

# Bounding Eigenvalues Cont

9/18/13  
Spectral

Thm (Courant-Fischer)  $A$   $n \times n$  sym with  
eigen  $\mu_1 \geq \dots \geq \mu_n$  then

$$\mu_k = \max_{\substack{S \subseteq \mathbb{R}^n \\ \dim(S) = k}} \min_{x \in S} \frac{x^T A x}{x^T x} = \min_{\substack{T \subseteq \mathbb{R}^n \\ \dim(T) = n-k+1}} \max_{x \in T} \frac{x^T A x}{x^T x}$$

eg  $\mu_1$  using  $S$   $Ax = \mu_1 x$

Suppose max is me & min is you

I would pick  $S = \langle x \rangle$  you would be forced to pick  $\alpha x$ .

pf For case  $S$ . ( $\geq$ )

pick  $S = \langle x_1, \dots, x_k \rangle$  let  $x = \sum c_i x_i$

$$\frac{x^T A x}{x^T x} = \frac{\sum \mu_i c_i^2}{\sum c_i^2} \leq \frac{\sum \mu_k c_i^2}{\sum c_i^2} = \mu_k$$

( $\Leftarrow$ ) To show  $\forall S \dim(S) = k$

$$\min_{x \in S} \frac{x^T A x}{x^T x} \leq \mu_k$$

pf Let  $T = \langle x_k, \dots, x_n \rangle \dim(T) = n - k + 1$

$\circ \Rightarrow T \cap S \neq \emptyset$

$$\min_{x \in S} \frac{x^T A x}{x^T x} \leq \min_{x \in S \cap T} \frac{x^T A x}{x^T x}$$

Thus  $x = \sum_{i=k}^n c_i x_i$

$$\frac{x^T A x}{x^T x} = \frac{\sum_k \mu_i c_i^2}{\sum c_i^2} \leq \frac{\sum_k \mu_k c_i^2}{\sum c_i^2} = \mu_k$$

□

~~Thus eigens for  $K_n$  are  $0, n, \dots, n$ .~~  
~~i.e.  $\forall x \perp \mathbf{1}$   $L_{K_n} x = nx$ .~~

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New notation:  $A \succeq 0$  if  $\forall x \ x^T A x \geq 0$   $A$  sym.

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In general  $A \preceq B$  if  $\forall x \ x^T A x \leq x^T B x$   
 i.e.  $B - A \succeq 0$

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note  $A \preceq B$  &  $B \preceq C$  then  $A \preceq C$

$$A \preceq B \Rightarrow A + C \preceq B + C$$

Sym  $A, B, C$ .

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Graphs  $G$  &  $H$   $G \preceq H$  if  $L_G \succeq L_H$

Claim  $G \preceq G \vee H$ .

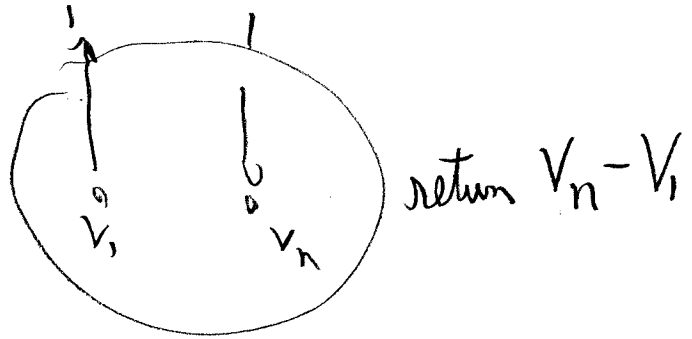
Thm  $A \preceq B \Rightarrow \lambda_k(A) \leq \lambda_k(B)$

$$\begin{aligned} \text{Pf (CF)} \quad \lambda_k(B) &= \min_{\dim(\mathcal{T})=k, \mathcal{T} \subseteq \mathbb{F}} \max_{x \in \mathcal{T}} \frac{x^T A x}{x^T x} \\ &\geq \min_{\dim(\mathcal{T})=k, \mathcal{T} \subseteq \mathbb{F}} \max_{x \in \mathcal{T}} \frac{x^T B x}{x^T x} \\ &= \lambda_k(B) \end{aligned}$$

$G$  graph &  $G_{1n} \equiv$  

Claim  $\alpha = \min_{\alpha} (G_{1n} \preceq \alpha G)$  iff  $ER_{1n} = \alpha$

Recall  $ER_{1n}$  test



ie solve  $LV = \begin{pmatrix} +1 \\ 0 \\ \vdots \\ 0 \\ -1 \end{pmatrix}$  return  $v_n - v_1$

$$\sum_m \begin{pmatrix} +1 \\ 0 \\ \vdots \\ 0 \\ -1 \end{pmatrix} \quad ER_{1n} = \sum_m^T L^+ \sum_m$$

potential flow  $Cd^T v = f$

$$ER_{1n} = f^T R f = (Cd^T v)^T R Cd^T v = v^T d^T C R C d^T v \\ = v^T L v$$

What is  $v^T L_{G_m} v$ ?  $= (v_n - v_1)^2 = ER_{1n}^2$

Claim  $G_m \preceq \alpha G \Rightarrow ER_m \preceq \alpha$

$f$  unit potential flow with potential  $v$

$$ER_m = f^T R f = v^T L v \geq \frac{v^T G_m v}{\alpha} = \frac{(ER_m)^2}{\alpha}$$

$$\Rightarrow \alpha \geq ER_m$$

(HW)  $ER_m \leq \alpha \Rightarrow G_m \preceq \alpha G$

# Upper bounding $\lambda_2$

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$$(CF) \quad \lambda_2 = \min_{v^T \mathbf{1} = 0} \frac{v^T L v}{v^T v}$$

$$\forall x \text{ st } x \perp \mathbf{1} \quad \lambda_2 \leq \frac{v^T L x}{x^T x}$$

\*8

Lower bding  $\lambda_2$  !

$K_n \equiv$  Complete graph

$$L_{K_n} = \begin{pmatrix} n-1 & & & -1 \\ & n-1 & & -1 \\ & & \ddots & -1 \\ -1 & & & n-1 \end{pmatrix}$$

$$\begin{pmatrix} n-1 & -1 \\ -1 & n-1 \\ \vdots & -1 \\ -1 & \vdots \end{pmatrix} \begin{pmatrix} 1 \\ -1 \\ 0 \\ \vdots \end{pmatrix} = \begin{pmatrix} n \\ -n \\ 0 \\ \vdots \end{pmatrix} = n \begin{pmatrix} 1 \\ -1 \\ 0 \\ \vdots \end{pmatrix}$$

$$x_1 = \begin{pmatrix} 1 \\ -1 \\ 0 \\ \vdots \end{pmatrix} \quad x_2 = \begin{pmatrix} 0 \\ 1 \\ -1 \\ \vdots \end{pmatrix} \quad \dots \quad x_{n-1} = \begin{pmatrix} 0 \\ \vdots \\ 1 \\ -1 \end{pmatrix}$$

$x_1, \dots, x_{n-1}$  are  $n-1$  indep eigenvector value  $n$ .

$$\lambda(K_n) = \{0, n, \dots, n\}$$

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Let  $P_n$  be unit weight path graph.

$$X = (-\frac{n}{2}, -\frac{n}{2}+1, \dots, +\frac{n}{2})$$

$$X^T L(P_n) X = \sum (x_i - x_{i+1})^2 = n-1$$

$$X^T X = 2 \sum_{i=1}^{\frac{n}{2}} i^2 \approx 2 \cdot \frac{(\frac{n}{2})^3}{3} = \frac{n^3}{12}$$

$$\lambda_2 \leq \frac{n-1}{\frac{n^3}{12}} \leq \frac{12}{n^2}$$

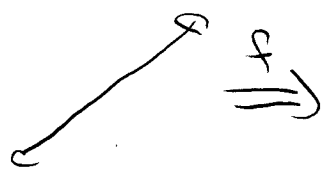
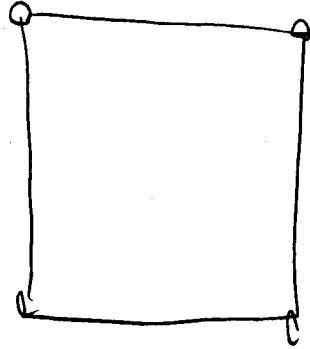
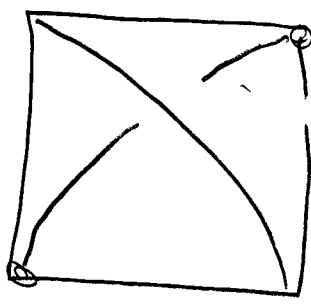
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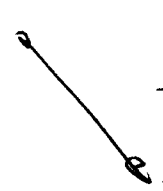
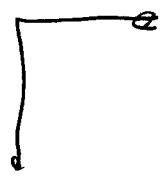
Def  $G$  (guest),  $H$  (host)

Path Embedding:  $f: E_G \rightarrow \text{Paths in } H$

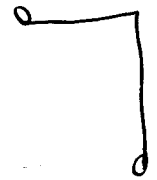
ex  $G = K_4$   $H = C_4$



$\Rightarrow$



$\Rightarrow$



$$\text{Congestion} = \max_{e \in H} |\{P \mid e \in P\}|$$

$$\text{Dilation} = \max_P \{|P|\}$$

In example: Cong = 3 & Dil = 2

## Mediant of Fractions

Fractions  $\left( \frac{a_1}{b_1}, \dots, \frac{a_n}{b_n} \right)$

$$\text{mediant} \equiv \frac{\sum a_i}{\sum b_i}$$

Claim:  $\frac{\sum a_i}{\sum b_i} \leq \max_i \frac{a_i}{b_i}$      $a_i > 0$      $b_i > 0$

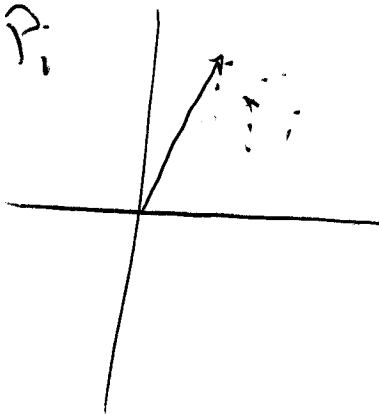
prelim proof

Consider points in  $\mathbb{R}^2$   $P_1 (b_1, a_1), \dots, P_n (b_n, a_n)$

Let  $\bar{P}$  average or center of mass.

$$\text{Slope}(P_i) = \frac{a_i}{b_i} \quad \text{Slope}(\bar{P}) = \frac{\sum a_i}{\sum b_i}$$

$$\text{Slope}(\bar{P}) \leq \max \text{Slope } P_i$$



# Laplacians as sums of Laplacians

$$E_{ij} = \begin{pmatrix} & i & j \\ \dots & & \\ 0 & 1 & -1 & \dots \\ \dots & & & \\ 0 & -1 & 1 & \dots \\ \dots & & & \end{pmatrix} \begin{matrix} i \\ j \end{matrix}$$

Def A is diagonal dominant if  $A_{ii} \geq \sum_{j \neq i} |A_{ij}|$

Claim  $L = L(G)$  then  $L = \sum_{(i,j) \in E} w_{ij} E_{ij}$

Claim A is SDD then  $A = \sum_{\substack{i \neq j \\ A_{ij} < 0}} -A_{ij} E_{ij} + \sum_{\substack{A_{ij} > 0 \\ i \neq j}} A_{ij} P_{ij} + D'$

$$P_{ij} = \begin{pmatrix} & i & j \\ \dots & & \\ & 1 & \\ \dots & & \\ & & 1 \\ \dots & & \end{pmatrix} \begin{matrix} i \\ j \end{matrix} \quad \& \quad D' \text{ is non-neg diagonal.}$$

Thm  $f: G \rightarrow H$  with congestion  $c$  & dilation  $d$

then  $cdL_H \geq L_G$

pt  $L(G) = G \quad L(H) = H$

$E_1, \dots, E_m$  edge subgraphs of  $G$

$P_1, \dots, P_m$  the path of  $H \quad \sum P_i \leq c \cdot H$

$$\frac{x^T G x}{x^T H x} \leq \frac{x^T \sum E_i x}{(1/c)x^T \sum P_i x} = \frac{c \sum x^T E_i x}{\sum x^T P_i x} \leq c \cdot \text{Max} \frac{x^T E_i x}{x^T P_i x}$$

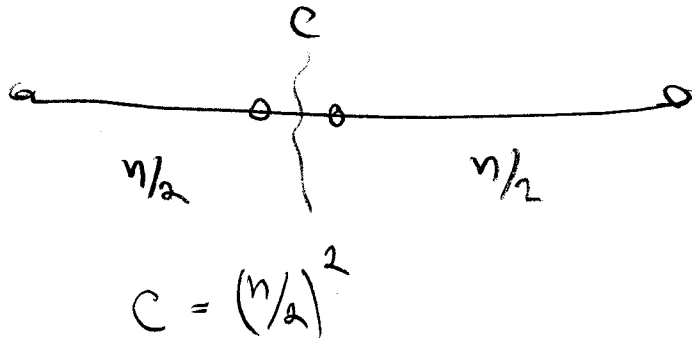
$$\leq c \cdot d$$

$$\text{ie } c \cdot d x^T H x \geq x^T G x$$

Back to  $\lambda_2$  for  $P_n$

f.o.  $K_n \xrightarrow{\text{path}} P_n$

Congestion



dilation  $\equiv n-1$

By Thm:  $\lambda_2(K_n) \leq c \cdot d \cdot \lambda_2(P_n)$

$$\frac{n}{c \cdot d} \leq \lambda_2(P_n)$$

$$\frac{n}{\binom{n}{2}(n-1)} = \frac{4}{n(n-1)} \leq \lambda_2(P_n)$$