

**15-410**

***“My other car is a cdr” -- Unknown***

**Exam #1**  
**Oct. 17, 2007**

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# Synchronization

## Checkpoint 2 –Wednesday, in cluster

- Reminders
  - context switch  $\neq$  mode switch
    - » Identify scenarios with one and not the other
  - context switch  $\neq$  interrupt
    - » Later it will be invoked in other circumstances

## Google “Summer of Code”

- <http://code.google.com/soc/>
- Hack on an open-source project
  - And get paid
  - And probably get recruited

## CMU SCS “Coding in the Summer”

# Synchronization

## Debugging advice

- Monday as I was buying lunch I received a fortune

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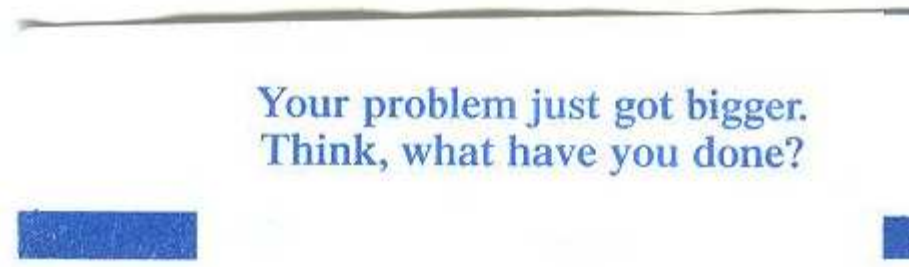


Image credit: Kartik Subramanian

# A Word on the Final Exam

## Disclaimer

- Past performance is not a guarantee of future results

## The course will change

- Up to now: “basics” - What you *need* for Project 3
- Coming: advanced topics
  - Design issues
  - Things you won't experience via implementation

## Examination will change to match

- More design questions
- Some things you won't have implemented (text useful!!)
- Still 3 hours, but more stuff (~100 points, ~7 questions)

# Outline

**Question 1**

**Question 2**

**Question 3**

**Question 4**

**Question 5**

# Q1 –Short Answer

## **Write pipe (also known as “write buffer”)**

- **Key concept: the part of a “modern” computer which makes it “modern”**
- **Popular but not as relevant to this course**
  - **The write side of a pipe**
  - **Some kind of write buffer which isn't a write pipe**

# Q1 –Short Answer

## Interrupt acknowledge

- **Best answers covered:**
  - **What it's for**
    - » **Sending device back to start of protocol (enabling it to assert another interrupt later)**
  - **When it happens**
    - » **When processor has acquired the information necessary to characterize and handle the interrupt**
  - **How it happens**
    - » **Processor sends a command (in our world, via an OUTB)**



# Q2 –Monitor Implementation

## Write some macros...

- `M_DECL()`, `M_INIT()`, etc.
- ...to support a “monitor style” of programming

## Getting started

- `work_setup()` needs to `thr_create()` a worker thread
  - Nobody else can...

## Locking issue

- Sometimes we need others to enter the monitor to progress us... `condition_wait()` will make that happen
- Sometimes we need others to *not* enter the monitor just yet... but `condition_wait()` will make that happen

# Q2 –Monitor Implementation

## Types and returning

- `M_RETURN(t,v)` –takes a type and a value
- There is a subtle locking problem here
  - What happens when I `M_RETURN(int,some_global_int)`?
  - `M_RETURN()` needs to accomplish two things
    - » Neither order will work
    - » So `M_RETURN()` needs to accomplish *three* things
      - That's what the type parameter is for

## Scoping

- A common `M_DECL()` mistake would mean each program could contain only one monitor.

# Q3 –Critical Section Algorithm

## The mission

- Evaluate a proposed critical-section algorithm in terms of whether it provides mutual exclusion, progress, and bounded waiting

## Terminology to watch out for

- Progress is about the *system*
- Bounded waiting is about a *particular victim*
  - Violating bounded waiting means “we can't write down a bound”
  - It does *not* mean “we can show there exists a small, bounded amount of unfairness” - strict FIFO behavior is *not* required, because it's much too hard

# Q3 –Critical Section Algorithm

## Mutual exclusion

- Pretty much everybody was able to show this was broken
- Some people lost some points for execution traces that were too terse (the loop is a key part of the story)

# Q3 –Critical Section Algorithm

## Progress

- **No!**
- **The key problem is that mutual exclusion is broken**
- **Two racing unlockers can leave the lock in a broken state**

Thread 2	Thread 1
T2 is done wanting	
Decide to appoint T1	
	T1 is done wanting
	Lock is available to all
Appoint T1	

**Now T1 goes on vacation to Belize...**

# Q3 –Critical Section Algorithm

## Progress

- **Not** progress violations
  - One thread might crash while holding the lock
  - One thread might never unlock the lock
    - » True, but not faults in the algorithm

## Another way to show progress isn't assured

- `set()` isn't atomic

## Other problems

- Bad execution traces which can't actually happen
- Explaining what the algorithm **wants** to do

# Q3 –Critical Section Algorithm

## Bounded waiting

- No!
  - Gee, this algorithm isn't so hot, is it?
- Key problem: `set ( )`

# Q4 –Deadlock

## Issues with the new cluster

- Description of resources (computers, servers, projector)
- Description of threads (OS, Networks)
- Deadlock? Yes/no/why?

## (A) –Can OS students deadlock?

- Observe: this is “Dining Philosophers”!
- Observe: the projector injects a subtle yet important property...

## (B) –Can Networks students deadlock?

- Can explain in terms of h&w or graph cycles
- Must state name of property and show it



# Q4 –Deadlock

## **(C) –Can mixture of students deadlock?**

- **Parts of a complete solution**
  - **Diagram of sufficient clarity**
  - **Event trace of sufficient clarity (clear text was accepted)**
  - **Explanation of why the situation, as diagrammed and traced, is classified the way it is**

## Q5 – Your Partner's Code

```
char *the_word(int num)
{
    char buf[8];

    switch (num % 4) {
        case 0: snprintf(buf, sizeof(buf), "zero"); break;
        case 1: snprintf(buf, sizeof(buf), "one"); break;
        case 2: snprintf(buf, sizeof(buf), "two"); break;
        case 3: snprintf(buf, sizeof(buf), "three"); break;
    }
    return (buf);
}
```

# Q5 –Your Partner's Code

## (A) –What's wrong with this picture?

- The “213 answer”: returning a pointer to “automatic storage”

## Claims difficult to support

- “Stack memory `disappears' when a function returns”
  - Set to zero...
  - Removed from address space...
  - Will cause a segmentation fault...
  - ...Unfortunately not true
- “`snprintf()` is not up to this job”
- “...the heap...”
- “`sizeof()` is evil”

# Q5 –Your Partner's Code

## “sizeof( ) is evil”

- There *are* times when sizeof() “doesn't do what you want”

```
void foo(char s[1024]) {  
    ... sizeof(s) ... // not 1024  
}
```

```
void bar(void) {  
    char *s;  
    s = malloc(512);  
    ... sizeof(s) ... // not 512  
}
```

# Q5 –Your Partner's Code

## “sizeof( ) is evil”

- There *are* times when sizeof() “doesn't do what you want”
- ...but it isn't *designed* to be wrong all the time!

## The problem isn't actually sizeof()

- The issue is that in C *some* things which look like arrays aren't
- Pointers can be used like arrays, but are pointer-sized
- Function parameters which look like arrays are actually pointers, and are pointer-sized
- Actual arrays (local or global) are actually arrays, and are array-sized

# Q5 –Your Partner's Code

## (B) What's wrong with the code –in context?

- Two possible answers
- For complete credit, the less-than-obvious one is better
  - There isn't another *thread* out there, but...

## Things to avoid

- “Some other thread...” - there are no other threads
- “The kernel...” - this code *is* the kernel
- Generally, avoid mysterious or missing actors

# Breakdown

<b>90% = 67.5</b>	<b>0 students</b>
<b>80% = 60.0</b>	<b>19 students</b>
<b>70% = 52.5</b>	<b>25 students (52 and up)</b>
<b>60% = 45.0</b>	<b>4 students</b>
<b>50% = 37.5</b>	<b>9 students</b>
<b>&lt;50%</b>	<b>3 students</b>

## Comparison

- Scores are lower than typical
  - Even if we correct for that person who clearly forgot to answer that one question

# Implications

## Further analysis will probably suggest a mild scaling

- Maybe something like 3-5 points

## Score below 70%?

- Figure out what happened
- Probably plan to do better on the final exam

## Warning...

- To pass the class you must demonstrate reasonable proficiency on exams (project grades alone are not sufficient)
- See syllabus