

Counting & FindingSpanning Trees

$G \equiv$ directed graph

$H \subseteq G$ is divergent ST with root r if

- 1) r can reach all of $V(G)$ in H .
 - 2) $\text{indegree}(v) = 1$ for $v \neq r$.
 - 3) $\text{indegree}(r) = 0$
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Def D is in-degree matrix of G if

$$D(i, j) = \begin{cases} \text{din}(i) & \text{if } i=j \\ -k = -\# \text{ edge from } i \text{ to } j & i \neq j \end{cases}$$

Lemma $G = (V, E)$ is a divergent tree iff

$$1) D(i, i) = \begin{cases} 0 & \text{if } i=r \\ 1 & \text{o.w} \end{cases}$$

$$2) \det \text{ of minor of } D_{r,r} \text{ (removing } r\text{th col \& row)} = 1$$

(\Rightarrow) 1) clear

2) After permuting row & col

D is upper Δ & $r=1$

$$\det(D_{r,r}) = 1$$

(\Leftarrow) suppose false.

$\therefore G$ must contain a disconnected component C .
Each node of C has in degree 1 and zero col sums.

Thus $\det(D(C)) = 0$.

$\Rightarrow \det(D(G)) = 0$ contra!

Thm # divergent ST rooted at $r \equiv \det(D_{r,r})$

Pf WLOG assume $r = V_1$

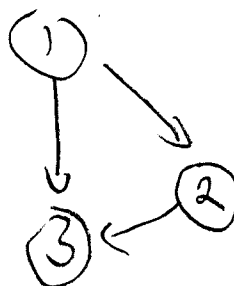
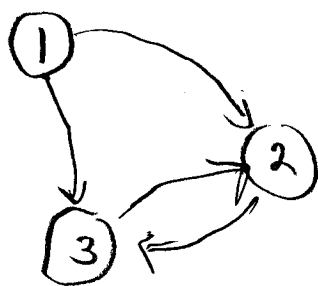
Note: Col sums of D are zero.

Start by replacing first col with $\begin{pmatrix} 1 \\ 0 \\ \vdots \\ 0 \end{pmatrix}$

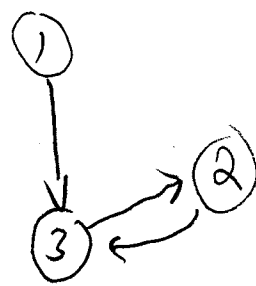
$$\det(D) = \det(D_{1,1})$$

Expand col 2 into subgraphs with fix edge into V_2 .

eg



G_1



G_2

$$D = \begin{pmatrix} 1 & -1 & -1 \\ 0 & 2 & -1 \\ 0 & -1 & 2 \end{pmatrix}$$

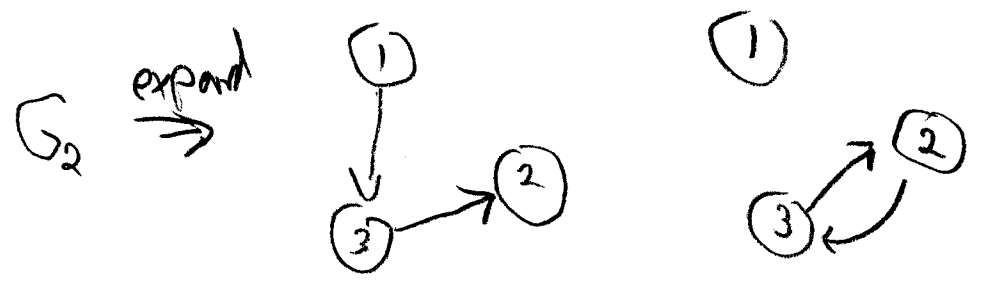
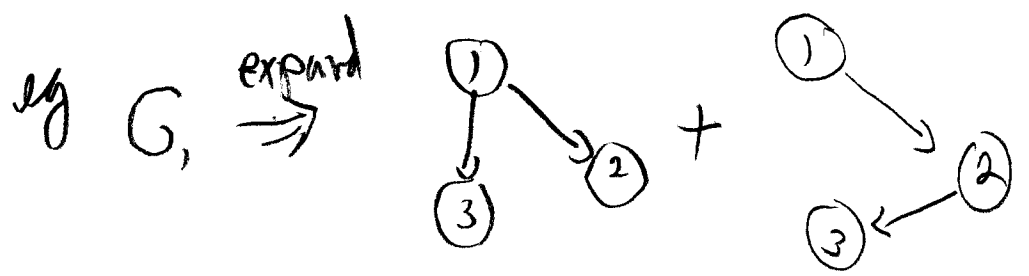
$$|D| = \begin{vmatrix} 1 & -1 & -1 \\ 0 & 1 & -1 \\ 0 & 0 & 2 \end{vmatrix} + \begin{vmatrix} 1 & 0 & -1 \\ 0 & 1 & -1 \\ 0 & -1 & 2 \end{vmatrix}$$

$G_1 \qquad G_2$

For matrix expand next col.

Repeat!

We get exact one matrix per indeg 1 subgraphs.



Each term in $\{1\}$ if \equiv divergent tree
 $\{0\}$ o.w.



The undirected Case

$G = (V, E)$ undirected

Let $G' = (V, E')$ where each edge $\{i, j\}$ is viewed as (i, j) & (j, i)

Claim $\#ST(G) = \#ST_r(G')$

$f: ST(G) \rightarrow ST_r(G')$



1-1 & onto

Note $D(G') \equiv \text{Laplacian of } G$