End of Semester Logistics ...

Final exam

- Monday, December 12th,
- 8:30-11:30, Hamerschlag B131
 - Cummulative, emphasis on 2nd half
 - Closed book, 2 pages of notes
- Study guide: bottom of syllabus page
- Review session:
 - Sunday, Dec 11th
 - 2pm 4pm
 - BH 255A (here)

End of Semester Logistics ...

Homework

- Solutions 1-6 online
- Solution 7 online tomorrow.
- Problem Set 8 due at midnight on Friday
- 711-6 canceled
- Late homework recieves a zero score once the solution sets have been posted.
- In calculating your final score, your lowest homework score will be dropped provided that all assignments have been submitted by the last day of classes.
- REMEMBER to note who you worked with!

End of Semester Logistics ...

Problem set 8

- Run 5 Blast searches with different parameter settings
- Record some results in Tables 1 & 2 (excel worksheet)
- Interpret in terms of Blast heuristics and Karlin Altschul stats

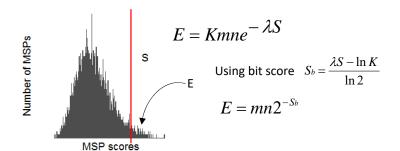
Recommendations:

- Run all five searches in one session
- Record results immediately
- Interpret results at your leisure

PLEASE PLEASE FILL OUT FACULTY COURSE EVALUATIONS

- How much information is available to distinguish between chance MSPs and MSPs in related sequences?
 - Information content of substitution matrices
 - Information content of alignments
- Which substitution matrix will maximize precision and recall?

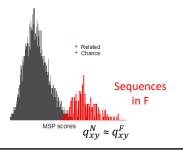
Recall: BLAST (Karlin-Altschul) Statistics

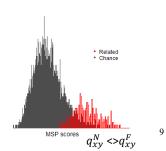


E = number of MSPs with scores > S.

<u>Maximal Segment Pair (MSP)</u>: an ungapped local alignment that cannot be improved by making it bigger or smaller.

- Each scoring matrix $S^N[x,y]$ has characteristic frequencies, q_{xy}^N
- Each protein family F has characteristic amino acid pair frequencies, q_{xy}^F
- The best results are obtained when q_{xy}^N is a good match for q_{xy}^F
- Given a query Q and a database sequence D_j both from family F, the alignment of Q and D_j will result in the highest MSP score when q_{xy}^N is a good match for q_{xy}^F
- In a database search with query Q, the best descrimination between chance MSPs and MSPs in alignments from related sequences in family F will be obtained when q^N_{xy} is a good match for q^F_{xy}.





The average score (in bits) per alignment position when using a PAM *Y* matrix to compare sequences in fact separated by n PAMs

(Calculated by simulation)

PAM matrix		Actual PAM distance n							
		40	80	120	160	200	240	280	320
	40	2.26	1:31	0.62	0.10	-0.30	-0.61	-0.86	-1.06
	80	2.14	1.44	0.92	0.53	0.23	-0.02	-0.21	-0.37
Υ	120	1.93	1.39	0.98	0.67	0.42	0.22	0.06	-0.07
•	160	1.71	1.28	0.95	0.70	0.50	0.33	0.20	0.09
	200	1.51	1.16	0.90	0.68	0.51	0.38	0.26	0.17
	240	1.32	1.05	0.82	0.65	0.51	0.39	0.29	0.21
	280	1.17	0.94	0.75	0.60	0.48	0.38	0.30	0.23
	320	1.03	0.84	0.68	0.56	0.46	0.37	0.30	0.24

Maxima highlighted in yellow

Best discrimination between related and chance MSPs :

Matrix divergence ~ Family divergence

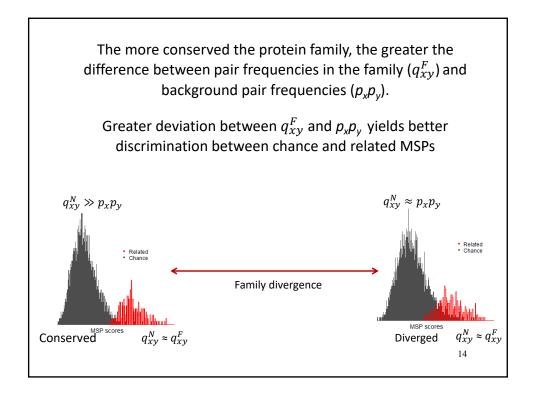
10

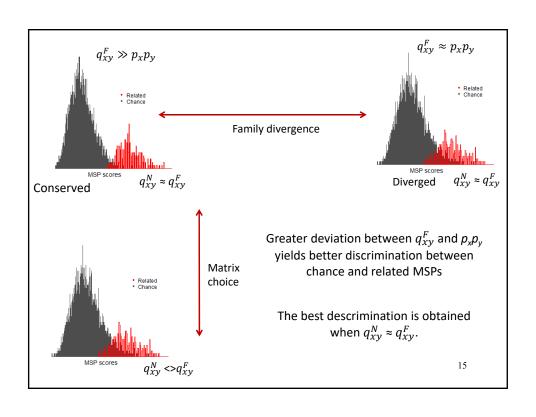
Altschul SF, J. Mol. Biol., 219, 555-565 (1991)

The averag <i>matr</i>	e score (ix to con	•			•		•	
		(Ca	lculated b	y simula	ation)			
DANA	Actual PAM distance n							
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						□ = Ef	ficiency ≥	94%
r#:-:		Score with PAM Y						
Effici	ency = –	Score with PAM n						
								11
				А	ltschul SE	J. Mol. Biol.	. 219. 555-	565 (1991

Choosing your scoring matrix

- BLAST will give reasonable accuracy as long as the empirical target frequencies do not deviate too far from the theoretical target frequencies
 - Use PAM40, BLOSUM62 & BLOSUM45, or BLOSUM62 & BLOSUM45





A warm-up thought experiment: How much information is available in a sequence of coin tosses to determine if the coin is fair or biased?

Alternate Hypothesis (H_A): Coin is biased

- $pr(H|H_A) = q$, $pr(T|H_A) = (1-q)$, where $q \neq 0.5$

Null Hypothesis (H_0) : : Coin is fair

 $- p(H|H_0) = p$, $p(T|H_0) = (1-p)$, where p = 0.5

- If q >>0.5 (e.g., q = 0.8), then a short series of coin tosses is sufficient to convince us that H_A is true.
- If $q \approx 0.5$ (e.g., q = 0.5001), then we require a much longer series of coin tosses is sufficient to convince us that $p(H) \neq 0.5$.

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Relative Entropy

Given two probability distributions, P and Q, defined on the same event space, $X=\{E_1, E_2, \dots E_N\}$

- $P = pr(X | \hat{H}_0) = \{p_1, p_2, ... p_N\}$
- $Q = pr(X | \hat{H}_A) = \{q_1, q_2, ... q_N\}$

the relative entropy or Kullback-Leibler Divergence

$$\mathcal{H} = \sum_{\mathbf{x}} q_i \log_2 \frac{q_i}{p_i}$$

is the expected information provided by each observation to discriminate in favor of hypothesis \hat{H}_a against hypothesis $\hat{H}_{0,}$ when \hat{H}_a is true.

Note: the KL Divergence is not symmetric and therefore not a distance.

Relative Entropy – coin toss example

Given two probability distributions, P and Q, defined on the same event space, $X=\{H,T\}$

- $P = pr(X | \hat{H}_0) = \{0.5, 0.5\}$
- $Q = pr(X | \hat{H}_{\Delta}) = \{q, (1-q)\}\)$, where $q \neq 0.5$

the relative entropy

$$\begin{split} \sum_{\{H.T\}} q_i \log_2 \frac{q_i}{p_i} \\ q \log_2 \frac{q}{p} + (1-q) \log_2 \frac{(1-q)}{p} \end{split}$$

is the expected information available per coss of the coin to discriminate in favor of hypothesis \hat{H}_a (biased coin) against hypothesis \hat{H}_{0} , (fair coin) if the coin is actually biased.

A warm-up thought experiment: How much information is available in a sequence of coin tosses to determine if the coin is fair or biased?

Alternate Hypothesis (H_A): Coin is biased

-
$$pr(H|H_{\Delta}) = q$$
, $pr(T|H_{\Delta}) = (1-q)$, where $q \neq 0.5$

Null Hypothesis (H_0) : : Coin is fair

$$- p(H|H_0) = p$$
, $p(T|H_0) = (1-p)$, where $p = 0.5$

- If q >>0.5 (e.g., q = 0.8), then a short series of coin tosses is sufficient to convince us that H_A is true.
- If $q \approx 0.5$ (e.g., q = 0.5001), then we require a much longer series of coin tosses is sufficient to convince us that $p(H) \neq 0.5$.

Relative Entropy – ungapped local alignments

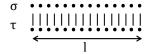
Alternate Hypothesis (H_A): σ and τ are related at N PAMs divergence.

- Amino acids x and y are aligned with frequency, q_{xy}^{N}
- Each alignment column is an observation, similar to a coin toss

Null Hypothesis (H_0): σ and τ are unrelated

- Amino acids x and y are aligned with background frequencies, $p_x p_y$
- If q_{xy}^N very different from $p_x p_y$, then only a few observations are sufficient to convince us that H_A is true (i.e., short alignment).
- If the family is diverged, $(q_{xy}^N \text{ more similar to } p_x p_y)$, then we require more observations (i.e., a longer alignment).

Relative Entropy – ungapped local alignments



Alternate Hypothesis (H_A): σ and τ are related at N PAMs divergence.

- Amino acids x and y are aligned with frequency, q_{xy}^{N}
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Null Hypothesis (H_0): σ and τ are unrelated

- Amino acids x and y are aligned with background frequencies, $p_x p_y$

Relative entropy
$$\mathcal{H}^N = \sum_{\{xy\}} q^N_{xy} \log_2 \frac{q^N_{xy}}{p_x p_y} = \sum_{\{xy\}} q^N_{xy} S^N[x,y]$$

gives the number of bits per position available to distinguish chance MSPs from MSPs in related sequences with N PAMs of divergence.

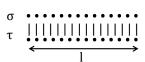
The *average* relative entropy of a substitution matrix is given in bits per position and can be calculated from S^N using the equation

$$\mathcal{H}^N = \sum_{\{xy\}} q_{xy}^N S^N[x, y]$$

BLO	SUM	F	PAM	Sequence			
	bits/site	bits/site		identity			
		20	2.95	83%			
		30	2.57				
		60	2.00	63%			
		70	1.60				
90	1.18	100	1.18	43%			
80	0.99	120	0.98	38%			
60	0.66	160	0.70	30%			
50	0.52	200	0.51	25%			
45	0.38	250	0.36	20%			

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An ungapped alignment of length I between two sequences separated N PAMs divergence contains $l\mathcal{H}^N$ bits of discriminatory information, on average.



How many bits of information are needed to find a related match in a database search? $E=m'n'2^{-\delta}$

$$S = \log_2 \frac{m'n'}{E}$$

Suppose we seek matches with E values no greater than E = 1. Then, we require $S \ge \log_2 m'n'$ bits. From this, we can estimate the minimum alignment length required to disinguish related from chance MSPs at N PAMS:

$$l\mathcal{H}^{N} = \log_{2} m'n'$$
$$l = \frac{\log_{2} m'n'}{\mathcal{H}^{N}}$$

Implications

The lower the relative entropy, \mathcal{H}^N , the longer the minimum alignment that is distinguishable from chance. $l=\frac{\log_2 m'n'}{\mathcal{H}^N}$

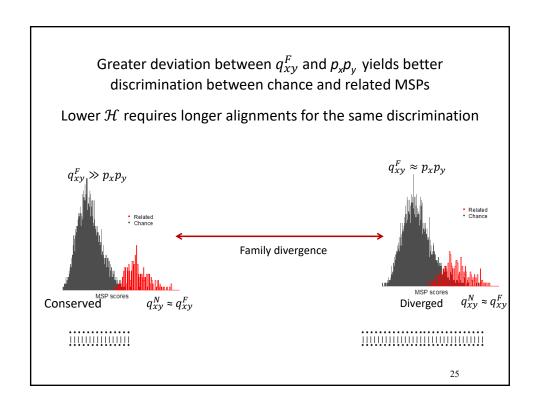
In a data base of length n = 50 billion, $\log_2 m'n' = 44$ bits are required. Since the alignment cannot be longer than the query, a query sequence must be at least $\frac{14}{2}$ 57 = 47, residues length $\frac{1}{2}$ 20 2006.

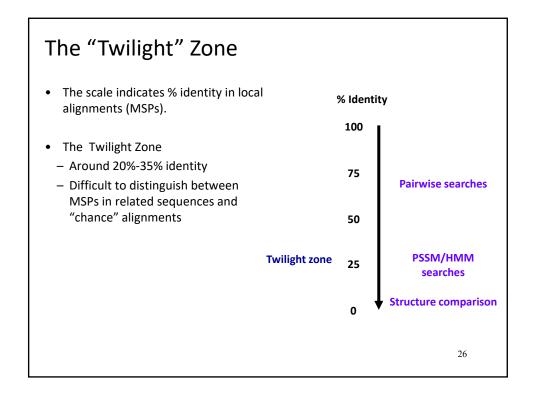
44/2.57 = 17 residues long at 30 PAMs 44/0.70 = 62 residues long at 160 PAMs 44/0.36 = 121 residues long at 250 PAMs

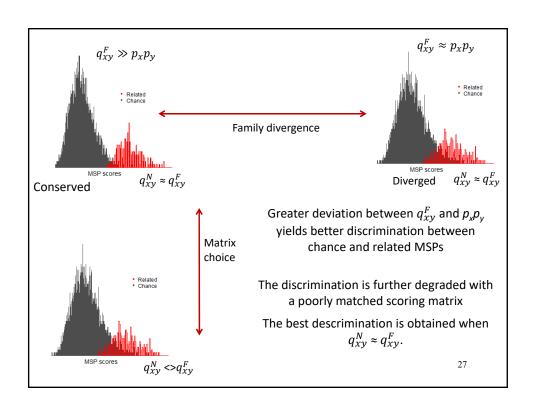
to distinguish significant HSP's from chance.

P.	AM	Seq Id		
30	2.57			
100	1.18	43 %		
120	0.98	38%		
160	0.70	30 %		
200	0.51	25%		
250	0.36	20 %		

Note that \mathcal{H}^N is an average over scoring matrix S^N . A shorter alignment may encode enough information if it contains many high-scoring pairs; alternatively, you may need a longer alignment if there are many low-scoring pairs.







Choosing your scoring matrix

- BLAST will give reasonable accuracy as long as the empirical target frequencies do not deviate too far from the theoretical target frequencies
 - Use PAM40, BLOSUM62 & BLOSUM45, or BLOSUM62 & BLOSUM45
- 2. The lower the relative entropy, *H*, the longer the minimum alignment that is distinguishable from chance.
- 3. If your query is short, you will only be able to find closely related matches.
 - Use PAM30

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DATABASE SEARCHING RECAP

Searching a sequence database

Input:

- query Q of length m
- database D=D1 D2 D3... D_N of length n

Search:

for j = 1 to N

- Find best local alignment of Q with Dj
- If "good alignment", add Dj to Results

Output: Results

PROBLEMS

- Too slow
- What is a "good" alignment?
- Which matrix should you use?
- · Which results are trustworthy?
- Can you find all related sequences in the database?

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Basic Local Alignment Search Tool

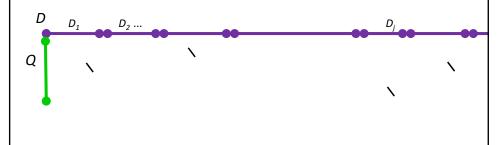
Altschul et al, 90

Construct hash table L

- Find all strings of length w that align with a w-mer in Q with score ≥ T Scan database D for hits - instances of words in L

Extend hits to find MSPs

If D_i contains MSP with score $S > S_T$ report D_i



Not all w-mers in Q will be included in L

Some w-mers not in W will be included in L

Suppose that w = 4 and T = 17 and word scores are calculated with BLOSUM 62

L L V L

LLV L

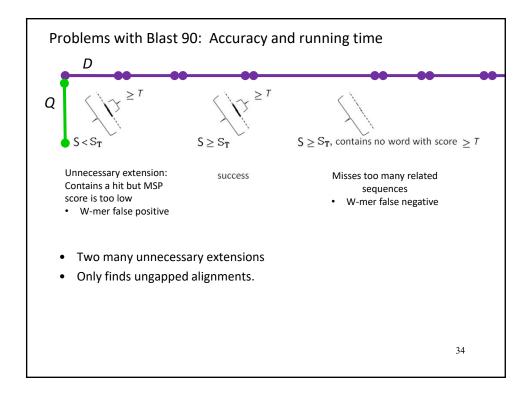
LLVL will not be included in L, because when aligned with itself the word score is lower than T

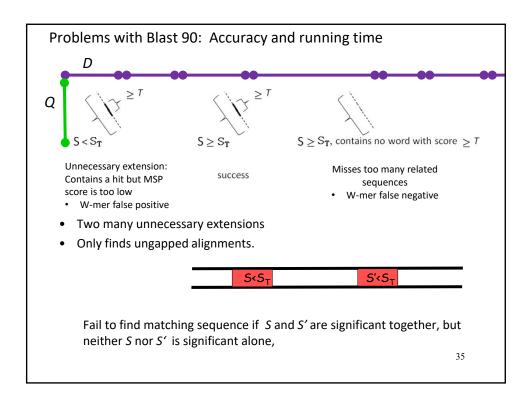
W D Y E

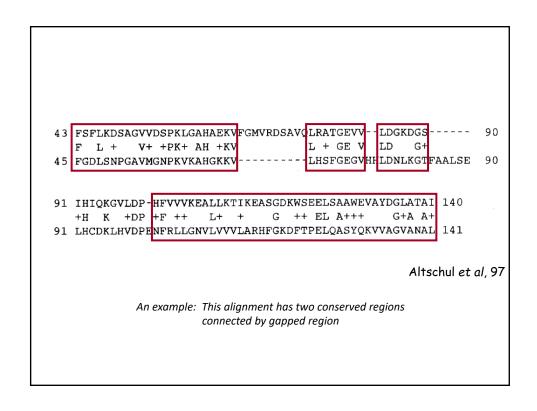
W E F E

11 2 3 5 21 > T

WEFE will be included in L because it aligns with a word in Q (WDYE) with a score greater than T

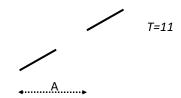






Two-Hit BLAST

- Reduce threshold T to obtain more hits
- Only trigger an ungapped extension if there are two hits on the same diagonal within distance A



- Misses fewer significant MSPs
- Fewer unnecessary extensions

Gapped, 2-hit Blast

Altschul et al, 97

- 1. Find hits of length w with similarity threshold T.
- 2. If D_J has

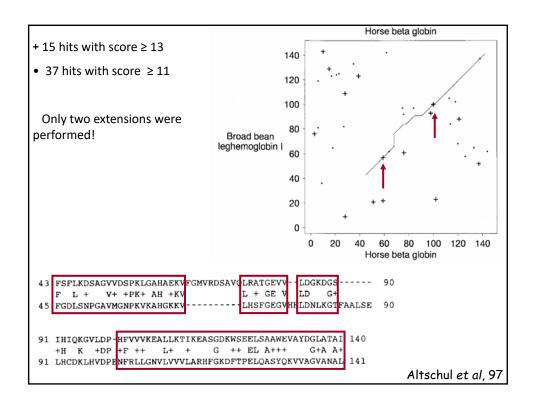
two hits

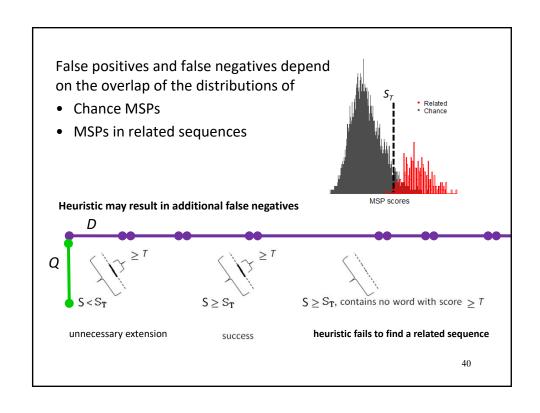
on same diagonal

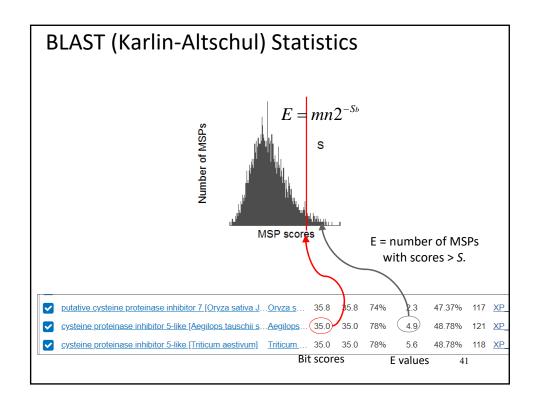
separated by a distance of at most A,

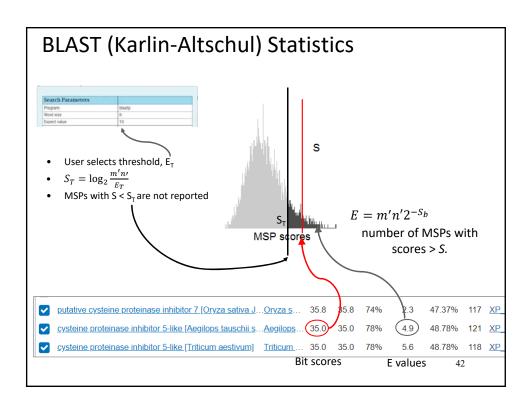
perform an ungapped extension to obtain MSP

- 3. If MSP score $S_1 > S_{g_i}$ perform a gapped extension with dynamic programming
- 4. If gapped extension score $S_2 > S_T$, report D_1 as a match.

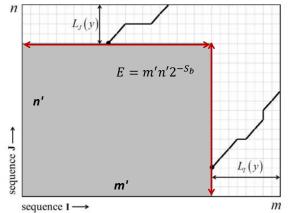








The distances m' and n' include an "edge correction"



Note: the NCBI papers show alignment matrices starting in the lower LH corner, not the upper LH corner as we use in this class.

"On average", an alignment must start within the gray box to accrue a score of at least S_{7} before reaching the end of the sequence

New finite-size correction for local alignment score distributions. Park, Sheetin, Ma, Madden, Spouge* *BMC Research Notes* 2012, 5:286 doi:10.1186/1756-0500-5-286