

More NP-Completeness

15-750
4/23/10

Graph Coloring

Through out $G = (V, E)$ undirected

Def A k -coloring is a fn $C: V \rightarrow \{0, 1, \dots, k-1\}$
s.t. the end-points of each edge have distinct colors.

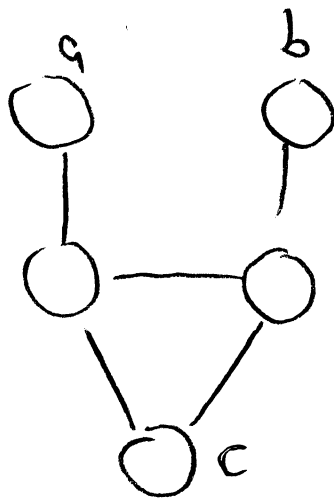
Def $3\text{-Color} \equiv \{G \mid G \text{ is } 3\text{-colorable}\}$

Thm 3-Color is NP-Complete

the color gadget:

$C \equiv 3\text{-coloring}$

Colors = $\{0, 1, 2\}$

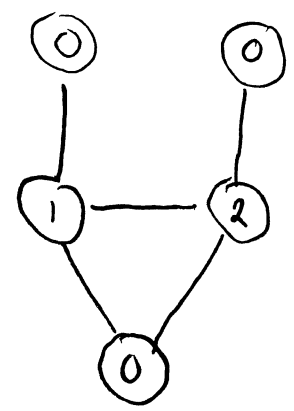


Critical Facts

1) $\forall c$ if $c(a) = c(b)$ then $c(a) = c(c)$

pf WLOG $c(a) = c(b) = 0$

Thus

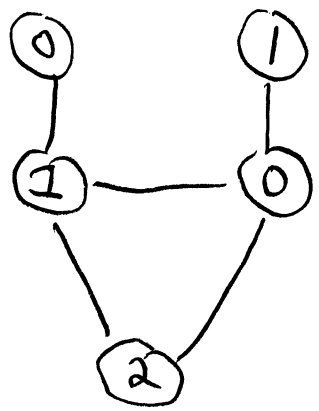
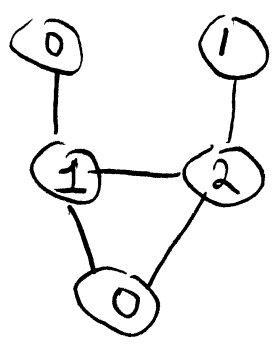


2) $\forall x \neq y, z \in \{0, 1, 2\} \exists c$ st

$$c(a) = x \wedge c(b) = y \wedge c(c) = z$$

WLOG $x = 0 \wedge y = 1$

WLOG $z \in \{0, 2\}$



Claim $3\text{-CNF} \leq_p 3\text{-Color}$

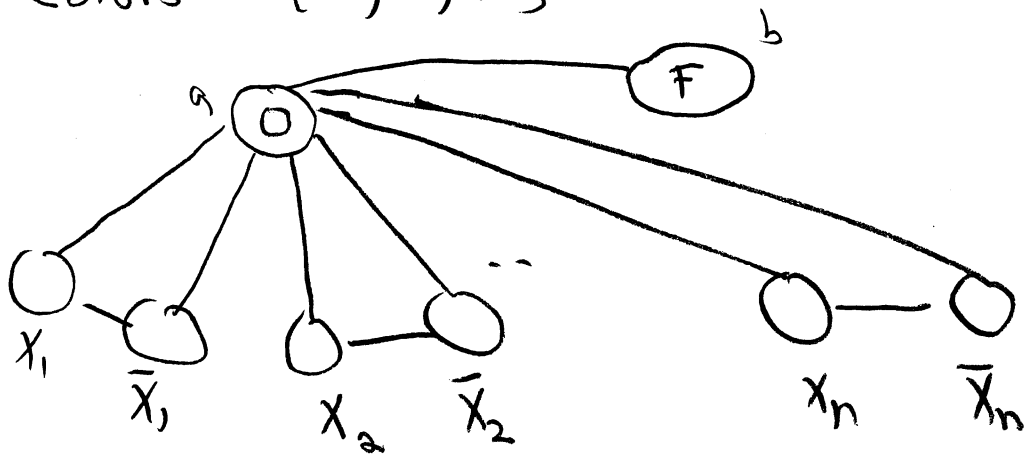
The construction of $f(\varphi)$

φ has clauses C_1, \dots, C_k

Variables x_1, \dots, x_n

A color wheel for each variable

Colors = $\{0, T, F\}$

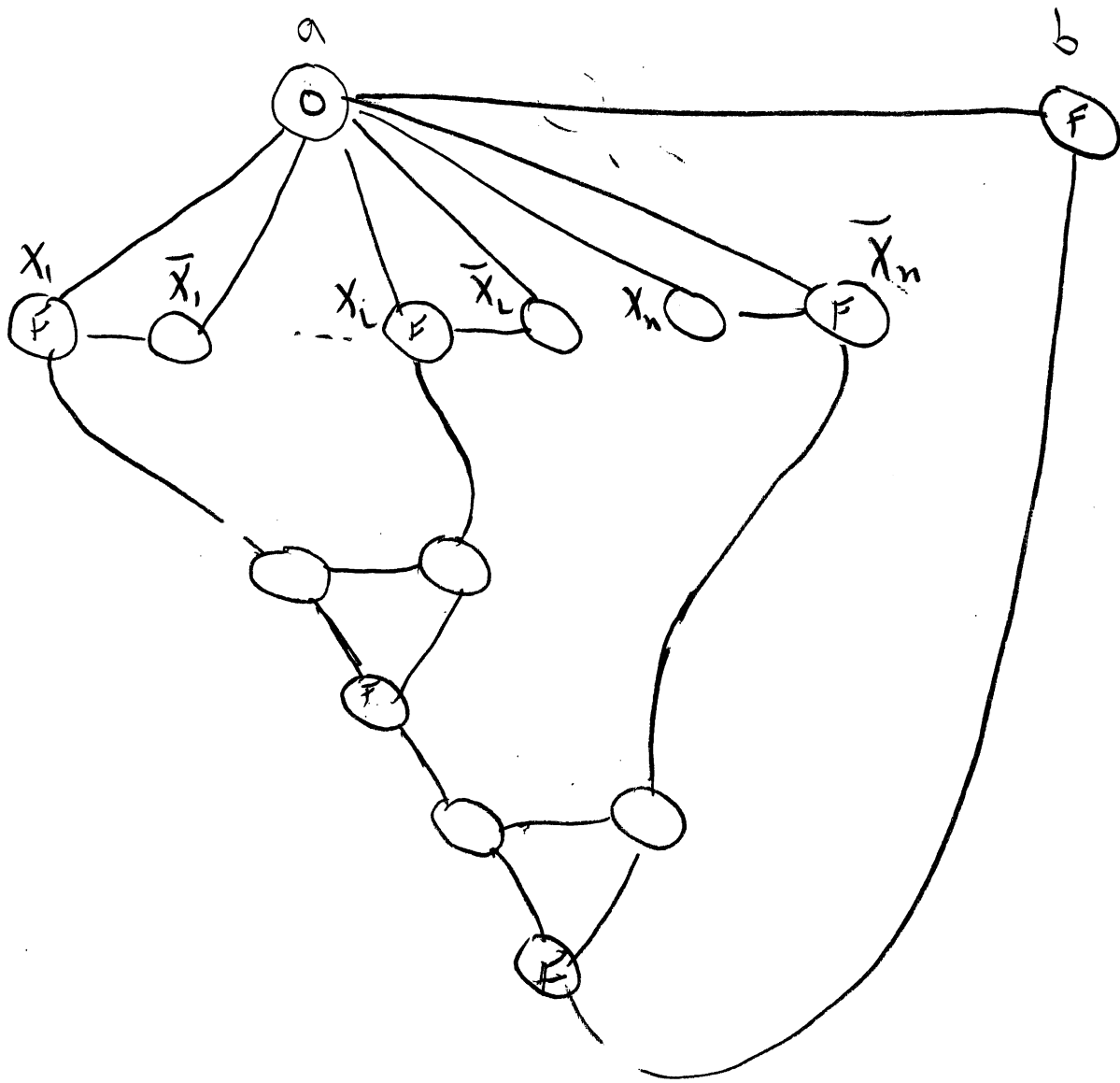


Our Color Wheels

WLOG $C(a) = 0$ & $C(b) = F$

Attaching 3-Clauses to Color Wheels

$$\text{eg } C_j = (x_1 \vee x_i \vee \bar{x}_n)$$



Add one for each clause

Claim φ is sat iff $G = f(\varphi)$ is 3-colorable

(\Rightarrow) 1) Use truth assignment to color color-wheels
with colors $\{0, T, F\}$

2) color gadgets

\Leftarrow view $C(A) = 0$ & $C(B) = F$ & remaining color as T .

Use colors of x_1, \dots, x_n as truth assignment.

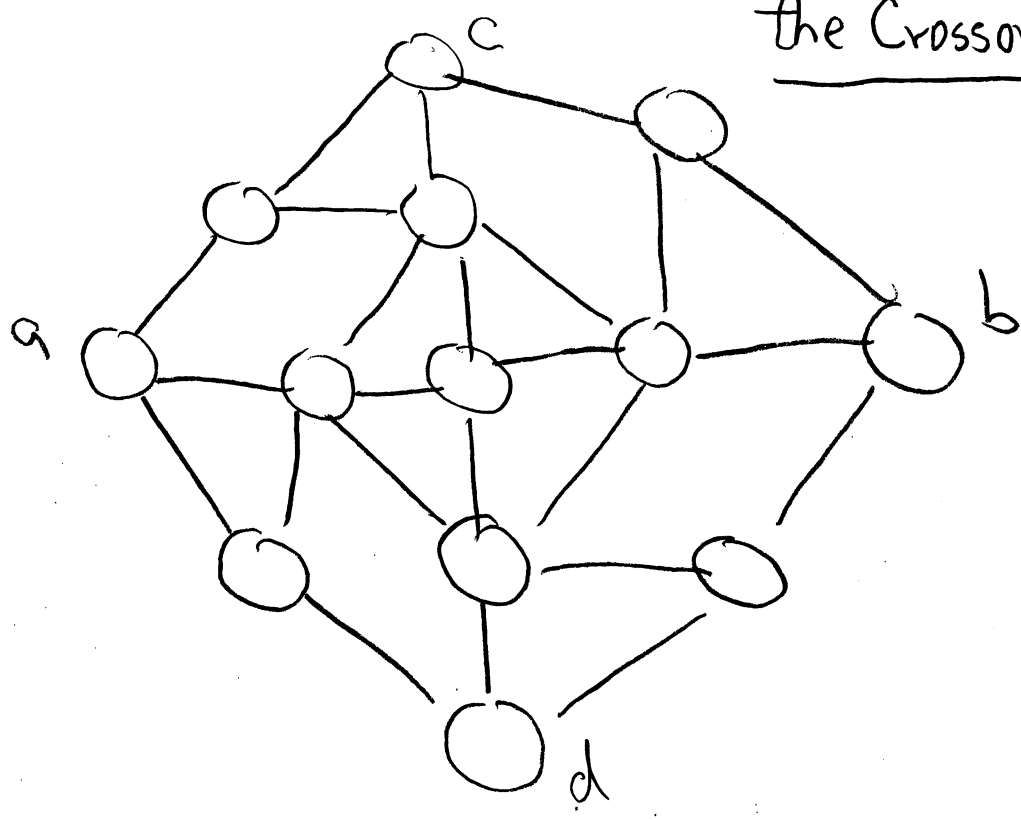
Def Planar- k -color = $\{ G \text{ planar} \mid G \text{ } k\text{-colorable} \}$

Thm Planar-4-color $\in P$

Thm Planar-3-color is NP-Complete

The gadget of all gadgets!

the Crossover gadget



1) IF C is a 3-coloring then

$$C(a) = C(b) \text{ \& } C(c) = C(d)$$

2) Given $x, y \in \{0, 1, 2\} \exists C$ st

$$x = C(a) = C(b) \text{ \& } y = C(c) = C(d)$$

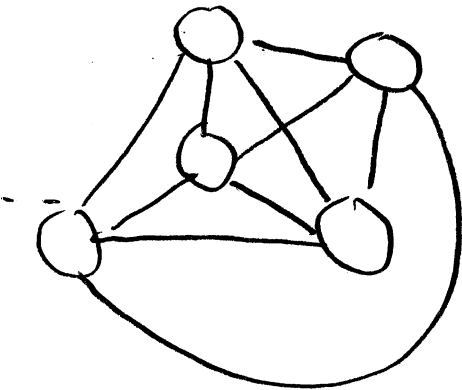
Given $G=(V,E)$ Goal construct $f(G)$

1) "Draw" G in plane with crossings

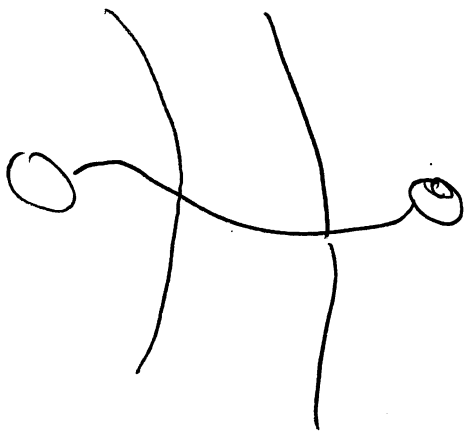
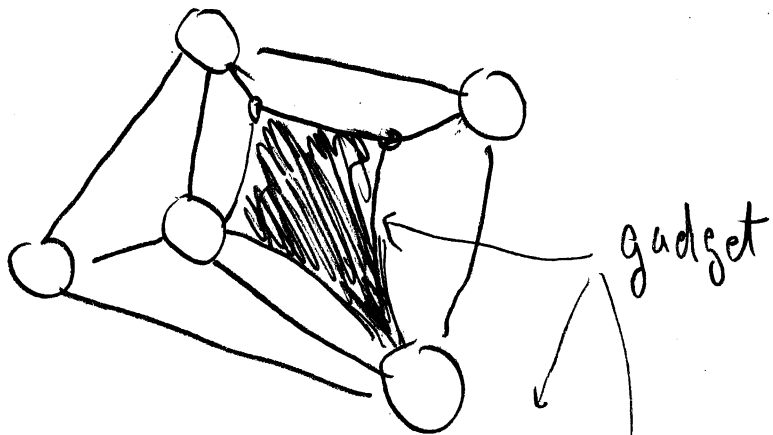
2) Insert a crossover gadget for each crossing

eg K_4

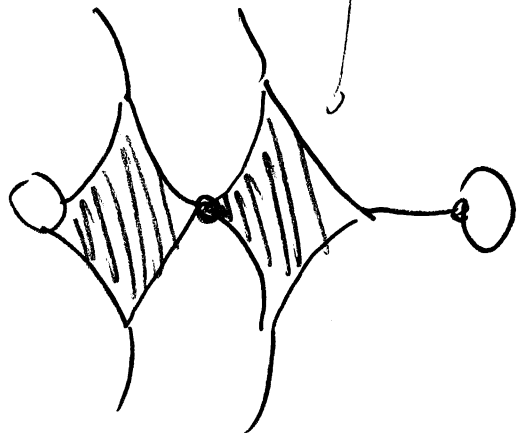
$G \equiv$



$f(G) \equiv$



$f \equiv$



Exact Cover

7

Input finite set X , Subsets $\{S_1, \dots, S_k\} = \mathcal{S}$

Question $\exists \mathcal{S}' \subseteq \mathcal{S}$ st. \mathcal{S}' is an exact cover

ie $\forall x \in X \exists ! s \in \mathcal{S}'$ s.t. $x \in s$

Thm Exact Cover is NP-Complete

Claim 3-Color \leq_p Exact Cover

$G = (V, E)$ is 3-color prob & $C = \{\text{red, blue, green}\}$

Def $N(u) = \{v \in V \mid (u, v) \in E\}$ $u \in V$

The Finite set X

For each $u \in V$ add four points

$$u, p_u^{\text{red}}, p_u^{\text{blue}}, p_u^{\text{green}}$$

For each $(u,v) \in E$, add six points

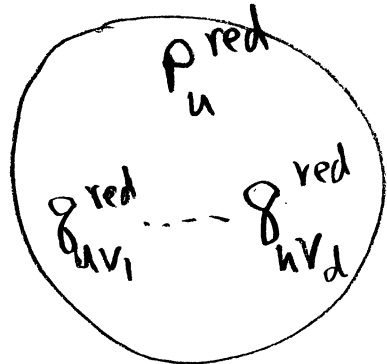
$$g_{uv}^{\text{red}}, g_{uv}^{\text{blue}}, g_{uv}^{\text{green}}, g_{vu}^{\text{red}}, g_{vu}^{\text{blue}}, g_{vu}^{\text{green}}$$

The subsets S

$\forall u \in V$ subsets $S_u^{\text{red}}, S_u^{\text{blue}}, S_u^{\text{green}}$

eg

$$S_u^{\text{red}} =$$



$$N(u) = \{v_1, \dots, v_d\}$$

Include $\{u, p_u^{\text{red}}\}, \{u, p_u^{\text{blue}}\}, \{u, p_u^{\text{green}}\}$

$$\left\{ \{g_{uv}^c, g_{vu}^{c'}\} \mid (u,v) \in E \wedge c \neq c' \right\}$$

Claim G is 3-colorable iff $f(G)$ has an exact cover

(\Rightarrow) Suppose $C: V \rightarrow \{\text{red, blue, green}\}$

proof by example! suppose

$(u, v) \in E$ $C(u) = \text{red}$ & $C(v) = \text{blue}$ $d(u) = 5$ & $d(v) = 3$

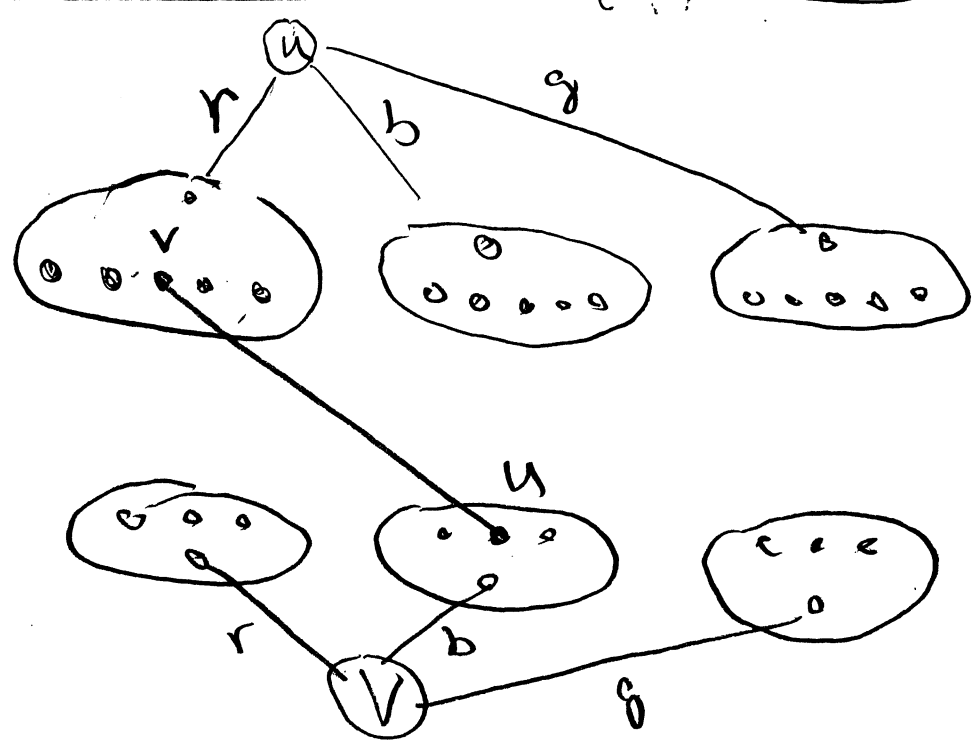
Include in cover

$\{u, P_u^{\text{red}}\} \{v, P_v^{\text{blue}}\}$

$\sum_u^c c' \neq \text{red}$

$\sum_v^c c' \neq \text{blue}$

$\{g_{uv}^{\text{red}}, g_{vu}^{\text{blue}}\}$



(\Leftarrow) The subsets $\{u, P_u^c\}$ give a coloring

More NP-Problems

10

Knap sack

Input Finite set S ; cost fn $w: S \rightarrow \mathbb{N}$
benefit fn $b: S \rightarrow \mathbb{N}$
int w, B

Question $\exists S' \subseteq S$ st

$$\sum_{a \in S'} w(a) \leq w \quad \& \quad \sum_{a \in S'} b(a) \geq B$$

Subset Sum

Input Finite set S ; $w: S \rightarrow \mathbb{N}$; int B

Question $\exists S' \subseteq S$ st

$$\sum_{a \in S'} w(a) = B$$

Partition Prob:

Input: Finite set S ; $w: S \rightarrow \mathbb{N}$

Question: $\exists S' \subseteq S$ st

$$\sum_{a \in S'} w(a) = \sum_{a \notin S'} w(a)$$

Partition \leq_p Subset Sum

$$\text{Set } B = \frac{1}{2} \sum_{a \in S} w(a)$$

Partition \leq_p Knapsack

$$\text{Set } b = W \text{ \& } W = B = \frac{1}{2} \sum_{a \in S} w(a)$$

Subset Sum \leq_p Partition

Set $S \subseteq \mathbb{N}$; $W: S \rightarrow \mathbb{N}$; B

Let $W = \sum_{a \in S} W(a) \wedge N \gg W$

New elements added to S a & b giving \bar{S}

$$W(a) = N - B \quad \& \quad W(b) = N - W + B$$

Claim S has subset sum iff \bar{S} has partition.

(\Rightarrow) Suppose S' is a subset sum
then $S' \cup \{a\}$ is a partition.

\Leftarrow Suppose \bar{S}' is a partition

if $a \in \bar{S}'$ then $b \notin \bar{S}'$

then $\bar{S}' - \{a\}$ is a subset sum.

Exact Cover \leq_p Subset Sum

pf Let (X, \mathcal{A}) be an exact-cover prob

We construct subset sum prob (\bar{X}, w) .

Set $\bar{X} = \mathcal{A}$ (Every subset is an element)

WLOG $X = \{0, 1, \dots, m-1\}$

for $x \in X$ $\#x = |\{A \in \mathcal{A} \mid x \in A\}|$

pick prime $p > \#x \forall x \in X$

$$\text{Def } w(A) = \sum_{x \in A} p^x \quad B = \sum_{x=0}^{m-1} p^x = \frac{p^m - 1}{p - 1}$$

Claim Exact Cover iff Subset Sum

By example

Suppose $A_1 = \{1, 3\}$ $A_2 = \{0, 3\}$

In p -ary notation

p^0 p^1 p^2 p^3 p^4 ... p^{m-1}

$$w(A_1) = (0, 1, 0, 1, 0, \dots, 0)$$

$$w(A_2) = (1, 0, 0, 1, 0, \dots, 0)$$

$$w(A_1) + w(A_2) = (1, 1, 0, 2, 0, \dots, 0)$$

Since $\#3 < p$ the 2 can not generate a carry.