

# Concurrency Assurance in Fluid

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Related reading: *Assuring and Evolving  
Concurrent Programs: Annotations and  
Policy*

15-413:  
Introduction to Software Engineering  
Jonathan Aldrich



## Find the Concurrency Bug!

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```
public class Logger {  
    private Filter filter;  
  
    public void setFilter(Filter newFilter)  
        throws SecurityException {  
        if (!anonymous) manager.checkAccess();  
        filter = newFilter;  
    }  
  
    public void log (LogRecord record) { ...  
        synchronized (this) {  
            if (filter != null && !filter.isLoggable(record))  
                return;  
        }  
    }  
}
```

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2

## PREfix: Language-Level Errors



- Error defined by language
  - Precise characterization of error
  - Any program that manifests that error is incorrect
  - Easy to define fully automated analysis
- Example: null pointer dereference
  - Occurs when `*p` is executed and `p == null`
  - Can be found by tracking which pointers may be null

## Concurrency Errors



- Example: data race condition
  - (Definition from Savage et al., *Eraser: A Dynamic Data Race Detector for Multithreaded Programs*)
  - Two threads access the same variable `v`
  - At least one access is a write
  - No explicit mechanism prevents the accesses from being simultaneous

## Concurrency Errors



- Example: data race condition
  - (Definition from Savage et al., *Eraser: A Dynamic Data Race Detector for Multithreaded Programs*)
  - Two threads access the same variable  $v$
  - At least one access is a write
  - No explicit mechanism prevents the accesses from being simultaneous
- Challenges
  - Difficult to check statically
    - How to tell if accesses can be simultaneous?
    - How to tell what synchronization mechanism is used?
  - Not always an error
    - Race may not affect correctness
- PREFIX approach will not work
  - Too many possibilities to explore, too many false positives

## Would Testing/Inspections Work?



## Would Testing/Inspections Work?



- Testing
  - Difficult because concurrency errors are non-deterministic
- Inspections
  - Concurrency errors are often non-local
    - Like errors that PREFIX finds
  - Require knowledge of programmer intent

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7

## Fluid: Models are missing



- **Programmer design intent** is missing
  - Not explicit in Java, C, C++, *etc*
    - What lock protects this object?
      - *This lock protects that state*
    - What is the actual extent of shared state of this object?
      - *This object is "part of" that object*
- **Adoptability**
  - Programmers: "Too difficult to express this stuff."
  - Fluid: Minimal **effort** — concise expression
    - Capture what programmers are **already thinking about**
    - No full specification
- **Incrementality**
  - Programmers: "I'm too busy; maybe after the deadline."
  - Fluid: **Payoffs** early and often
    - Direct programmer utility — *negative marginal cost*
    - Increments of payoff for increments of effort

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8

## Reporting Code–Model Consistency



### Tool analyzes model/code consistency

- No model  $\Rightarrow$  no assurance
- Identify likely model sites

### Three classes of results



- Code–model consistency



- Code–model inconsistency



- Informative — Request for annotation

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9

## BoundedFIFO



```
public class BoundedFIFO {
  //@aggregate [] into Instance
  //@unshared
  LoggingEvent[] buf;

  //@ lock BufLock is this protects Instance

  int numElts = 0, first = 0, next = 0, size;

  public BoundedFIFO(int size) { ... }

  //@ requires BufLock
  public LoggingEvent get() {
    if(numElts == 0) return null;
    LoggingEvent r = buf[first];
    if(++first == size) first = 0;
    numElts--;
    return r;
  }

  //@ requires BufLock
  public void put(LoggingEvent o) {
    if(numElts != size) {
      buf[next] = o;
      if(++next == size) next = 0;
      numElts++;
    }
  }

  //@ requires BufLock
  public int getMaxSize() {
    return size;
  }

  // no annotation required
  public synchronized void resize(int
    newSize) { ... }
  ...
}
```

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10

## BoundedFIFO Client



```
public class FIFOClient {
    private final BoundedFIFO fifo = ...;
    ...
    public void putter(LoggingEvent e) {
        synchronized(fifo) {
            while(fifo.isFullO) {
                try { fifo.wait(); }
                catch(InterruptedException ie) {}
            }
            fifo.put(e);
            if(fifo.wasEmptyO) fifo.notify();
        }
    }

    public LoggingEvent getter() {
        synchronized(fifo) {
            LoggingEvent e;
            while(fifo.length() == 0) {
                try { fifo.wait(); }
                catch(InterruptedException ie) {}
            }
            e = fifo.get();
            if(fifo.wasFull()) fifo.notify();
            return e ;
        }
    }

    public int length() {
        synchronized(fifo) { return fifo.length(); }
    }
}
```

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11

## Logger Revisited



```
/** @lock FilterLock is this protects filter */
public class Logger {
    private Filter filter;

    public void setFilter(Filter newFilter)
        throws SecurityException {
        if (!anonymous) manager.checkAccess();
        filter = newFilter;
    }

    public void log (LogRecord record) { ...
        synchronized (this) {
            if (filter != null && !filter.isLoggable(record))
                return;
        }
    }
}
```

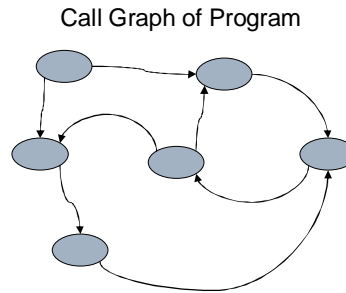
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12

# How Incrementality Works



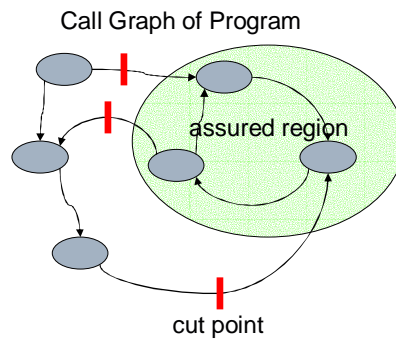
- Incrementality
  - When you annotate a portion of the program, you'll get immediate feedback on concurrency correctness
  - Incremental benefit for each unit of cost
- How can one provide incremental benefit with mutual dependencies?



# How Incrementality Works



- Incrementality
  - When you annotate a portion of the program, you'll get immediate feedback on concurrency correctness
  - Incremental benefit for each unit of cost
- How can one provide incremental benefit with mutual dependencies?
- Cut points
  - Method annotations partition call graph
  - Can assure property of a subgraph
  - Assurance is *contingent* on accuracy of trusted cut point method annotations



## Questions for Evaluating Tools

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- What class of errors does the tool find?
  - And can that class be found with other techniques?
- Can the tool miss errors?
- How many false errors does it report?
- Can I run it on part of a system?
- How much manual effort is required?
- Does it find errors across procedure boundaries?
- Does it scale to large systems?