

15-410

“...[insert witty quote here]...”

Exam #1
Mar. 14, 2005

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Synchronization

Checkpoint 2 – Friday, in cluster

- **Reminder: context switch \neq interrupt**
 - **Later other things will invoke it too**

Upcoming events

- **15-412 – switched to Fall semester!**
- **Summer internship with SCS Facilities?**

A Word on the Final Exam

Disclaimer

- Past performance is not a guarantee of future results

The course will change

- Up to now: “basics”
 - What you need for Project 3
- Coming: advanced topics
 - Design issues
 - Things you won't experience via implementation

Examination will change to match

- More design questions
- Some things you won't have implemented

Outline

Question 1

Question 2

Question 3

Question 4

Question 5

Q1 – Short Answer

User mode vs. kernel mode

- Some instructions reserved for kernel
- All memory is (in some sense) available to kernel
- All hardware devices available to kernel
- Interrupts, traps enter kernel mode

Less strong

- Kernel stacks are smaller (OS design choice)
- Kernel schedules threads (usually; if they're k threads...)

Q1 – Definitions

runnable

- “...if it can be run” - not very deep
- “right answer” has two parts
 - No impediment to running / eligible for scheduling
 - » No I/O in progress; not stunned by deschedule()
 - Not currently running on any CPU

Q1 – Definitions

progress

- We were looking for: critical-section progress requirement
- Less strong: “doing useful work”

bus lock

- Almost always refers to *memory* bus, during atomic operation (Test&Set, Exchange, etc.)
- Less strong: some notion of interrupt queueing

Q2 – Deadlock (“ClusterF”)

Part A – Prevention

- Recite the “Four Ways to Forgiveness”
- Not right: describing some dynamic scheme

Part B – Analysis of Prevention Approaches

- Should be clear that API *requires* mutex, non-preemption
- Banning hold & wait
 - Changes protocol, not API – fine
 - Risks starvation – better than deadlock
- Banning circular wait
 - “Always go clockwise” doesn't work: consider *two* circuits...
 - Acquiring links in numerical order (vs. *circuit* order) works
 - » requires changing protocol – fine

Q3 – iSpend Interrupt Handling

Key concept

- Trap gates \Rightarrow *each* interrupt can arrive once
- Stack must hold *sum* of all stack requirements, not max

Q4 – Porting traceback() to P1

What's missing in the P1 environment?

- Convenient pointer-validation approach!
 - No `msync()`, no `/proc`, no `write()`

How would you approximate?

- Basic idea: pointers outside 0...16M are not valid
- Refinements
 - Stack frames probably don't occupy “special” memory
 - » Video RAM, kernel code, ...
 - Stack frames should probably be in stack region(s)
 - » ...various approaches...

Q5 – Critical-section Algorithm

Algorithm – Dekker (see text), broken subtly

Result: doesn't guarantee *bounded waiting*

- The “back off if we're both trying to enter” part is broken
 - One party has an execution pattern which backs off *always*
 - Then other party backs off *never*
- Work through this... see a course staff member if further hints necessary

Popular near-miss

- Showing bounded waiting is broken, calling it progress

Summary

90%	=	67.5	7	students
80%	=	60.0	13	students
70%	=	52.5	10	students
60%	=	45.0	11	students
<60%			10	students

Summary

90% = 67.5 7 students

80% = 60.0 13 students

70% = 52.5 10 students

60% = 45.0 11 students

<60% 10 students

But...

- **Most popular score was 52.0!!!**

Summary (tweaked)

89% = 67.0 9 students

80% = 60.0 11 students

69% = 52.0 16 students

60% = 45.0 5 students

<60% 10 students

Implications

Score below 52?

- Figure out what happened
- Probably plan to do better on the final exam

Score below 35?

- Something went *drastically* wrong
- Passing the final exam may be a serious challenge
- To pass the class you must demonstrate some proficiency on exams (project grades alone are not sufficient)