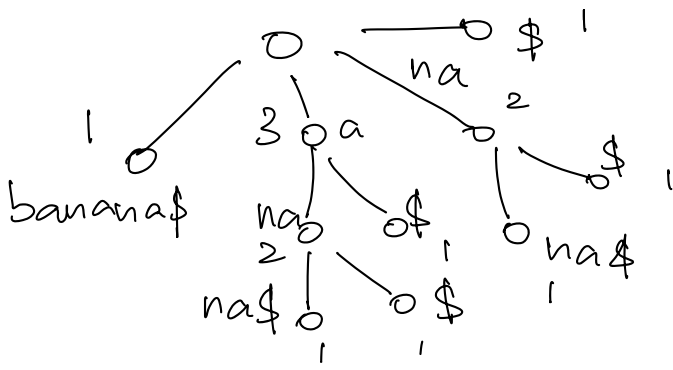


Recall: last time we talked about suffix & suffix tree.



Given string $s = s[0]s[1] \dots s[n-1]$ over Σ
 Suffix: $S_i = s[i:n] = s[i]s[i+1] \dots s[n-1]$

It helps solve problems:

- pattern matching: find P to T
- longest common substring

Now, the question is:

How to construct suffix tree?

naive solution: create the "complete" tree then compress.

more efficient solution: suffix array.

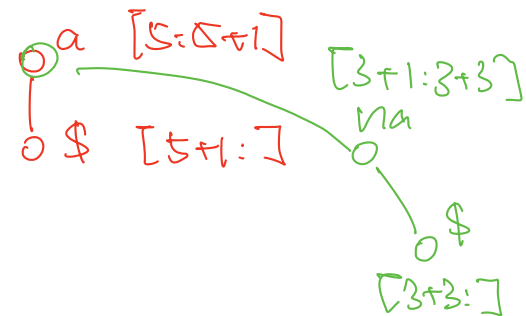
↳ Def: An array of integers $SA[0 \dots n-1]$ that contains a permutation of $[0, n-1]$ such that $SA[0] < SA[1] < \dots < SA[n-1]$

i.e. $SA[i] = j \rightarrow$ "the i -th smallest suffix starts at pos. j in the

original string.
 Example: b a n a n a \$
 0 1 2 3 4 5 6

sorted suffix strings ($\$ < \Sigma$)

length of common prefix	idx	Suffix
0	6	\$ → shared
1	5	a \$ → not shared
3	3	a na \$ → not shared
0	1	a na na \$ → shared with next
0	0	b a n a n a \$
0	4	n a \$
2	2	n a na \$

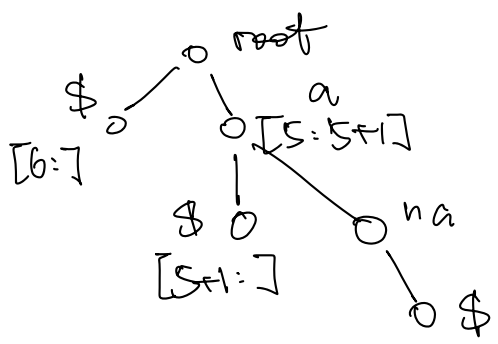


suffix array:

6 5 3 1 0 4 2

LCP: - - -

2. add nodes in order

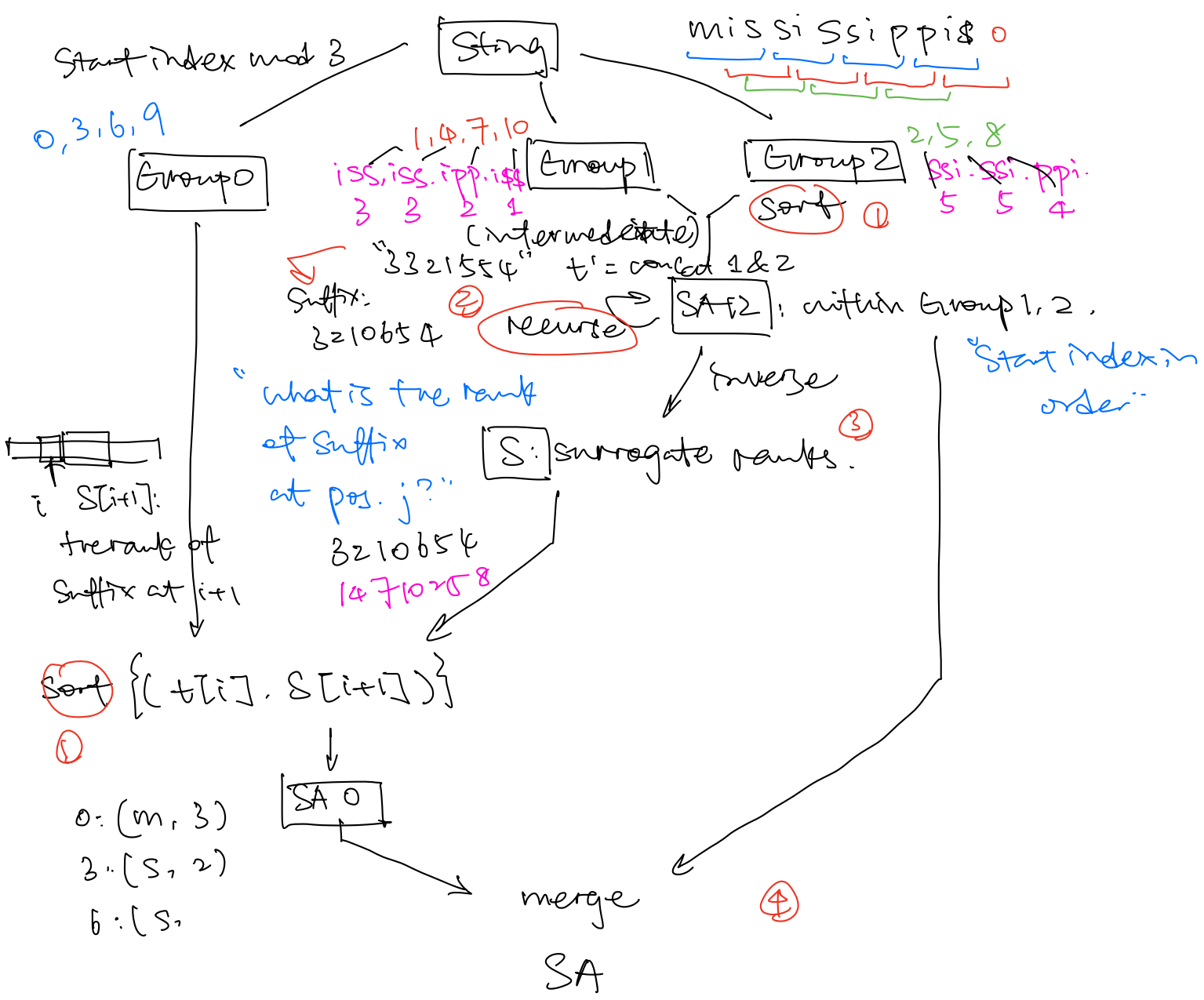


$O(n)$

Suffix tree $\leftrightarrow$ Suffix array
 $O(n)$

String $\xleftarrow{O(n)}$?

Skew Algorithm / "PS3"
(intuition)



1. Sort - Radix Sort

1. Sorting Group 0. $\{(char, int)\}$.

problem: sort a list of length n of pairs (x_i, y_i)

where $0 \leq x_i, y_i \leq n-1$.

algo: 1. create n buckets. put the pairs in buckets by y_i .

2. process buckets in increasing order, putting the new pairs into new buckets by x_i .

3. read off the sorted list

eg: $(3,1)(1,5)(5,5)(5,3)(3,1)(1,5)(5,0)$

Step 1:

0	:	$(5,0)$
1	:	$(3,1)$ $(3,1)$
2	:	
3	:	$(5,3)$
4	:	
5	:	$(1,5)$ $(5,5)$ $(1,5)$
6	:	

Step 2:

0	:	
1	:	$(1,5)$ $(1,5)$
2	:	$(3,1)$ $(3,1)$
3	:	
4	:	
5	:	$(5,0)$ $(5,3)$ $(5,5)$
6	:	

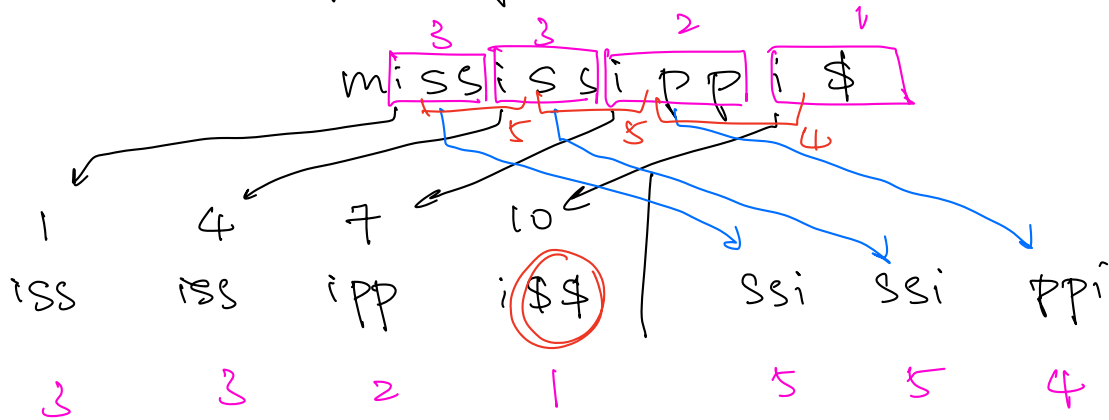
Sorted

2. Sort 3-charac list

$\{(s[i], s[i+1], s[i+2])\}$

2. encoded $t \rightarrow$ suffix array

claim: t 's suffix array is SA12.



$$\begin{array}{llll}
 \varphi. & i \bmod 3 = 1 & j \bmod 3 = 0 & S_i = (t[i] \cdot S_{i+1}) \\
 & 2 & 0 & (S[i], S[i+1], \\
 & & = & S_{i+2}) \\
 & & & \uparrow \\
 & & & \text{Group 1}
 \end{array}
 \quad
 \begin{array}{ll}
 S_j = (t[j] \cdot S_{j+1}) \\
 (S[j], S[j+1], \\
 S_{j+2}) \\
 \uparrow \\
 \text{Group 2}
 \end{array}$$

Runtime $T(n) = O(n) + T(\lceil 2n/3 \rceil)$
 recursive calls. $\uparrow$
 (2).

$2n/3 \rightarrow$ There can be no more lexicographic names than characters in $S[2]$, so the alphabet size in a recursive call never exceeds the size of the string.