

Greedy Elim-1-2

Goal 1) Use GE to remove all degree 1 & 2 nodes.
2) But keep spanning tree.

Greedy Elim

Input: $G = (V, E, w)$ & ST T of G

Output: $\hat{G} = (\hat{V}, \hat{E}, \hat{w})$ & ST $\hat{T}$ of $\hat{G}$

1) $\hat{G} := G$; $\hat{E}_T = E_T$

2) while $\exists v$ of degree ≤ 2 do

if $\text{degree}(v) = 1$ then remove v .

if $\text{degree}(v) = 2$ $u_1 \xrightarrow{w_1} v \xrightarrow{w_2} u_2$

$$w \leftarrow \frac{w_1 \cdot w_2}{w_1 + w_2}$$

— tree edges
~ non-tree edges

Case 1 $u_1 \text{---} v \text{---} u_2 \Rightarrow u_1 \overset{w}{\text{---}} u_2$ (2 samples)

Case 2 $u_1 \text{---} v \text{---} u_2 \Rightarrow u_1 \overset{w+w_3}{\text{---}} u_2$

Case 3 $u_1 \text{ --- } v \text{ --- } u_2 \Rightarrow u_1 \xrightarrow{w+w_3} u_2$

Claim $\hat{T}$ is ST of $\hat{G}$ & $|\text{nontree}(\hat{G})| \leq |\text{nontree}(G)|$

if $K = |\text{nontree}(G)| = K$ then $|\hat{V}| \leq 2K$ & $|\hat{E}| \leq 3K$

pf

Let $n_1 = \# \text{leaves in } \hat{T}$
 $n_2 = \# \text{degree 2 node in } \hat{T}$
 $n_3 = \# \text{ " } \geq 3 \text{ " "}$

note $n_3 \leq n_1$ thus $n = n_1 + n_2 + n_3 \leq 2n_1 + n_2 \leq 2K$

$\hat{G}$ has $2K-1$ tree edges & K nontree edges

$$m \leq 3K$$

Claim $\text{Stretch}_T(e) \geq \text{Stretch}_{\hat{T}}(e)$

Gaussian Elim Details

Suppose $A \xrightarrow{GE} G$

ie $L_k \cdots L_1 A L_1^T \cdots L_k^T = \begin{pmatrix} I & 0 \\ 0 & G \end{pmatrix}$

$$A = L_1^{-1} \cdots L_k^{-1} \begin{pmatrix} I & 0 \\ 0 & G \end{pmatrix} L_k^{-T} \cdots L_1^{-T}$$

Good Solve $Ax = b$

[project & solve] $\begin{pmatrix} I & 0 \\ 0 & G \end{pmatrix} L_k^{-1T} \cdots L_1^{-1T} x = L_k^{-1} \cdots L_1^{-1} b = b'$

solve $\begin{pmatrix} I & 0 \\ 0 & G \end{pmatrix} y = b' \Rightarrow Gy_G = b'_G$

$$y = L_k^{-1T} \cdots L_1^{-1T} x$$

[expand] $x = L_1^T \cdots L_k^T y$

Building Precondition Chain

Def $C = \{G = G_1, H_1, G_2, \dots, G_d\}$ chains of graphs.

$K = \{K_1, \dots, K_{d-1}\}$ list of numbers

$\{C, K\}$ is a good Precon Chain if $\exists \{\mu_1, \dots, \mu_d\}$ st.

$$1) G_i \leq H_i \leq K_i G_i$$

$$2) G_{i+1} = \text{Greedy Elim}(H_i)$$

$$3) |G_i| = \mu_i$$

$$4) \mu_1, \mu_2 \leq m$$

$$5) \mu_i / \mu_{i+1} > \lceil c \sqrt{K_i} \rceil \quad \forall i$$

$$6) K_i \geq K_{i+1}$$

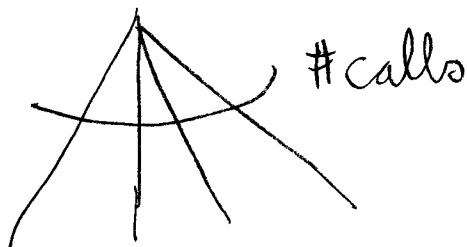
$$7) \mu_d \text{ bdd.}$$

Determining Work

2-Level

Probs

$$G_i x = b$$



$$G_{i+1} x = b_{i+1}$$

~~A~~

$$G_{i+1} x = b_{i+1}$$

~~A~~

#subcalls.

$$\text{Level}_i \text{ work} \equiv O(|G_i| \cdot \# \text{calls})$$

#calls
Matrix-Vector products

$$\text{Level}_{i+1} \text{ work} \equiv$$

$$O(\# \text{calls} \cdot |G_{i+1}| \cdot \# \text{subcalls})$$

$$\text{Goal: } \frac{\text{Level}_{i+1} \text{ work}}{\text{Level}_i \text{ work}} \leq C < 1$$

$$\frac{\# \text{ calls} \cdot |G_{i+1}| \cdot \# \text{ sub calls}}{|G_i| \cdot \# \text{ calls}} \leq \# \text{ calls} \cdot \frac{|G_{i+1}|}{|G_i|} \quad (6)$$

$$< \frac{\# \text{ calls}}{c_r \sqrt{K_i}} = \frac{\sqrt{K_i}}{c_r \sqrt{K_i}} < \frac{1}{c_r} < 1$$