

# Low-Rank PSD Approximation in Input-Sparsity Time

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# Low Rank Approximation

- A is an  $n \times n$  matrix
- A is typically well-approximated by low rank matrix
  - E.g., high rank because of noise
- **Goal:** find a low rank matrix approximating A
  - Easy to store in factored form
  - Data more interpretable

# What is a Good Low Rank Approximation?

## Singular Value Decomposition (SVD)

Any matrix  $A = U \Sigma V^T$

- $U$  has orthonormal columns
- $\Sigma$  is diagonal with non-increasing positive entries down the diagonal
- $V^T$  has orthonormal rows
  
- Rank-k approximation:  $A_k = U_k \Sigma_k V_k^T$

$$\begin{pmatrix} \mathbf{A} \end{pmatrix} = \begin{pmatrix} \mathbf{U}_k \end{pmatrix} \begin{pmatrix} \Sigma_k \end{pmatrix} \begin{pmatrix} \mathbf{V}_k^T \end{pmatrix} + \begin{pmatrix} \mathbf{E} \end{pmatrix}$$

$$A_k = \operatorname{argmin}_{\text{rank } k \text{ matrices } B} |A-B|_F$$

$$(|C|_F = (\sum_{i,j} C_{i,j}^2)^{1/2})$$

Computing  $A_k$  exactly is expensive

# Approximate Low Rank Approximation

- [CW13] output a rank  $k$  matrix  $A'$ , so that with probability  $> 2/3$ ,

$$\|A - A'\|_F \leq (1 + \varepsilon) \|A - A_k\|_F,$$

in  $\text{nnz}(A) + n \cdot \text{poly}(k/\varepsilon)$  time

# Structure-Preserving Low Rank Approximation

- Let  $A$  be an arbitrary  $n \times n$  matrix
- Instead of just finding a rank- $k$  matrix  $A'$  for which  $\|A-A'\|_F$  is small, suppose we also require  $A'$  to be positive semidefinite (PSD)
  - $A'$  is symmetric and all eigenvalues are non-negative
- Covariance matrices, kernel matrices, Laplacians are PSD
  - Approximate them for efficiency
- Roundoff errors may make a PSD matrix non-PSD
  - We do not assume  $A$  is PSD but want  $A'$  to be PSD

# Structure-Preserving Low Rank Approximation

- **Goal:** output a PSD rank- $k$  matrix  $A'$  for which  $\|A-A'\|_F$  is small
- Can assume  $A$  is symmetric
  - $A = A^{\text{sym}} + A^{\text{asym}}$ , where  $A_{i,j}^{\text{sym}} = \frac{A_{i,j}+A_{j,i}}{2}$  and  $A_{i,j}^{\text{asym}} = \frac{A_{i,j}-A_{j,i}}{2}$
  - $\|A - A'\|_F^2 = \|A^{\text{sym}} - A'\|_F^2 + \|A^{\text{asym}}\|_F^2$
  - Compute  $A^{\text{sym}}$  in  $\text{nnz}(A)$  time
- What is the best PSD rank- $k$  approximation  $A_{k,+}$  to  $A$ ?
- **Lemma:**  $A_{k,+}$  is obtained by zeroing out all but the top  $k$  *positive* eigenvalues in eigendecomposition of  $A$ 
  - If  $A = U D U^T$ , then  $A_{k,+} = U D_{k,+} U^T$
  - If  $A$  has fewer than  $k$  positive eigenvalues, zero out all except these eigenvalues

# Our Result for PSD Low Rank Approximation

- **(PSD low rank approximation result):** In  $\text{nnz}(A) + n \text{ poly}(k/\epsilon)$  time, can find a PSD rank- $k$   $A'$  so that

$$\|A - A'\|_F \leq (1 + \epsilon) \|A - A_{k,+}\|_F$$

- Previous work
  - “Nystrom method” based on uniform sampling requires incoherence assumptions on  $A$
  - [GM] Weaker bound  $\|A - A'\|_F \leq \|A - A_{k,+}\|_F + \epsilon \|A - A_{k,+}\|_*$  where  $\|\cdot\|_*$  is nuclear norm
  - [WLZ] Running time at least  $n^2 k / \epsilon$  and  $A'$  has a larger rank  $k/\epsilon$

# Our Result for PSD Column Subset Selection

- **(PSD Column Subset Selection):** In  $\text{nnz}(A)\log n + n \text{poly}(k(\log n)/\epsilon)$  time, find a subset  $C$  of  $O(k/\epsilon)$  columns of  $A$  so that  $A' = CUC^T$  is rank- $k$ , PSD, and

$$\|A - A'\|_F \leq (1 + \epsilon) \|A - A_{k,+}\|_F$$

- Column subsets preserve sparsity, interpretability
- Most previous results require incoherence assumptions or achieve weaker guarantees in terms of the nuclear norm
- [WLZ] Get  $O\left(\frac{k}{\epsilon}\right)$  columns but running time at least  $n^2 k/\epsilon$  and  $\text{rank}(A') = k/\epsilon$

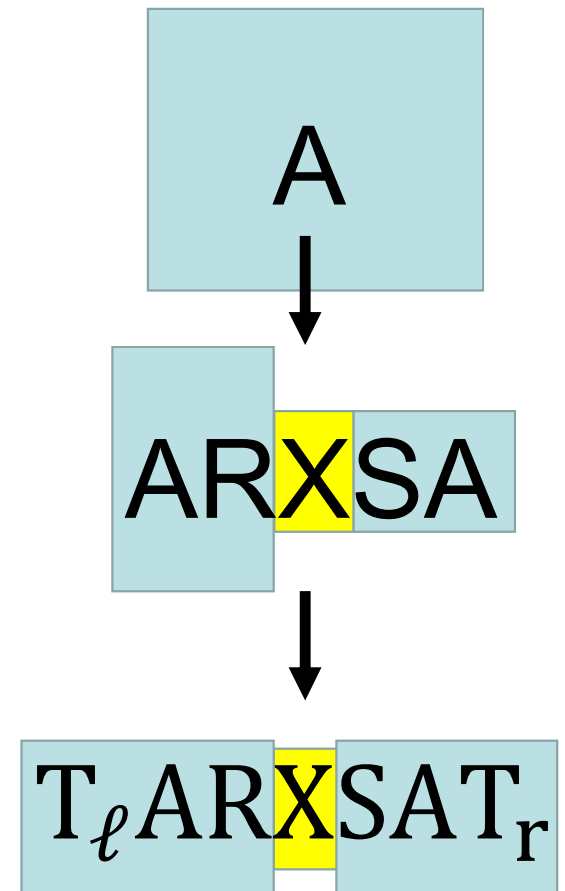


# Talk Outline

- Low Rank PSD Approximation
  - Pitfalls of Usual Approach
  - SF( $\epsilon$ , k) Property
  - Error Term Lemmas
  - Solving a Small Problem Quickly
- Low Rank PSD Column Subset Selection

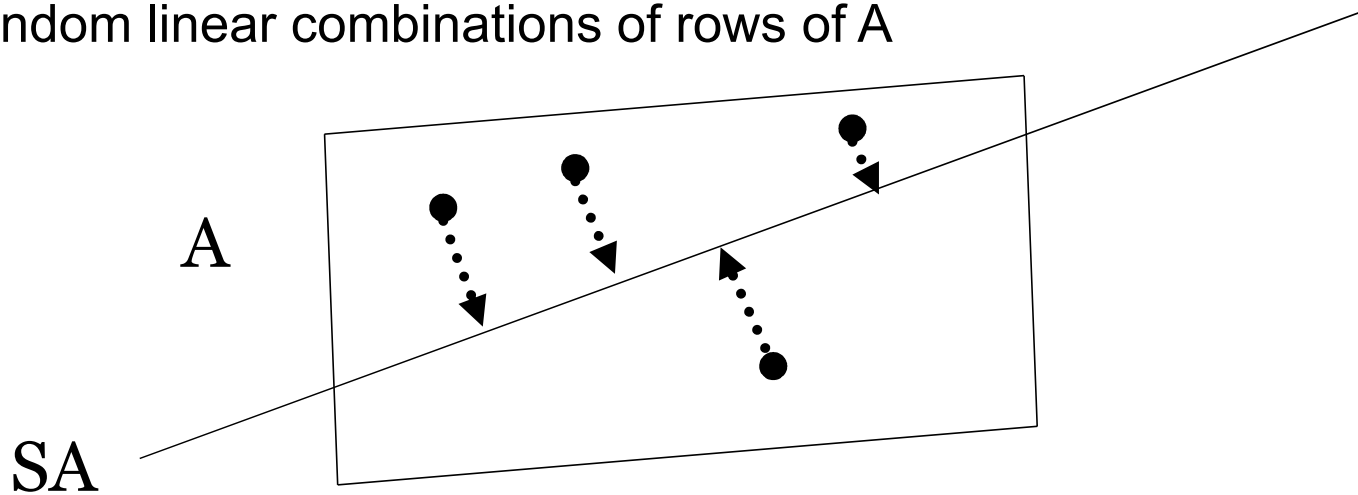
# Pitfalls of Usual Approach

- [CW09,HMT,BWZ] Choose random matrices R, S
  - R has  $k/\epsilon$  columns and S has  $k/\epsilon$  rows
  - $A' = \operatorname{argmin}_{\text{rank-}k \text{ ARXSA}} \|ARXSA - A\|_F^2$  and  $\|A-A'\|_F \leq (1+\epsilon) \|A-A_k\|_F$
  - Compute AR and SA in  $\text{nnz}(A)$  time
- To solve for A', solve
  - $\min_{\text{rank-}k X} \|T_\ell ARXSAT_r - T_\ell AT_r\|_F^2$ , where  $T_\ell, T_r$  are random
  - [FT]  $\text{poly}(k/\epsilon)$  time
- Need  $R = S^T$  for A' to be symmetric, but analysis requires  $\|x^T SAR\|_2 = \Theta(\|x^T SA\|_2)$  for all x !



# SF( $\epsilon, k$ ) Property

- Given  $n \times n$  input matrix  $A$
- Compute  $S \cdot A$  using a sketching matrix  $S$  with  $k/\epsilon \ll n$  rows.  $S \cdot A$  takes random linear combinations of rows of  $A$



- Let  $P$  be the  $n \times n$  projection matrix onto the row span of  $SA$
- [S,CW,MM,NN,...]  $|A(I - P)|_F^2 \leq (1 + \epsilon)|A - A_k|_F^2$
- $P$  is **SF( $\epsilon, k$ )** if  $|A(I - P)|_2^2 \leq \frac{\epsilon}{k}|A - A_k|_F^2$  (recall  $|B|_2 = \sup_x |Bx|_2/|x|_2$ )

# SF( $\epsilon$ , $k$ ) Property

- SF( $\epsilon$ ,  $k$ ) property implies usual theory

$$|A(I - P)|_F^2$$

$$= |A_k(I - P)|_F^2 + |(A - A_k)(I - P)|_F^2 \text{ by Pythagorean theorem}$$

$$\leq k \cdot |A_k(I - P)|_2^2 + |(A - A_k)(I - P)|_F^2 \text{ since } A_k \text{ has rank } k$$

$$\leq k \cdot |A(I - P)|_2^2 + |(A - A_k)(I - P)|_F^2 \text{ since } |A(I - P)x|_2 \geq |A_k(I - P)x|_2 \text{ for all } x$$

$$\leq k \left(\frac{\epsilon}{k}\right) |A - A_k|_F^2 + |(A - A_k)(I - P)|_F^2 \text{ by SF}(\epsilon, k) \text{ property}$$

$$\leq \epsilon |A - A_k|_F^2 + |A - A_k|_F^2 \text{ since projections don't increase norms}$$

# A Basic Lemma

- **Lemma:** For symmetric  $A, B$  with  $(A-B)B = 0$ , and projection matrix  $P$ ,

$$|A - PBP|_F^2 = |A - B|_F^2 + |B - PBP|_F^2 + 2\text{Tr}(A - B)(I - P)BP$$

# SF( $\epsilon$ , k) Projections Give Good Solutions

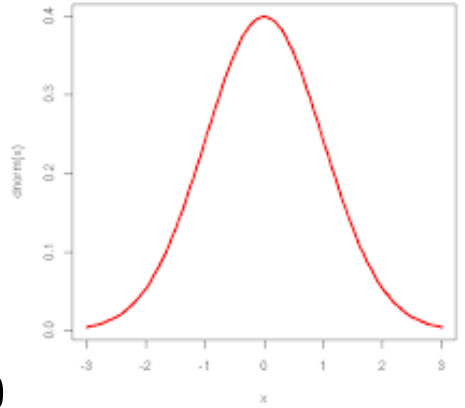
- **Lemma:** If P is SF( $\epsilon$ , k) for A, then  $|A - PA_{k,+}P|_F \leq (1 + O(\epsilon)) |A - A_{k,+}|_F$
- Proof: Since  $(A - A_{k,+})A_{k,+} = 0$ , can apply the basic lemma:

$$|A - PA_{k,+}P|_F^2 = |A - A_{k,+}|_F^2 + \underbrace{|A_{k,+} - PA_{k,+}P|_F^2}_{\text{Use SF}(\epsilon, k) \text{ property of } P} + 2\text{Tr}(A - A_{k,+})(I - P)A_{k,+}P$$

Use SF( $\epsilon$ , k) property of P to show these terms are  $O(\epsilon)|A - A_{k,+}|_F^2$

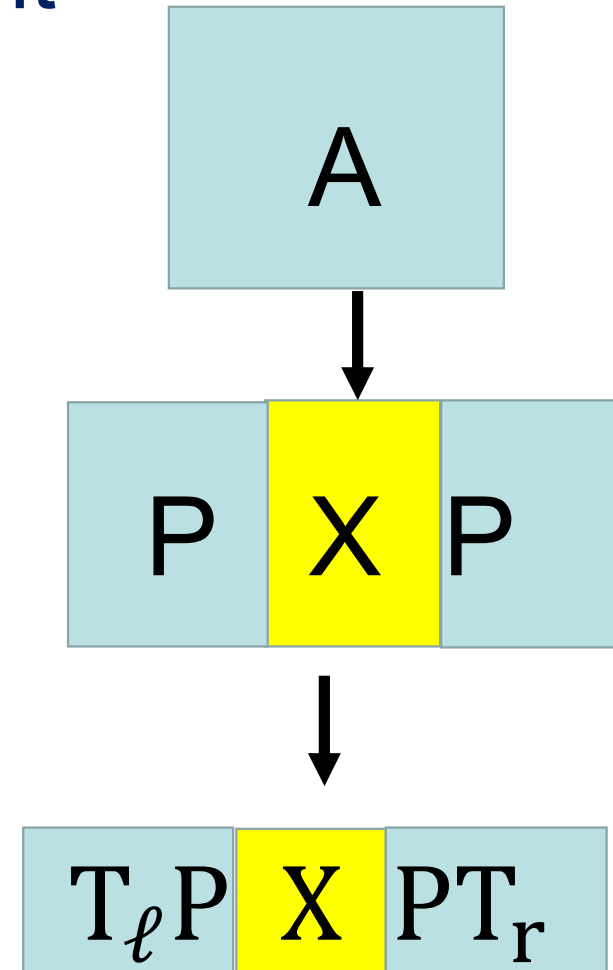
# Finding an $SF(\epsilon, k)$ Projection

- [CEMMP] implies if  $S$  is  $\text{poly}(k/\epsilon) \times n$  i.i.d. Gaussian, Fast Hadamard Transform, or a Sparse Embedding Matrix, then  $P = (SA)^{-1}SA$  is  $SF(\epsilon, k)$

$$\left[ \begin{array}{c} \text{density} \\ 0.4 \\ 0.3 \\ 0.2 \\ 0.1 \\ 0.0 \\ -3 \\ -2 \\ -1 \\ 0 \\ 1 \\ 2 \\ 3 \end{array} \right], P \cdot H \cdot D, \left[ \begin{array}{ccccc} 1 & 0 & 0 & -1 & 0 & 1 \\ 0 & 0 & 1 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 1 & 0 \\ 0 & -1 & 0 & 0 & 0 & 0 \end{array} \right]$$


# Completing the Argument

- If  $P$  is  $SF(\epsilon, k)$  for  $A$ , then  
 $A' = \operatorname{argmin}_{\text{rank-}k \text{ PSD } PXP} |A - PXP|_F$  satisfies  
 $|A' - A|_F \leq (1 + \epsilon)|A - A_{k,+}|_F$
- We can also find  $P = (SA)^{-1}SA$  in factored form in  $\text{nnz}(A)$  time
- How to solve for  $A'$ ?
  - Multiply by small random matrices  $T_\ell, T_r$ ,  
solve  $A' = \operatorname{argmin}_{\text{rank-}k \text{ PSD } PXP} |T_\ell A T_r - T_\ell P X P T_r|_F$
- **Tiny problem but is it in polynomial time?**
  - We show how to solve it up to a  $(1 + \epsilon)$  –factor quickly





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- Low Rank PSD Column Subset Selection

# Low Rank Column Subset Selection

- [CEMMP], together with a composition lemma, gives a  $P = (SA)^{-1}SA$  where  $S$  is a sampling and rescaling matrix, such that  $P$  is  $SF(\epsilon, k)$  for  $A$  and  $SA$  can be found in  $\text{nnz}(A)\log n + n \text{poly}\left(\frac{k}{\epsilon}\right)$  time
- $S$  samples  $O\left(\frac{k}{\epsilon}\right)$  rows of  $A$
- Our earlier lemma implies there is a rank- $k$  PSD solution  $A' = PXP$
- Our earlier procedure finds  $X$

# Other Results

- **Symmetric Matrices**

- Analogous results for  $\text{nnz}(A)$  time for finding a rank- $k$  symmetric approximation to  $A$ , and for column/row subset selection

- **Low Rank Approximation with Tail Guarantee**

- Let  $t = 2k/\epsilon$
- A PSD rank- $k$  matrix  $A'$  can be found in  $\text{nnz}(A) + n \text{ poly}\left(\frac{k}{\epsilon}\right)$  time with

$$\|A - A'\|_F^2 \leq \|A - A_{k,+}\|_F^2 + \|A_{t+k} - A_t\|_F^2$$

# Conclusion and Open Questions

- First rank-k PSD approximation of an arbitrary matrix  $A$  in  $\text{nnz}(A)$  time
- Optimal  $O\left(\frac{k}{\epsilon}\right)$  columns/rows for rank-k PSD subset selection, in  $\text{nnz}(A) \log n$  time
- Similar results for symmetric approximations
- Should be able to improve the time for column/row subset selection to  $\text{nnz}(A)$  using known techniques
- High-level question – quickly find low rank approximations with additional structure, such as being PSD

# What is $A_{k,+}$ ?

- $A = U D U^T$  is eigendecomposition
- $\|A - A_{k,+}\|_F^2 = \|D - U^T A_{k,+} U\|_F^2$ , where  $U^T A_{k,+} U$  is PSD and rank  $k$
- Let  $D_{k,+}$  be the best rank- $k$  PSD approximation to  $D$
- $\|D - U^T A_{k,+} U\|_F^2 \geq \|D - D_{k,+}\|_F^2 = \|UDU^T - UD_{k,+}U^T\|_F^2 = \|A - UD_{k,+}U^T\|_F^2$
- $A_{k,+} = UD_{k,+}U^T$

*But what is  $D_{k,+}$ , the best rank  $k$  PSD approximation to diagonal matrix  $D$ ?*

$$\left\| \begin{bmatrix} D_+ & 0 \\ 0 & D_- \end{bmatrix} - \begin{bmatrix} Y_t & Y_r \\ Y_l & Y_b \end{bmatrix} \right\|_F^2 \geq \|D_+ - Y_t\|_F^2 + \|D_- - Y_b\|_F^2$$

$D$ 
 $D_{k,+}$

*$D_{k,+}$  is diagonal matrix with top  $k$  non-negative eigenvalues of  $A$*

# Solving the Small Problem

- How to solve for  $A' = \operatorname{argmin}_{\text{rank-}k \text{ PSD PXP}} |T_\ell A T_r - T_\ell P X P T_r|_F$ ?
- Can find  $A'$  minimizing this up to a  $1+\epsilon$  factor
- Write  $SA = R\Sigma W^T$  in its SVD, so  $P = (SA)^-SA = WW^T$
- $A' = \operatorname{argmin}_{\text{rank-}k \text{ PSD PXP}} |T_\ell A T_r - (T_\ell W)W^T X W (W^T T_r)|_F$
- Write  $(T_\ell W) = U_\ell \Sigma_\ell V_\ell^T$  and  $(W^T T_r) = U_r \Sigma_r V_r^T$  in their SVD, use that all their singular values are  $1 \pm \epsilon$ , and after some algebra,

$$X = \left[ \frac{M+M^T}{2} \right]_{k,+} \quad \text{where } M = V_\ell \Sigma_\ell^{-1} U_\ell^T T_\ell A T_r V_r \Sigma_r^{-1} U_r^T$$