

# 15-410

*“My computer is 'modern'!”*

## Synchronization #1

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**Dave Eckhardt**

# Notice

## Me vs. OSC Chapter 6

- I will cover 6.3 much more than the text does...
  - ...even more than the previous edition did...
  - This is a good vehicle for understanding race conditions

## Me vs. OS:P+P Chapter 5

- Philosophically very similar
- Examples and focus are different

## Not in the book

- “Atomic sequences vs. voluntary de-scheduling”
  - “Sim City” example

## Textbook recommended!

- We will spend ~4 lectures on one chapter (~7 on two)
- This is important stuff
  - Getting a “second read” could be very useful

# Outline

**An intrusion from the “real world”**

**Two fundamental operations**

**Three necessary critical-section properties**

**Two-process solution**

**N-process “Bakery Algorithm”**

# Mind your P's and Q's

**Imagine you wrote this code:**

```
choosing[i] = true;  
number[i] =  
    max(number[0], number[1], ...) + 1;  
choosing[i] = false;
```

# Mind your P's and Q's

**Imagine you wrote this code:**

```
choosing[i] = true;  
number[i] =  
    max(number[0], number[1], ...) + 1;  
choosing[i] = false;
```

**Imagine what is sent out over the memory bus is:**

```
number[i] = 11;  
choosing[i] = false;
```

**Is that ok?**

# Mind your P's and Q's

**Imagine you wrote this code:**

```
choosing[i] = true;  
number[i] =  
    max(number[0], number[1], ...) + 1;  
choosing[i] = false;
```

**How about this??**

```
choosing[i] = false;  
number[i] = 11;
```

**Is my computer broken???**

- “Computer Architecture for \$200, Dave”...

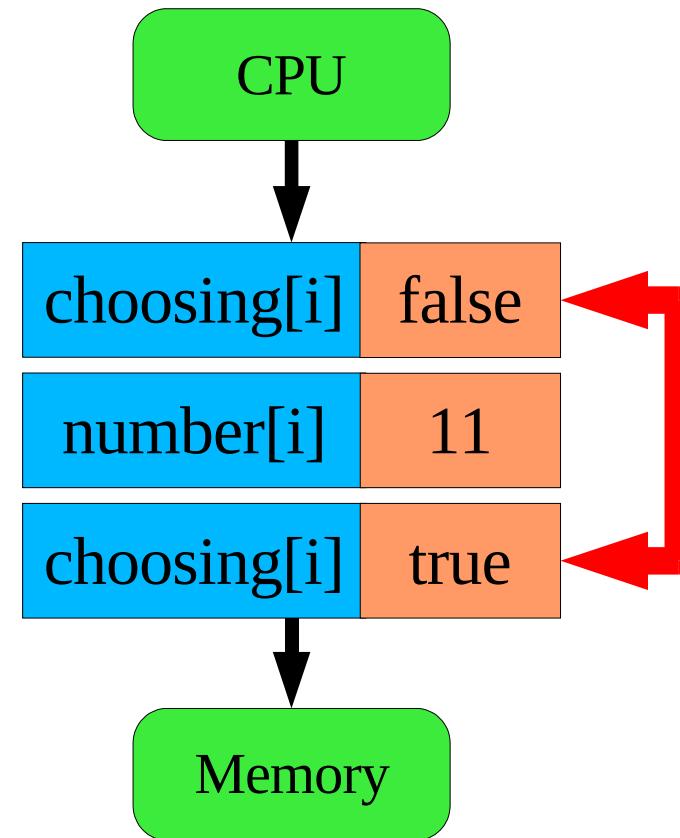
# Is my computer broken?!

**No, your computer is  
“modern”**

- Processor “write pipe” queues memory stores
- ...*and* coalesces “redundant” writes!

**Crazy?**

- Not if you're pounding out pixels!



# My Computer is Broken?!

## Magic “memory barrier” instructions available...

- ...stall processor until write pipe is empty

## Ok, now I understand

- Probably not!
  - <http://www.cs.umd.edu/~pugh/java/memoryModel/>
    - » see “Double-Checked Locking is Broken” Declaration
  - See also “release consistency”

## Textbook mutual exclusion algorithm memory model

- ...is “what you expect” (pre-“modern”)
- Ok to use simple model for homework, exams, P2
  - But it's not right for multi-processor Pentium-4 systems...

# Synchronization Fundamentals

## Two fundamental operations

- Atomic instruction sequence
- Voluntary de-scheduling

## Multiple implementations of each

- Uniprocessor vs. multiprocessor
- Special hardware vs. special algorithm
- Different OS techniques
- Performance tuning for special cases

## Be *very clear* on features, differences

- The two operations are more “opposite” than “the same”

# Synchronization Fundamentals

**Multiple client abstractions use the two operations**

**Textbook prefers**

- “Critical section”, semaphore, monitor

**Very relevant**

- Mutex/condition variable (POSIX pthreads)
- Java “synchronized” keyword (3 flavors)

# Synchronization Fundamentals

## Two Fundamental operations

- ▷ Atomic instruction sequence
- Voluntary de-scheduling

# Atomic Instruction Sequence

## Problem domain

- *Short* sequence of instructions
- Nobody else may interleave same sequence
  - or a “related” sequence
- “Typically” nobody is competing

# Non-interference

## Multiprocessor simulation (think: “Sim City”)

- Coarse-grained “turn” (think: hour)
- Lots of activity within each turn
- Think: M:N threads, M=objects, N=#processors

**Most cars don't interact in a game turn...**

- Must model those that do
- So street intersections can't generally be “processed” by multiple cars at the same time

# Commerce

<i>Customer 0</i>	<i>Customer 1</i>
<b>cash = store-&gt;cash;</b>	<b>cash = store-&gt;cash;</b>
<b>cash += 50;</b>	<b>cash += 20;</b>
<b>wallet -= 50;</b>	<b>wallet -= 20;</b>
<b>store-&gt;cash = cash;</b>	<b>store-&gt;cash = cash;</b>

Should the store call the police?

Is deflation good for the economy?

# Commerce – Observations

## Instruction sequences are “short”

- Ok to “mutually exclude” competitors (make them wait)

## Probability of collision is “low”

- Many non-colliding invocations per second
  - (lots of stores in the city)
- ***Must not*** use an expensive anti-collision approach!
  - “Just make a system call” is ***not*** an acceptable answer
- Common (non-colliding) case must be fast

# Synchronization Fundamentals

## Two Fundamental operations

- Atomic instruction sequence
- ⇒ Voluntary de-scheduling

# Voluntary De-scheduling

## Problem domain

- “Are we there yet?”
- “Waiting for Godot”

## Example - “Sim City” disaster daemon

```
while (date < 1906-04-18) cwait(date);
while (hour < 5) cwait(hour);
for (i = 0; i < max_x; i++)
  for (j = 0; j < max_y; j++)
    wreak_havoc(i,j);
```

# Voluntary De-scheduling

## Anti-atomic

- We **want** to be “maximally interleaved against”

## Running and making others wait is **wrong**

- Wrong for them – we won't be ready for a while
- Wrong for us – we can't be ready until **they** progress

## We don't **want** exclusion

## We **want** others to run - they **enable** us

## CPU **de**-scheduling is an OS service!

# Voluntary De-scheduling

## Wait pattern

```
LOCK WORLD
while (! (ready = scan_world())){
    UNLOCK WORLD
    WAIT_FOR(progress_event)
    LOCK WORLD
}
```

## Your partner-competitor will

```
SIGNAL(progress_event)
```

# Standard Nomenclature

## “Traditional CS” code skeleton / naming

```
do {  
    entry section  
    critical section:  
        ...computation on shared state...  
    exit section  
    remainder section:  
        ...private computation...  
} while (1);
```

# Standard Nomenclature

## What's muted by this picture?

- What's *in* that critical section?
  - Quick atomic sequence?
  - Need for a long sleep?

## For now...

- Pretend critical section is a brief atomic sequence
- Study the entry/exit sections

# Three Critical Section Requirements

## *Mutual Exclusion*

- At most one thread is executing each critical section

## *Progress*

- Choosing protocol must have bounded time
  - Common way to fail: choosing next entrant cannot wait for non-participants

## *Bounded waiting*

- Cannot wait forever once you begin entry protocol
- ...bounded number of entries by others
  - not necessarily a bounded number of *instructions*

# Notation For 2-Process Protocols

## Assumptions

- Multiple threads (1 CPU with timer, or multiple CPU's)
- Shared memory, but no locking/atomic instructions

**Thread  $i$  = “us”**

**Thread  $j$  = “the other thread”**

**$i, j$  are *thread-local* variables**

- $\{i,j\} = \{0,1\}$
- $j == 1 - i$

**This notation is “odd”**

- But it *may well appear in an exam question*

# Idea #1 - “Taking Turns”

```
int turn = 0;  
  
while (turn != i)  
    continue;  
...critical section...  
turn = j;
```

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```

**Mutual exclusion – yes (make sure you see it)**

# Idea #1 - “Taking Turns”

```
int turn = 0;  
  
while (turn != i)  
    continue;  
...critical section...  
turn = j;
```

**Mutual exclusion – yes (make sure you see it)**

**Progress - *no***

- *Strict* turn-taking is fatal
- If T0 never tries to enter, T1 will wait forever
  - Violates the “depends on non-participants” rule

# Idea #2 - “Registering Interest”

```
boolean want[2] = {false, false};
```

```
want[i] = true;  
while (want[j])  
    continue;  
...critical section...  
want[i] = false;
```

# Mutual Exclusion (Intuition)

<i>Thread 0</i>	<i>Thread 1</i>
<b>want[0] = true;</b>	
<b>while (want[1]) ;</b>	
<b>...enter...</b>	<b>want[1] = true;</b>
	<b>while (want[0]) ;</b>
	<b>while (want[0]) ;</b>
<b>want[0] = false;</b>	<b>while (want[0]) ;</b>
	<b>...enter...</b>

# Mutual Exclusion (Intuition)

<i>Thread 0</i>	<i>Thread 1</i>
<b>want[0] = true;</b>	
<b>while (want[1]) ;</b>	
<b>...enter...</b>	<b>want[1] = true;</b>
	<b>while (want[0]) ;</b>
	<b>while (want[0]) ;</b>
<b>want[0] = false;</b>	<b>while (want[0]) ;</b>
	<b>...enter...</b>

How about progress?

# Failing “Progress”

<i>Thread 0</i>	<i>Thread 1</i>
<b>want[0] = true;</b>	
	<b>want[1] = true;</b>
<b>while (want[1]) ;</b>	
	<b>while (want[0]) ;</b>

It works for every ***other*** interleaving!

# “Peterson's Solution” (1981)

(“Taking turns when necessary”)

```
boolean want[2] = {false, false};  
int turn = 0;  
  
want[i] = true;  
turn = j;  
while (want[j] && turn == j)  
    continue;  
...critical section...  
want[i] = false;
```

# Proof Sketch of Exclusion

**Assume contrary: two threads in critical section**

**Both in c.s. implies  $\text{want}[i] == \text{want}[j] == \text{true}$**

**Thus both while loops exited because “turn != j”**

**Cannot have ( $\text{turn} == 0 \ \&\& \text{turn} == 1$ )**

- So one exited first

**w.l.o.g., T0 exited first because “turn ==1” failed**

- So  $\text{turn} == 0$  before  $\text{turn} == 1$
- So T1 had to set  $\text{turn} == 0$  before T0 set  $\text{turn} == 1$
- So T0 could not see  $\text{turn} == 0$ , could *not* exit loop first!

# Proof Sketch Hints

**want[i] == want[j] == true**

“want[]” fall away, focus on “turn”

**turn[] vs. loop exit...**

What really happens here?

<i>Thread 0</i>	<i>Thread 1</i>
<b>turn = 1;</b>	<b>turn = 0;</b>
<b>while (turn == 1);</b>	<b>while (turn == 0);</b>

# Bakery Algorithm

## More than two processes?

- Generalization based on bakery/deli counter
  - Get monotonically-increasing ticket number from dispenser
  - Wait until monotonically-increasing “now serving” == you
    - » You have lowest number  $\Rightarrow$  all people with smaller numbers have already been served

## Multi-process version

- Unlike “reality”, two people can get the same ticket number
- Sort by “ticket number with tie breaker”:
  - (ticket number, process number) tuple

# Bakery Algorithm

## Phase 1 – Pick a number

- Look at all presently-available numbers
- Add 1 to highest you can find

## Phase 2 – Wait until you hold *lowest* number

- Not strictly true: processes may have same number
- Use process-id as a tie-breaker
  - (ticket 7, process 99) > (ticket 7, process 45)
- Your turn when you hold lowest (t,pid)

# Bakery Algorithm

```
boolean choosing[n] = { false, ... };  
int number[n] = { 0, ... } ;
```

# Bakery Algorithm

## Phase 1: Pick a number

```
choosing[i] = true;
```

```
number[i] =  
    max(number[0], number[1], ...) + 1;
```

```
choosing[i] = false;
```

**Worst case: everybody picks same number!**

**But at least *next wave* of arrivals will pick a larger number...**

# Bakery Algorithm

## Phase 2: Sweep “proving” we have lowest number

```
for (j = 0; j < n; ++j) {  
    while (choosing[j])  
        continue;  
    while ((number[j] != 0) &&  
        ((number[i], i) > (number[j], j)))  
        continue;  
    }  
    ...critical section...  
    number[i] = 0;
```

# Summary

**Memory is *weird***

**Two fundamental operations - understand!**

- *Brief exclusion* for atomic sequences
- *Long-term yielding* to get what you want

**Three necessary critical-section properties**

**Understand these “exclusion algorithms” (which are also race-condition parties)**

- Two-process solution
- N-process “Bakery Algorithm”