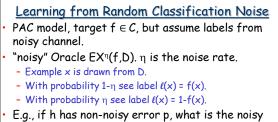


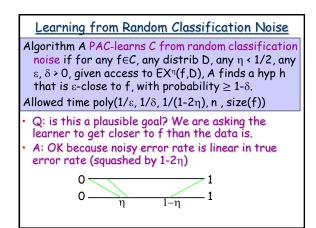
# Learning when there is no perfect hypothesis

- Hoeffding/Chernoff bounds: minimizing training error will approximately minimize true error: just need  $O(1/\varepsilon^2)$  samples versus  $O(1/\varepsilon)$ .
- What about polynomial-time algorithms? Seems harder.
- Given data set S, finding apx best conjunction is NP-hard.
  Can do other things, like minimize hinge-loss, but may be a big gap wrt error rate ("0/1 loss").
- One way to make progress: make assumptions on the "noise" in the data. E.g., Random Classification Noise model.



- E.g., if h has non-noisy error p, what is the noisy error rate?
  - $p(1-\eta) + (1-p)\eta = \eta + p(1-2\eta)$ .





## **Notation**

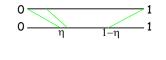
- Use "Pr[...]" for probability with respect to non-noisy distribution.
- Use "Pr<sub>η</sub>[...]" for probability with respect to noisy distribution.

## Learning OR-functions (assume monotone) Let's assume noise rate $\eta$ is known. Say $p_i = \Pr[f(x)=0 \text{ and } x_i=1]$ Any h that includes all $x_i$ such that $p_i=0$ and no $x_i$ such that $p_i > \varepsilon/n$ is good. So, just need to estimate $p_i$ to $\pm \frac{\varepsilon}{2n}$ .

- Rewrite as  $p_i = \Pr[f(x)=0|x_i=1] \times \Pr[x_i=1]$ .
- $2^{nd}$  part unaffected by noise (and if tiny, can ignore  $x_i).$  Define  $q_i$  as  $1^{st}$  part.
- Then  $\Pr_{\eta}[\ell(x)=0|x_i=1] = q_i(1-\eta)+(1-q_i)\eta = \eta+q_i(1-2\eta).$
- So, enough to approx LHS to  $\pm O\left(\frac{\epsilon}{2n}(1-2\eta)\right)$ .

# Learning OR-functions (assume monotone)

• If noise rate not known, can estimate with smallest value of  $\Pr_{\eta}[\ell(x)=0|x_i=1]$ .



#### Generalizing the algorithm

Basic idea of algorithm was:

- See how can learn in non-noisy model by asking about probabilities of certain events with some "slop".
- Try to learn in noisy model by breaking events into:
  - Parts predictably affected by noise.
  - Parts unaffected by noise.
- Let's formalize this in notion of "statistical query" (SQ) algorithm. Will see how to convert any SQ alg to work with noise.

## The Statistical Query Model

- No noise.
- Algorithm asks: "what is the probability a labeled example will have property χ? Please tell me up to additive error τ."
  - Formally,  $\chi{:}X\times\{0,1\}\to\{0,1\}.$  Must be poly-time computable.  $\tau\ge 1/\text{poly}(...).$
  - Let  $P_{\chi} = \Pr[\chi(x,f(x))=1].$
  - World responds with  $\mathsf{P}'_\chi\in[\mathsf{P}_\chi\text{-}\tau,\,\mathsf{P}_\chi\text{+}\tau].$
  - [can extend to [0,1]-valued or vector-valued  $\boldsymbol{\chi}$ ]
- May repeat poly(...) times. Can also ask for unlabeled data. Must output h of error  $\leq \epsilon.$  No  $\delta$  in this model.

## The Statistical Query Model

- Examples of queries:
- What is the probability that x<sub>i</sub>=1 and label is negative? - What is the error rate of my current hypothesis h?  $[\chi(x,\ell)=1 \text{ iff } h(x) \neq \ell]$
- Get back answer to  $\pm \tau$ . Can simulate from  $\approx 1/\tau^2$  examples. [That's why need  $\tau \ge 1/\text{poly}(...)$ .]
- To learn OR-functions, ask for Pr[x<sub>i</sub>=1 and f(x)=0] with  $\tau = \frac{\epsilon}{2n}$ . Produce OR of all x<sub>i</sub> s.t.  $P'_{\chi} \le \frac{\epsilon}{2n}$ .

# The Statistical Query Model

- Many algorithms can be simulated with statistical queries:
  - Perceptron: ask for E[f(x)x :  $h(x) \neq f(x)$ ] (formally define vector-valued  $\chi = f(x)x$  if  $h(x) \neq f(x)$ , and 0 otherwise. Then divide by Pr[ $h(x) \neq f(x)$ ].)
  - Hill-climbing type algorithms: what is error rate of h? What would it be if I made this tweak?
- Properties of SQ model:
  - Can automatically convert to work in presence of classification noise.
  - Can give a nice characterization of what can and cannot be learned in it.

## <u>SQ-learnable $\Rightarrow$ (PAC+Noise)-learnable</u>

Given query  $\chi_{\text{,}}$  need to estimate from noisy data. Idea:

- Break into part predictably affected by noise, and part unaffected.
- Estimate these parts separately.
- Can draw fresh examples for each query or estimate many queries from same sample if VCDim of query space is small.
- Running example:  $\chi(x, \ell)=1$  iff  $x_i=1$  and  $\ell=0$ .

## How to estimate $Pr[\chi(x,f(x))=1]$ ?

- Let CLEAN =  $\{x : \chi(x,0) = \chi(x,1)\}$
- Let NOISY = { $x : \chi(x,0) \neq \chi(x,1)$ }

- What are these for " $\chi(x,\ell)=1$  iff  $x_i=1$  and  $\ell=0$ "? Now we can write:

- $\Pr[\chi(x,f(x))=1] = \Pr[\chi(x,f(x))=1 \text{ and } x \in CLEAN] + \Pr[\chi(x,f(x))=1 \text{ and } x \in NOISY].$
- Step 1: first part is easy to estimate from noisy data (easy to tell if  $x \in CLEAN$ ).
- What about the 2<sup>nd</sup> part?

#### How to estimate $Pr[\chi(x,f(x))=1]$ ?

- Let  $CLEAN = \{x : \chi(x,0) = \chi(x,1)\}$
- Let NOISY = { $x : \chi(x,0) \neq \chi(x,1)$ } - What are these for  $\chi(x,\ell)=1$  iff  $x_i=1$  and  $\ell=0$ ?
- Now we can write: -  $Pr[\chi(x,f(x))=1] = Pr[\chi(x,f(x))=1 \text{ and } x \in CLEAN] + Pr[\chi(x,f(x))=1 \text{ and } x \in NOISY].$
- Can estimate Pr[x ∈ NOISY].
- Also estimate  $P_{\eta} \equiv \Pr_{\eta}[\chi(x,\ell)=1 \mid x \in \text{NOISY}].$
- Want  $P \equiv Pr[\chi(x, f(x))=1 | x \in NOISY].$
- Write  $P_{\eta} = P(1-\eta) + (1-P)\eta = \eta + P(1-2\eta)$ .

- Just need to estimate  $P_{\eta}$  to additive error  $\tau(1-2\eta)$ .
- If don't know η, can have "guess and check" wrapper around entire algorithm.

#### <u>Characterizing what's learnable using</u> <u>SQ algorithms</u>

- Key tool: Fourier analysis of boolean functions.
- Sounds scary but it's a cool idea!
- Let's think of functions from {0,1}<sup>n</sup> →{-1,1}.
  View function f as a vector of 2<sup>n</sup> entries:
- $\left(\sqrt{D[000]}f(000),\sqrt{D[001]}f(001),\dots,\sqrt{D[x]}f(x),\dots\right)$
- What is (f, f)? What is (f, g)?
- What is an orthonormal basis? Will see connection to SQ algs next time...