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15-826: Multimedia Databases and Data Mining

Lecture #26: Graph mining - patterns

Christos Faloutsos

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Must-read Material

- [Graph mining textbook] Deepayan Chakrabarti and Christos Faloutsos [*Graph Mining: Laws, Tools and Case Studies*](#), Morgan Claypool, 2012
 - Part I (patterns)

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Must-read Material

- Michalis Faloutsos, Petros Faloutsos and Christos Faloutsos, On Power-Law Relationships of the Internet Topology, SIGCOMM 1999.
- R. Albert, H. Jeong, and A.-L. Barabasi, Diameter of the World Wide Web Nature, 401, 130-131 (1999).
- Reka Albert and Albert-Laszlo Barabasi Statistical mechanics of complex networks, Reviews of Modern Physics, 74, 47 (2002).
- Jure Leskovec, Jon Kleinberg, Christos Faloutsos Graphs over Time: Densification Laws, Shrinking Diameters and Possible Explanations, KDD 2005, Chicago, IL, USA

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Must-read Material (cont'd)

- D. Chakrabarti and C. Faloutsos, Graph Mining: Laws, Generators and Algorithms, in ACM Computing Surveys, 38 (1), 2006

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Main outline

- Introduction
- Indexing
- Mining
 - Graphs – patterns
 - Graphs – generators and tools
 - Association rules
 - ...

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Outline


- ➔ • Introduction – Motivation
- Problem#1: Patterns in graphs
- Problem#2: Scalability
- Conclusions

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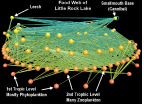


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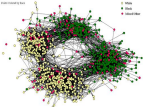
Graphs - why should we care?



Friendship Network
[Moody '01]



Food Web
[Martinez '91]



Internet Map
[lumeta.com]


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Graphs - why should we care?

- IR: bi-partite graphs (doc-terms)

D_1
 \dots
 D_N



T_1
 \dots
 T_M

- web: hyper-text graph
- ... and more:

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Graphs - why should we care?

- 'viral' marketing
- web-log ('blog') news propagation
- computer network security: email/IP traffic and anomaly detection
-

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Outline

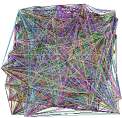


- Introduction – Motivation
- ➡ • Problem#1: Patterns in graphs
 - Static graphs
 - Weighted graphs
 - Time evolving graphs
- Problem#2: Scalability
- Conclusions

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Problem #1 - network and graph mining

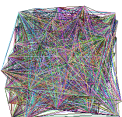



- What does the Internet look like?
- What does FaceBook look like?
- What is 'normal'/'abnormal'?
- which patterns/laws hold?

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Problem #1 - network and graph mining

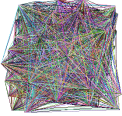



- What does the Internet look like?
- What does FaceBook look like?
- What is 'normal'/'abnormal'?
- which patterns/laws hold?
 - To spot **anomalies** (rarities), we have to discover **patterns**


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Problem #1 - network and graph mining



- What does the Internet look like?
- What does FaceBook look like?
- What is 'normal'/'abnormal'?
- which patterns/laws hold?
 - To spot **anomalies** (rarities), we have to discover **patterns**
 - **Large** datasets reveal patterns/anomalies that may be invisible otherwise...



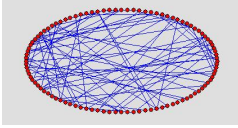
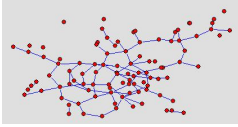
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Are real graphs random?

- random (Erdos-Renyi) graph – 100 nodes, avg degree = 2
- before layout
- after layout
- No obvious patterns

(generated with: pajek
<http://vlado.fmf.uni-lj.si/pub/networks/pajek/>)

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Graph mining

- Are real graphs random?

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Laws and patterns

- Are real graphs random?
- A: NO!!
 - Diameter
 - in- and out- degree distributions
 - other (surprising) patterns
- So, let's look at the data

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Solution# S.1

- Power law in the degree distribution [SIGCOMM99]

internet domains

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Solution# S.1

- Power law in the degree distribution [SIGCOMM99]

internet domains

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Solution# S.1

- Q: So what?

internet domains

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Solution# S.1

- Q: So what?
- A1: # of two-step-away pairs:
= friends of friends (F.O.F.)
internet domains

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Solution# S.1

Gaussian trap

- Q: So what?
- A1: # of two-step-away pairs: $O(d_{\max}^2) \sim 10M^2$
internet domains

$\sim 0.8PB \rightarrow$
a data center(!)

DCO @ CMU

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Solution# S.1

Gaussian trap

- Q: So what?
- A1: # of two-step-aware inter-
 $\sim 10M^2$
 $\sim 0.8PB \rightarrow$
a data center(!)

**Such patterns \rightarrow
New algorithms**

-0.82

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Solution# S.2: Eigen Exponent E

Eigenvalue

Exponent = slope
 $E = -0.48$
May 2001

$Ax = \lambda x$

- A2: power law in the eigenvalues of the adjacency matrix

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Solution# S.2: Eigen Exponent E

Eigenvalue

Exponent = slope
 $E = -0.48$
May 2001

- [Mihail, Papadimitriou '02]: slope is $\frac{1}{2}$ of rank exponent

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But:

How about graphs from other domains?

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More power laws:

- web hit counts [w/ A. Montgomery]

Count (log scale)

Web Site Traffic

Zipf

ebay

users

sites

in-degree (log scale)

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epinions.com

- who-trusts-whom [Richardson + Domingos, KDD 2001]

count

Original graph

R-MAT graph

trusts-2000-people user

(out) degree

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And numerous more

- # of sexual contacts
- Income [Pareto] – ‘80-20 distribution’
- Duration of downloads [Bestavros+]
- Duration of UNIX jobs (‘mice and elephants’)
- Size of files of a user
- ...
- ‘Black swans’

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 - degree, diameter, eigen,
 - Triangles
 - Weighted graphs
 - Time evolving graphs

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Solution# S.3: Triangle ‘Laws’




- Real social networks have a lot of triangles

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Solution# S.3: Triangle ‘Laws’




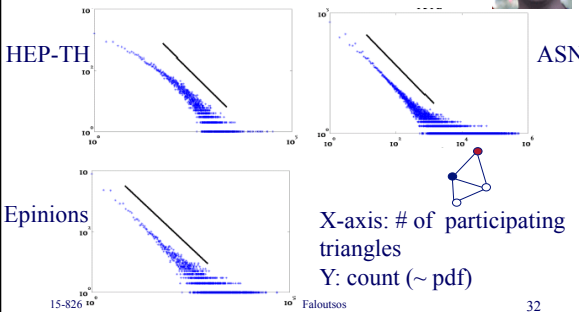
- Real social networks have a lot of triangles
 - Friends of friends are friends
- Any patterns?

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Triangle Law: #S.3

[Tsourakakis ICDM 2008]

HEP-TH ASN

Epinions


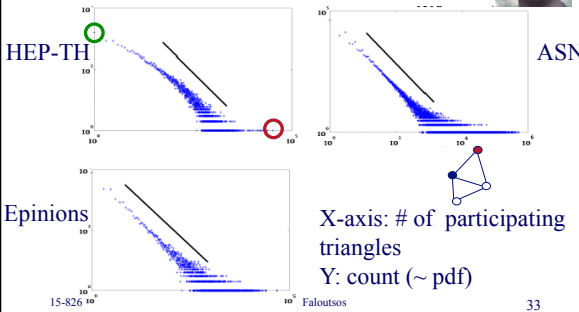
X-axis: # of participating triangles
Y: count (~ pdf)

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Triangle Law: #S.3

[Tsourakakis ICDM 2008]

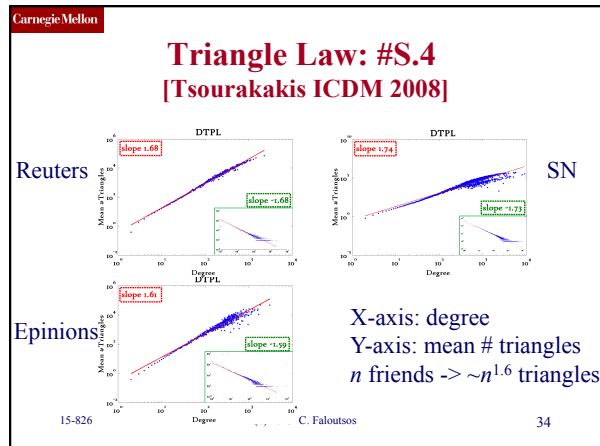



HEP-TH ASN

Epinions

X-axis: # of participating triangles
Y: count (~ pdf)

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Triangle Law: Computations [Tsourakakis ICDM 2008]

But: triangles are expensive to compute
(3-way join; several approx. algos)
Q: Can we do that quickly?

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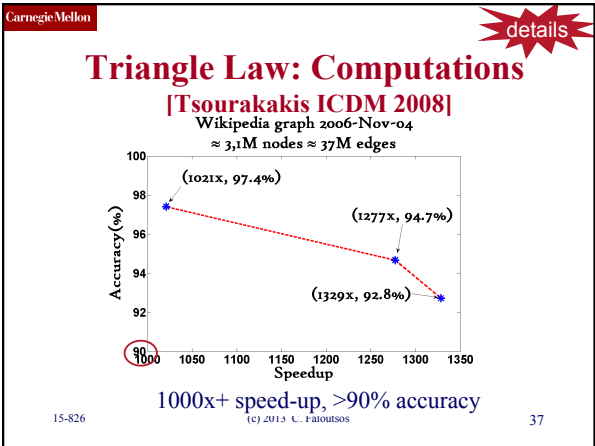
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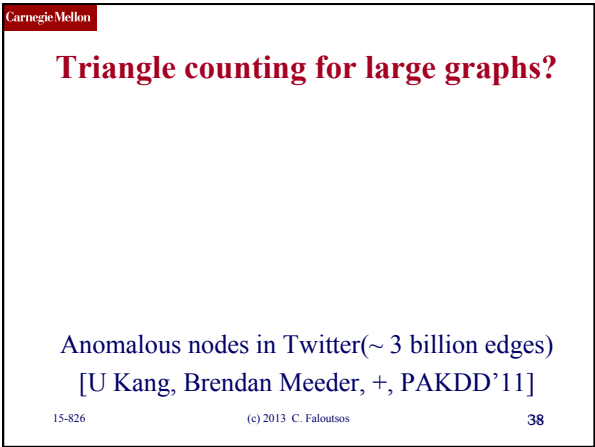
Triangle Law: Computations [Tsourakakis ICDM 2008]

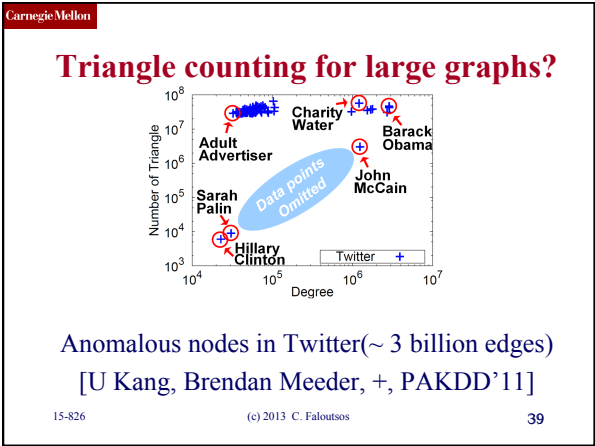
But: triangles are expensive to compute
(3-way join; several approx. algos)
Q: Can we do that quickly?
A: Yes!

#triangles = $\frac{1}{6} \sum (\lambda_i^3)$
(and, because of skewness (S2) ,
we only need the top few eigenvalues!

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Any other ‘laws’?

Yes!

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Any other ‘laws’?

Yes!

- Small diameter (~ constant!) –
 - six degrees of separation / ‘Kevin Bacon’
 - small worlds [Watts and Strogatz]

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Any other ‘laws’?

- Bow-tie, for the web [Kumar+ ‘99]
- IN, SCC, OUT, ‘tendrils’
- disconnected components

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Any other 'laws'?

- power-laws in communities (bi-partite cores) [Kumar+, '99]

Log(count)

2:3 core (m:n core)

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Any other 'laws'?

- "Jellyfish" for Internet [Tauro+ '01]
- core: ~clique
- ~5 concentric layers
- many 1-degree nodes

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EigenSpokes

B. Aditya Prakash, Mukund Seshadri, Ashwin Sridharan, Sridhar Machiraju and Christos Faloutsos: *EigenSpokes: Surprising Patterns and Scalable Community Chipping in Large Graphs*, PAKDD 2010, Hyderabad, India, 21-24 June 2010.

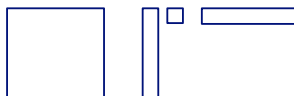
Useful for fraud detection!

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EigenSpokes

- Eigenvectors of adjacency matrix
 - equivalent to singular vectors (symmetric, undirected graph)

$$A = U\Sigma U^T$$


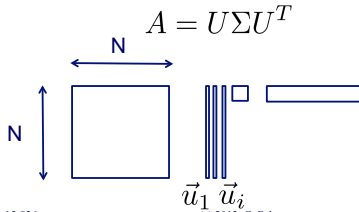
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EigenSpokes

details

- Eigenvectors of adjacency matrix
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$$A = U\Sigma U^T$$


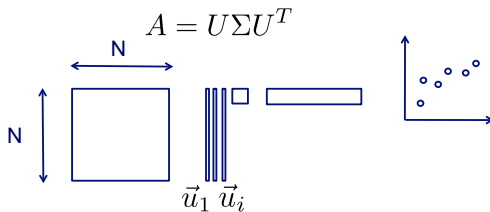
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EigenSpokes

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- Eigenvectors of adjacency matrix
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$$A = U\Sigma U^T$$


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EigenSpokes

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- Eigenvectors of adjacency matrix
 - equivalent to singular vectors (symmetric, undirected graph)

$$A = U\Sigma U^T$$

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EigenSpokes

details

- Eigenvectors of adjacency matrix
 - equivalent to singular vectors (symmetric, undirected graph)

$$A = U\Sigma U^T$$

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EigenSpokes

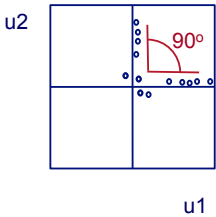
- EE plot:
 - Scatter plot of scores of u_1 vs u_2
 - One would expect
 - Many points @ origin
 - A few scattered ~randomly

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EigenSpokes

- EE plot:
- Scatter plot of scores of u_1 vs u_2
- One would expect
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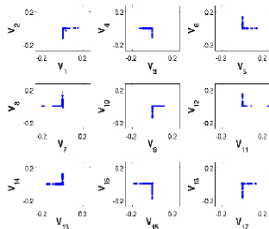


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EigenSpokes - pervasiveness

- Present in mobile social graph
 - across time and space
- Patent citation graph

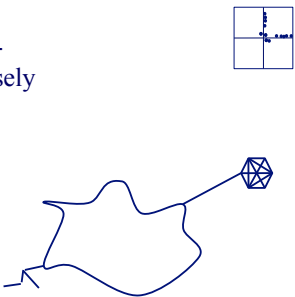


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EigenSpokes - explanation

Near-cliques, or near-bipartite-cores, loosely connected

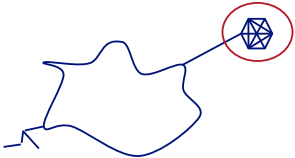



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EigenSpokes - explanation

Near-cliques, or near-bipartite-cores, loosely connected

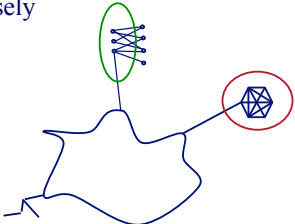




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EigenSpokes - explanation

Near-cliques, or near-bipartite-cores, loosely connected

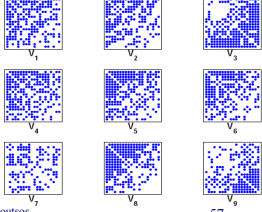
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EigenSpokes - explanation

Near-cliques, or near-bipartite-cores, loosely connected

spy plot of top 20 nodes



So what?

- Extract nodes with high scores
- high connectivity
- Good “communities”

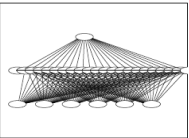
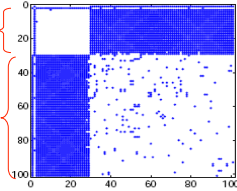
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Bipartite Communities!

patents from same inventor(s)
'cut-and-paste' bibliography!

magnified bipartite community

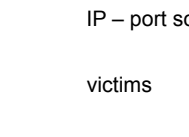
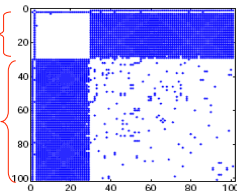
Useful for fraud detection!

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Bipartite Communities!

IP – port scanners
victims

Useful for fraud detection!


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➡




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Observations on weighted graphs?

- A: yes - even more 'laws'!



M. McGlohon, L. Akoglu, and C. Faloutsos
Weighted Graphs and Disconnected Components: Patterns and a Generator.
SIG-KDD 2008

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Observation W.1: Fortification

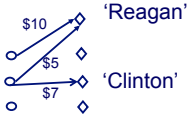
Q: How do the weights of nodes relate to degree?

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Observation W.1: Fortification

More donors, more \$?




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Observation W.1: fortification: Snapshot Power Law

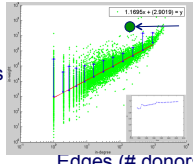
- Weight: super-linear on in-degree
- exponent 'iw': $1.01 < iw < 1.26$

**More donors,
even more \$**



In-weights (\$)

Orgs-Candidates




Edges (# donors)

e.g. John Kerry, \$10M received, from 1K donors

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Problem: Time evolution

- with Jure Leskovec (CMU -> Stanford)
- and Jon Kleinberg (Cornell – sabb. @ CMU)





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T.1 Evolution of the Diameter

- Prior work on Power Law graphs hints at **slowly growing diameter**:
 - diameter $\sim O(\log N)$
 - diameter $\sim O(\log \log N)$
- What is happening in real data?




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T.1 Evolution of the Diameter

- Prior work on Power Law graphs hints at **slowly growing diameter**:
 - diameter $\sim O(\log N)$
 - diameter $\sim O(\log \log N)$
- What is happening in real data?
- Diameter **shrinks** over time

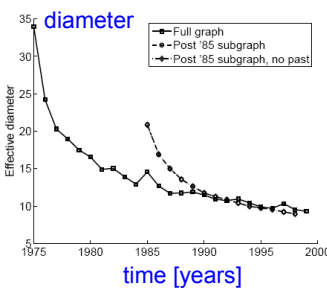


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T.1 Diameter – “Patents”

- Patent citation network
- 25 years of data
- @1999
 - 2.9 M nodes
 - 16.5 M edges



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T.2 Temporal Evolution of the Graphs

- $N(t)$... nodes at time t
- $E(t)$... edges at time t
- Suppose that
 - $N(t+1) = 2 * N(t)$
- Q: what is your guess for
 - $E(t+1) = ? 2 * E(t)$

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T.2 Temporal Evolution of the Graphs

- $N(t)$... nodes at time t
- $E(t)$... edges at time t
- Suppose that
 - $N(t+1) = 2 * N(t)$
- Q: what is your guess for
 - $E(t+1) = ? 2 * E(t)$
- A: over-doubled!
 - But obeying the “Densification Power Law”

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T.2 Densification – Patent Citations

- Citations among patents granted
- @1999
 - 2.9 M nodes
 - 16.5 M edges
- Each year is a datapoint

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More on Time-evolving graphs

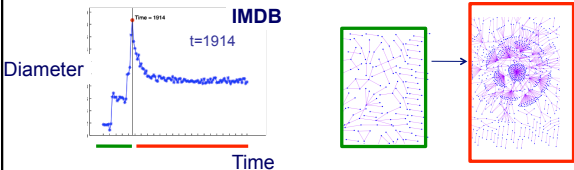
M. McGlohon, L. Akoglu, and C. Faloutsos
Weighted Graphs and Disconnected Components: Patterns and a Generator.
 SIG-KDD 2008

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[Gelling Point]

- Most real graphs display a gelling point
- After gelling point, they exhibit typical behavior. This is marked by a spike in diameter.



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Observation T.3: NLCC behavior

Q: How do NLCC's emerge and join with the GCC?

('NLCC' = non-largest conn. components)

- Do they continue to grow in size?
- or do they shrink?
- or stabilize?



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Observation T.3: NLCC behavior

Q: How do NLCC's emerge and join with the GCC?

('NLCC' = non-largest conn. components)

- Do they continue to grow in size?
- or do they shrink?
- or stabilize?



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Observation T.3: NLCC behavior

Q: How do NLCC's emerge and join with the GCC?

('NLCC' = non-largest conn. components)

YES - Do they continue to grow in size?

YES - or do they shrink?

YES - or stabilize?

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Observation T.3: NLCC behavior

- After the gelling point, the GCC takes off, but NLCC's remain ~constant (actually, **oscillate**).

CC size

Time-stamp

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Timing for Blogs

- with Mary McGlohon (CMU->Google)
- Jure Leskovec (CMU->Stanford)
- Natalie Glance (now at Google)
- Mat Hurst (now at MSR)

[SDM'07]

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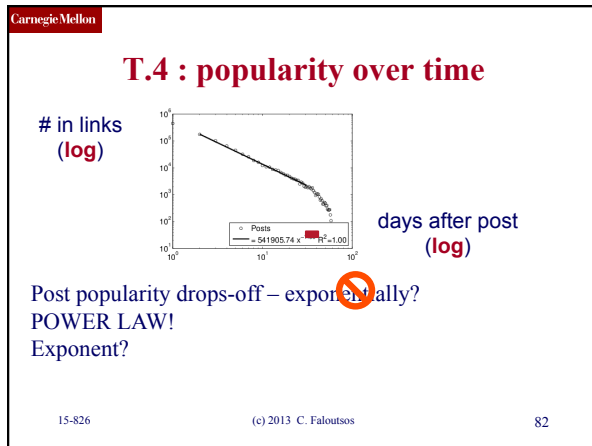
T.4 : popularity over time

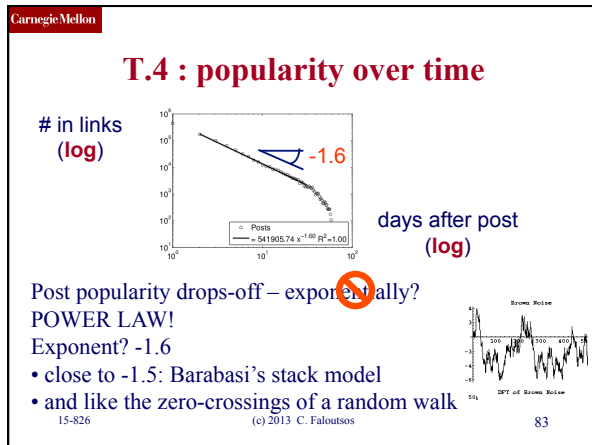
in links

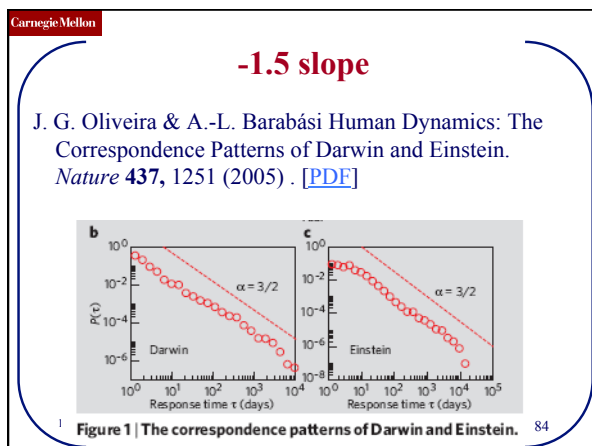
lag: days after post

Post popularity drops-off – exponentially?

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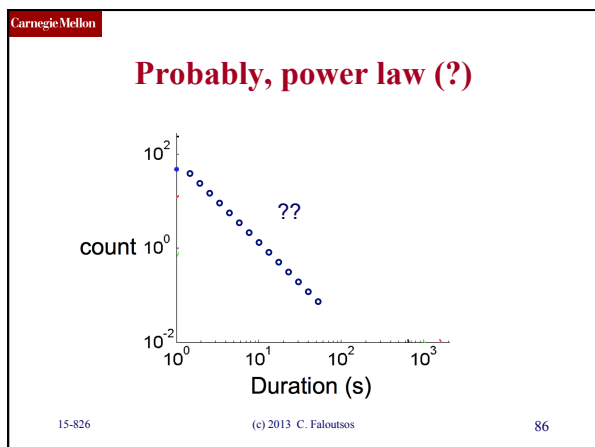
T.5: duration of phonecalls

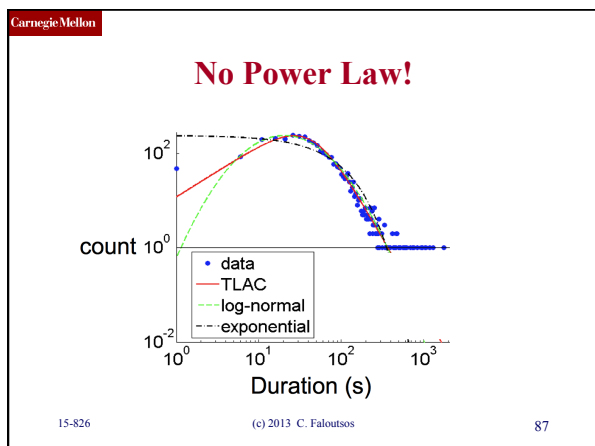
*Surprising Patterns for the Call
Duration Distribution of Mobile
Phone Users*

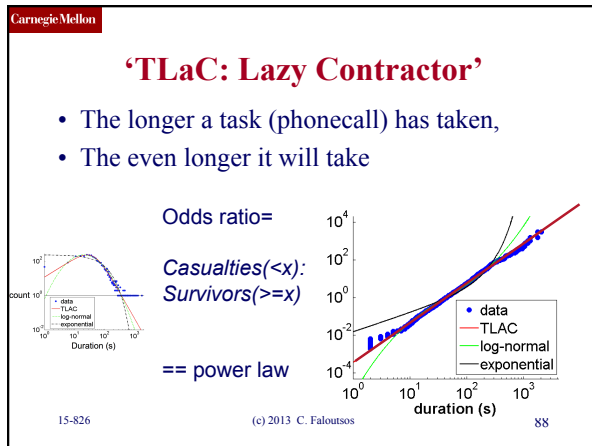


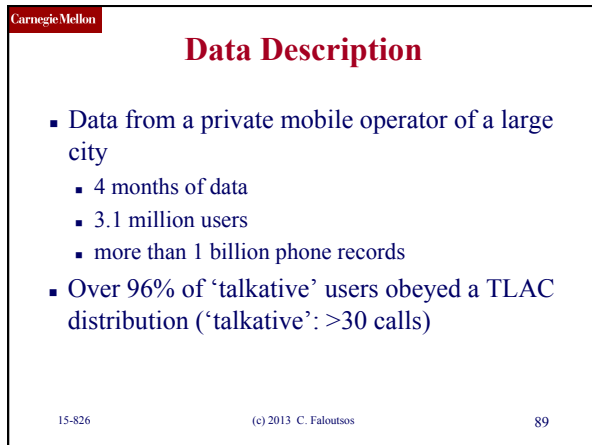
Pedro O. S. Vaz de Melo, Leman
Akoglu, Christos Faloutsos, Antonio
A. F. Loureiro
PKDD 2010

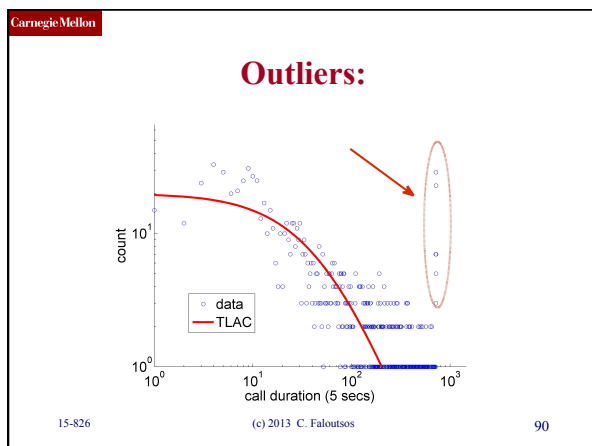
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Outline




- Introduction – Motivation
- Problem#1: Patterns in graphs
- ➡ • Problem#2: Scalability -PEGASUS
- Conclusions


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Scalability




- Google: > 450,000 processors in clusters of ~2000 processors each [Barroso, Dean, Hölzle, “Web Search for a Planet: The Google Cluster Architecture” IEEE Micro 2003]
- Yahoo: 5Pb of data [Fayyad, KDD’07]
- Problem: machine failures, on a daily basis
- How to parallelize data mining tasks, then?
- A: map/reduce – hadoop (open-source clone)
<http://hadoop.apache.org/>



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Outline – Algorithms & results



	Centralized	Hadoop/ PEGASUS
Degree Distr.	old	old
Pagerank	old	old
Diameter/ANF	old	HERE
Conn. Comp	old	HERE
Triangles	done	HERE
Visualization	started	

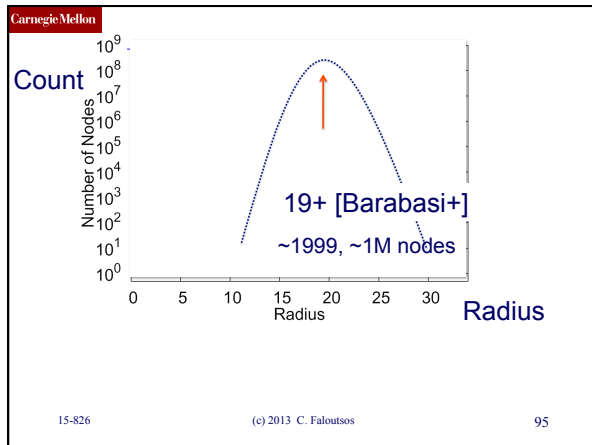
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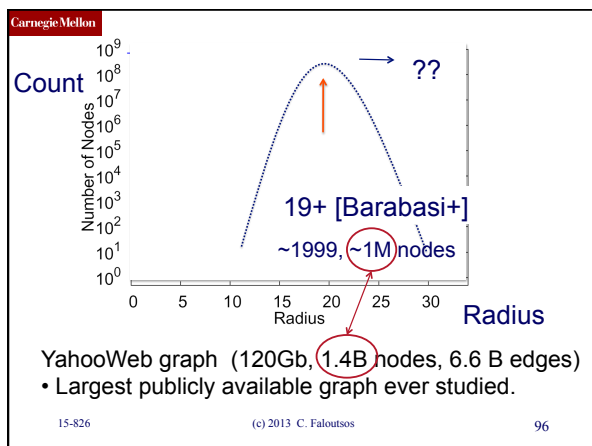
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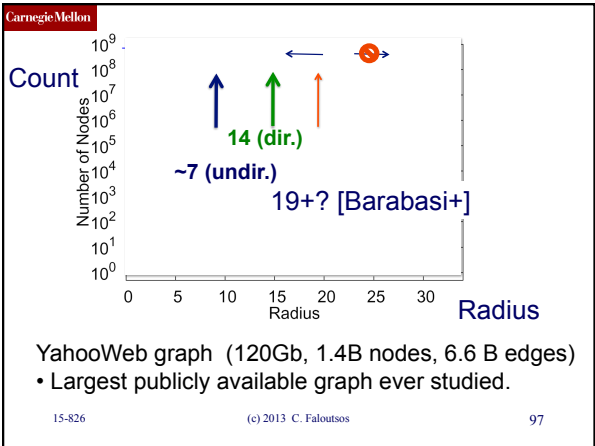
HADI for diameter estimation

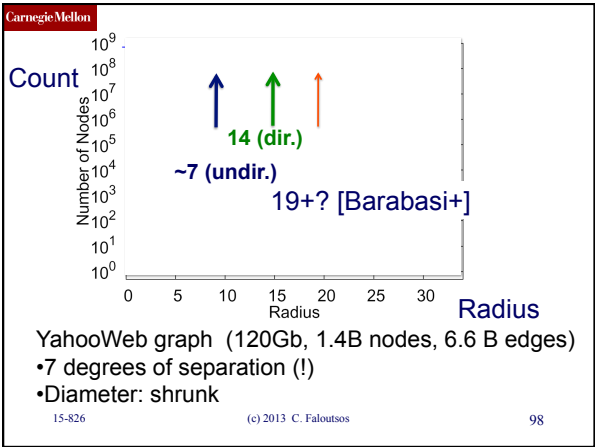
- *Radius Plots for Mining Tera-byte Scale Graphs* U **Kang**, Charalampos Tsourakakis, Ana Paula Appel, Christos Faloutsos, Jure Leskovec, SDM'10
- Naively: diameter needs $O(N^2)$ space and up to $O(N^3)$ time – **prohibitive** ($N \sim 1B$)
- Our HADI: linear on E ($\sim 10B$)
 - Near-linear scalability wrt # machines
 - Several optimizations \rightarrow 5x faster

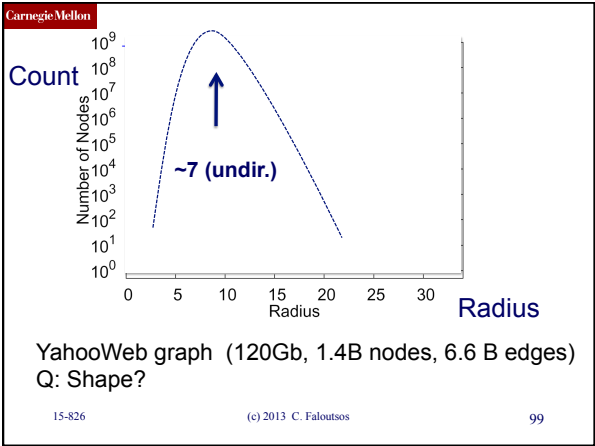
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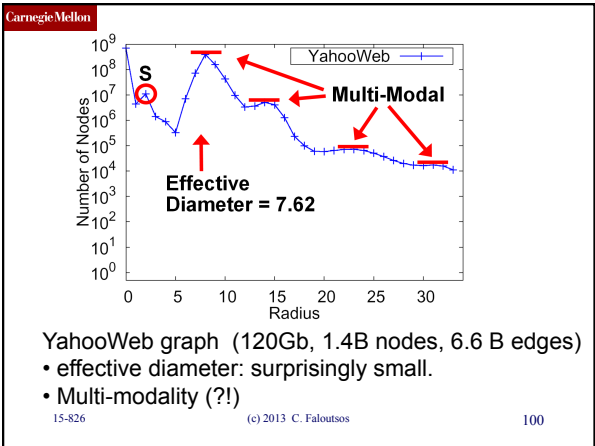


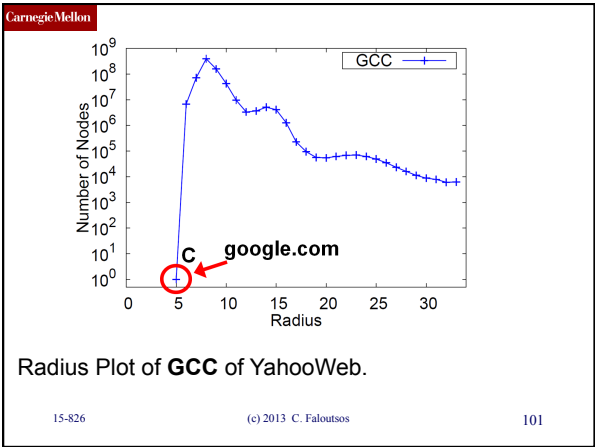


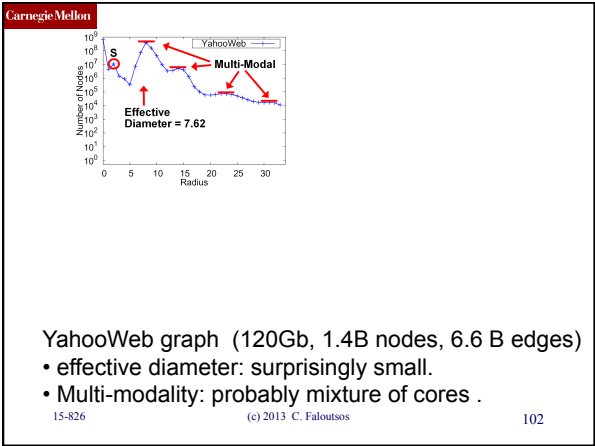












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Conjecture:

EN

~7

DE

BR

YahooWeb graph (120Gb, 1.4B nodes, 6.6 B edges)

- effective diameter: surprisingly small.
- Multi-modality: probably mixture of cores .

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Conjecture:

~7

YahooWeb graph (120Gb, 1.4B nodes, 6.6 B edges)

- effective diameter: surprisingly small.
- Multi-modality: probably mixture of cores .

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Run time in hours

7.6x faster

5.1x

3.8x

3.2x

Legend: HADI-plain, HADI-BSE, HADI-BL, HADI-OPT

Data: KR-2B, KR-1.1B, ER-2B, ER-1.1B

Running time - Kronecker and Erdos-Renyi Graphs with billions edges.

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Outline – Algorithms & results

	Centralized	Hadoop/ PEGASUS
Degree Distr.	old	old
Pagerank	old	old
Diameter/ANF	old	HERE
→ Conn. Comp	old	HERE
Triangles		HERE
Visualization	started	

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Generalized Iterated Matrix
Vector Multiplication (GIMV)

[PEGASUS: A Peta-Scale Graph Mining System - Implementation and Observations.](#)
U Kang, Charalampos E. Tsourakakis,
and Christos Faloutsos.
[\(ICDM\)](#) 2009, Miami, Florida, USA.
Best Application Paper (runner-up).

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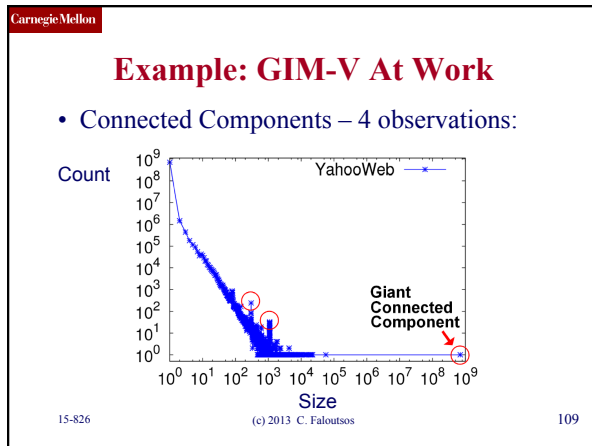
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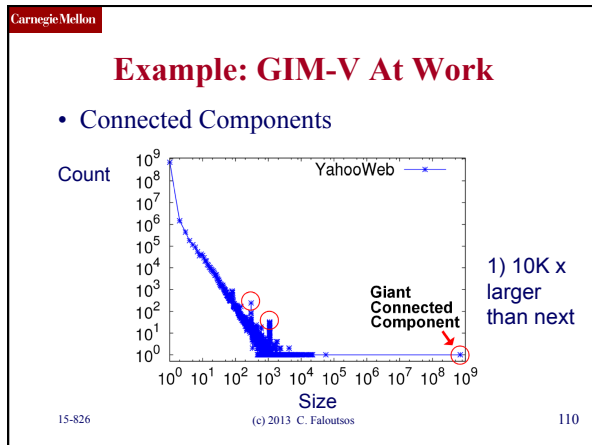
Generalized Iterated Matrix
Vector Multiplication (GIMV)

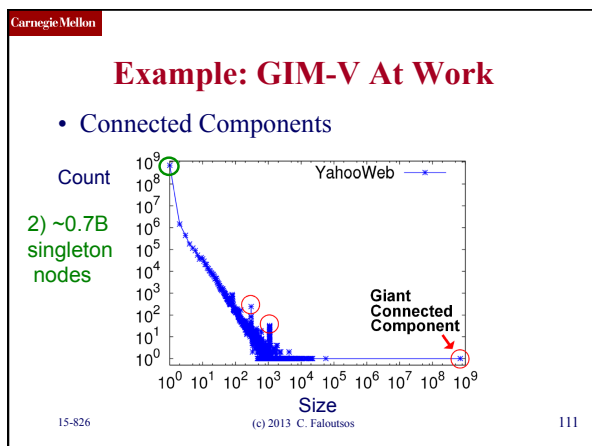
- PageRank
- proximity (RWR)
- Diameter
- Connected components
- (eigenvectors,
- Belief Prop.
- ...)

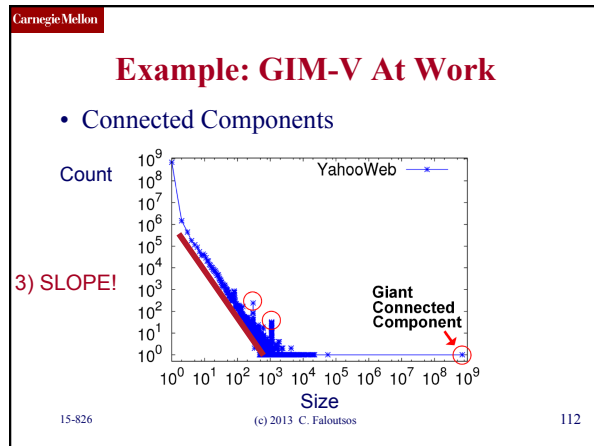
Matrix – vector
Multiplication
(iterated)

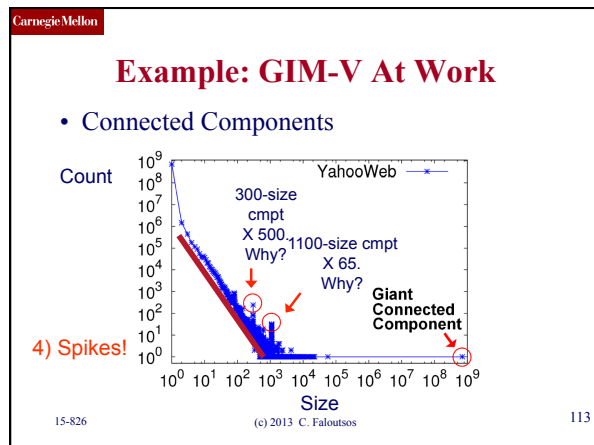
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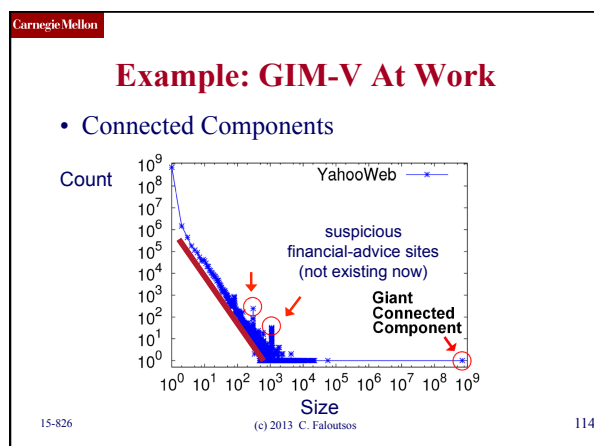


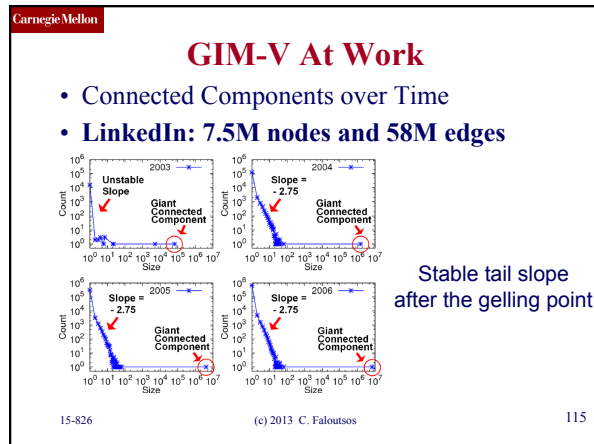













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Outline



- Introduction – Motivation
- Problem#1: Patterns in graphs
- DELETE
- Problem#2: Scalability
- ➡ • Conclusions

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OVERALL CONCLUSIONS – low level:

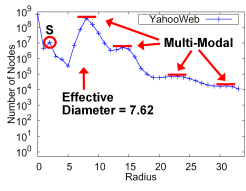
- Several new **patterns** (fortification, shrinking diameter, triangle-laws, conn. components, etc)
- New **tools**:
 - anomaly detection (OddBall), belief propagation, immunization
- **Scalability**: PEGASUS / hadoop

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OVERALL CONCLUSIONS – high level

- **BIG DATA:** Large datasets reveal patterns/ outliers that are invisible otherwise



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(Project info)


www.cs.cmu.edu/~pegasus




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McGlohon,
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