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## Announcements

- Bring ear buds or headphones to lab!

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## Representing and Manipulating Data

### Last Unit

- How to represent data as a sequence of bits
- How to interpret bit representations
- Use of levels of abstraction in representing more complex information (music, pictures) using simpler building blocks (numbers)

### This Unit

- How sequences of bits are implemented using electrical signals, and manipulated by circuits
- Use of levels of abstraction in designing more complex computer components from simpler components

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## Foundations

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Boolean logic is the logic of digital circuits

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## Implementing Bits

- Computers compute by manipulating electricity according to specific rules.
- We associate electrical signals inside the machine with bits. Any electrical device with two distinct states (e.g. on/off switch, two distinct voltage or current levels) could implement our bits.
- The rules are implemented by electrical circuits.

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## Conceptualizing bits and circuits

- **ON or 1** : **True**
- **OFF or 0** : **False**
- **circuit behavior:**  
expressed in *Boolean logic* or *Boolean algebra*

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## Boolean Logic (Algebra)

Computer circuitry works based on **Boolean Logic** (Boolean Algebra) : operations on **True (1)** and **False (0)** values.

A	B	$A \wedge B$ (A AND B) (conjunction)	$A \vee B$ (A OR B) (disjunction)	A	$\neg A$ (NOT A) (negation)
0	0	0	0	0	1
0	1	0	1	1	0
1	0	0	1		
1	1	1	1		

- A and B in the table are Boolean variables, AND and OR are operations (also called functions).

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## Boolean Logic & Truth Tables

- Example:** You can think of  $A \wedge B$  below as  
*15110 is fun and 15110 is useful*  
 where A stands for the statement *15110 is fun*,  
 B stands for the statement *15110 is useful*.

A	B	$A \wedge B$ (A AND B) (conjunction)	$A \vee B$ (A OR B) (disjunction)	A	$\neg A$ (NOT A) (negation)
0	0	0	0	0	1
0	1	0	1	1	0
1	0	0	1		
1	1	1	1		

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## Logic gates

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the basic elements of digital circuits

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## Logic Gates

- A gate is a physical device that implements a Boolean operator by performing basic operations on electrical signals.
- Nowadays, gates are built from transistors.

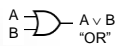


physical picture of gates

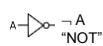
Physical behavior of circuits is beyond the scope of our course.



"AND"



"OR"



"NOT"

logical picture of gates

## AND, OR, NOT Gates

A	B	$A \wedge B$ (A AND B) (conjunction)	$A \vee B$ (A OR B) (disjunction)
0	0	0	0
0	1	0	1
1	0	0	1
1	1	1	1

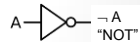
A	$\neg A$ (NOT A) (negation)
0	1
1	0



"AND"



"OR"



"NOT"

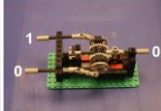
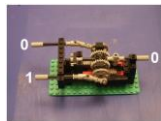
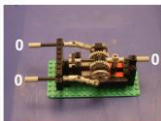
Truth tables define the input - output behavior of logic gates.

## A Mechanical Implementation

Push-pull logic AND gate

For an input pushed-in lever represents 1

For an output pushed-in lever represents 0



Source:  
randomwraith.com  
by Martin Howard

## Combinational circuits

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combinations of logic gates

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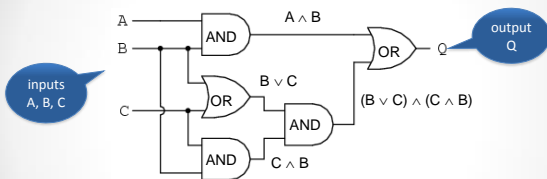
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## Combinational Circuits

The logic states of inputs at any given time determine the state of the outputs.



What is Q?  $(A \wedge B) \vee ((B \vee C) \wedge (C \wedge B))$

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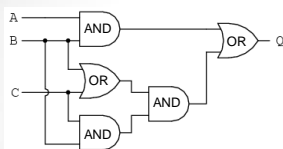
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## Truth Table of a Circuit



$Q = (A \wedge B) \vee ((B \vee C) \wedge (C \wedge B))$

A	B	C	Q
0	0	0	
0	0	1	
0	1	0	
0	1	1	
1	0	0	
1	0	1	
1	1	0	
1	1	1	

How do I know that there should be 8 rows in the truth table?

Describes the relationship between inputs and outputs of a device

[http://www.allaboutcircuits.com/vol\\_4/chpt\\_7/6.html](http://www.allaboutcircuits.com/vol_4/chpt_7/6.html)

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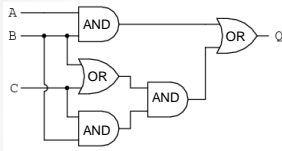
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## Truth Table of a Circuit



$$Q = (A \wedge B) \vee ((B \vee C) \wedge (C \wedge B))$$

A	B	C	Q
0	0	0	0
0	0	1	
0	1	0	
0	1	1	
1	0	0	
1	0	1	
1	1	0	
1	1	1	

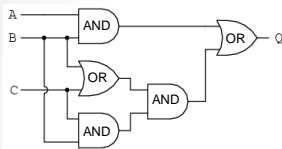
Describes the relationship between inputs and outputs of a device

[http://www.allaboutcircuits.com/vol\\_4/chpt\\_7/6.html](http://www.allaboutcircuits.com/vol_4/chpt_7/6.html)

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[illegible]

## Truth Table of a Circuit



$$Q = (A \wedge B) \vee ((B \vee C) \wedge (C \wedge B))$$

A	B	C	Q
0	0	0	0
0	0	1	0
0	1	0	
0	1	1	
1	0	0	
1	0	1	
1	1	0	
1	1	1	

Describes the relationship between inputs and outputs of a device

[http://www.allaboutcircuits.com/vol\\_4/chpt\\_7/6.html](http://www.allaboutcircuits.com/vol_4/chpt_7/6.html)

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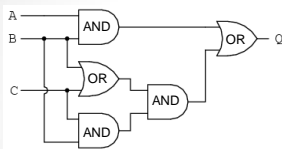
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## Truth Table of a Circuit



$$Q = (A \wedge B) \vee ((B \vee C) \wedge (C \wedge B))$$

A	B	C	Q
0	0	0	0
0	0	1	0
0	1	0	0
0	1	1	
1	0	0	
1	0	1	
1	1	0	
1	1	1	

Describes the relationship between inputs and outputs of a device

[http://www.allaboutcircuits.com/vol\\_4/chpt\\_7/6.html](http://www.allaboutcircuits.com/vol_4/chpt_7/6.html)

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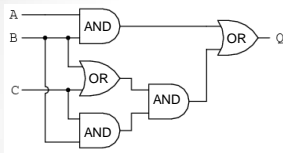
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## Truth Table of a Circuit



$$Q = (A \wedge B) \vee ((B \vee C) \wedge (C \wedge A))$$

A	B	C	Q
0	0	0	0
0	0	1	0
0	1	0	0
0	1	1	1
1	0	0	
1	0	1	
1	1	0	
1	1	1	

Describes the relationship between inputs and outputs of a device

[http://www.allaboutcircuits.com/vol\\_4/chpt\\_7/6.html](http://www.allaboutcircuits.com/vol_4/chpt_7/6.html)

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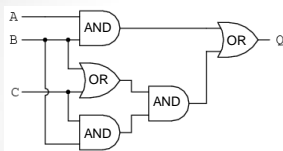
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## Truth Table of a Circuit



$$Q = (A \wedge B) \vee ((B \vee C) \wedge (C \wedge A))$$

A	B	C	Q
0	0	0	0
0	0	1	0
0	1	0	0
0	1	1	1
1	0	0	0
1	0	1	
1	1	0	
1	1	1	

Describes the relationship between inputs and outputs of a device

[http://www.allaboutcircuits.com/vol\\_4/chpt\\_7/6.html](http://www.allaboutcircuits.com/vol_4/chpt_7/6.html)

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## Describing Behavior of Circuits

- Boolean expressions
- Circuit diagrams
- Truth tables

} Equivalent notations

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## Manipulating circuits

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Boolean algebra and logical equivalence

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## Why manipulate circuits?

- The design process
  - simplify a complex design for easier manufacturing, faster or cooler operation, ...
- Boolean algebra helps us find another design guaranteed to have same behavior

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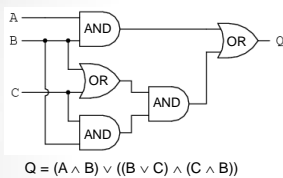
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## Logical Equivalence



A	B	C	Q
0	0	0	0
0	0	1	0
0	1	0	0
0	1	1	1
1	0	0	0
1	0	1	0
1	1	0	1
1	1	1	1

Can we come up with a simpler circuit implementing the same truth table?  
Simpler circuits are typically cheaper to produce, consume less energy etc.

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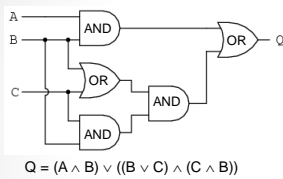
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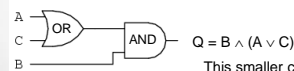
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## Logical Equivalence



A	B	C	Q
0	0	0	0
0	0	1	0
0	1	0	0
0	1	1	1
1	0	0	0
1	0	1	0
1	1	0	1
1	1	1	1



This smaller circuit is logically equivalent to the one above: they have the same truth table. By using laws of Boolean Algebra we convert a circuit to another equivalent circuit.

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## Laws for the Logical Operators $\wedge$ and $\vee$ (Similar to $\times$ and $+$ )

- Commutative:  $A \wedge B = B \wedge A$        $A \vee B = B \vee A$
- Associative:  $A \wedge B \wedge C = (A \wedge B) \wedge C = A \wedge (B \wedge C)$   
 $A \vee B \vee C = (A \vee B) \vee C = A \vee (B \vee C)$
- Distributive:  $A \wedge (B \vee C) = (A \wedge B) \vee (A \wedge C)$   
 $A \vee (B \wedge C) = (A \vee B) \wedge (A \vee C)$
- Identity:  $A \wedge 1 = A$        $A \vee 0 = A$
- Dominance:  $A \wedge 0 = 0$        $A \vee 1 = 1$
- Idempotence:  $A \wedge A = A$        $A \vee A = A$
- Complementation:  $A \wedge \neg A = 0$        $A \vee \neg A = 1$
- Double Negation:  $\neg \neg A = A$

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## Laws for the Logical Operators $\wedge$ and $\vee$ (Similar to $\times$ and $+$ )

- Commutative:  $A \wedge B = B \wedge A$        $A \vee B = B \vee A$
- Associative:  $A \wedge B \wedge C = (A \wedge B) \wedge C = A \wedge (B \wedge C)$   
 $A \vee B \vee C = (A \vee B) \vee C = A \vee (B \vee C)$
- Distributive:  $A \wedge (B \vee C) = (A \wedge B) \vee (A \wedge C)$   
 $A \vee (B \wedge C) = (A \vee B) \wedge (A \vee C)$
- Identity:  $A \wedge 1 = A$        $A \vee 0 = A$

.....

The A's and B's here are schematic variables! You can instantiate them with any expression that has a Boolean value:

$$\underbrace{(x \vee y)}_A \wedge \underbrace{z}_{B \wedge A} = z \wedge \underbrace{(x \vee y)}_A \quad (\text{by commutativity})$$

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## Applying Properties for $\wedge$ and $\vee$

Showing  $(x \wedge y) \vee ((y \vee z) \wedge (z \wedge y)) = y \wedge (x \vee z)$

Commutativity  $A \wedge B = B \wedge A$

$$(x \wedge y) \vee ((z \wedge y) \wedge (y \vee z))$$

Distributivity  $A \wedge (B \vee C) = (A \wedge B) \vee (A \wedge C)$

$$(x \wedge y) \vee (z \wedge y \wedge y) \vee (z \wedge y \wedge z)$$

Associativity, Commutativity, Idempotence

$$(x \wedge y) \vee ((z \wedge y) \vee (y \wedge z))$$

Commutativity, idempotence  $A \wedge A = A$

$$(y \wedge x) \vee (y \wedge z)$$

Distributivity (backwards)  $(A \wedge B) \vee (A \wedge C) = A \wedge (B \vee C)$

$$(y \wedge (x \vee z))$$

**Conclusion:**

$$(x \wedge y) \vee ((y \vee z) \wedge (z \wedge y)) = y \wedge (x \vee z)$$

## Extending the system

...

more gates and DeMorgan's laws

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## More gates (NAND, NOR, XOR)

A	B	A nand B	A nor B	A xor B
0	0	1	1	0
0	1	1	0	1
1	0	1	0	1
1	1	0	0	0

- nand ("not and"):  $A \text{ nand } B = \text{not } (A \text{ and } B)$

$$\begin{matrix} A \\ B \end{matrix} \rightarrow \neg(A \wedge B)$$

- nor ("not or"):  $A \text{ nor } B = \text{not } (A \text{ or } B)$

$$\begin{matrix} A \\ B \end{matrix} \rightarrow \neg(A \vee B)$$

- xor ("exclusive or"):

$$A \text{ xor } B = (A \text{ and not } B) \text{ or } (B \text{ and not } A)$$

$$\begin{matrix} A \\ B \end{matrix} \rightarrow A \oplus B$$

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## A curious fact

- Functional Completeness of NAND and NOR
  - Any logical circuit can be implemented using NAND gates only
- Same applies to NOR

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## DeMorgan's Law

Nand:  $\neg(A \wedge B) = \neg A \vee \neg B$

Nor:  $\neg(A \vee B) = \neg A \wedge \neg B$

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## DeMorgan's Law

Nand:  $\neg(A \wedge B) = \neg A \vee \neg B$

`if not (x > 15 and x < 110): ...`

is logically equivalent to

`if (not x > 15) or (not x < 110): ...`

Nor:  $\neg(A \vee B) = \neg A \wedge \neg B$

`if not (x < 15 or x > 110): ...`

is logically equivalent to

`if (not x < 15) and (not x > 110): ...`

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## A circuit for parity checking

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Boolean expressions and circuits

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## A Boolean expression that checks parity

- 3-bit odd parity checker F: an expression that should be true when the count of 1 bits is odd: when 1 or 3 of the bits are 1s.

$$P = (\neg A \wedge \neg B \wedge C) \vee (\neg A \wedge B \wedge \neg C) \vee (A \wedge \neg B \wedge \neg C) \vee (A \wedge B \wedge C)$$

A	B	C	P
0	0	0	0
0	0	1	1
0	1	0	1
0	1	1	0
1	0	0	1
1	0	1	0
1	1	0	0
1	1	1	1

There are specific methods for obtaining canonical Boolean expressions from a truth table, such as writing it as a disjunction of conjunctions or as a conjunction of disjunctions. See the bonus slide at the end.

Note we have four subexpressions above each of them corresponding to exactly one row of the truth table where P is 1.

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## The circuit

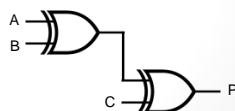
3-bit odd parity checker

$$P = (\neg A \wedge \neg B \wedge C) \vee (\neg A \wedge B \wedge \neg C) \vee (A \wedge \neg B \wedge \neg C) \vee (A \wedge B \wedge C)$$

A	B	C	P
0	0	0	0
0	0	1	1
0	1	0	1
0	1	1	0
1	0	0	1
1	0	1	0
1	1	0	0
1	1	1	1

logically equivalent

$$P = (A \oplus B) \oplus C$$



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## Summary

You should be able to:

- Identify basic gates
- Describe the behavior of a gate or circuit using Boolean expressions, truth tables, and logic diagrams
- Transform one Boolean expression into another given the laws of Boolean algebra

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## Bonus slide Using Minterms to Construct a Boolean Function from a Truth Table

As presented by Alvarado et. al. in *CS for All*:

1. Write down the truth table for the Boolean function that you are considering
2. Delete all the rows from the truth table where the value of the function is 0
3. For each remaining row create a "minterm" as follows:
  - a. For each variable that has a 1 in that row write the name of the variable.  
If the input variable is 0 in that row, write the variable with a negation symbol.
  - b. Take their conjunction (AND them together)
4. Combine all of the minterms using OR (take their disjunction)

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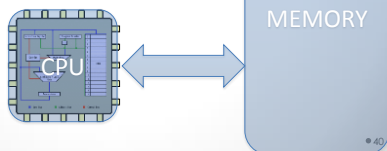
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## Next Time

- How circuits are combined to form a computer
  - Von Neumann architecture revisited
  - Fetch – Decode - Execute Cycle



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