition Graph

**Combinatorial explosion** of system states renders explicit model construction infeasible.



## Exponential Growth of ...

- ... global state space in number of concurrent components.
- ... memory states in memory size.

**Feasibility of model checking inherently tied to handling state explosion.**

# Combating State Explosion



- **Binary Decision Diagrams** can be used to represent state transition systems more efficiently.  
→ **Symbolic Model Checking 1992**
- **Semantic techniques** for alleviating state explosion:
  - Partial Order Reduction.
  - Abstraction.
  - Compositional reasoning.
  - Symmetry.
  - Cone of influence reduction.
  - Semantic minimization.

# Model Checking since 1981



1981	Clarke / Emerson: CTL Model Checking Sifakis / Quielle	$10^5$
1982	EMC: Explicit Model Checker Clarke, Emerson, Sistla	
1990	Symbolic Model Checking Burch, Clarke, Dill, McMillan	$10^{100}$
1992	SMV: Symbolic Model Verifier McMillan	
1998	Bounded Model Checking using SAT Biere, Clarke, Zhu	$10^{1000}$
2000	Counterexample-guided Abstraction Refinement Clarke, Grumberg, Jha, Lu, Veith	

**1990s: Formal Hardware Verification in Industry: Intel, IBM, Motorola, etc.**

# Model Checking since 1981



1981 Clarke / Emerson: CTL Model Checking  
Sifakis / Quielle

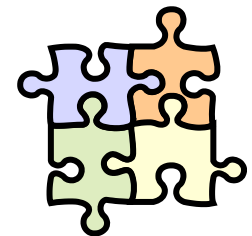
1982 EMC: Explicit Model Checker  
Clarke, Emerson, Sistla

1990 Symbolic Model Checking  
Burch, Clarke, Dill, McMillan

1992 SMV: Symbolic Model Verifier  
McMillan

1998 **Bounded Model Checking** using SAT  
Biere, Clarke, Zhu

2000 **Counterexample-guided Abstraction Refinement**  
Clarke, Grumberg, Jha, Lu, Veith



**CBMC**



**MAGIC**

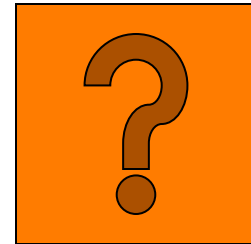
# Grand Challenge: **Model Check Software !**

What makes **Software Model Checking**  
different ?



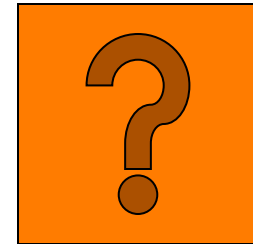


# What Makes Software Model Checking Different ?



- Large/unbounded base types: `int`, `float`, `string`
- User-defined types/classes
- Pointers/aliasing + unbounded #'s of heap-allocated cells
- Procedure calls/recursion/calls through pointers/dynamic method lookup/overloading
- Concurrency + unbounded #'s of threads

# What Makes Software Model Checking Different ?



- Templates/generics/include files
- Interrupts/exceptions/callbacks
- Use of secondary storage: files, databases
- Absent source code for: libraries, system calls, mobile code
- Esoteric features: continuations, self-modifying code
- Size (e.g., MS Word = 1.4 MLOC)

# Grand Challenge: Model Check Software !

**Early attempts in the 1980s failed to scale.**

**2000s: renewed interest / demand:**

**Java Pathfinder:** NASA Ames

**SLAM:** Microsoft

**Bandera:** Kansas State

**BLAST:** Berkeley

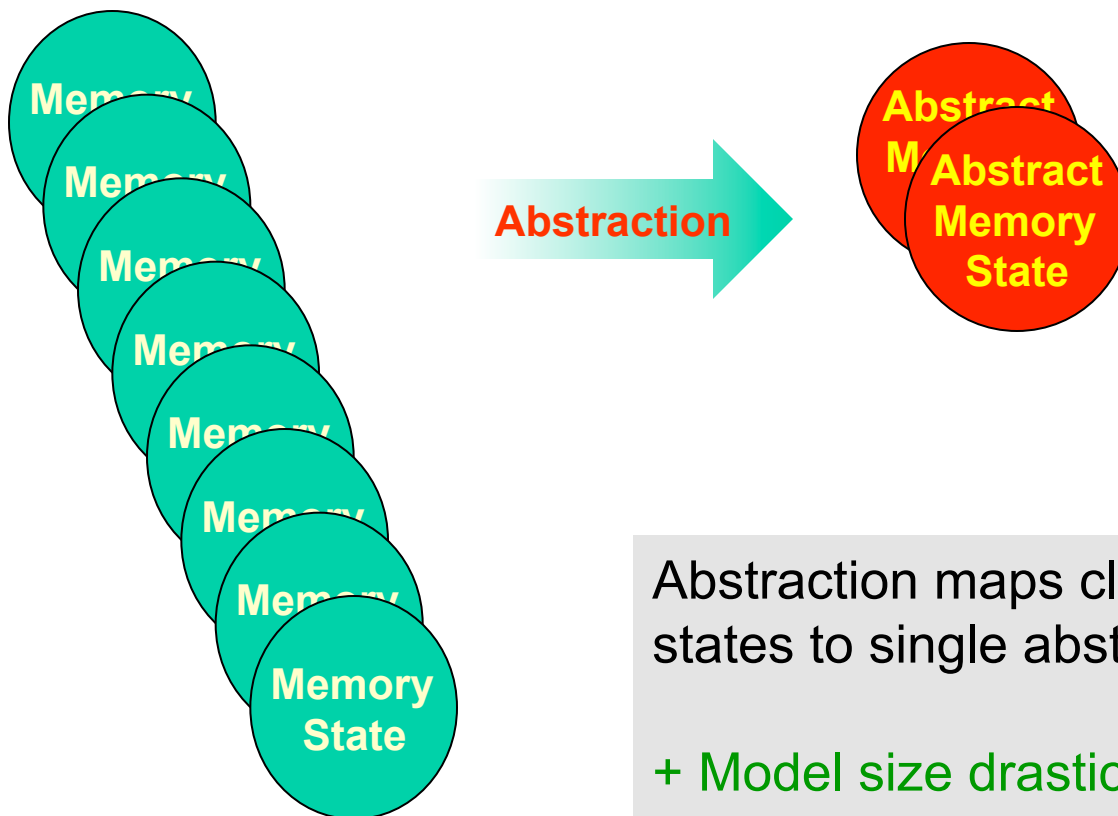
...

*SLAM to be shipped to Windows device driver developers.*

In general, these tools are unable to handle **complex data structures** and **concurrency**.

# The MAGIC Tool:

## Counterexample-Guided Abstraction Refinement



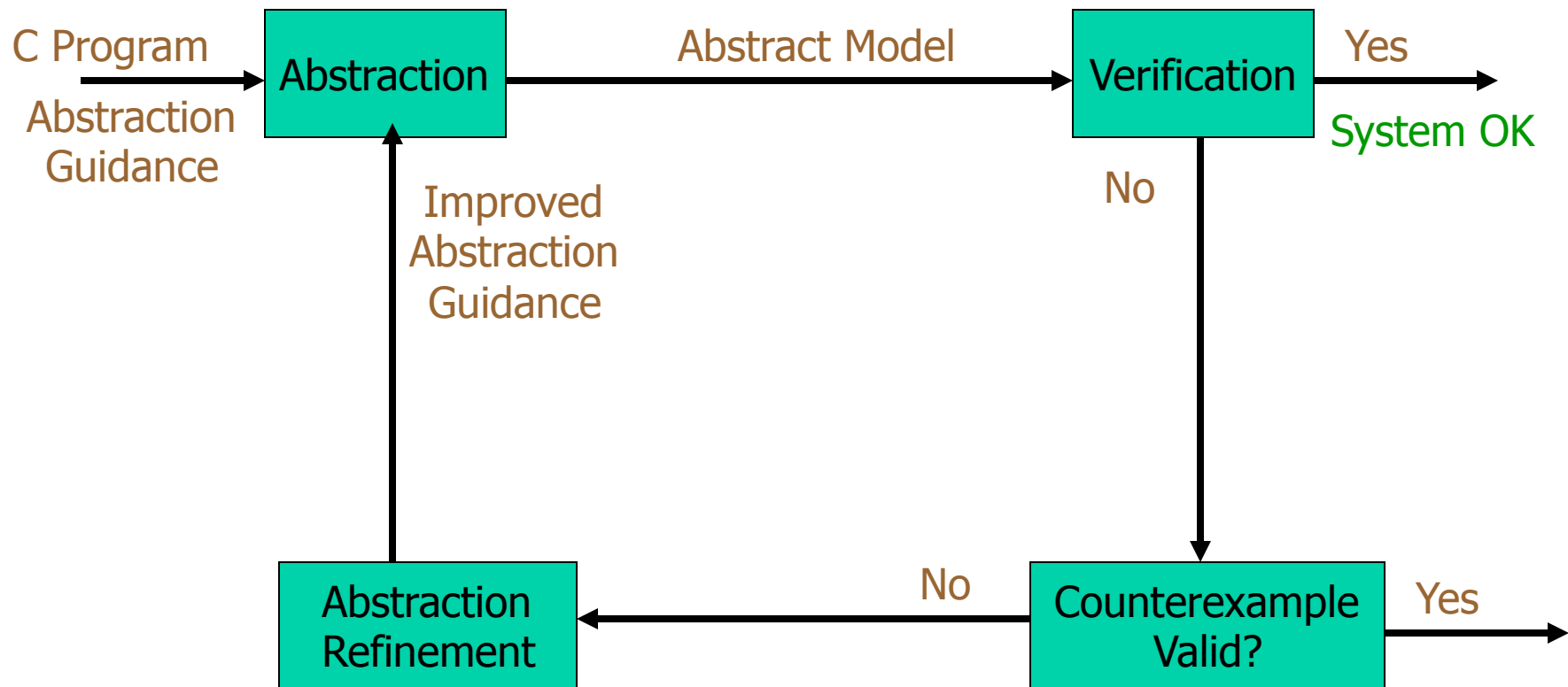
Abstraction maps classes of similar memory states to single abstract memory states.

+ Model size drastically reduced.

- Invalid counterexamples possible.

# The MAGIC Tool:

## Counterexample-Guided Abstraction Refinement



# CBMC: Embedded Systems Verification

The screenshot shows the CBMC-GUI interface. The main window displays a C program named 'while.c'. The code defines three integer arrays: 'table0' with values {0xf324, 0}, 'table1' with values {0xec26, 0x626e, 0}, and 'table2' with value {0}. A pointer array 'tables' is defined with elements 'table0', 'table1', and 'table2'. The 'main' function takes an 'unsigned index' and a pointer '\*p'. It checks if 'index > 2', sets 'index=0' if true, and then assigns 'p=tables[index]'. A red highlight is under the line 'while(\*p!=0) p++;'. Below the code editor is a 'Watch' window with tabs for 'Output', 'Errors', 'Watch', and 'Debug'. The 'Watch' window shows a table of variables and their values:

Name	Value
index	1 (0000000000)
c::tables	{ c::table0,
c::table2	{ 0 }
c::table1	{ 60454, 2
c::table0	{ 62244,
p	c::table

The status bar at the bottom indicates 'while.1.log'.

- Method: **Bounded Model Checking**
- Implemented **GUI** to facilitate tech transfer
- Applications:
  - Part of train controller from GE
  - Cryptographic algorithms (DES, AES, SHS)
  - C Models of ASICs provided by nVidia

# Case Study:

## Verification of MicroC/OS

- **Real-Time Operating System**
  - About 6000 lines of C code
  - Used in commercial *embedded systems*
    - UPS, Controllers, Cell-phones, ATMs
- Required **mutual exclusion** in the kernel
  - `OS_ENTER_CRITICAL()` and `OS_EXIT_CRITICAL()`
- **MAGIC** and **CBMC**:
  - Discovered **one unknown bug** related to the locking discipline
  - Discovered three more bugs
  - Verified that **no similar bugs are present**

