A Problematic Program

John C. Reynolds
Carnegie Mellon University
and
Microsoft Research Cambridge

5 February 2008 — Dagstuhl
(Joint work with Josh Berdine)

(Research partially supported by National Science Foundation Grant CCF-054102l.)

A Problematic Program

It is widely believed that two concurrent processes that both mutate the same location may cause a potential race condition unless all mutations occur within critical regions associated with the same resource.

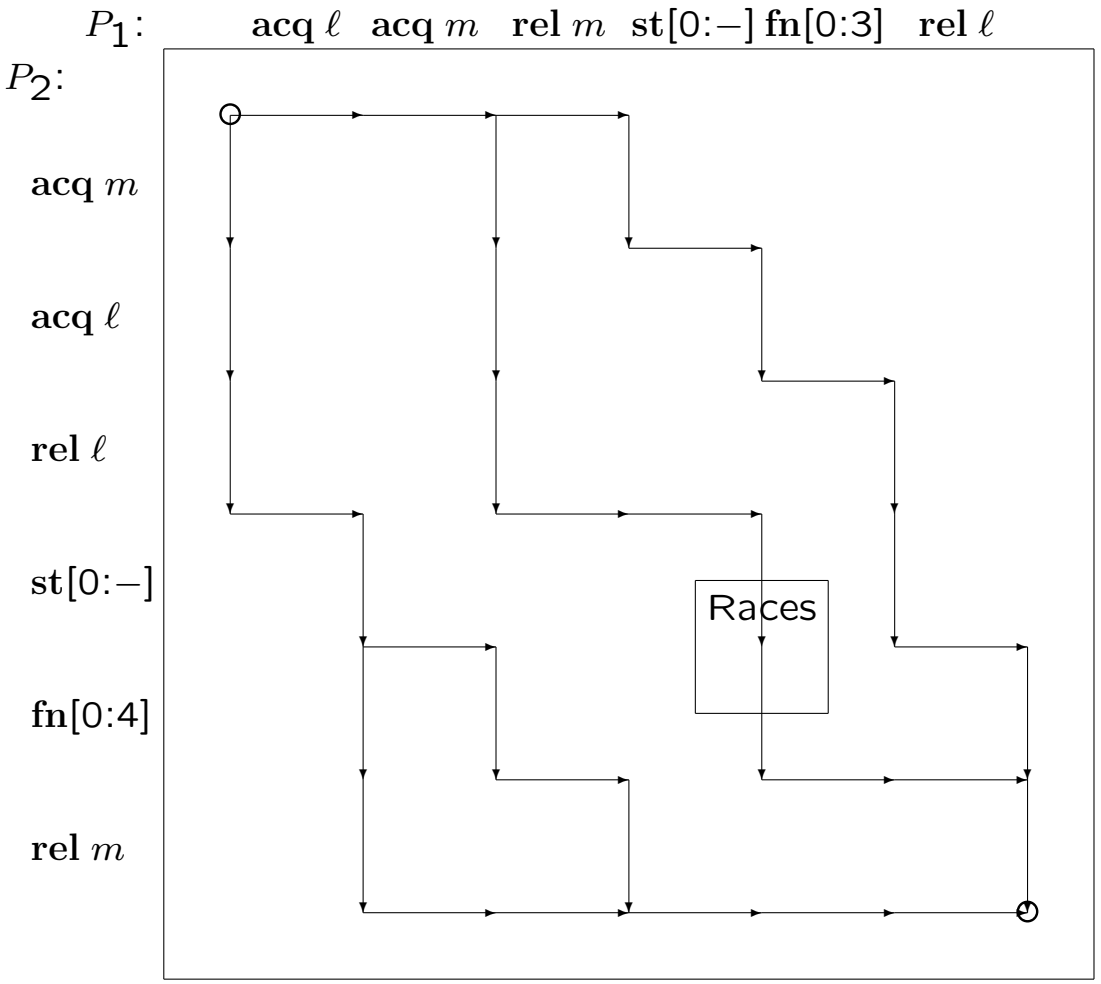
In fact, the following program cannot cause a race condition:

$$\begin{array}{l}\mathbf{resource}\ \ell\ \mathbf{in}\ \mathbf{resource}\ m\ \mathbf{in}\\ \qquad (\mathbf{with}\ \ell\ \mathbf{do}\ ((\mathbf{with}\ m\ \mathbf{do}\ \mathbf{skip}); [0] := 3))\\ \qquad \parallel (\mathbf{with}\ m\ \mathbf{do}\ ((\mathbf{with}\ \ell\ \mathbf{do}\ \mathbf{skip}); [0] := 4)).\end{array}$$

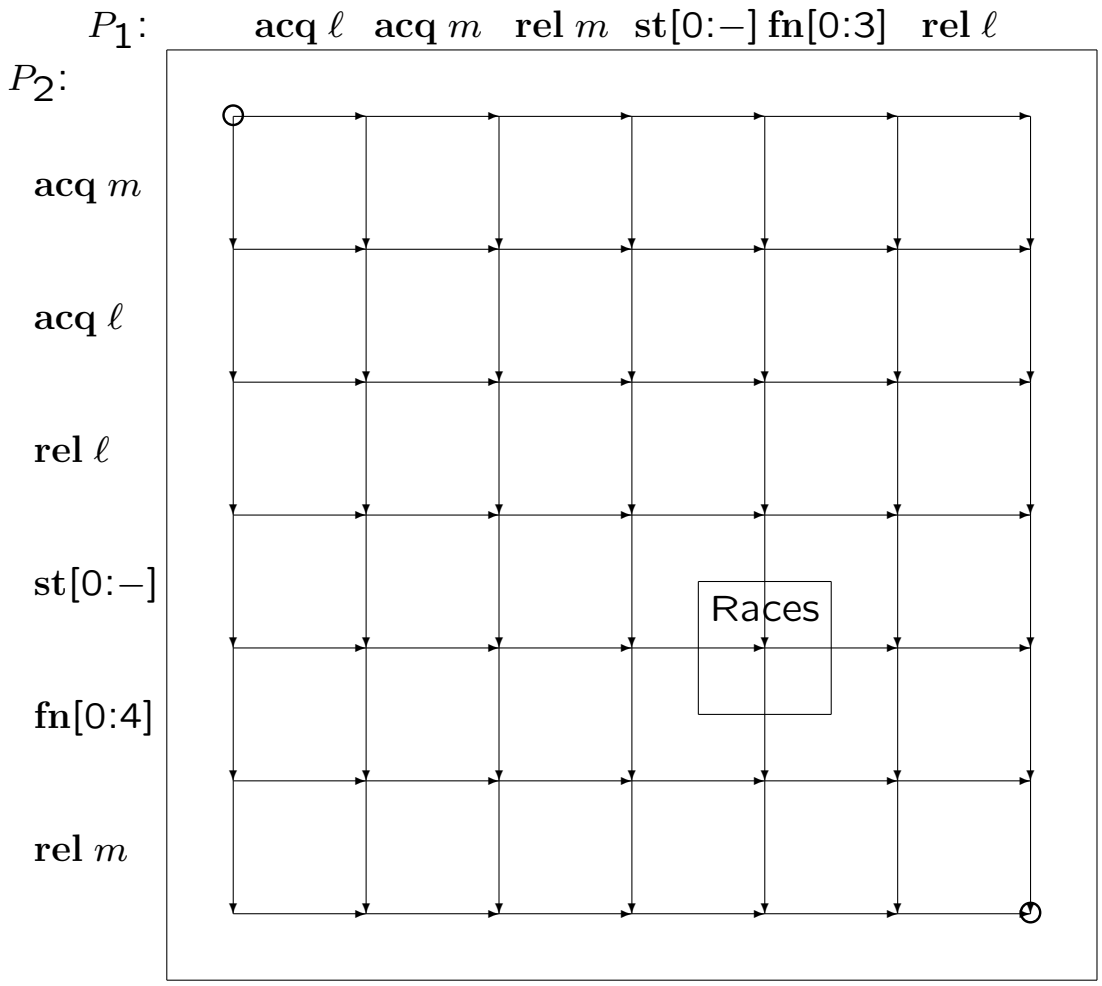
(although it can deadlock).

Interleavings as Paths

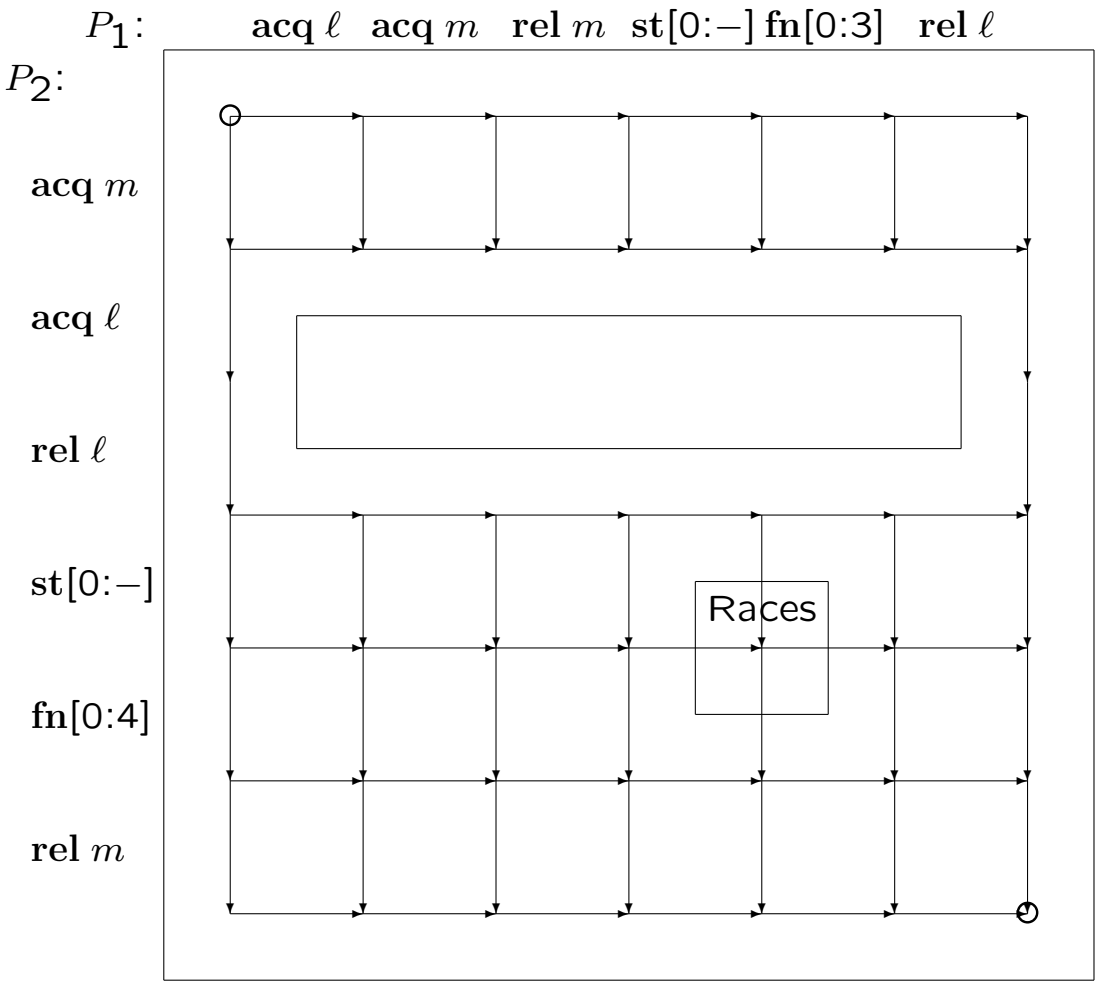
$$(\textbf{with } \ell \textbf{ do } ((\textbf{with } m \textbf{ do skip}); [0] := 3))$$
$$\| \; (\textbf{with } m \textbf{ do } ((\textbf{with } \ell \textbf{ do skip}); [0] := 4)).$$



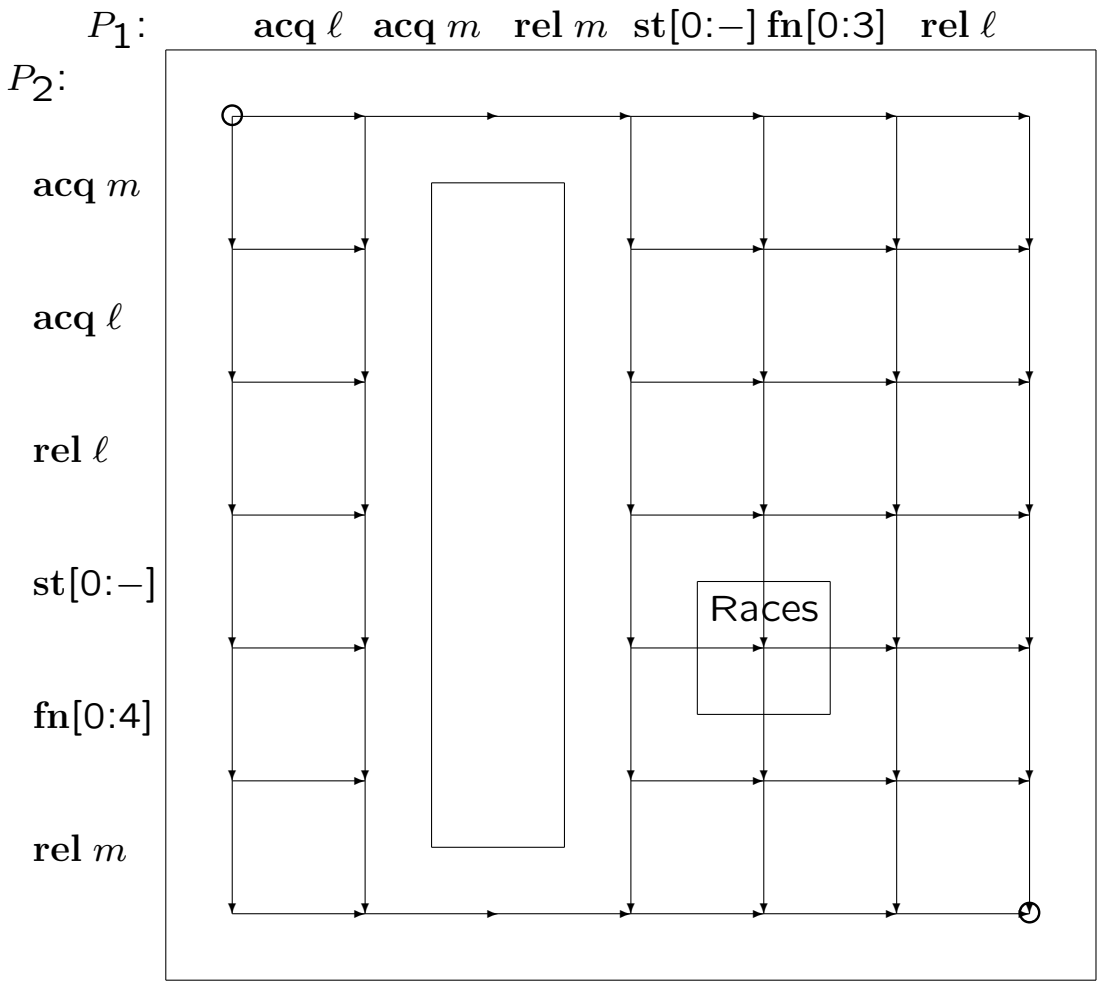
All Possible Interleavings



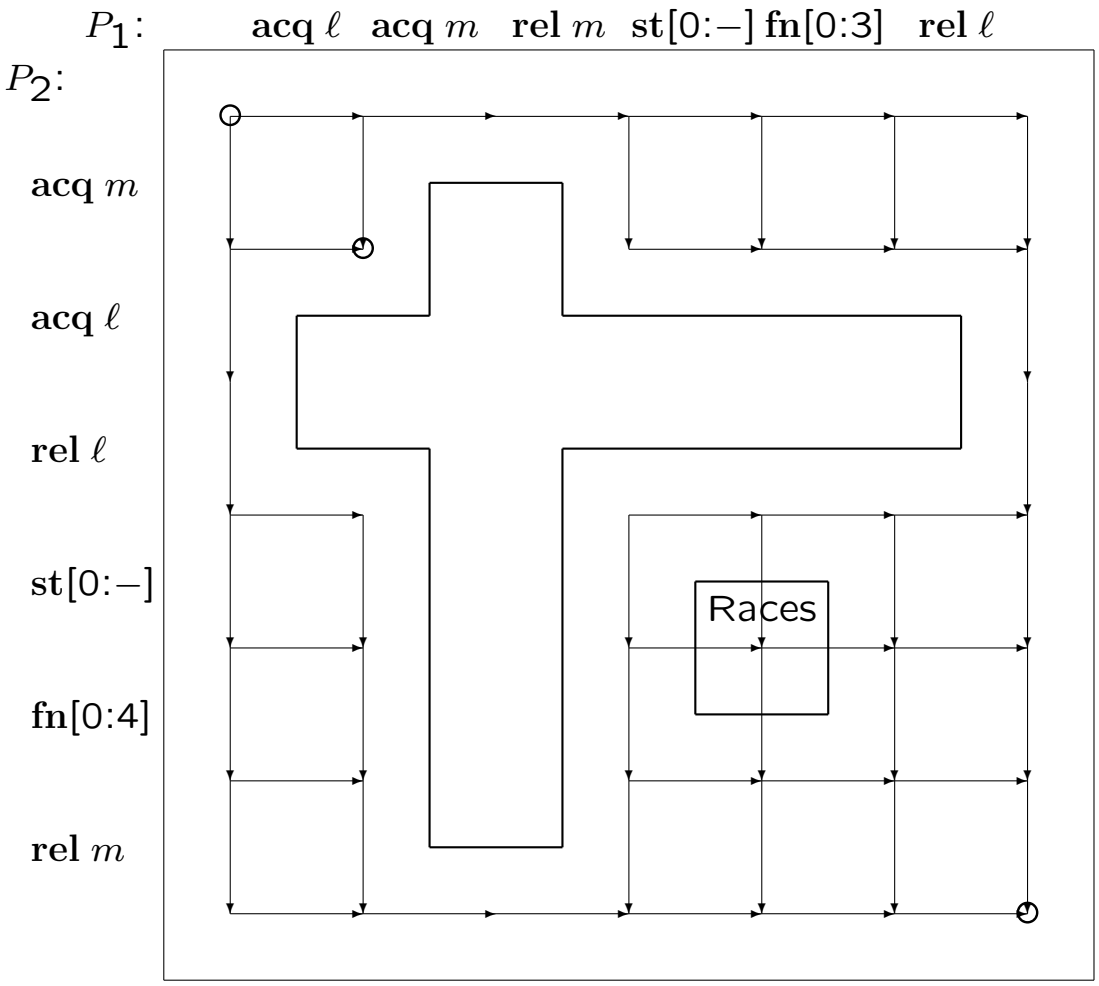
Exclusion by ℓ



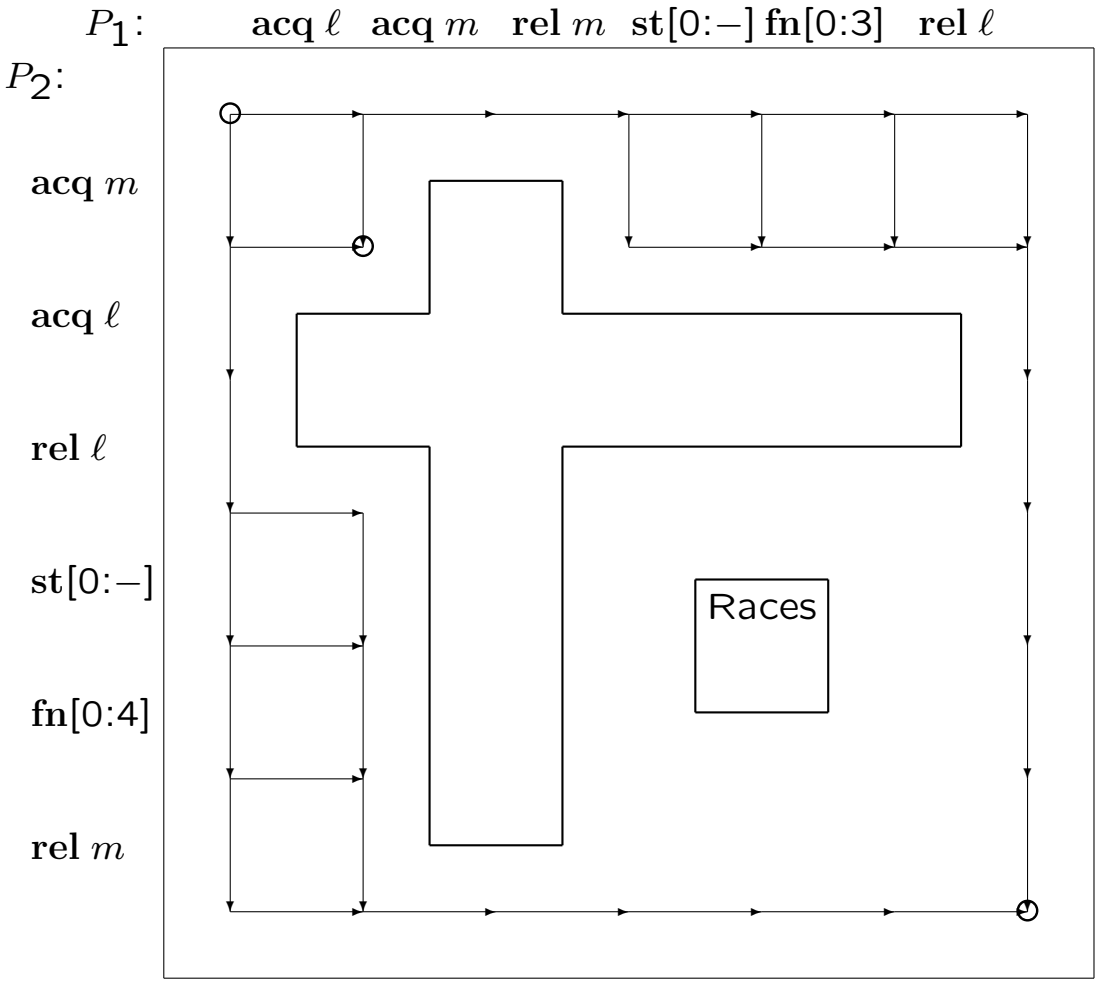
Exclusion by m



The Combined Exclusion



Excluding Unreachable Nodes



How to Prove it: Use an Auxiliary Variable

$\{0 \mapsto -\}$

$\mathbf{resource}\ \ell\ \mathbf{in}\ \mathbf{resource}\ m\ \mathbf{in}$

$\quad (\mathbf{with}\ \ell\ \mathbf{do}\ ((\mathbf{with}\ m\ \mathbf{do}\ p := 0); [0] := 3))$

$\quad \| (\mathbf{with}\ m\ \mathbf{do}\ ((\mathbf{with}\ \ell\ \mathbf{do}\ p := 1); [0] := 4))$

$\{0 \mapsto -\}$

The Resource Invariants

Let

$$R_\ell = \textbf{if } p = 0 \textbf{ then } 0 \mapsto - \textbf{ else emp}$$
$$R_m = \textbf{if } p = 0 \textbf{ then emp else } 0 \mapsto -$$

Then

$R_\ell * R_m$

iff $\textbf{if } p = 0 \textbf{ then } 0 \mapsto - * \textbf{emp else emp} * 0 \mapsto -$

iff $\textbf{if } p = 0 \textbf{ then } 0 \mapsto - \textbf{ else } 0 \mapsto -$

iff $0 \mapsto -$

and

$$R_\ell * (p = 0 \wedge \textbf{emp}) \text{ iff } 0 \mapsto - * (p = 0 \wedge \textbf{emp})$$
$$R_m * (p \neq 0 \wedge \textbf{emp}) \text{ iff } 0 \mapsto - * (p \neq 0 \wedge \textbf{emp})$$

Thus

$$
\begin{array}{l}
\{R_\ell * R_m\} \\
\{0 \mapsto -\} \\
p := 0 \\
\{0 \mapsto - * (p = 0 \wedge \mathbf{emp})\} \\
\{R_\ell * R_m * (p = 0 \wedge \mathbf{emp})\}
\end{array}
$$

$$
\begin{array}{l}
\{R_\ell\} \\
\mathbf{with}\ m\ \mathbf{do}\ p := 0; \\
\{R_\ell * (p = 0 \wedge \mathbf{emp})\} \\
\{0 \mapsto - * (p = 0 \wedge \mathbf{emp})\} \\
[0] := 3 \\
\{0 \mapsto - * (p = 0 \wedge \mathbf{emp})\} \\
\{R_\ell * (p = 0 \wedge \mathbf{emp})\} \\
\{R_\ell\}
\end{array}
$$

$\{\mathbf{emp}\}$
$\mathbf{with}\ \ell\ \mathbf{do}\ ((\mathbf{with}\ m\ \mathbf{do}\ p := 0)\ ;\ [0] := 3)$
$\{\mathbf{emp}\}$

and similarly

$\{\mathbf{emp}\}$
$\mathbf{with}\ m\ \mathbf{do}\ ((\mathbf{with}\ \ell\ \mathbf{do}\ p := 1)\ ;\ [0] := 4)$
$\{\mathbf{emp}\}$

So

$\{\mathbf{emp} * \mathbf{emp}\}$
$\mathbf{with}\ \ell\ \mathbf{do}\ ((\mathbf{with}\ m\ \mathbf{do}\ p := 0)\ ;\ [0] := 3)$
$\|\ \mathbf{with}\ m\ \mathbf{do}\ ((\mathbf{with}\ \ell\ \mathbf{do}\ p := 1)\ ;\ [0] := 4)$
$\{\mathbf{emp} * \mathbf{emp}\}$

and finally

$\{0 \mapsto -\}$

$\{R_\ell * R_m\}$

$\mathbf{resource}\ \ell\ \mathbf{in}\ \mathbf{resource}\ m\ \mathbf{in}$

$\quad (\mathbf{with}\ \ell\ \mathbf{do}\ ((\mathbf{with}\ m\ \mathbf{do}\ p := 0); [0] := 3))$

$\quad \| (\mathbf{with}\ m\ \mathbf{do}\ ((\mathbf{with}\ \ell\ \mathbf{do}\ p := 1); [0] := 4))$

$\{R_\ell * R_m\}$

$\{0 \mapsto -\}$

Note that the resources ℓ and m each have half permission for the variable p (in the sense of Bornat).