



10-607 Computational Foundations for Machine Learning

Recursion, Induction, &
Dynamic Programming

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Plan

Recursion

- Recursion
- Practical considerations
- Proof by induction

Divide and Conquer

- Sorting example

Dynamic Programming

- Overlapping subproblems
- Examples



Factorial

Complete the recursive function for factorial!

Note: A recursive function calls itself

```
def factorial(int n):  
    if _____:   
        return _____  
    return _____
```

$$0! = 1$$

$$1! = 1$$

$$2! = 2$$

$$3! = 6$$

$$4! = 24$$

Recursion



<https://www.cs.cmu.edu/~112/notes/notes-recursion-part1.html>

Template

```
def recursiveFunction():  
    if (this is the base case):  
        do something non-recursive  
    else:  
        do something recursive
```

Recursion Examples

List sum

Problem: sum all of the numbers in a given list

```
def listSum(L):
```

Base Case: the list is empty, so the sum is 0

```
if (len(L) == 0):
```

```
    return 0
```

```
else:
```

Recursive Case: assume we already know the sum of the entire list

after the first element. Add that sum to the first element.

```
return L[0] + listSum(L[1:])
```

```
print(listSum([2,3,5,7,11])) # 28
```

Recursion Examples

Power

Problem: raise the number base to the given exponent

```
def power(base, expt):  
    # assume expt is non-negative integer  
    if (expt == 0):  
        return 1  
    else:  
        return base * power(base, expt-1)
```

```
print(power(2,5)) # 32
```

$$2^{10}$$

$$2 \cdot 2^9$$

Recursion Examples

Power – Multiple base cases!

```
def power(base, expt):  
    # This version allows for negative exponents  
    # It still assumes that expt is an integer, however.  
    if (expt == 0):  
        return 1  
    elif (expt < 0): # new recursive case!  
        return 1.0 / power(base, abs(expt))  
    else:  
        return base * power(base, expt-1)  
  
print(power(2,5)) # 32  
print(power(2,-5)) # 1/32 = 0.03125
```

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Iterative vs Recursive

Factorial

Iterative



```
def factorial(n):  
    result = 1  
    for i in range(2, n + 1):  
        result *= i  
    return result  
  
print(factorial(5))
```

Recursive



```
def factorial(n):  
    if (n < 2):  
        return 1  
    else:  
        return n * factorial(n - 1)  
  
print(factorial(5))
```

How does this look in memory?

Memory



```
long factorial(int n) {  
    if (n == 0) {  
        return 1;  
    }  
  
    return n * factorial(n-1);  
}  
  
int main() {  
    int n = 4;  
  
    long result;  
    result = factorial(n);  
  
    cout << result << endl;  
}
```

How does this look in memory?

Memory



```
long factorial(int n) {  
    if (n == 0) {  
        return 1;  
    }  
  
    return n * factorial(n-1);  
}
```

```
int main() {  
    int n = 4;   
    long result = 0;  
    result = factorial(n);  
  
    cout << result << endl;  
}
```

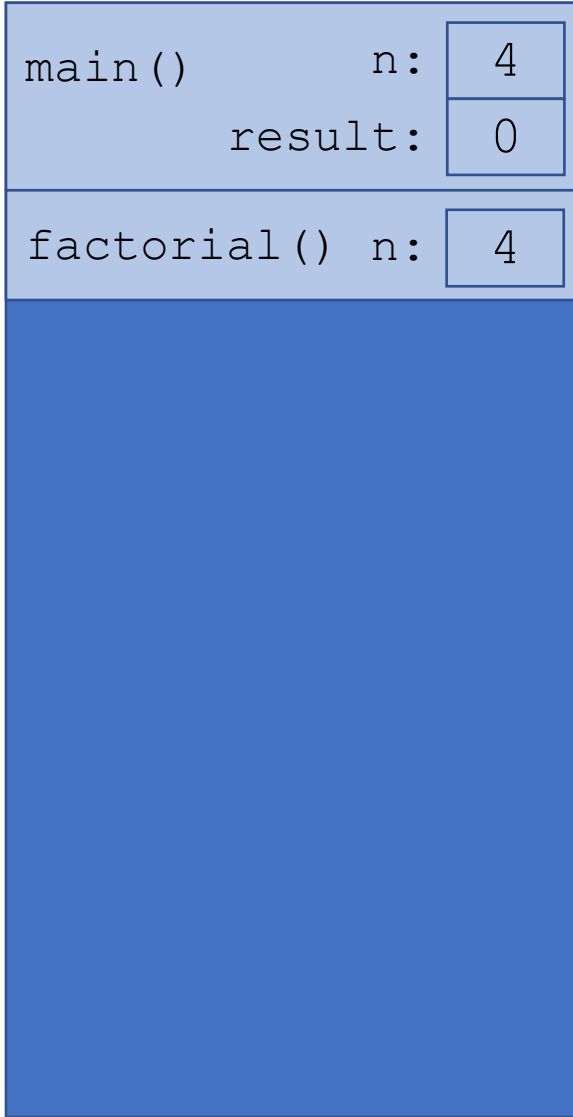
The “stack” part of memory is a
stack

Function call: push

Return: pop

“Stack” part of memory is a stack

Memory



```
long factorial(int n) {  
    if (n == 0) {  
        return 1;  
    }  
  
    return n * factorial(n-1);  
}  
  
int main() {  
    int n = 4;  
  
    long result = 0;  
    result = factorial(n);  
  
    cout << result << endl;  
}
```

“Stack” part of memory is a stack

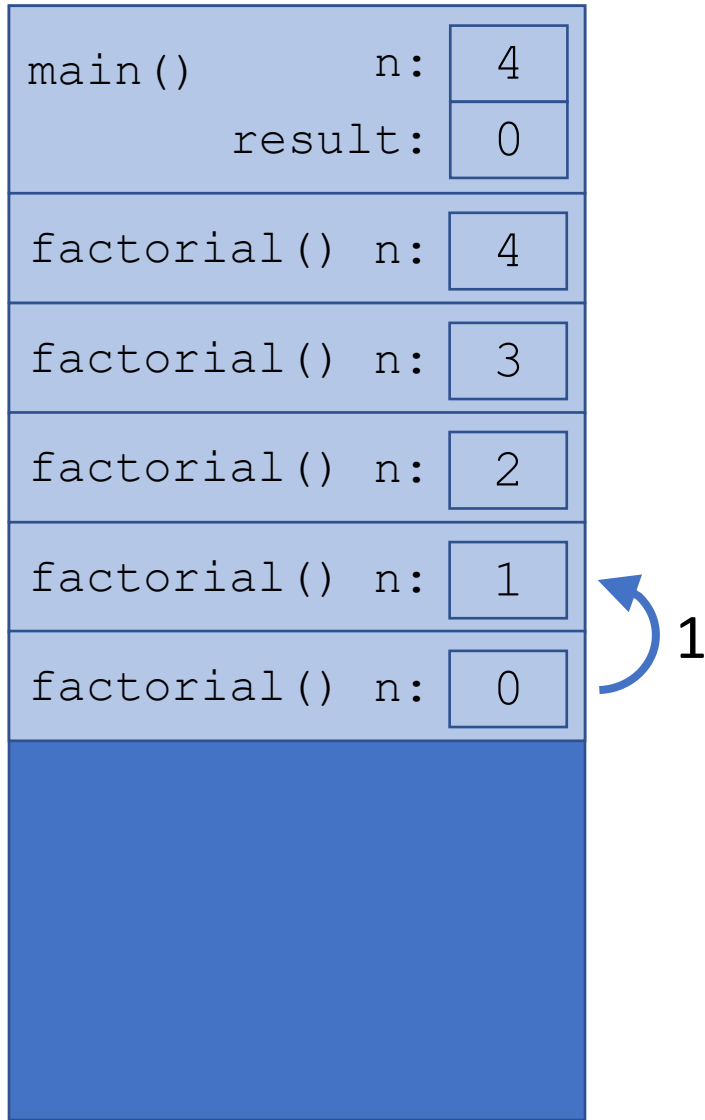
Memory

main()	n:	4
	result:	0
factorial()	n:	4
factorial()	n:	3
factorial()	n:	2
factorial()	n:	1
factorial()	n:	0

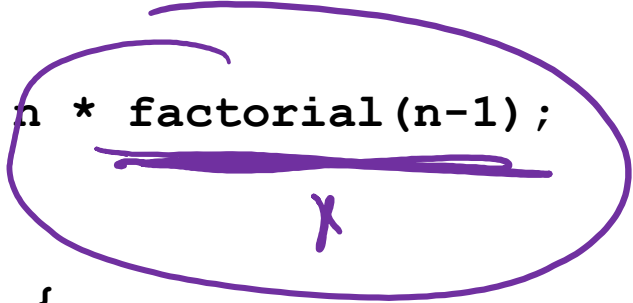
```
long factorial(int n) {  
    if (n == 0) {  
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    }  
  
    return n * factorial(n-1);  
}  
  
int main() {  
    int n = 4;  
  
    long result = 0;  
    result = factorial(n);  
  
    cout << result << endl;  
}
```

“Stack” part of memory is a stack

Memory

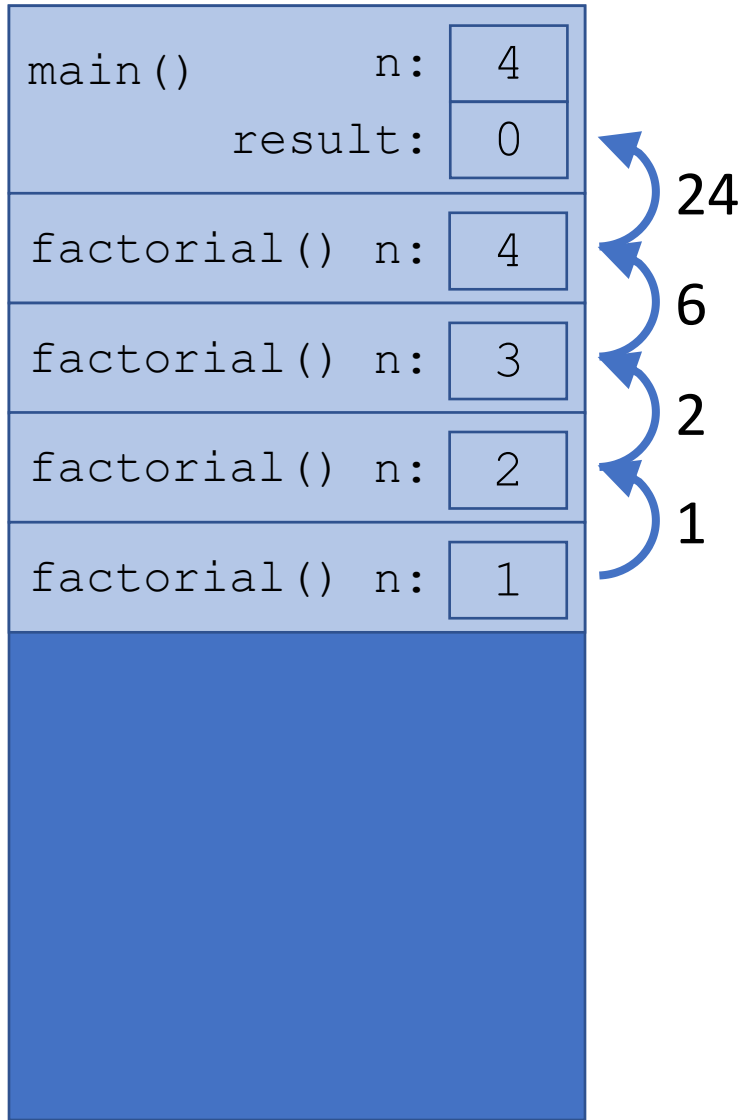


```
long factorial(int n) {  
    if (n == 0) {  
        return 1;  
    }  
    return n * factorial(n-1);  
}  
  
int main() {  
    int n = 4;  
  
    long result = 0;  
    result = factorial(n);  
  
    cout << result << endl;  
}
```



“Stack” part of memory is a stack

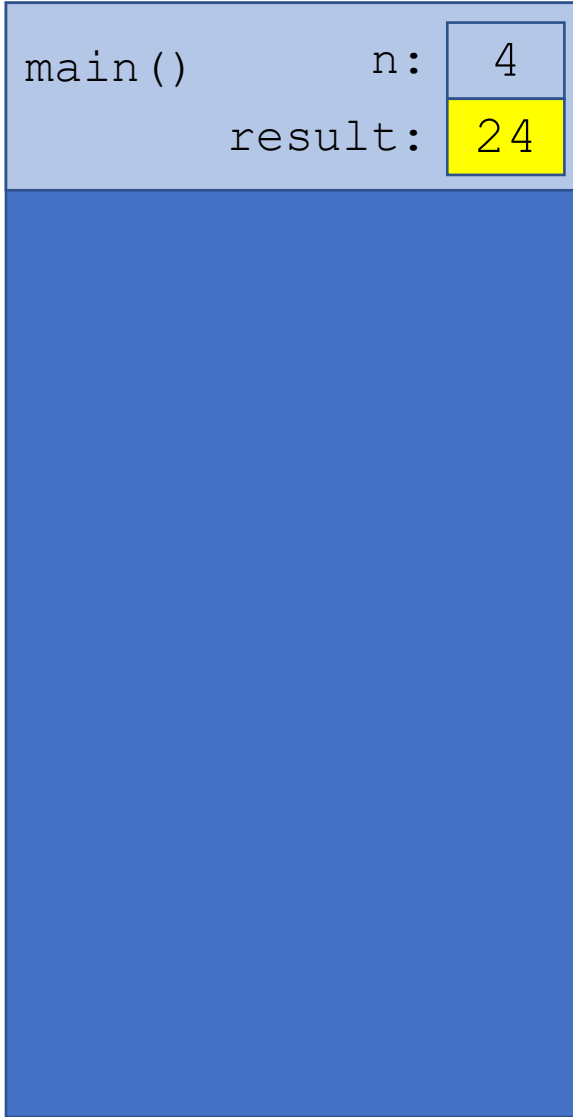
Memory



```
long factorial(int n) {  
    if (n == 0) {  
        return 1;  
    }  
    return n * factorial(n-1);  
}  
  
int main() {  
    int n = 4;  
  
    long result = 0;  
    result = factorial(n);  
  
    cout << result << endl;  
}
```

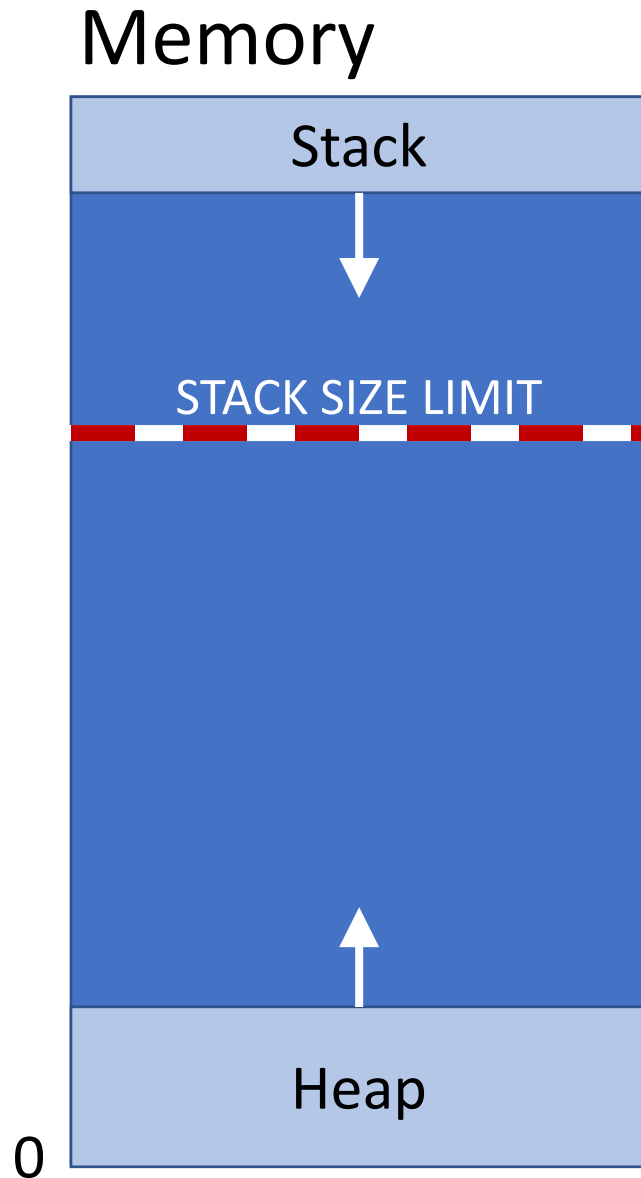

“Stack” part of memory is a stack

Memory



```
long factorial(int n) {  
    if (n == 0) {  
        return 1;  
    }  
  
    return n * factorial(n-1);  
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int main() {  
    int n = 4;  
  
    long result = 0;  
    result = factorial(n);  
  
    cout << result << endl;  
}
```

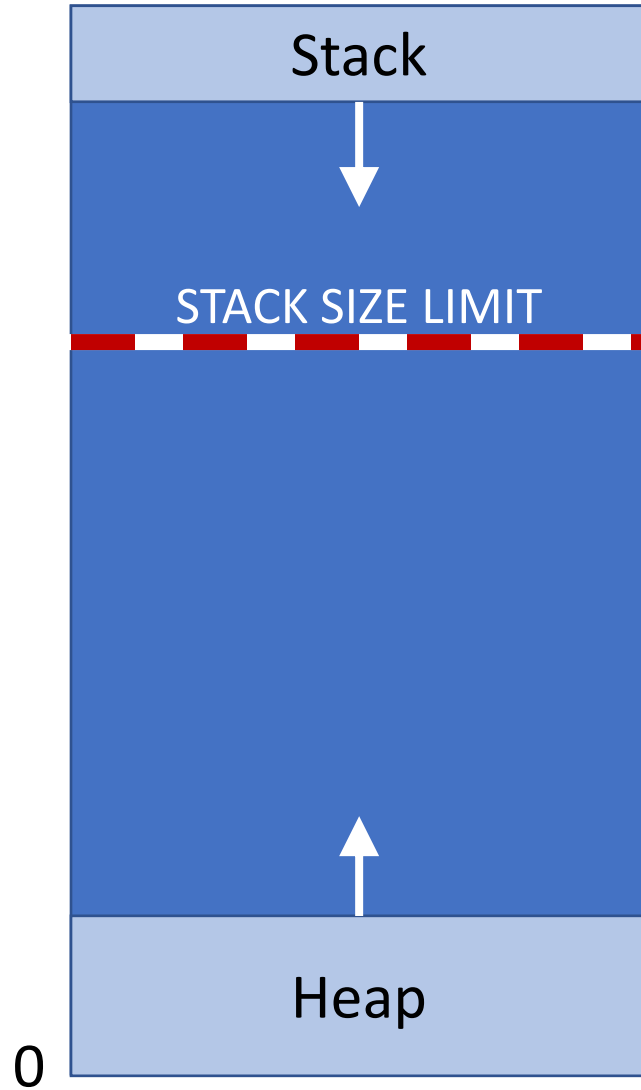
What can go wrong?



```
long factorial(int n) {  
    if (n == 0) {  
        return 1;  
    }  
  
    return n * factorial(n-1);  
}  
  
int main() {  
    int n = 4;  
  
    long result = 0;  
    result = factorial(n);  
  
    cout << result << endl;  
}
```

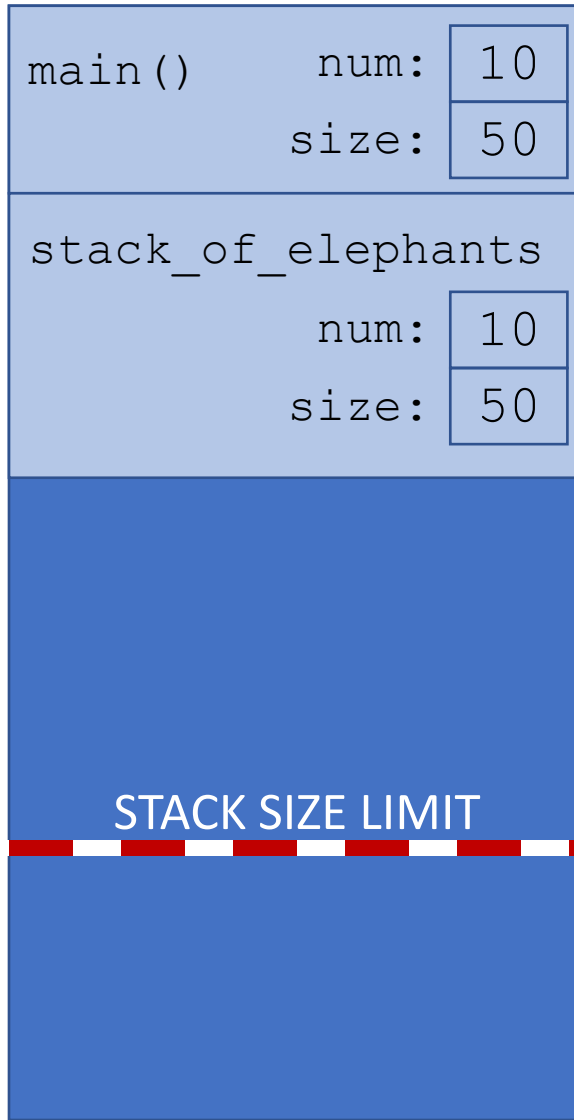
Demo: Stack of elephants

Memory



Demo: Stack of elephants

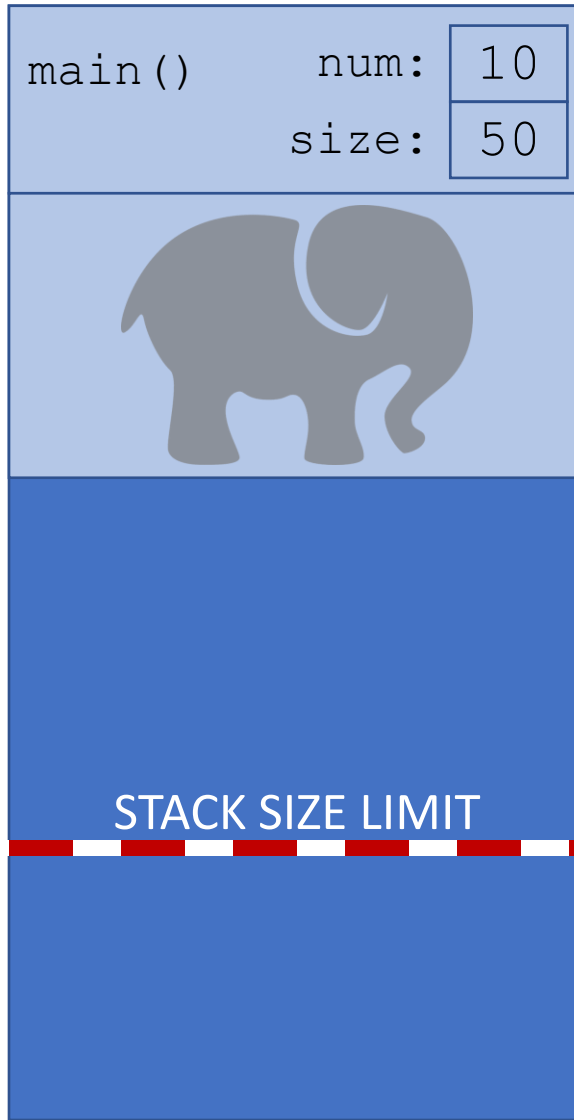
Memory



```
void stack_of_elephants(  
    int num, int size) {  
    if (num == 0) {  
        return;  
    }  
  
    int elephant_array[size]   
  
    stack_of_elephants(num-1, size);  
}  
  
int main() {  
    int num = 10;   
    int size = 50;   
    stack_of_elephants(num, size)  
}
```

Demo: Stack of elephants

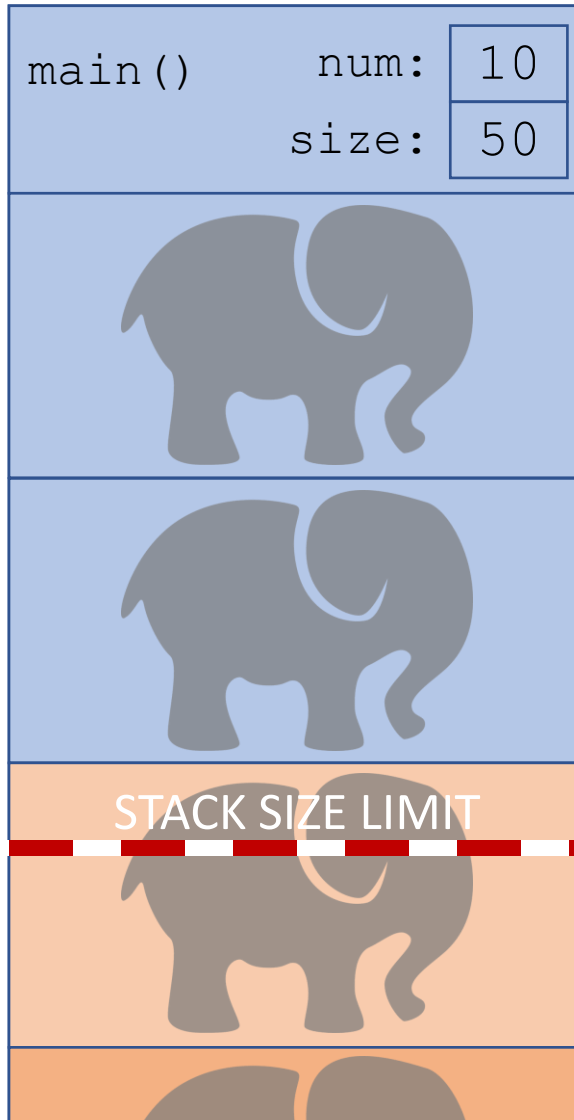
Memory



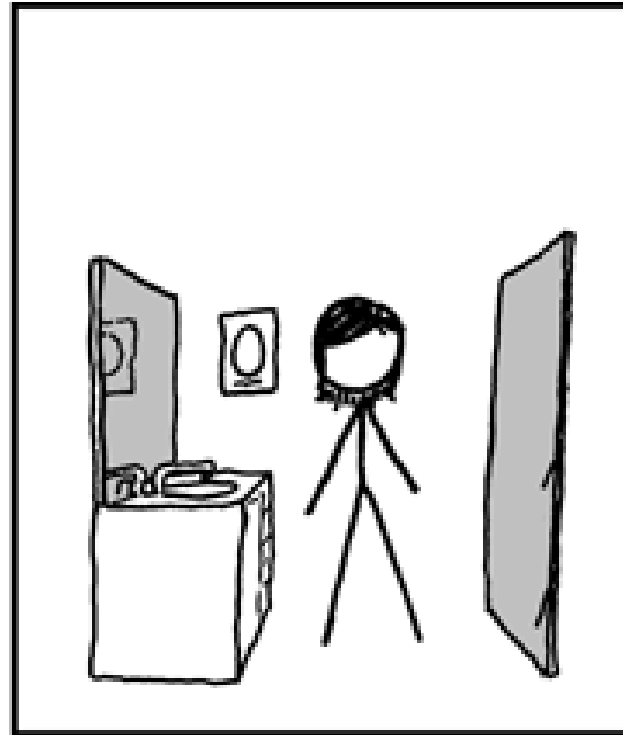
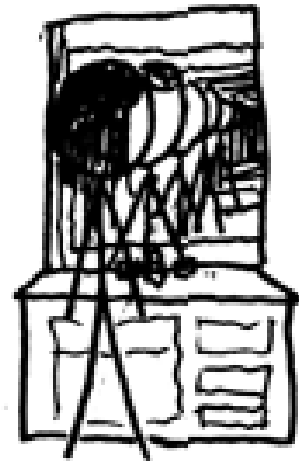
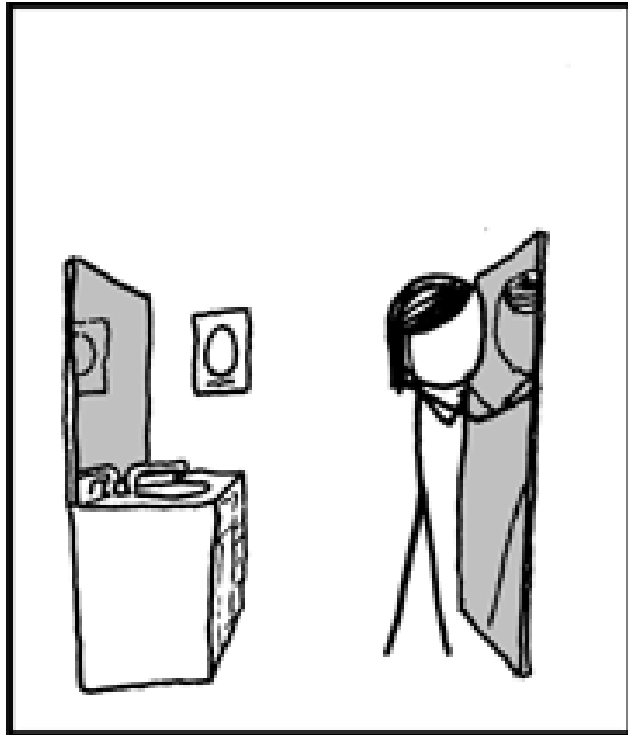
```
void stack_of_elephants(  
    int num, int size) {  
    if (num == 0) {  
        return;  
    }  
  
    int elephant_array[size]  
  
    stack_of_elephants(num-1, size);  
}  
  
int main() {  
    int num = 10;  
    int size = 50;  
    stack_of_elephants(num, size)  
}
```

Demo: Stack of elephants

Memory



```
void stack_of_elephants(  
    int num, int size) {  
    if (num == 0) {  
        return;  
    }  
  
    int elephant_array[size]  
  
    stack_of_elephants(num-1, size);  
}  
  
int main() {  
    int num = 10;  
    int size = 50; 50000  
    stack_of_elephants(num, size)  
}
```



A fatal error has occurred.

0x00000539 0x4641494C

0x4F4F5053 0x786B6364

Press Ctrl+Alt+Delete to
restart the universe.

Plan

Recursion

- ✓ ■ Recursion
- ✓ ■ Practical considerations
- Proof by induction

Divide and Conquer

- Sorting example

Dynamic Programming

- Overlapping subproblems
- Examples

Proof by Induction

Template

- Prove proposition for base case, e.g., $n=0$
- Inductive step
 - Assume proposition holds for $n=k$
 - Prove proposition holds for $n=k+1$



Proof by Induction

$$f(n) = n(n+1)/2$$

Example: Prove sum of integers 1 to n is $n(n+1)/2$

$$\xrightarrow{k+1} (k+1)(k+1+1)/2$$

Base $n=1$

$$f(1) = 1(1+1)/2 = 1 \quad \checkmark$$

Assume $n=k$ hold $f(k)$

Prove $n=k+1$ is correct

$$f(k+1) = \underline{f(k)} + k+1 \quad \leftarrow$$

$$= k(k+1)/2 + k+1$$

$$= (k(k+1) + 2(k+1))/2$$

$$= (k+1)(k+2)/2$$

$\checkmark$

Exercise 1

Use induction to prove that for $f(n) = n^2$ in $O(2^n)$

$$n^2 \leq 1 \cdot 2^n$$

$$C=1$$

$$n_0 = 5$$

Base $n=5$: $5^2 \leq 2^5$
 $25 \leq 32 \quad \checkmark$

Assume $n=k$: $k^2 \leq 2^k$

Prove $n=k+1$:

$$(k+1)^2 \leq 2^{(k+1)}$$

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Divide and Conquer

Merge sort

Detailed code: [15-112 Merge Sort](#)

```
def mergesort(L):  
    if (len(L) < 2):  
        return L  
    else:  
        mid = len(L)//2  
        left = mergesort(L[:mid])  
        right = mergesort(L[mid:])  
        return merge(left, right)  
  
print(mergesort([1, 5, 3, 4, 2, 0]))
```

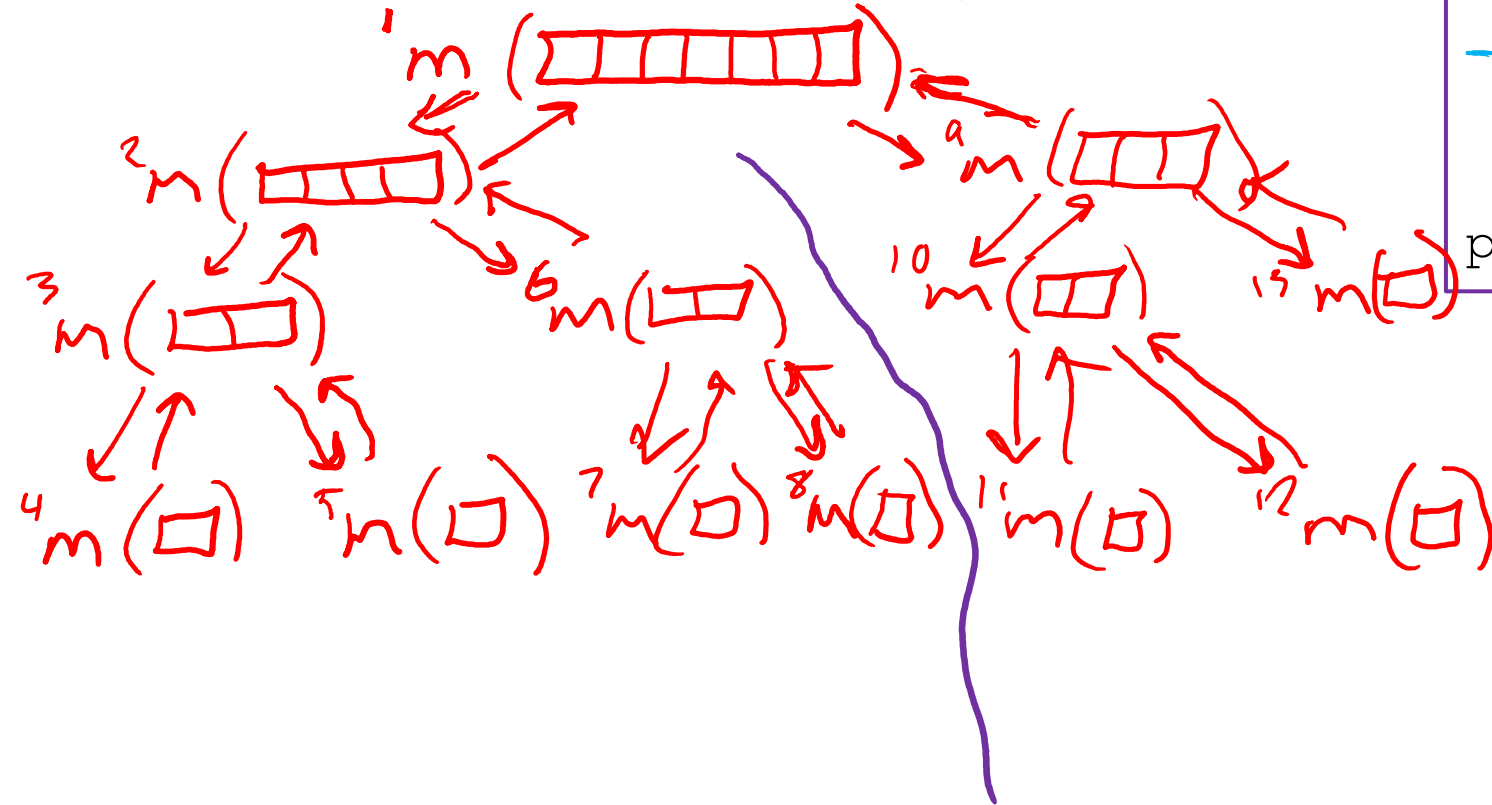
Divide and Conquer

$m()$ →

Merge sort

Detailed code: [15-112 Merge Sort](#)

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def mergesort(L):  
    if (len(L) < 2):  
        return L  
    else:  
        mid = len(L)//2  
        left = mergesort(L[:mid])  
        → right = mergesort(L[mid:])  
        return merge(left, right)  
  
print(mergesort([1,5,3,4,2,0]))
```



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Reminder Binary Search

Graph of recursive calls

¹ S ([] [] [] [] [] [])



² S ([] [] [])



³ S ([] [])



⁴ S ([])

```
binarySearch2(data, key, first, last) {  
    // TODO base case  
    int mid = (first + last) // 2;  
    if (data[mid] == key)  
        return true;  
    else if (data[mid] > key)  
        return binarySearch2(data, key, first, mid-1);  
    else  
        return binarySearch2(data, key, mid+1, last)  
}
```


Exercise 2: Fibonacci Naïve

Notebook

Fibonacci

N Value



0 1

1 1

2 2 = 1+1

3 3 = 1+2

4 5 = 2+3

5 8 = 3+5

6 13 = 5+8

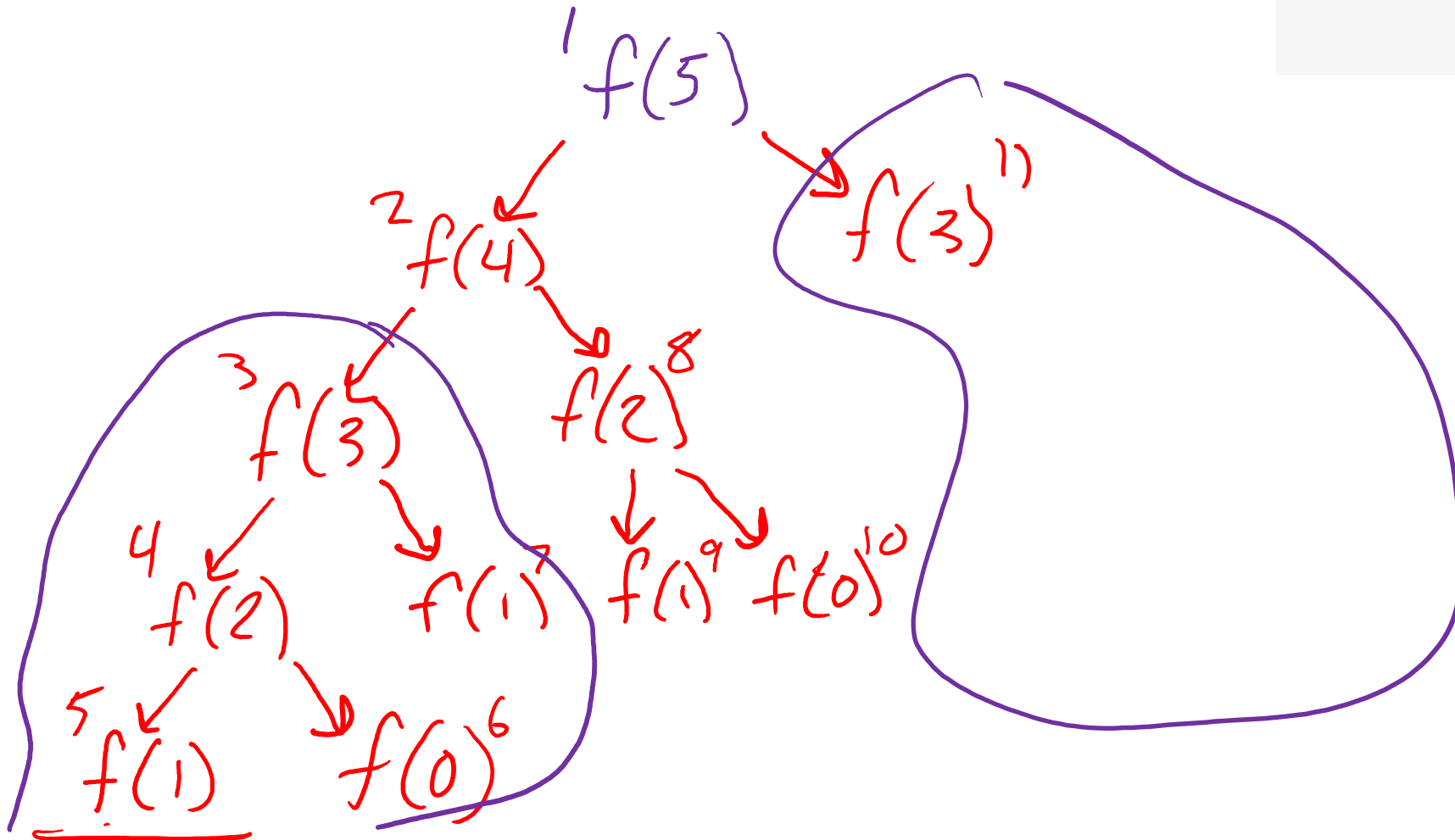
7 21 = 8+13

$$f(n) = f(n-2) + f(n-1)$$

Exercise 3: Fibonacci Naïve

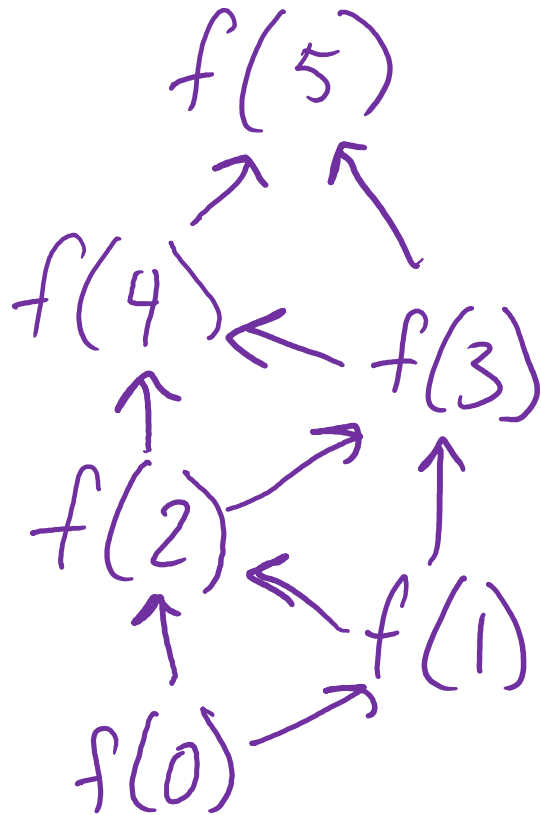
Graph of recursive calls

```
def fib(n):  
    if n < 2:  
        return 1  
    else:  
        return fib(n-1) + fib(n-2)
```



Fibonacci Dynamic Programming

Graph with overlapping subproblems



Fibonacci Dynamic Programming

Notebook

Dynamic Programming

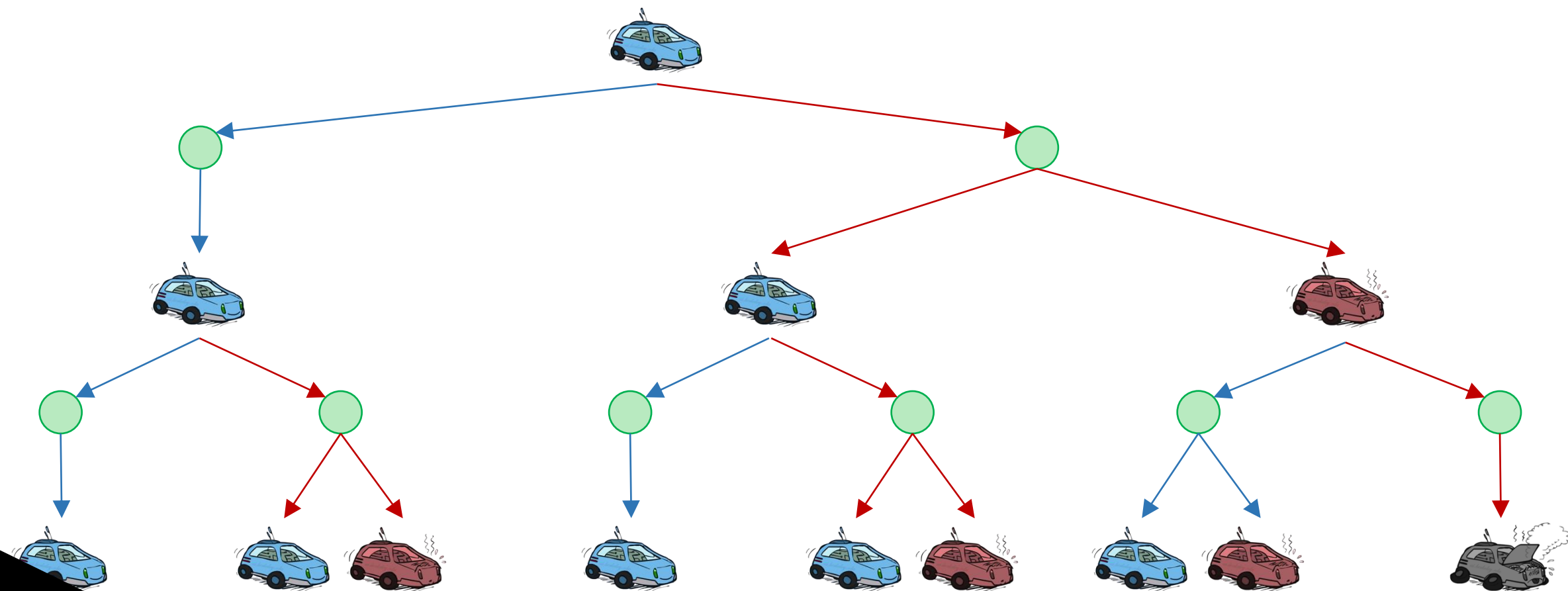
Efficient way to solve problems that can take a long time

- Dividing them into a set of smaller problems
- Optimal solutions of each subproblem are stored and used to find the optimal solution for the larger problem
- Unlike divide and conquer, these problems can overlap.

Examples from ML

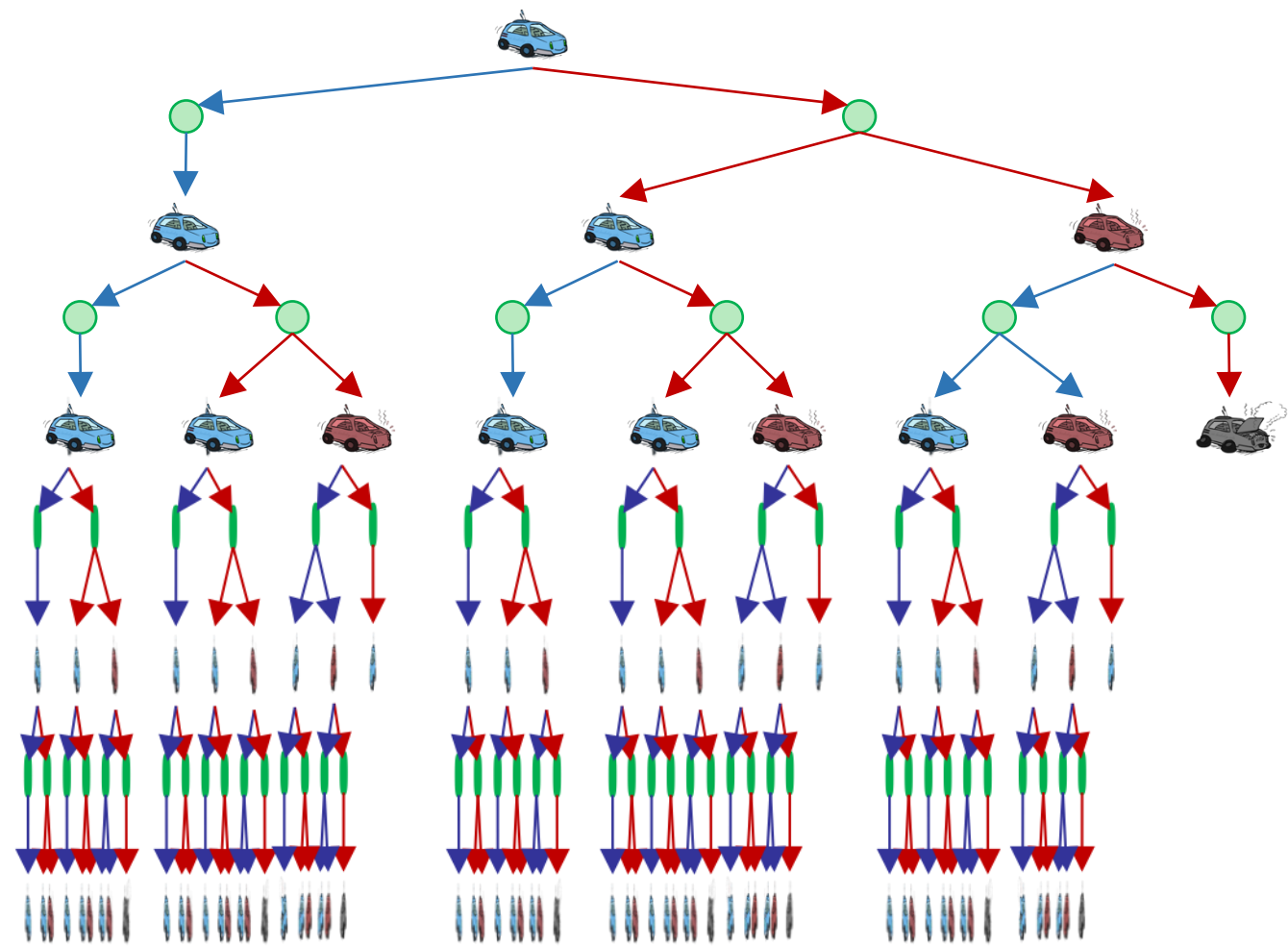
Dynamic Programming

MDP: Racing Search Tree



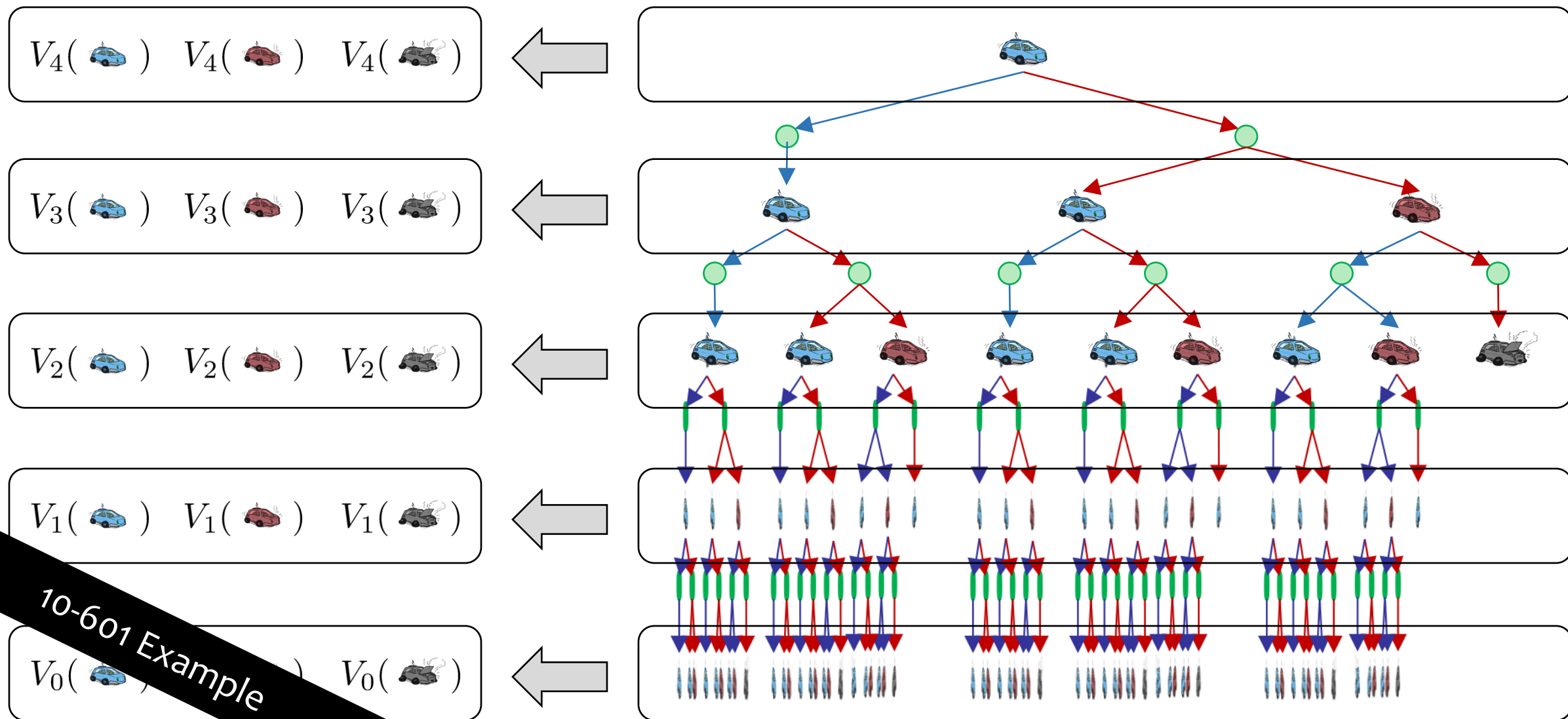
10-601 Example

MDP: Racing Search Tree



10-601 Example

MDP: Racing Tree Example



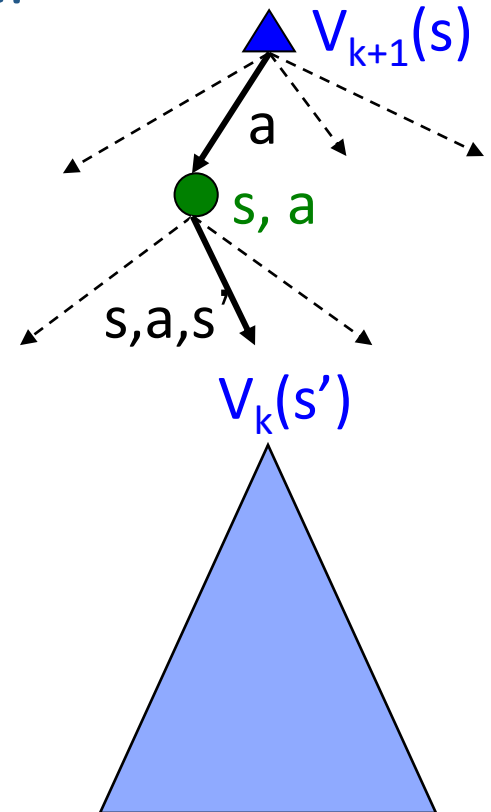
MDP: Value Iteration

Start with $V_0(s) = 0$: no time steps left means an expected reward sum of zero

Given vector of $V_k(s)$ values, do one ply of expectimax from each state:

$$V_{k+1}(s) \leftarrow \max_a \sum_{s'} T(s, a, s') [R(s, a, s') + \gamma V_k(s')]$$

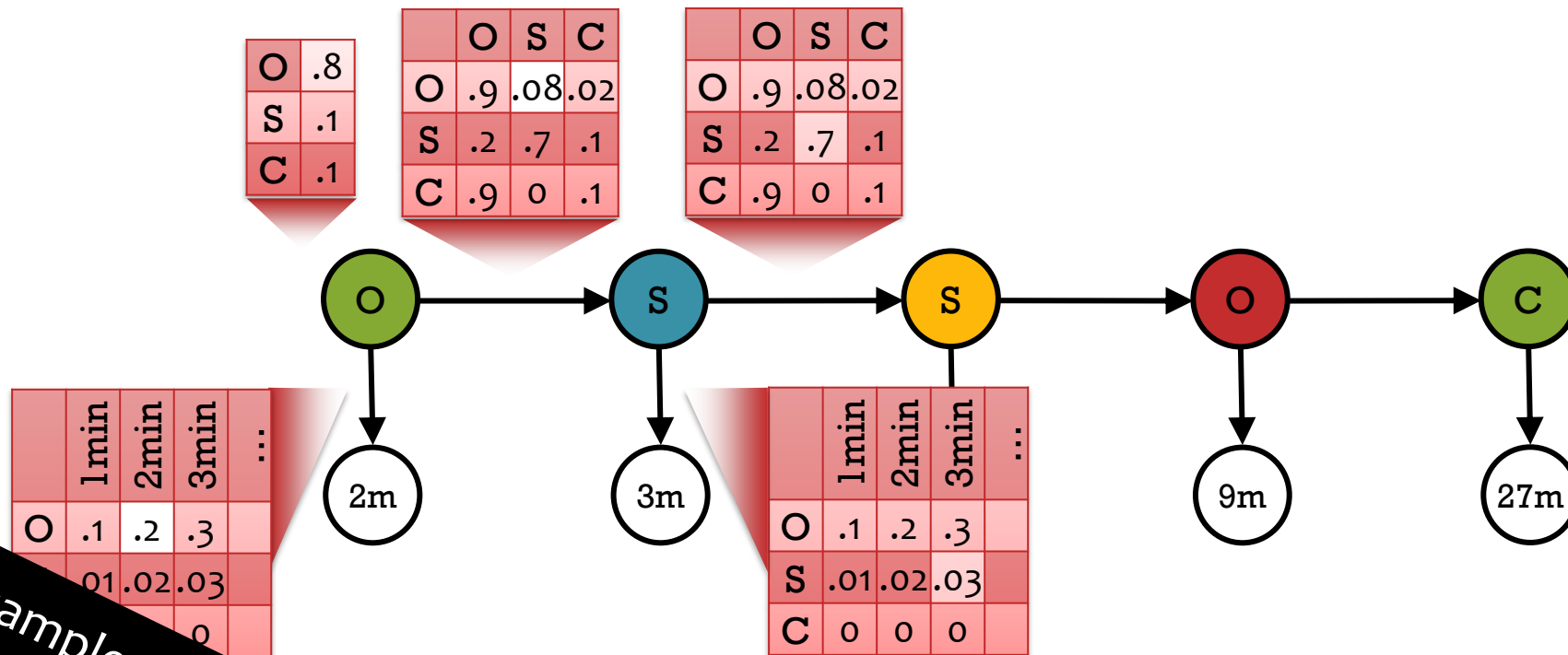
Repeat until convergence



Hidden Markov Model

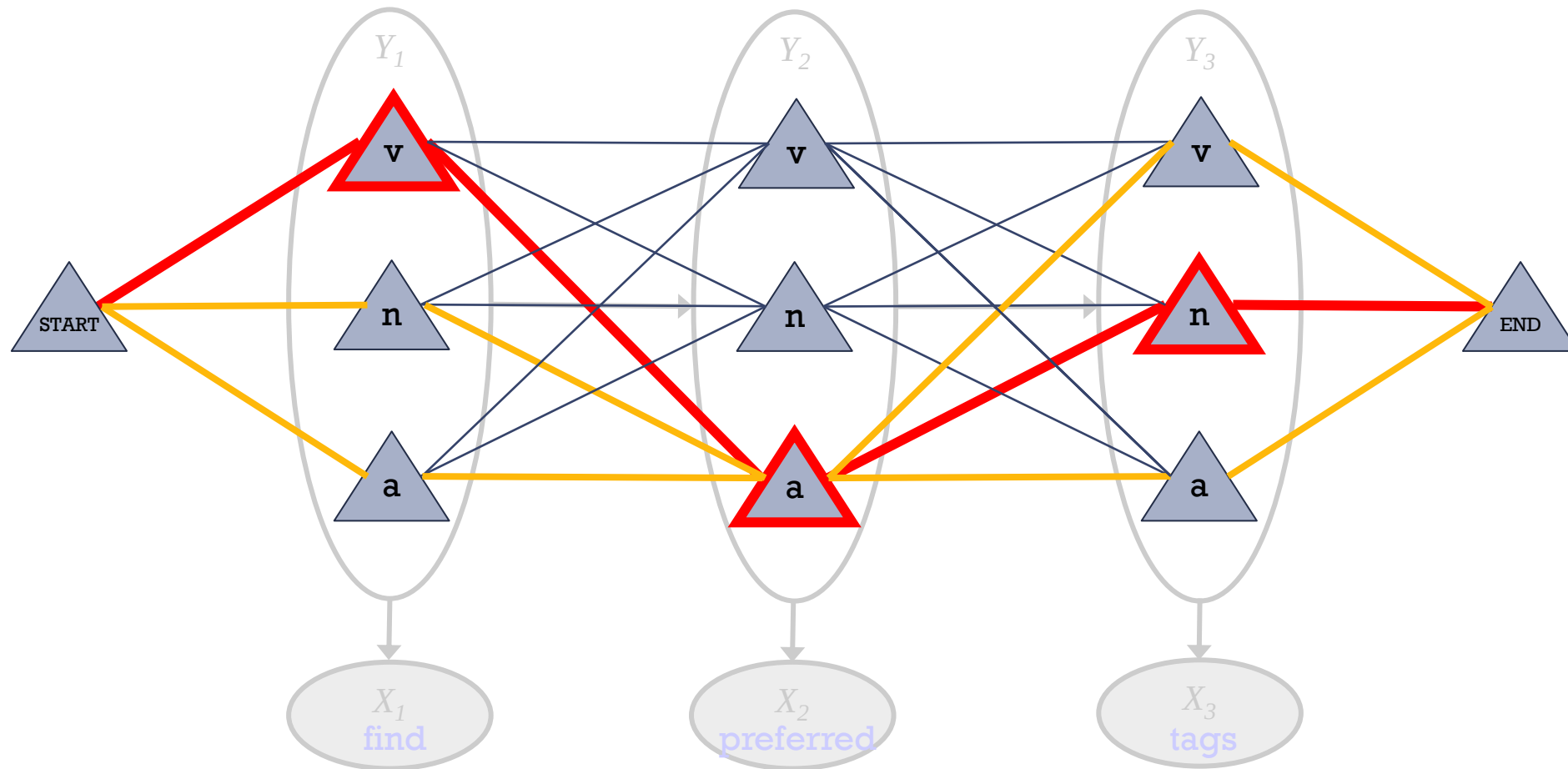
A Hidden Markov Model (HMM) provides a joint distribution over the the tunnel states / travel times with an assumption of dependence between adjacent tunnel states.

$$p(o, s, s, o, c, 2m, 3m, 18m, 9m, 27m) = (.8 * .08 * .2 * .7 * .03 * \dots)$$



10-601 Example

Forward-Backward Algorithm: Finds Marginals



- So $p(v \ a \ n) = (1/Z) * \text{product weight of one path}$

Marginal probability $p(Y_2 = a)$

$= (1/Z) * \text{total weight of all paths through}$

