

UNIT 4B

Iteration: Sorting

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1

Sorting

The collage illustrates sorting in different contexts:

- Windows 'Sort' dialog box:** Shows options to sort by 'DESCRIPTION' (Ascending/Descending) and 'Then by' (Ascending/Descending).
- Music Library Table:** A table with columns 'Name', 'Artist', and 'Time'. The 'Time' column is highlighted, and a mouse cursor is over the 'Time' header.
- YouTube Search Results:** Shows search results for 'amd'. The 'Search options' dropdown is open, showing 'Result type' (All, Videos, Channels, Playlists) and 'Sort by' (Relevance, Upload date, View count, Rating). The 'Relevance' and 'Upload date' options are circled.

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2

Insertion Sort Outline

```
def isort(datalist):  
    result = []  
    for value in datalist:  
        # insert value in its  
        # proper place in result  
    return result
```

insert function

```
datalist.insert(position, value)
```

```
>>> a = [10, 30, 20]  
>>> a  
[10, 30, 20]  
>>> a.insert(0, 15)  
>>> a  
[15, 10, 30, 20]
```

insert function (cont'd)

```
>>> a.insert(2, 50)
>>> a
[15, 10, 50, 30, 20]
>>> a.insert(5, 40)
>>> a
[15, 10, 50, 30, 20, 40]
```

The insert function just inserts at the given index. It doesn't assume the list is sorted.

Insertion Sort, Refined

```
def isort(datalist):
    result = []
    for value in datalist:
        # compute place to insert
        result.insert(place, value)
    return result
```

How do we find the right place to insert?

gr_index

Compute the index of first element greater than item

```
def gr_index(datalist, item):  
    # precondition: datalist is sorted!  
    index = 0  
    while index < len(datalist) and datalist[index] < item:  
        index = index + 1  
    return index
```

Testing gr_index

```
>>> a = [10, 20, 30, 40, 50]  
>>> a  
[10, 20, 30, 40, 50]  
>>> gr_index(a, 3)  
0  
>>> gr_index(a, 14)  
1  
>>> gr_index(a, 37)  
3  
>>> gr_index(a, 99)  
5
```

Insertion Sort, Complete

```
def isort(datalist):  
    result = [ ]  
    for value in datalist:  
        place = gr_index(result, value)  
        result.insert(place, value)  
    return result
```

Debugging Insertion Sort

```
def isort(list):  
    result = [ ]  
    print(result)      # for debugging  
    for val in list:  
        place = gindex(result, val)  
        result.insert(place, val)  
        print(result)  # for debugging  
    return result
```

Testing isort

```
>>> isort([3, 1, 4, 1, 5, 9, 6, 2])  
[]  
[3]  
[1, 3]  
[1, 3, 4]  
[1, 1, 3, 4]  
[1, 1, 3, 4, 5]  
[1, 1, 3, 4, 5, 9]  
[1, 1, 3, 4, 5, 6, 9]  
[1, 1, 2, 3, 4, 5, 6, 9]  
=> [1, 1, 2, 3, 4, 5, 6, 9]
```

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11

Can We Do Better?

- isort doesn't change its input list.
- Instead it makes a new list, called result.
- This takes twice as much memory.
- Can we write a destructive ("in place") version of the algorithm that doesn't use extra memory?

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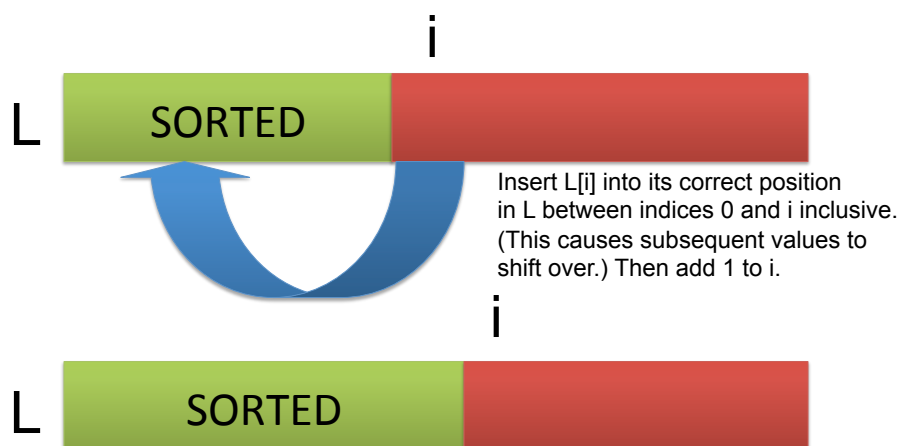
12

Destructive (In Place) Insertion Sort

Given a list L of length n , $n > 0$.

1. Set $i = 1$.
2. While i is not equal to n , do the following:
 - a. Insert $L[i]$ into its correct position in L between indices 0 and i inclusive.
 - b. Add 1 to i .
3. Return the list L which will now be sorted.

General Idea: Any one Iteration



Same basic picture!

Insertion Sort in Python ("in place")

```
def isort(datalist):  
    i = 1  
    while i < len(datalist):  
        x = datalist.pop(i)  
        index = 0  
        while index < i and datalist[index] < x:  
            index = index + 1  
        datalist.insert(index, x)  
        i = i + 1  
    return datalist
```

pop: remove the item at
position i in list
and store it in x

Look familiar?
This is essentially
gr_index, except it
stops at index i rather
than scanning the
whole list!

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15

Look Closer at Insertion Sort

Given a list L of length n, $n > 0$.

1. Set $i = 1$.
2. While i is not equal to n , do the following:

Precondition for each iteration: $L[0..i)$ is sorted

a. Insert $L[i]$ into its correct position in L between
index 0 and index i inclusive.

b. Add 1 to i .

Postcondition for each iteration: $L[0..i)$ is sorted

3. Return the list L which will now be sorted.

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16

Look Closer at Insertion Sort

Given a list L of length n , $n > 0$.

1. Set $i = 1$.
2. While i is not equal to n , do the following:
Loop invariant: $L[0..i)$ is sorted
 - a. Insert $L[i]$ into its correct position in L between index 0 and index i inclusive.
 - b. Add 1 to i .
3. Return the list L which will now be sorted.

A **loop invariant** is a condition that is true at the start and end of each iteration of a loop.

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17

Reasoning with the Loop Invariant

The loop invariant is true at the end of each iteration, including the last iteration. After the last iteration, when we go to step 3:

$L[0..i)$ is sorted (from the last iteration)

AND

i is equal to n (due to the while loop terminating)

These 2 conditions imply that $L[0..n)$ is sorted, but this range is the entire list, so the list must always be sorted when we return our final answer!

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18