

Graph Sparsifiers

Goal Input: $G = (V, E, w)$ where $|E| \gg |V|$

Output: $H = (V, E', w')$ where $|E'| \approx |V|$

s.t. H similar G .

Question: What is similar?

Ans: Cut preserving. $\forall S \subseteq V$

$$|\{w_{ij} \mid i \in S \& j \notin S\}| \approx |\{w'_{ij} \mid i \in S \& j \notin S\}|$$

Spectral Preserving

$$(1-\epsilon)G \preceq H \preceq (1+\epsilon)G$$

HW Show Spectral \Rightarrow Cut Preserving!

Random Sampling

$$G = (V, E, W)$$

edge	e_1	\dots	e_m	$\sum p_i = 1$
prob	p_1	\dots	p_m	

Sample: 1) Pick an edge e_i using probs p_i
 2) Return $(1/p_i)e_i$

K-Sample: 1) Pick Sample₁, ..., Sample_k (replacement)
 2) $H = \{ \text{Sample}_1 + \dots + \text{Sample}_k \}$
 Return H/k

Question: What p_1, \dots, p_m , and k work?

Sample: can be viewed as matrix!
 K-Sample: Sum of matrix!

Chernoff Bounds

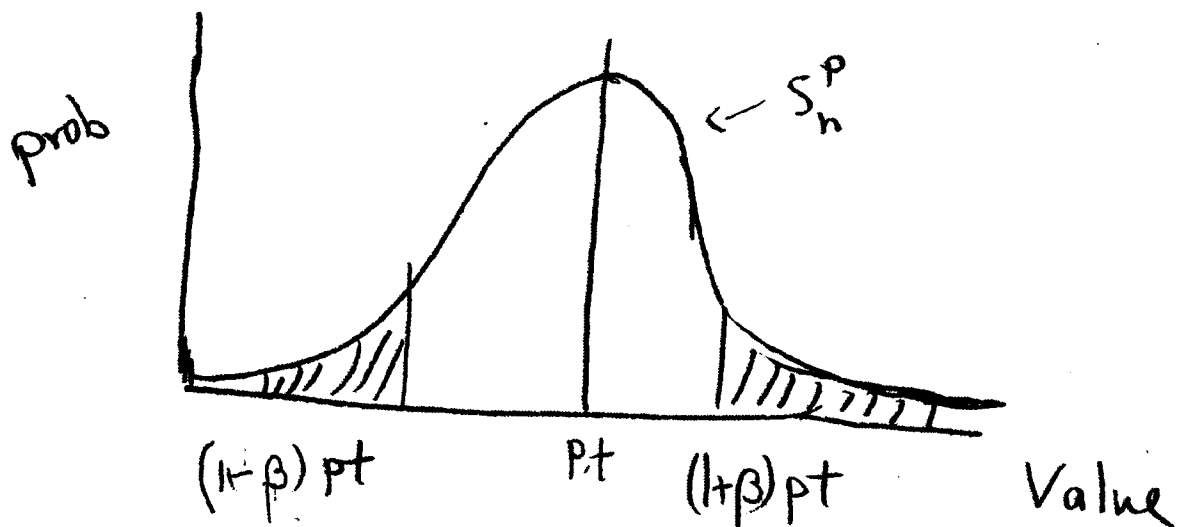
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Let X_1, \dots, X_t be independent 0/1 random variables

Assume $\text{Prob}(X_i=1) = p$

The binomial random variable is

$$S_n^p = X_1 + \dots + X_t$$



$$\text{Expect}(S_t^p) = \sum E(X_i) = p \cdot t$$

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Thm $\text{Prob}(S_t^P < (1-\beta)pt) < e^{-\beta^2 pt/2} \quad \forall 0 \leq \beta < 1$

Thm $\text{Prob}(S_t^P > (1+\beta)pt) < e^{-\beta^2 pt/2} \quad \forall 0 \leq \beta < 1$

Ahlsvede-Winter Thm

Note $\min_{\mu} (A \leq \mu B)$ is $\sigma(A/B)$

(AW Thm) Let Y_1, \dots, Y_k be iid random SPS_d $n \times n$ matrices
 at 1) $E(Y_i) = Z$ & 2) $Y_i \leq \mu Z$ then

$$P_c \left[(1-\epsilon)Z \leq \sum_{i=1}^k \frac{Y_i}{k} \leq (1+\epsilon)Z \right] \geq 1 - 2n e^{-c\epsilon^2 k/\mu}$$

$C \equiv 1/4$ works!

Lets do some examples

$$1) Z = I \text{ \& } E_i = \begin{pmatrix} 0 & & & 0 \\ & \ddots & & \\ & & 1 & \\ & & & \ddots \\ 0 & & & & 0 \\ \vdots & & & & \\ 0 & & & & 0 \end{pmatrix}$$

$Y \equiv$ pick random $i \in [1, \dots, n]$ return $n E_i$

$$E(Y) = \sum_{i=1}^n P(i) n E_i = \sum E_i = I = Z$$

$$\sigma(n E_i / I) = n \sigma(E_i / I) = n$$

Goal $2n e^{-c \epsilon^2 K/n} < \delta$

pick $K = c'n \log n$

$$2n e^{-c \epsilon^2 K/n} = 2n (e^{\log n})^{-c'} \approx n^{c''-1}$$

Ex 2) $Z = LG = L$ $H_{ij} = \begin{pmatrix} 1 & -1 \\ -1 & 1 \end{pmatrix} = H_e$

$$L = \sum_{e \in E} H_e \quad (\text{unweighted case})$$

$Y \equiv$ pick random $e \in E$, return $m H_e$

$$E(Y) = \sum_{i=1}^m p(i) m H_{e_i} = \sum H_{e_i} = L$$

$$V(m H_e / L) = m V(H_e / L) = m E R_e \text{ in } G. \quad (*)$$

If e is an isthmus then $E R_e = 1$

Thus $\mu = m$ and AW will require $O(m \log m)$ samples.

Picking edges uniformly is not good!

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Let's try picking proportional to ER_e !

Question: What is $\sum_{e \in E} ER_e$ (G connected)

Claim $\sum_{e \in E} ER_e = n-1$ (Trees done!)

Foster: $\sum_{e \in E} w(e) ER_e = n-1$

(We will prove later.)

$Y \equiv$ pick an edge e with prob = $\frac{w(e)ER_e}{n-1} = p_e$

return $\frac{w(e)}{p_e} H_e = \frac{(n-1)w(e)}{w(e)ER_e} H_e = \frac{(n-1)}{ER_e} H_e$

$E(Y) = \sum_{e \in E} p_e \left(\frac{w(e)}{p_e} \right) H_e = \sum_{e \in E} w_e H_e = L$

$$\nabla \left(\frac{n-1}{ER_e} H_e/L \right) = \frac{(n-1)}{ER_e} \nabla (H_e/L) = \frac{n-1}{ER_e} ER_e = n-1$$

$$\nabla (Y/L) \leq n$$

Thus AW gives concentration at

$$K = c' n \log n / \epsilon^2$$

Note $L_k = \sum_{i=1}^k Y_i / k$ is a Laplacian $\leq k$ edges.

Thm (SS) If H is obtained from G

by sampling proportional to $w_e ER_e$ $O(n \log n)$

times then WHP $(1-\epsilon)G \preceq H \preceq (1+\epsilon)G$.

Questions

1) \exists linear #edge spectral sparsifiers?
(yes)

2) Can we find $O(n \log n)$ spectral sparse quickly?
(yes)

3) Proof of AW?
(to do!)

4) Do we need exact ER's?
(no)

5) Proof of Foster Thm?
(to do!)

What if we use upper estimates of ER_e !

Given $P_e \geq ER_e \forall e \in E$

Set $t = \sum w_e P_e \geq n-1$

$P'_e = \frac{w_e P_e}{t}$ (probabilities)

$Y \equiv$ Pick e with prob P'_e ; return $\frac{w_e H_e}{P'_e} = \frac{t}{P_e} H_e$

Clear $E(Y) \equiv L$

$$\nabla(Y_i/L) = \max_e \nabla\left(\frac{t}{P_e} H_e/L\right) =$$

$$= \max_e \left(\frac{t}{P_e} ER_e\right) \leq t \quad (P_e \geq ER_e)$$

AW \Rightarrow Concentration: for $k = O(t \log t)$

Thus $P_e \geq ER_e$ work but may get more edges!