## 15-750:Algorithms in the Real World

## Data Compression:

Linear Transform Coding
(for both lossless and lossy compression)

## 5. Linear Transform Coding

Goal: Transform the data into a form that is easily compressible (through lossless or lossy compression)

Select a set of linear basis functions $\phi_{i}$ that span the space

- sin, cos, spherical harmonics, wavelets, ...


## Example: Cosine Transform



## Other Transforms

Polynomial:


Wavelet (Haar):


## How to Pick a Transform

## Goals:

- Decorrelate the data
- Low coefficients for many terms
- Basis functions that can be ignored from the perception point-of-view


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## Quantization (lossy)

## Scalar Quantization

Quantize regions of values into a single value
E.g. Drop least significant bit (Can be used to reduce \# of bits for a pixel)

Q: Why is this lossy?
Many-to-one mapping

Two types

- Uniform: Mapping is linear
- Non-uniform: Mapping is non-linear


## Scalar Quantization


uniform

non uniform

Q: Why use non-uniform?
Error metric might be non-uniform.
E.g. Human eye sensitivity to specific color regions

Can formalize the mapping problem as an optimization problem

## Vector Quantization

Mapping a multi-dimensional space into a smaller set of messages


Encode
Decode

## Vector Quantization (VQ)

What do we use as vectors?

- Color (Red, Green, Blue)
- Can be used, for example to reduce 24bits/pixel to 8bits/pixel
- Used in some monitors to reduce data rate from the CPU (colormaps)
- K consecutive samples in audio
- Block of K pixels in an image

How do we decide on a codebook

- Typically done with clustering

VQ most effective when the variables along the dimensions of the space are correlated

## Vector Quantization: Example



## Case Study: JPEG

A nice example since it uses many techniques:

- Transform coding (Cosine transform)
- Scalar quantization
- Residual coding
- Run-length coding
- Huffman or arithmetic coding


## 15-750: Algorithms in the Real World

Algorithms for coding
(Error Correcting Codes)

$$
\begin{aligned}
& \text { Welc**e } t^{*} t^{*} \text { is clas* } o^{*} c^{*} d^{*} n g \\
& \quad y^{*} u \text { a** in } f^{*} r \text { a } f^{*} n \text { rid*! }
\end{aligned}
$$

What do these sentences say?

## Why did this work?

## Redundancy!

Codes are clever ways of judiciously adding redundancy to enable recovery under "noise".

## General Model


codeword' (c')
decoder
message or error
"Noise" introduced by the channel:

- changed fields in the codeword vector (e.g. a flipped bit).
- Called errors
- missing fields in the codeword vector (e.g. a lost byte).
- Called erasures

How the decoder deals with errors and/or erasures?

- detection (only needed for errors)
- correction


## Applications

Numerous applications:
Some examples

- Storage: Hard disks, cloud storage, NAND flash...
- Wireless: Cell phones, wireless links,
- Satellite and Space: TV, Mars rover, ...

Reed-Solomon codes are by far the most used in practice.

Low density parity check codes (LDPC) codes used for 4G (and 5G) communication and NAND flash

## Block Codes

symbols (e.g., bits)

message I

message 2

Other kind: convolutional codes (we won't cover it)...

## Block Codes


message or error

- Each message and codeword is of fixed size
- Notation:

"length of the code"
length of the codeword

C = "code" = set of codewords

## Simple Examples

3-Repetition code: $k=1, n=3$

| Message |  | Codeword |
| :--- | :--- | :---: |
| 0 | -> | 000 |
| 1 | -> | 111 |

- How many erasures can be recovered?
- How many errors can be detected?
- Up to how many errors can be corrected?

Errors are much harder to deal with than erasures.
Why?

