

Dataflow Analysis (part 2)

15-411/15-611 Compiler Design

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Today

- Dataflow Analysis
 - reaching definitions
 - liveness
 - available expressions
 - very busy expressions
- Framework

A sample program

```
int fib10(void) {
    int n = 10;
    int older = 0;
    int old = 1;
    int result = 0;
    int i;

    if (n <= 1) return n;
    for (i = 2; i<n; i++) {
        result = old + older;
        older = old;
        old = result;
    }
    return result;
}
```

```
1:      n <- 10
2:      older <- 0
3:      old <- 1
4:      result <- 0
5:      if n <= 1 goto 14
6:      i <- 2
7:      if i > n goto 13
8:      result <- old + older
9:      older <- old
10:     old <- result
11:     i <- i + 1
12:     JUMP 7
13:     return result
14:     return n
```

Simple Constant Propagation

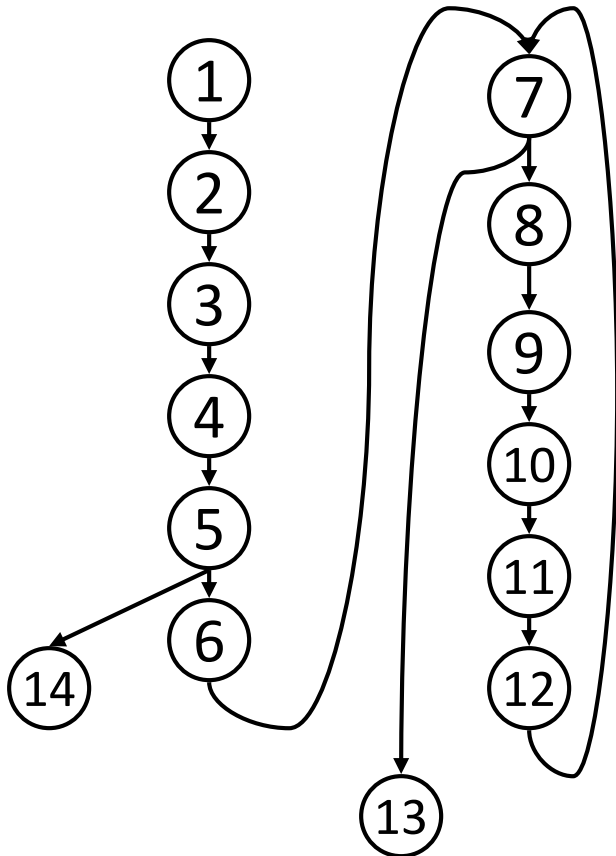
- Can we do SCP?
- How do we recognize it?
- What aren't we doing?
- Metanote:
 - keep opts simple!
 - Use combined power

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Reaching Definitions

A definition of variable v at program point d reaches program point u if there exists a path of control flow edges from d to u that does not contain a definition of v .



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```

Reaching Definitions (ex)

- 1 reaches 5, 7, and 14

14, Really?

Meta-notes:

- (almost) always conservative
- only know what we know
- Keep it simple:
 - What opt(s), if run before this would help
 - What about:

```
1: x ← 0
2: if (false) x ← -1
3: ... x ...
```

 - Does 1 reach 3?
 - What opt changes this?

```
1:      n ← 10
2:      older ← 0
3:      old ← 1
4:      result ← 0
5:      if n ≤ 1 goto 14
6:      i ← 2
7:      if i > n goto 13
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Calculating Reaching Definitions

A definition of variable v at program point d reaches program point u if there exists a path of control flow edges from d to u that does not contain a definition of v .

- Build up RD stmt by stmt
- Stmt s , “ $d: v \leftarrow x \text{ op } y$ ”, generates d
- Stmt s , “ $d: v \leftarrow x \text{ op } y$ ”, kills all other defs(v)

Or,

- $\text{Gen}[s] = \{ d \}$
- $\text{Kill}[s] = \text{defs}(v) - \{ d \}$

Gen and kill for each stmt

	Gen	kill
1: n ← 10	1	
2: older ← 0	2	9
3: old ← 1	3	10
4: result ← 0	4	8
5: if n ≤ 1 goto 14		
6: i ← 2	6	11
7: if i > n goto 13		
8: result ← old + older	8	4
9: older ← old	9	2
10: old ← result	10	3
11: i ← i + 1	11	6
12: JUMP 7		
13: return result		
14: return n		

we determine the defs that reach a node by using:

- control flow information
- gen and kill info

Computing $in[n]$ and $out[n]$

- $In[n]$: the set of defs that reach the beginning of node n
- $Out[n]$: the set of defs that reach the end of node n

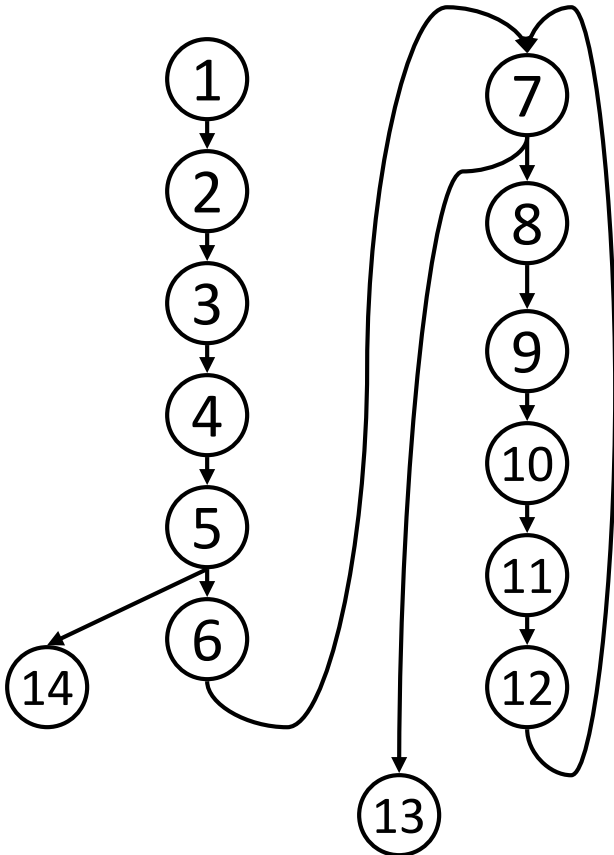
$$in[n] = \bigcup_{p \in pred[n]} out[p]$$

$$out[n] = gen[n] \cup (in[n] - kill[n])$$

- Initialize $in[n]=out[n]=\{\}$ for all n
- Solve iteratively

pred[n]?

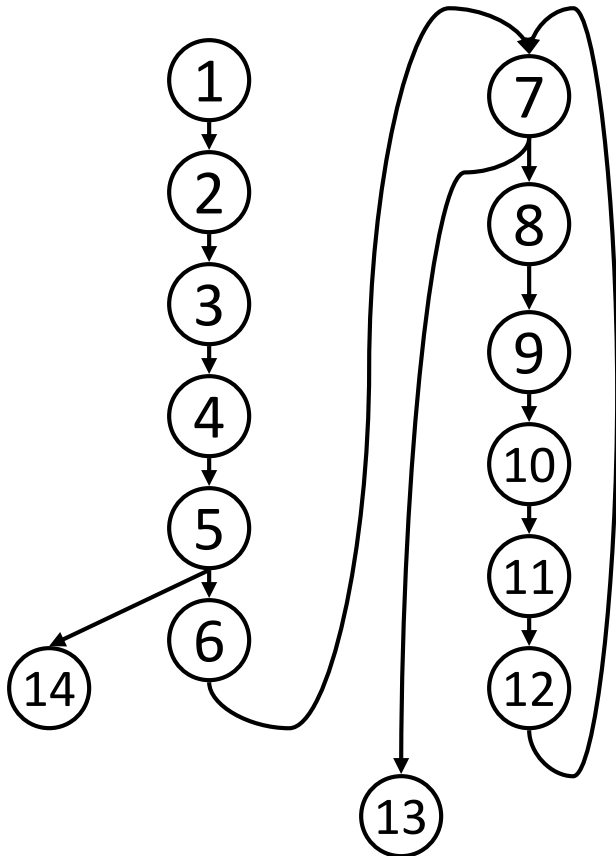
- Pred[n] are all nodes that can directly reach n in the control flow graph.
- E.g., $\text{pred}[7] = \{ 6, 12 \}$



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What order to eval nodes?

- Does it matter?
- Lets do: 1,2,3,4,5,14,6,7,13,8,9,10,11,12



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Example:

- Order: 1,2,3,4,5,14,6,7,13,8,9,10,11,12

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9: older ← old	9	2	1-3, 6, 8	1, 3, 6, 8, 9
10: old ← result	10	3	1, 3, 6, 8, 9	1, 6, 8-10
11: i ← i + 1	11	6	1, 6, 8-10	1, 8-11
12: JUMP 7			1, 8-11	1, 8-11
13: return result			1-4, 6, 8-11	1-4, 6, 8-11
14: return n			1-4	1-4

Example (pass 2)

- Order: 1,2,3,4,5,14,6,7,13,8,9,10,11,12

$$in[n] = \bigcup_{p \in pred[n]} out[p] \quad out[n] = gen[n] \cup (in[n] - kill[n])$$

	Gen	kill	in	out
1: n ← 10	1			1
2: older ← 0	2	9	1	1, 2
3: old ← 1	3	10	1, 2	1, 2, 3
4: result ← 0	4	8	1-3	1-4
5: if n ≤ 1 goto 14			1-4	1-4
6: i ← 2	6	11	1-4	1-4, 6
7: if i > n goto 13			1-4, 6, 8-11	1-4, 6, 8-11
8: result ← old + older	8	4	1-4, 6, 8-11	1-3, 6, 8-11
9: older ← old	9	2	1-3, 6, 8-11	1, 3, 6, 8-11
10: old ← result	10	3	1, 3, 6, 8, 9	1, 6, 8-10
11: i ← i + 1	11	6	1, 6, 8-10	1, 8-11
12: JUMP 7			1, 8-11	1, 8-11
13: return result			1-4, 6, 8-11	1-4, 6, 8-11
14: return n			1-4	1-4

Example (pass 2)

- Order: 1,2,3,4,5,14,6,7,13,8,9,10,11,12

$$in[n] = \bigcup_{p \in pred[n]} out[p] \quad out[n] = gen[n] \cup (in[n] - kill[n])$$

	Gen	kill	in	out
1: n ← 10	1			1
2: older ← 0	2	9	1	1, 2
3: old ← 1	3	10	1, 2	1, 2, 3
4: result ← 0	4	8	1-3	1-4
5: if n ≤ 1 goto 14			1-4	1-4
6: i ← 2	6	11	1-4	1-4, 6
7: if i > n goto 13			1-4, 6, 8-11	1-4, 6, 8-11
8: result ← old + older	8	4	1-4, 6, 8-11	1-3, 6, 8-11
9: older ← old	9	2	1-3, 6, 8-11	1, 3, 6, 8-11
10: old ← result	10	3	1, 3, 6, 8-11	1, 6, 8-11
11: i ← i + 1	11	6	1, 6, 8-10	1, 8-11
12: JUMP 7			1, 8-11	1, 8-11
13: return result			1-4, 6, 8-11	1-4, 6, 8-11
14: return n			1-4	1-4

Example (pass 2)

- Order: 1,2,3,4,5,14,6,7,13,8,9,10,11,12

$$in[n] = \bigcup_{p \in pred[n]} out[p] \quad out[n] = gen[n] \cup (in[n] - kill[n])$$

	Gen	kill	in	out
1: n ← 10	1			1
2: older ← 0	2	9	1	1, 2
3: old ← 1	3	10	1, 2	1, 2, 3
4: result ← 0	4	8	1-3	1-4
5: if n ≤ 1 goto 14			1-4	1-4
6: i ← 2	6	11	1-4	1-4, 6
7: if i > n goto 13			1-4, 6, 8-11	1-4, 6, 8-11
8: result ← old + older	8	4	1-4, 6, 8-11	1-3, 6, 8-11
9: older ← old	9	2	1-3, 6, 8-11	1, 3, 6, 8-11
10: old ← result	10	3	1, 3, 6, 8-11	1, 6, 8-11
11: i ← i + 1	11	6	1, 6, 8-11	1, 8-11
12: JUMP 7			1, 8-11	1, 8-11
13: return result			1-4, 6, 8-11	1-4, 6, 8-11
14: return n			1-4	1-4

Example (pass 2)

- Order: 1,2,3,4,5,14,6,7,13,8,9,10,11,12

$$in[n] = \bigcup_{p \in pred[n]} out[p] \quad out[n] = gen[n] \cup (in[n] - kill[n])$$

	Gen	kill	in	out
1: n ← 10	1			1
2: older ← 0	2	9	1	1, 2
3: old ← 1	3	10	1, 2	1, 2, 3
4: result ← 0	4	8	1-3	1-4
5: if n ≤ 1 goto 14			1-4	1-4
6: i ← 2	6	11	1-4	1-4, 6
7: if i > n goto 13			1-4, 6, 8-11	1-4, 6, 8-11
8: result ← old + older	8	4	1-4, 6, 8-11	1-3, 6, 8-11
9: older ← old	9	2	1-3, 6, 8-11	1, 3, 6, 8-11
10: old ← result	10	3	1, 3, 6, 8-11	1, 6, 8-11
11: i ← i + 1	11	6	1, 6, 8-11	1, 8-11
12: JUMP 7			1, 8-11	1, 8-11
13: return result			1-4, 6, 8-11	1-4, 6, 8-11
14: return n			1-4	1-4

An Improvement: Basic Blocks

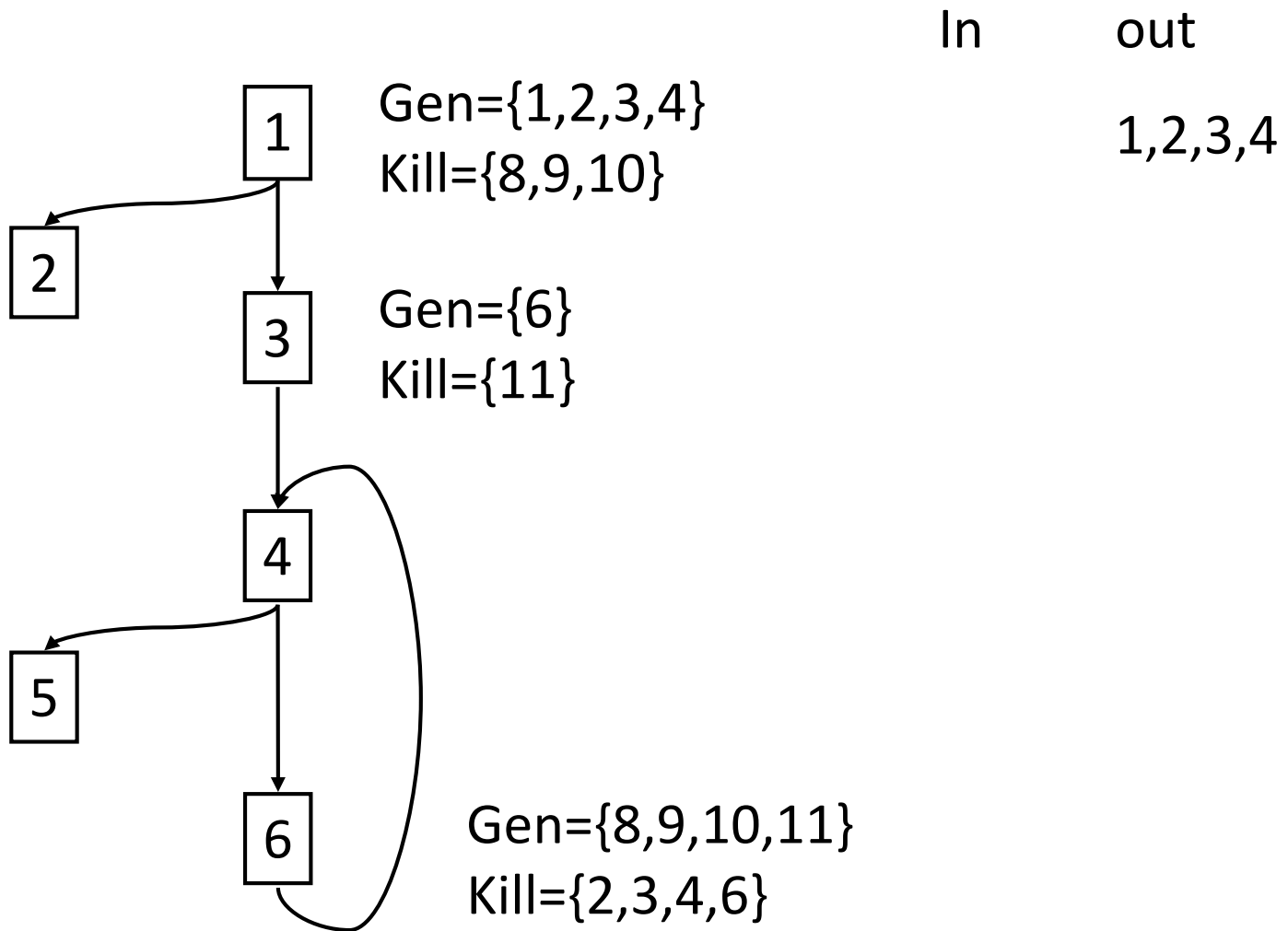
- No need to compute this one stmt at a time
- For straight line code:
 - $\text{In}[s1; s2] = \text{in}[s1]$
 - $\text{Out}[s1; s2] = \text{out}[s2]$
- Combine the gen and kill sets into one per BB.

		Gen	kill
• $\text{Gen}[\text{BB}] = \{2, 3, 4, 5\}$	1: $i \leftarrow 1$	1	8, 4
	2: $j \leftarrow 2$	2	
• $\text{Kill}[\text{BB}] = \{1, 8, 11\}$	3: $k \leftarrow 3 + i$	3	11
	4: $i \leftarrow j$	4	1, 8
	5: $m \leftarrow i + k$	5	

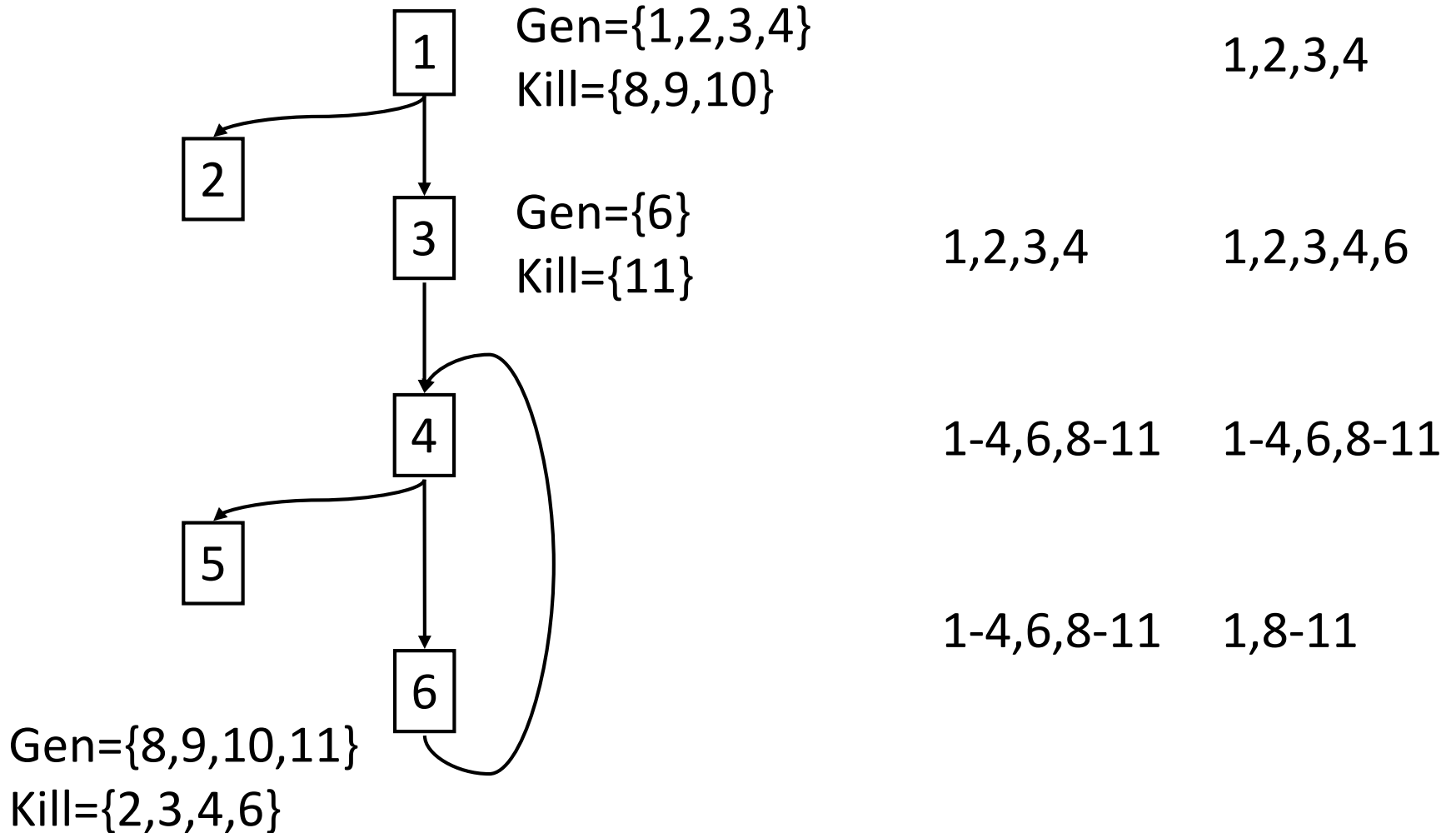
BB sets

		Gen	kill
	1: n ← 10	1	
	2: older ← 0	2	9
1	3: old ← 1	3	10
	4: result ← 0	4	8
	5: if n ≤ 1 goto 14		1, 2, 3, 4 8, 9, 10
3	6: i ← 2	6	11 6 11
4	7: if i > n goto 13		
	8: result ← old + older	8	4
	9: older ← old	9	2
6	10: old ← result	10	3
	11: i ← i + 1	11	6
	12: JUMP 7		8-11 2-4, 6
5	13: return result		
2	14: return n		

BB sets



BB sets



Forward Dataflow

- Reaching definitions is a forward dataflow problem:
It propagates information from the predecessors of a node to the node
- Defined by:
 - Basic attributes: (gen and kill)
 - Transfer function: $F_{bb} \quad out[n] = gen[n] \cup (in[n] - kill[n])$
 - Meet operator: union $in[n] = \bigcup_{p \in pred[n]} out[p]$
 - Set of values (a lattice, in this case powerset of program points)
 - Initial values for each node b
- Solve for fixed point solution

How to implement?

- Values?
- Gen?
- Kill?
- F_{bb} ?
- Order to visit nodes?
- When are we done?
 - In fact, do we know we terminate?

Implementing RD

- Values: bits in a bit vector
- Gen: 1 in each position generated, otherwise 0
- Kill: 0 in each position killed, otherwise 1
- F_{bb} : $out[b] = gen[b] \mid (in[b] \ \& \ kill[b])$
- Init $in[b]=out[b]=0$

- When are we done?
- What order to visit nodes? Does it matter?

RD Worklist algorithm

Initialize: $\text{in}[B] = \text{out}[b] = \emptyset$

Initialize: $\text{in}[\text{entry}] = \emptyset$

Work queue, $W =$ all Blocks in topological order

while ($|W| \neq 0$) {

 remove b from W

$\text{old} = \text{out}[b]$

$\text{in}[b] = \{\text{over all } \text{pred}(p) \in b\} \cup \text{out}[p]$

$\text{out}[b] = \text{gen}[b] \cup (\text{in}[b] - \text{kill}[b])$

 if ($\text{old} \neq \text{out}[b]$) $W = W \cup \text{succ}(b)$

Storing Rd information

- Use-def chains: for each use of var x in s , a list of definitions of x that reach s

1: $n \leftarrow 10$	1	
2: $older \leftarrow 0$	1	1, 2
3: $old \leftarrow 1$	1, 2	1, 2, 3
4: $result \leftarrow 0$	1-3	1-4
5: $if\ n \leq 1\ goto\ 14$	1-4	1-4
6: $i \leftarrow 2$	1-4	1-4, 6
7: $if\ i > n\ goto\ 13$	1-4, 6, 8-11	1-4, 6, 8-11
8: $result \leftarrow old + older$	1-4, 6, 8-11	1-3, 6, 8-11
9: $older \leftarrow old$	1-3, 6, 8-11	1, 3, 6, 8-11
10: $old \leftarrow result$	1, 3, 6, 8-11	1, 6, 8-11
11: $i \leftarrow i + 1$	1, 6, 8-11	1, 8-11
12: JUMP 7	1, 8-11	1, 8-11
13: $return\ result$	1-4, 6	1-4, 6
14: return n	1-4	1-4

Storing Rd information

- Use-def chains: for each use of var x in s, a list of definitions of x that reach s

1: n <- 10	1	
2: older <- 0	1	1, 2
3: old <- 1	1, 2	1, 2, 3
4: result <- 0	1-3	1-4
5: if n <= 1 goto 14	1-4	1-4
6: i <- 2	1-4	1-4, 6
7: if i > n goto 13	1-4, 6, 8-11	1-4, 6, 8-11
8: result <- old + older	1-4, 6, 8-11	1-3, 6, 8-11
9: older <- old	1-3, 6, 8-11	1, 3, 6, 8-11
10: old <- result	1, 3, 6, 8-11	1, 6, 8-11
11: i <- i + 1	1, 6, 8-11	1, 8-11
12: JUMP 7	1, 8-11	1, 8-11
13: return result	1-4, 6	1-4, 6
14: return n	1-4	1-4

Storing Rd information

- Use-def chains: for each use of var x in s , a list of definitions of x that reach s

1: $n \leftarrow 10$	1	
2: $older \leftarrow 0$	1	1, 2
3: $old \leftarrow 1$	1, 2	1, 2, 3
4: $result \leftarrow 0$	1-3	1-4
5: $if\ n \leq 1\ goto\ 14$	1-4	1-4
6: $i \leftarrow 2$	1-4	1-4, 6
7: $if\ i > n\ goto\ 13$	1-4, 6, 8-11	1-4, 6, 8-11
8: $result \leftarrow old + older$	1-4, 6, 8-11	1-3, 6, 8-11
9: $older \leftarrow old$	1-3, 6, 8-11	1, 3, 6, 8-11
10: $old \leftarrow result$	1, 3, 6, 8-11	1, 6, 8-11
11: $i \leftarrow i + 1$	1, 6, 8-11	1, 8-11
12: JUMP 7	1, 8-11	1, 8-11
13: $return\ result$	1-4, 6	1-4, 6
14: return n	1-4	1-4

Storing Rd information

- Use-def chains: for each use of var x in s , a list of definitions of x that reach s

1: $n \leftarrow 10$	1	
2: $older \leftarrow 0$	1	1, 2
3: $old \leftarrow 1$	1, 2	1, 2, 3
4: $result \leftarrow 0$	1-3	1-4
5: $if\ n \leq 1\ goto\ 14$	1-4	1-4
6: $i \leftarrow 2$	1-4	1-4, 6
7: $if\ i > n\ goto\ 13$	1-4, 6, 8-11	1-4, 6, 8-11
8: $result \leftarrow old + older$	1-4, 6, 8-11	1-3, 6, 8-11
9: $older \leftarrow old$	1-3, 6, 8-11	1, 3, 6, 8-11
10: $old \leftarrow result$	1, 3, 6, 8-11	1, 6, 8-11
11: $i \leftarrow i + 1$	1, 6, 8-11	1, 8-11
12: JUMP 7	1, 8-11	1, 8-11
13: $return\ result$	1-4, 6	1-4, 6
14: $return\ n$	1-4	1-4

Constant Folding + DCE

1: n ← 10	1	
2: older ← 0	1	1, 2
3: old ← 1	1, 2	1, 2, 3
4: result ← 0	1-3	1-4
5: if 10 ← 1 goto 14	1-4	1-4
6: i ← 2	1-4	1-4, 6
7: if i > 10 goto 13	1-4, 6, 8-11	1-4, 6, 8-11
8: result ← old + older	1-4, 6, 8-11	1-3, 6, 8-11
9: older ← old	1-3, 6, 8-11	1, 3, 6, 8-11
10: old ← result	1, 3, 6, 8-11	1, 6, 8-11
11: i ← i + 1	1, 6, 8-11	1, 8-11
12: JUMP 7	1, 8-11	1, 8-11
13: return result	1-4, 6	1-4, 6
14: return 10	1-4	1-4

Better Constant Propagation

- What about:

```
x <- 1
```

```
if (y > z)
```

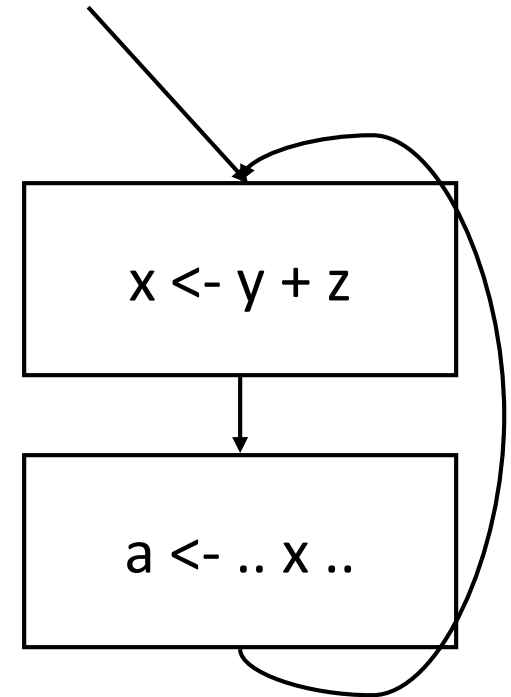
```
    x <- 1
```

```
a <- x
```

- How might you solve this with SSA?

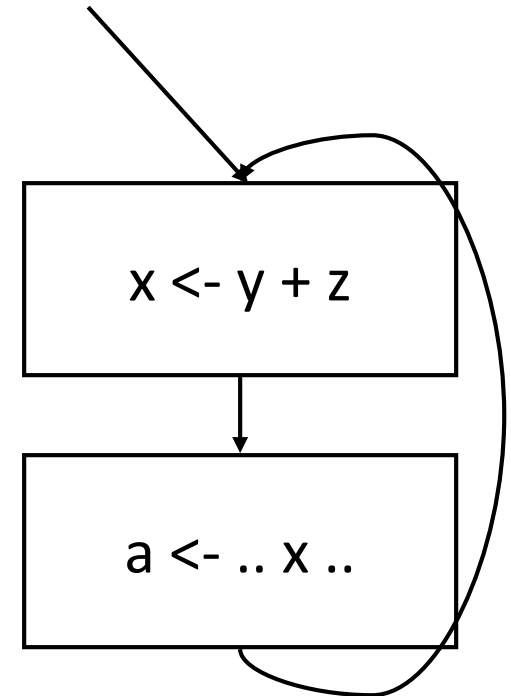
Loop Invariant Code Motion

- When can expression be moved out of a loop?



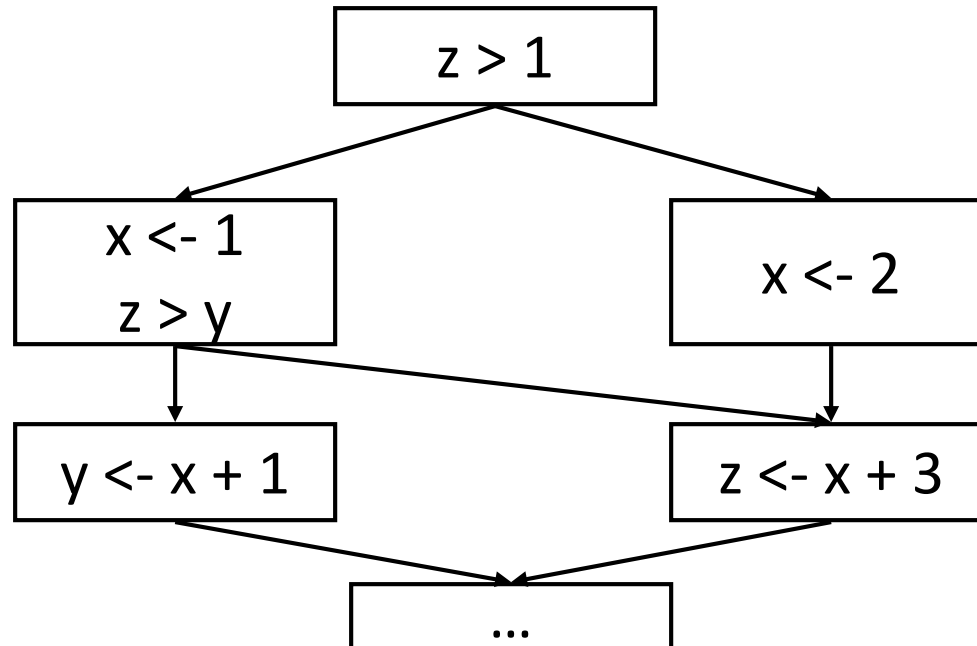
Loop Invariant Code Motion

- When can expression be moved out of a loop?
- When all reaching definitions of operands are outside of loop, expression is loop invariant
- Use ud-chains to detect



Def-use chains are valuable too

- Def-use chain: for each definition of var x , a list of all uses of that definition
- Computed from liveness analysis, a backward dataflow problem
- Def-use is NOT symmetric to use-def



Liveness (def-use chains)

- A variable x is live-out of a stmt s if x can be used along some path starting a s , otherwise x is dead.

Liveness as a dataflow problem

- This is a backwards analysis
 - A variable is live out if used by a successor
 - Gen: For a use: indicate it is live coming into s
 - Kill: Defining a variable v in s makes it dead before s (unless s uses v to define v)
 - Lattice is just live (top) and dead (bottom)
- Values are variables
- $In[n]$ = variables live before n
= $(out[n] - kill[n]) \cup gen[n]$
- $Out[n]$ = variables live after n
= $\bigcup_{s \in succ(n)} In[s]$

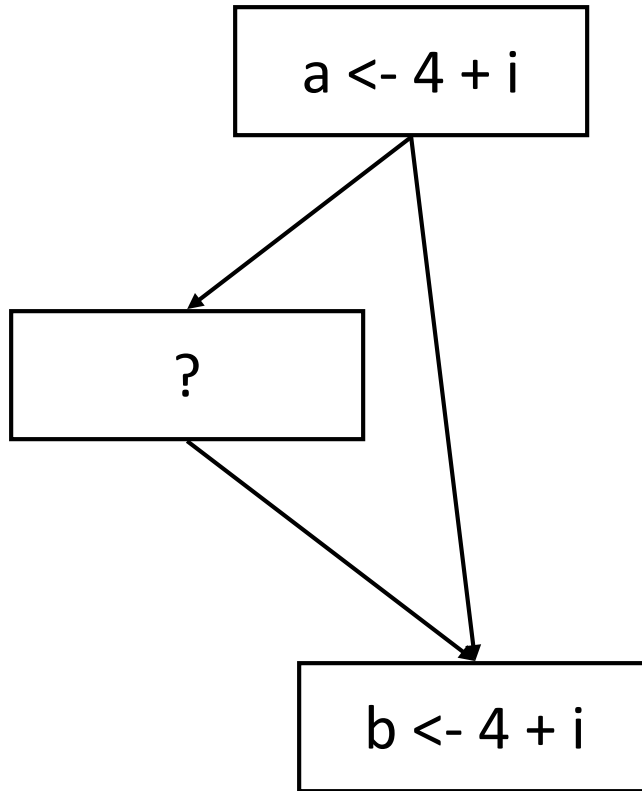
Dead Code Elimination

- Code is dead if it has no effect on the outcome of the program.
- When is code dead?

Dead Code Elimination

- Code is dead if it has no effect on the outcome of the program.
- When is code dead?
 - When the definition is dead, and
 - When the instruction has no side effects
- So:
 - run liveness
 - Construct def-use chains
 - Any instruction which has no users and has no side effects can be eliminated

When can we do CSE?



Available Expressions

- $X+Y$ is “available” at statement S if
 - $x+y$ is computed along every path from the start to S
AND
 - neither x nor y is modified after the last evaluation of $x+y$

$a \leftarrow b+c$

$b \leftarrow a-d$

$c \leftarrow b+c$

$d \leftarrow a-d$

Available Expressions

- $X+Y$ is “available” at statement S if
 - $x+y$ is computed along every path from the start to S
AND
 - neither x nor y is modified after the last evaluation of $x+y$

$a \leftarrow b+c$

$b \leftarrow a-d$

$c \leftarrow b+c$ ← $b+c$ Not available, since b redefined

$d \leftarrow a-d$ ← $a-d$ is available

Computing Available Expressions

- Forward or backward?
- Values?
- Lattice?
- $\text{gen}[b] =$
- $\text{kill}[b] =$
- $\text{in}[b] =$
- $\text{out}[b] =$
- initialization?

Computing Available Expressions

- Forward
- Values: all expressions
- Lattice: available, not-avail
- $\text{gen}[b]$ = if b evals expr e and doesn't define variables used in e
- $\text{kill}[b]$ = if b assigns to x , then all exprs using x are killed.
- $\text{out}[b] = (\text{in}[b] - \text{kill}[b]) \cup \text{gen}[b]$
- $\text{in}[b]$ = what to do at a join point?
- initialization?

Computing Available Expressions

- Forward
- Values: all expressions
- Lattice: available, not-avail
- $gen[b]$ = if b evals expr e and doesn't define variables used in e
- $kill[b]$ = if b assigns to x, $exprs(x)$ are killed
 $out[b] = (in[b] - kill[b]) \cup gen[b]$
- $in[b]$ = An expr is avail only if avail on ALL edges, so:
 $in[b] = \cap$ over all $p \in pred(b)$, $out[p]$
- Initialization
 - All nodes, but entry are set to ALL avail
 - Entry is set to NONE avail

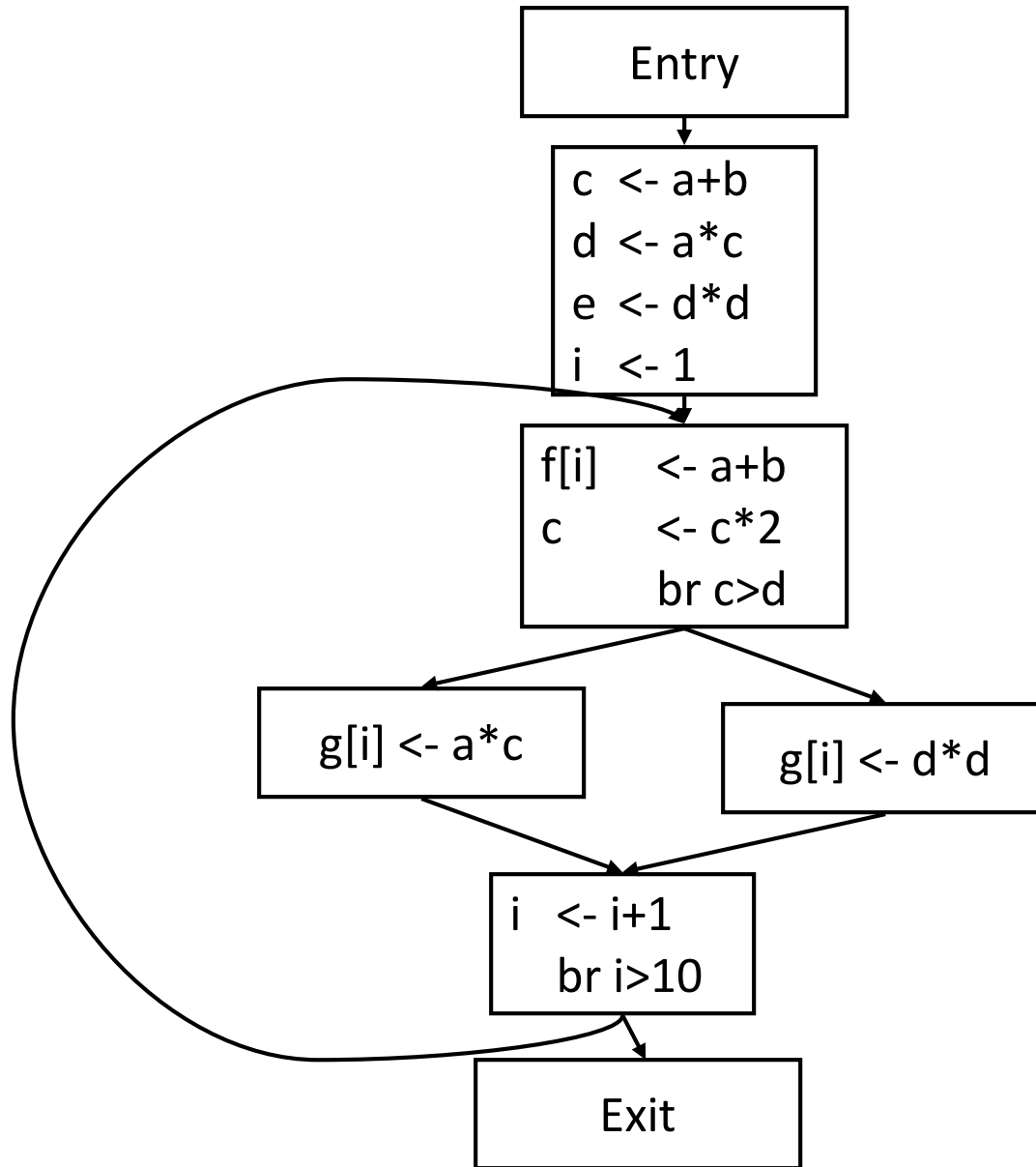
Constructing Gen & Kill

Stmt s	Gen	Kill
$t \leftarrow x \text{ op } y$	$\{x \text{ op } y\}\text{-kill}[s]$	$\{\text{exprs containing } t\}$
$t \leftarrow M[a]$	$\{M[a]\}\text{-kill}[s]$	
$M[a] \leftarrow b$		
$f(a, \dots)$		$\{M[x] \text{ for all } x\}$
$t \leftarrow f(a, \dots)$		

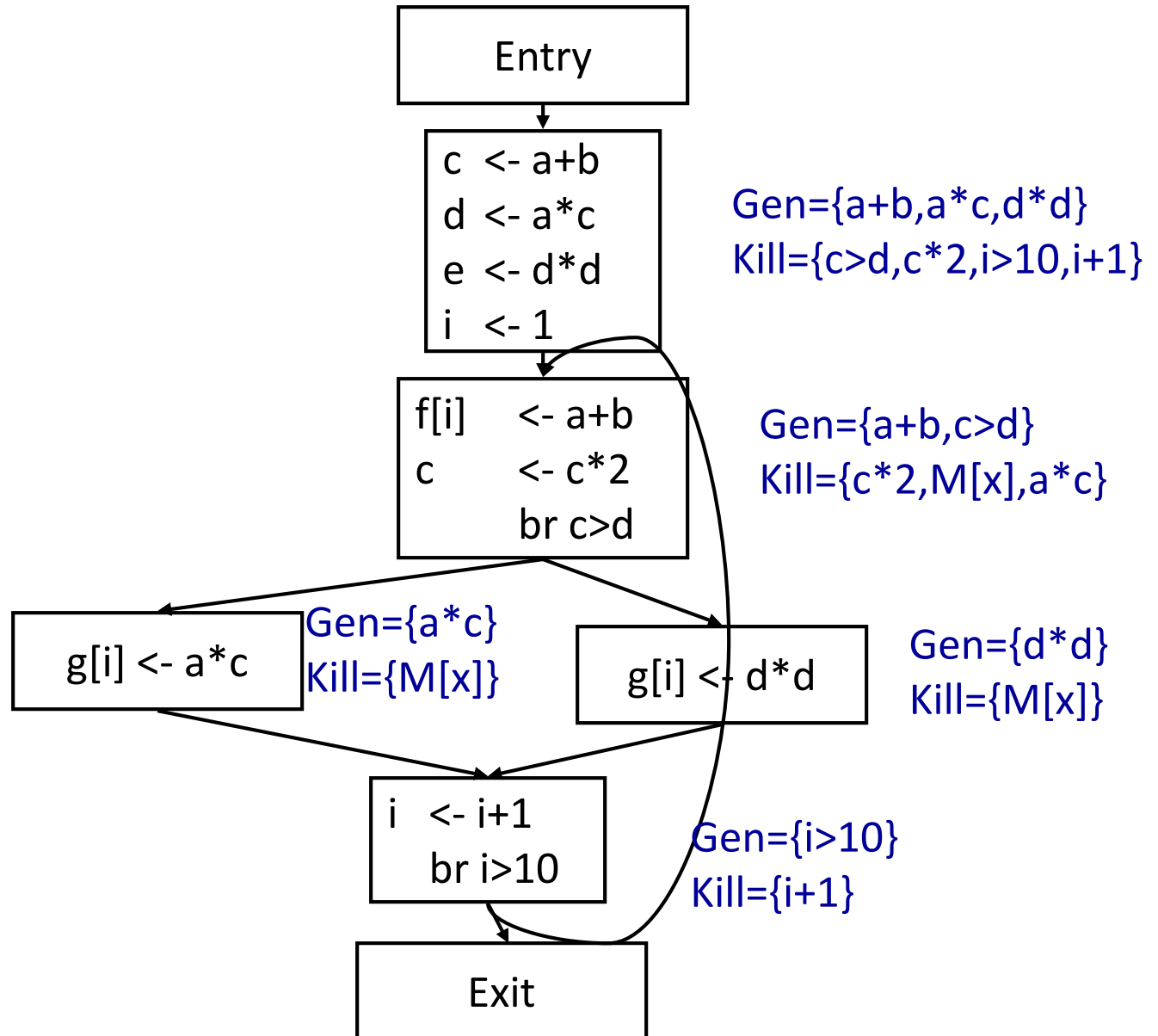
Constructing Gen & Kill

Stmt s	Gen	Kill
$t \leftarrow x \text{ op } y$	$\{x \text{ op } y\}$ -kill[s]	{exprs containing t}
$t \leftarrow M[a]$	$\{M[a]\}$ -kill[s]	{exprs containing t}
$M[a] \leftarrow b$	{}	{for all x, M[x]}
$f(a, \dots)$	{}	{for all x, M[x]}
$t \leftarrow f(a, \dots)$	{}	{exprs containing t for all x, M[x]}

Example

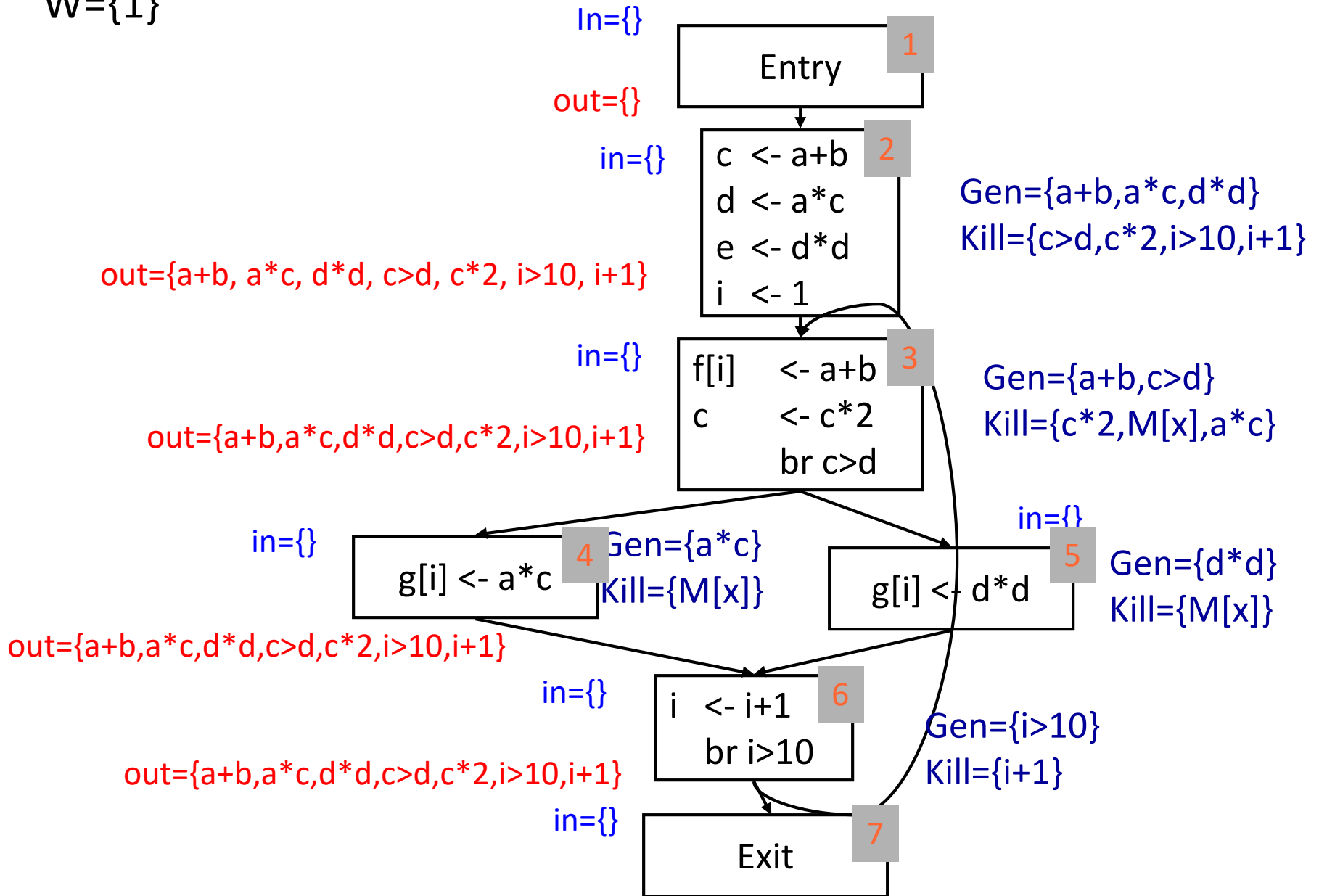


Example



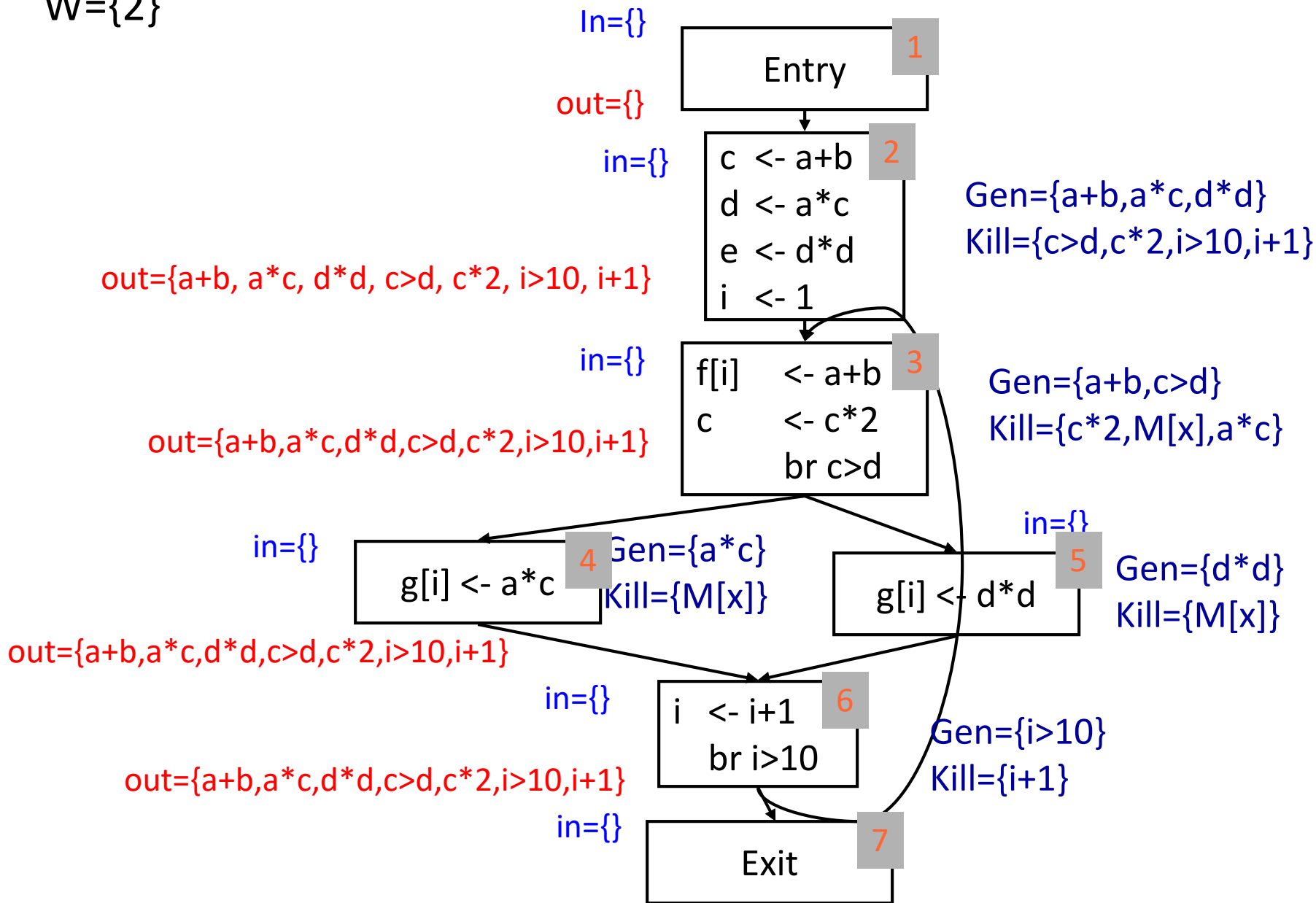
Example

$W=\{1\}$



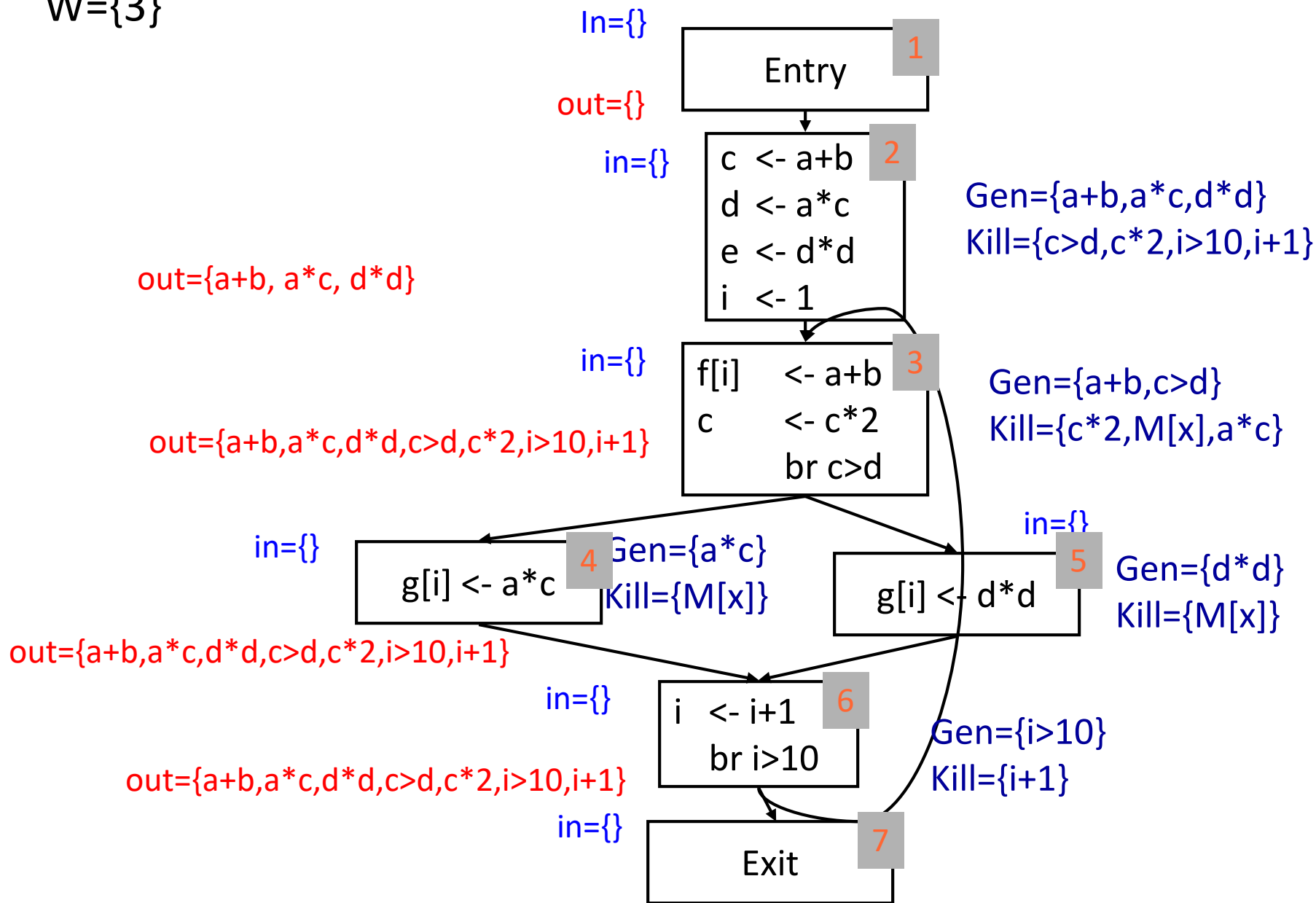
Example

$W=\{2\}$



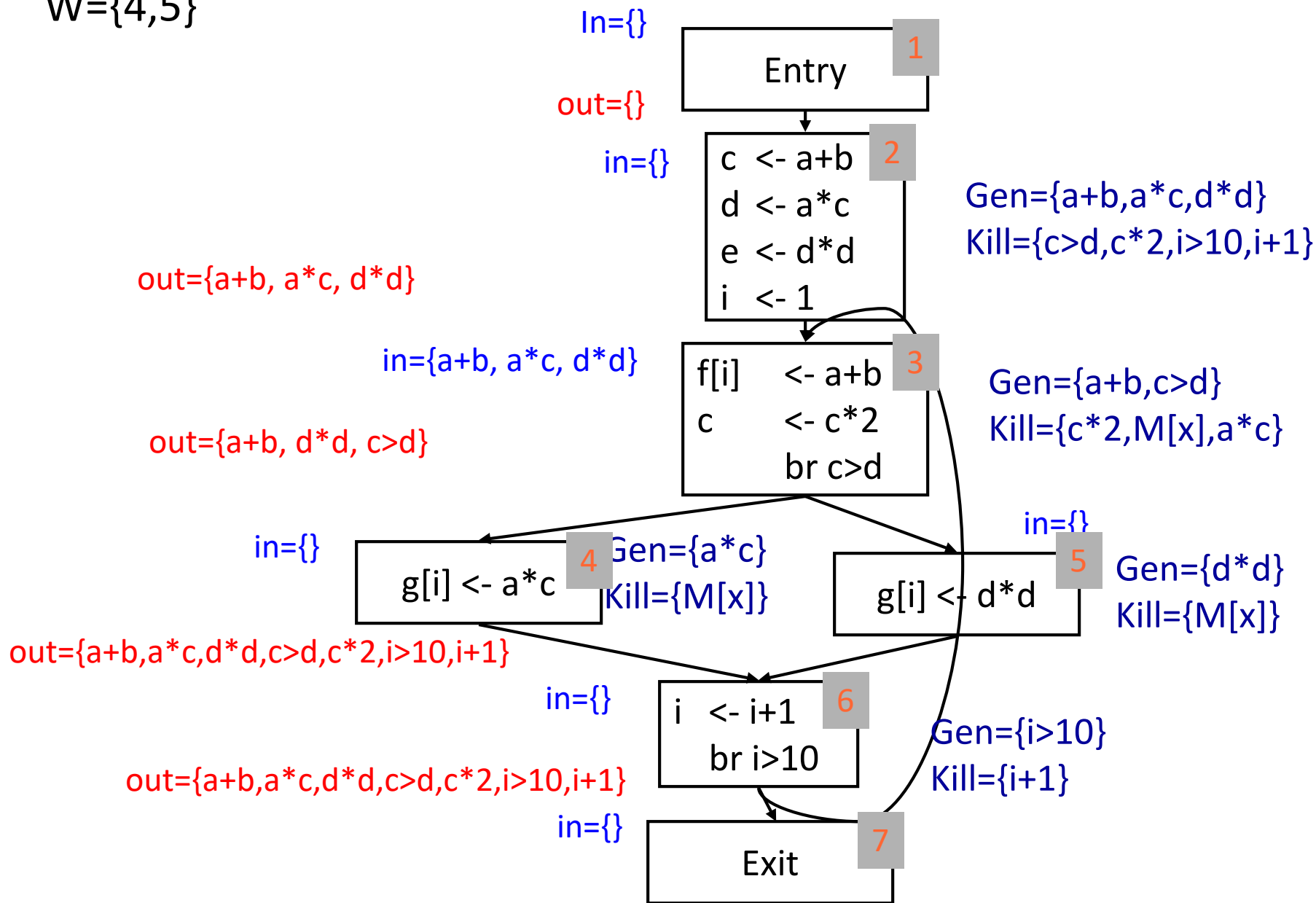
Example

$W=\{3\}$



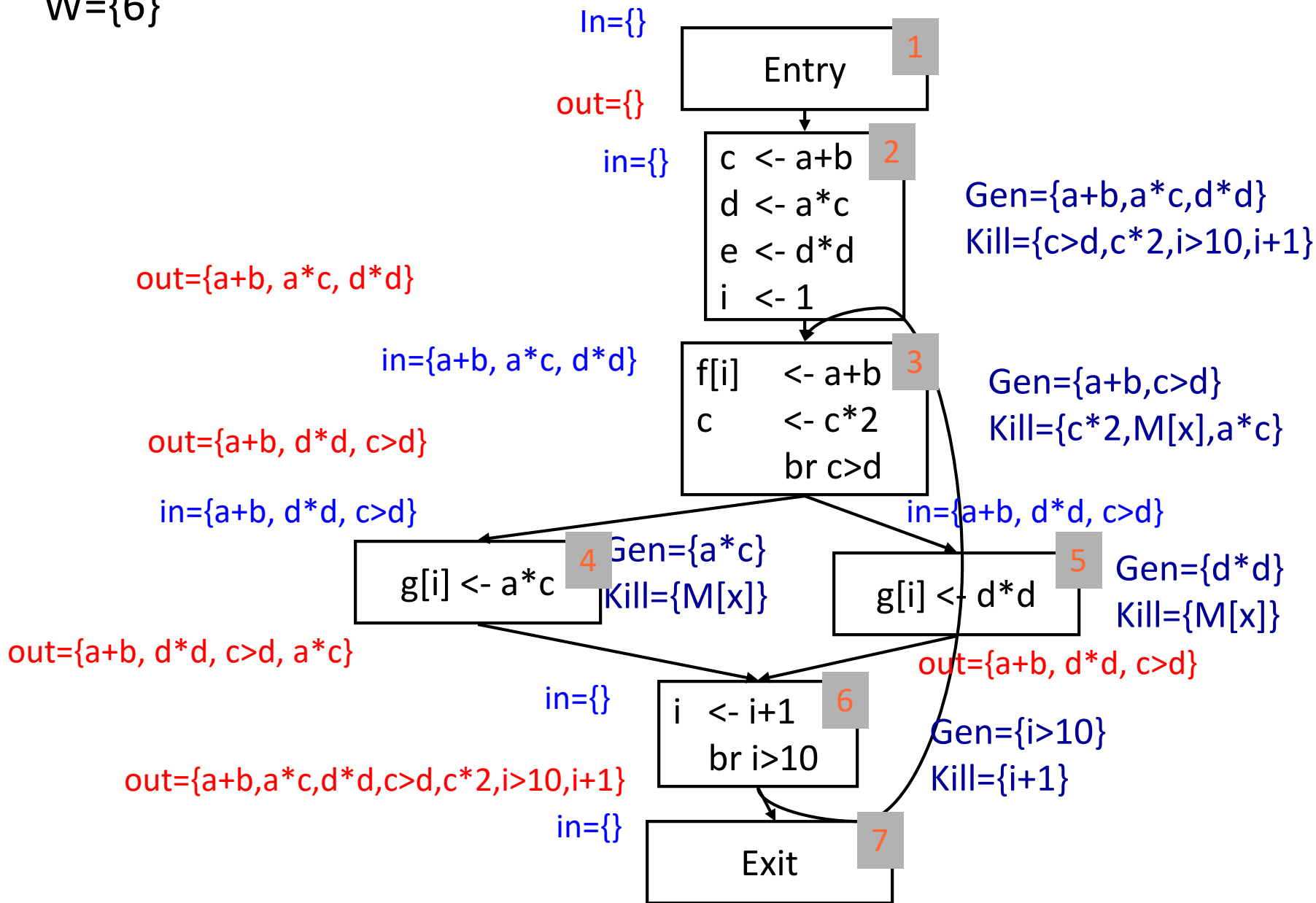
Example

$W=\{4,5\}$



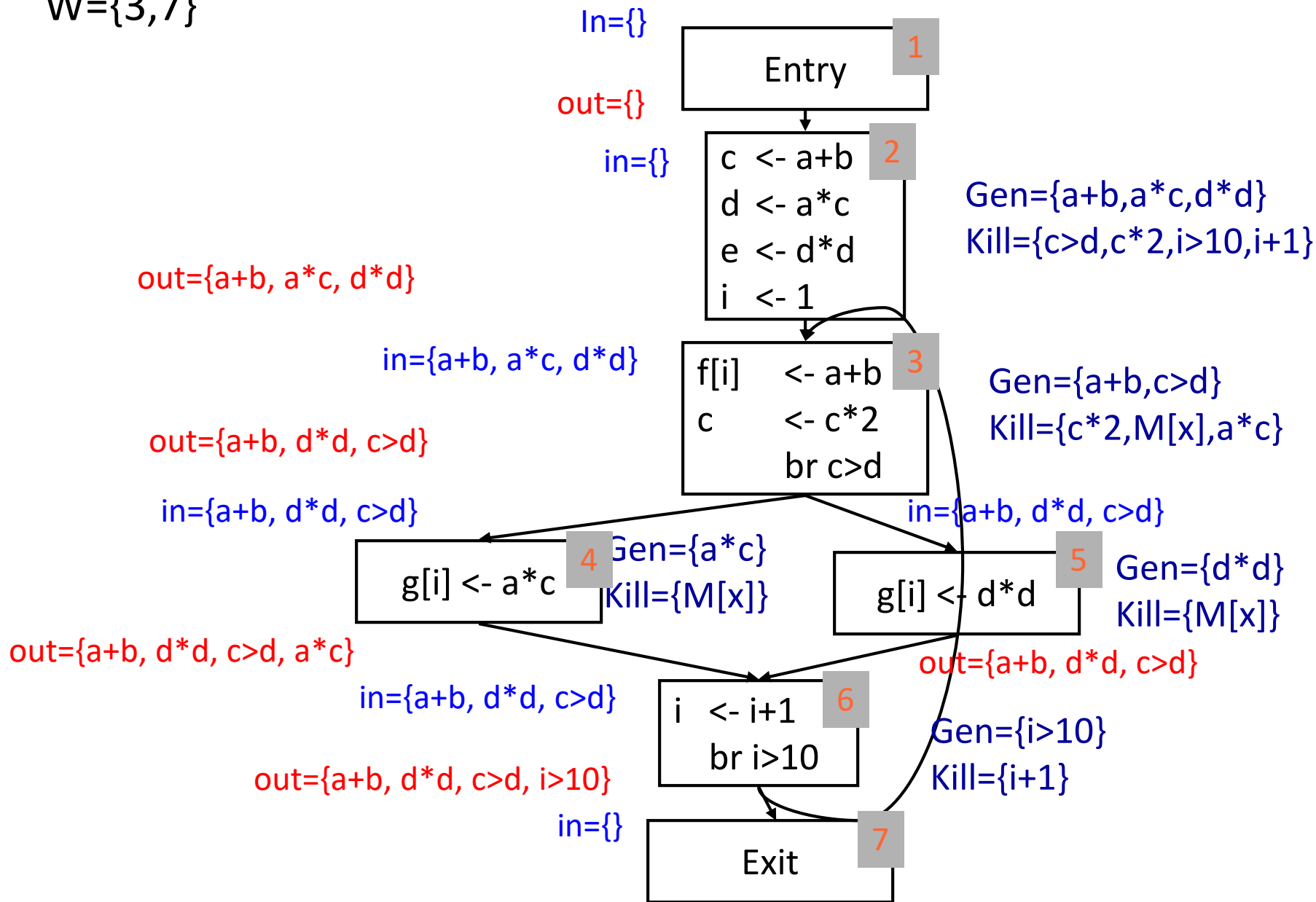
Example

$W=\{6\}$



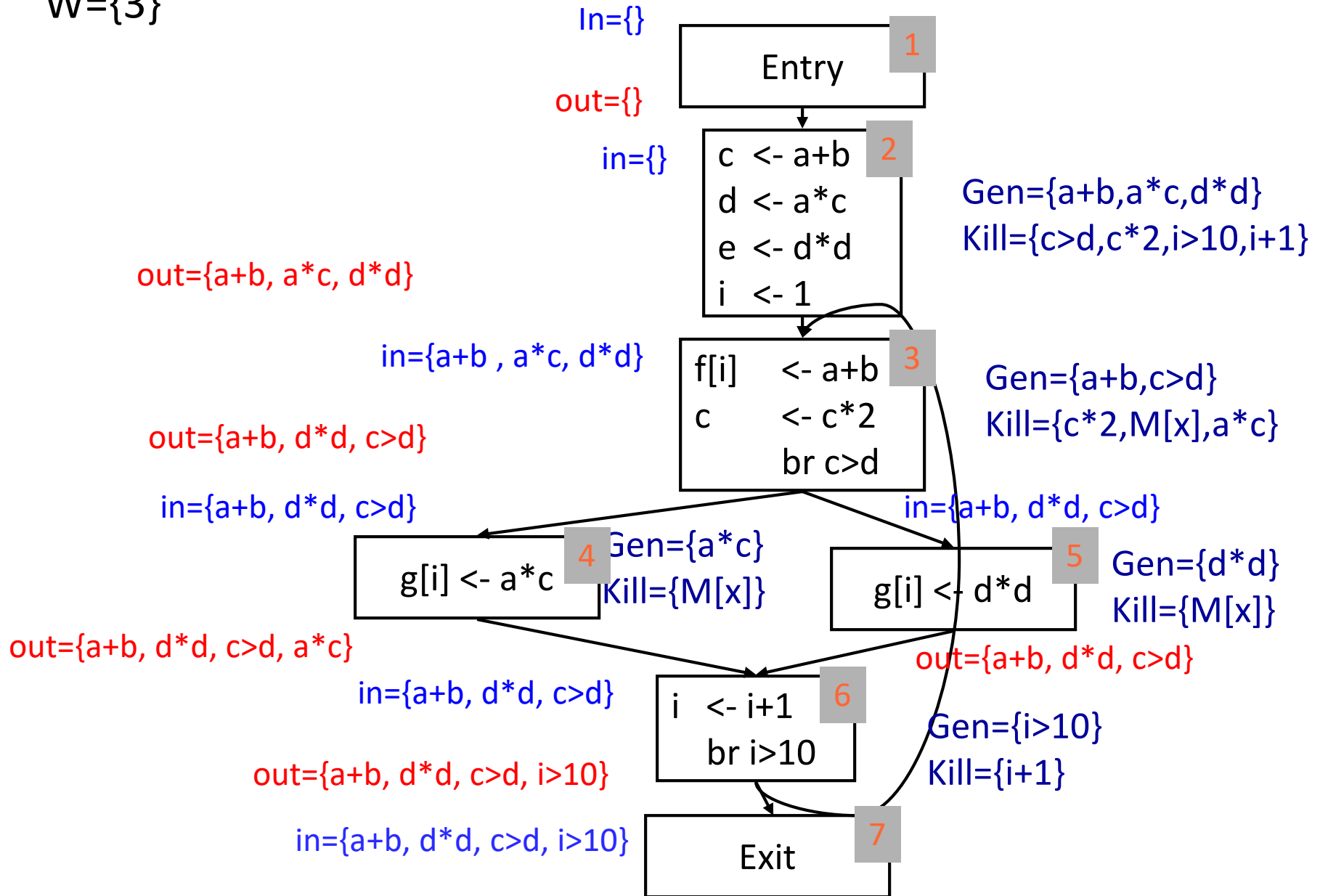
Example

$W=\{3,7\}$

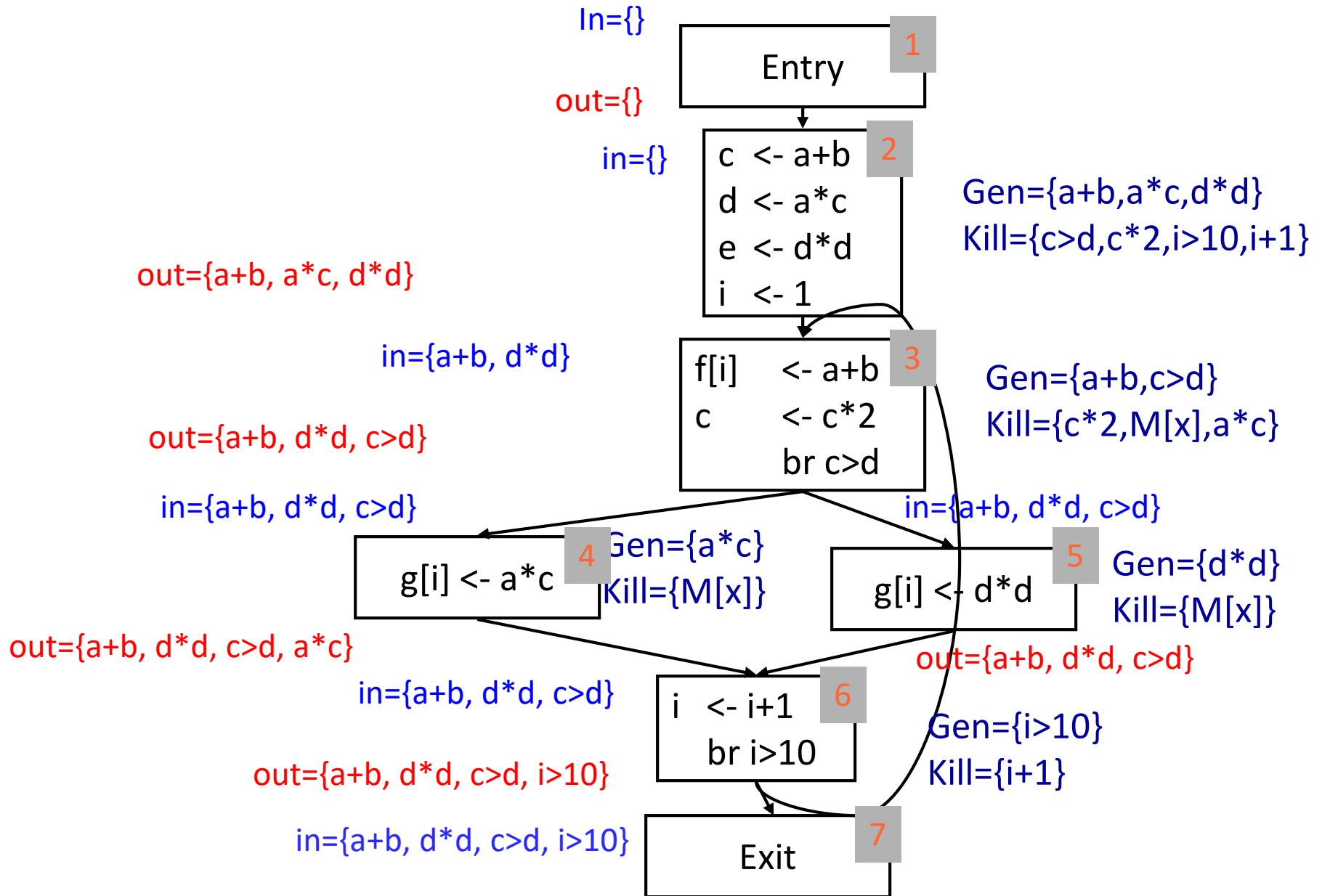


Example

$W=\{3\}$



Example

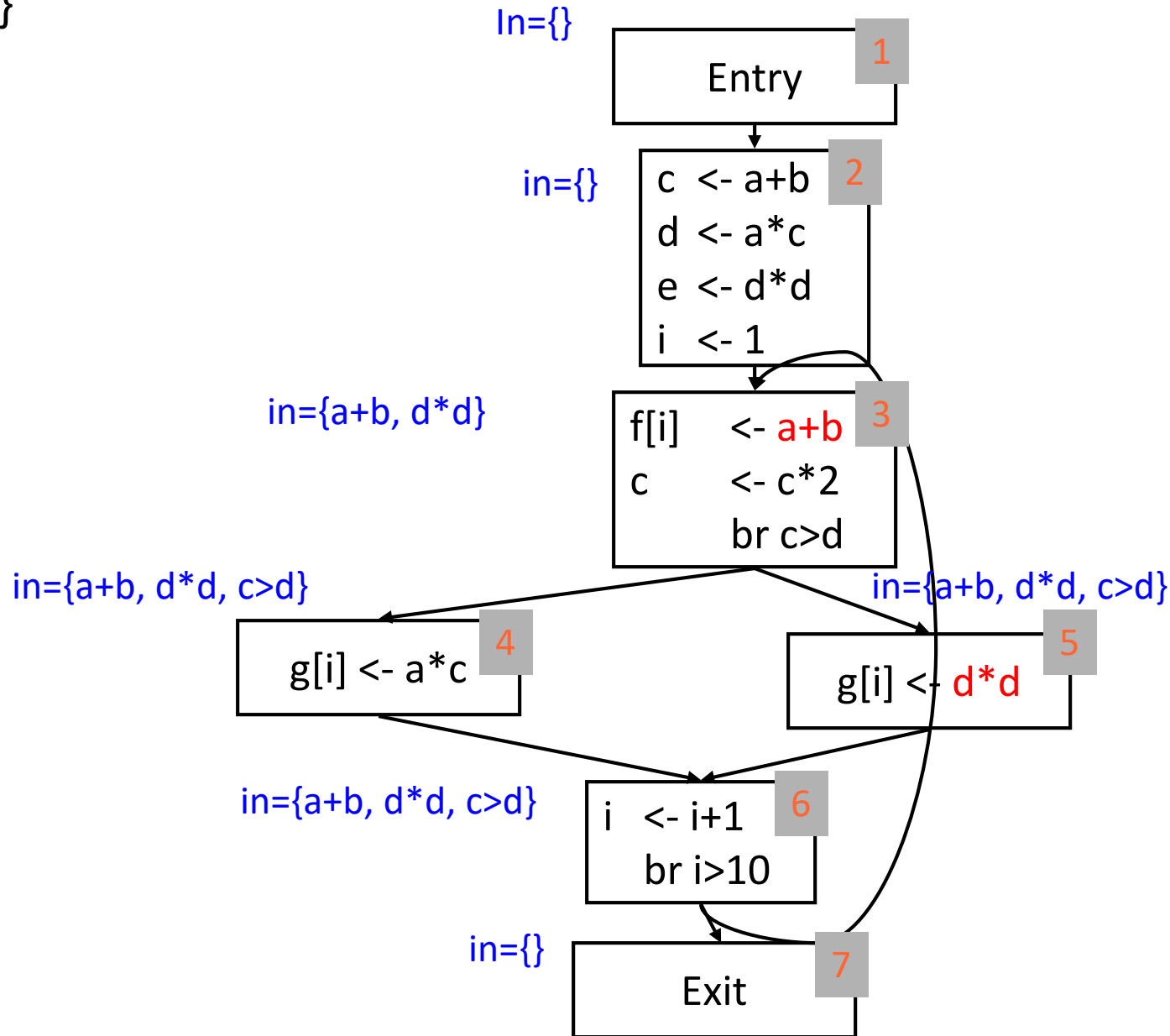


CSE

- Calculate Available expressions
- For every stmt in program
 - If expression, $x \text{ op } y$, is available {
 - Compute reaching expressions for “ $x \text{ op } y$ ”
at this stmt
 - foreach stmt in RE of the form $t \leftarrow x \text{ op } y$
 - rewrite at: $t' \leftarrow x \text{ op } y$
 - $t \leftarrow t'$
- replace “ $x \text{ op } y$ ” in stmt with t'

Find x op y available

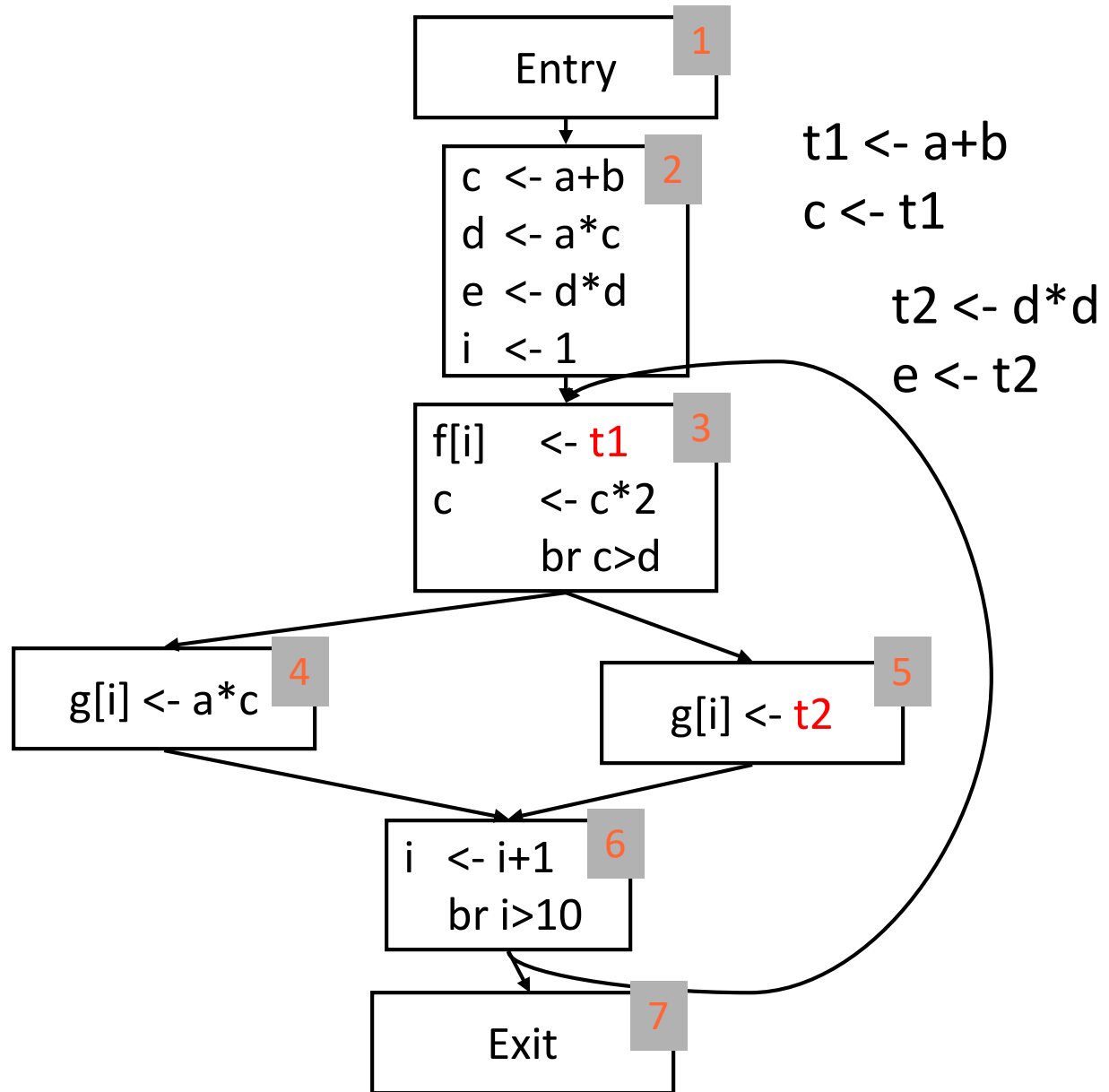
$W=\{3\}$



Calculating Reaching Expressions

- Could be dataflow problem, but not needed enough, so ...
- To find RE for “ $x \text{ op } y$ ” at stmt S
 - traverse cfg backward from S until
 - reach $t \leftarrow x + y$ (& put into RE)
 - reach definition of x or y

Example



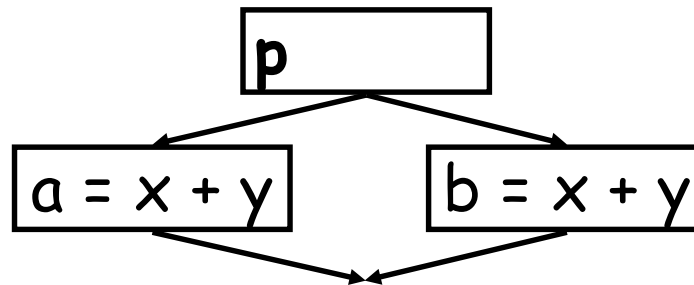
Dataflow Summary

	Union (may)	intersection (must)
Forward	Reaching defs	Available exprs
Backward	Live variables	very busy exprs

Later in course we look at bidirectional dataflow

Very Busy Expressions

- A Backward, Must data flow analysis
- An expression e is *very busy at point p* if On every path from p , e is evaluated before the value of e is changed
- Optimization
 - Can hoist very busy expression computation



Forward Must Data Flow Algorithm

$\text{Out}(s) = \text{Gen}(s)$ for all statements s

$W = \{\text{all statements}\}$

Repeat

 Take s from W

$\text{In}(s) = \bigcap_{s' \in \text{pred}(s)} \text{Out}(s')$

$\text{Temp} = \text{Gen}(s) \cup (\text{In}(s) - \text{Kill}(s))$

 If $(\text{temp} \neq \text{Out}(s))$ {

$\text{Out}(s) = \text{temp}$

$W = W \cup \text{succ}(s)$

 }

Until $W = \emptyset$

Forward May Data Flow Algorithm

$\text{Out}(s) = \text{Gen}(s)$ for all statements s

$W = \{\text{all statements}\}$

Repeat

 Take s from W

$\text{In}(s) = \bigcup_{s' \in \text{pred}(s)} \text{Out}(s')$

$\text{Temp} = \text{Gen}(s) \cup (\text{In}(s) - \text{Kill}(s))$

 If $(\text{temp} \neq \text{Out}(s))$ {

$\text{Out}(s) = \text{temp}$

$W = W \cup \text{succ}(s)$

 }

Until $W = \emptyset$

Backward May Data Flow Algorithm

$In(s) = Gen(s)$ for all statements s

$W = \{\text{all statements}\}$ (worklist)

Repeat

 Take s from W

$Out(s) = \bigcup_{s' \in succ(s)} In(s')$

$Temp = Gen(s) \cup (Out(s) - Kill(s))$

 If ($temp \neq In(s)$) {

$In(s) = temp$

$W = W \cup pred(s)$

 }

Until $W = \emptyset$

Backward Must Data Flow Algorithm

$In(s) = Gen(s)$ for all statements s

$W = \{\text{all statements}\}$ (worklist)

Repeat

 Take s from W

$Out(s) = \bigcap_{s' \in succ(s)} In(s')$

$Temp = Gen(s) \cup (Out(s) - Kill(s))$

 If ($temp \neq In(s)$) {

$In(s) = temp$

$W = W \cup pred(s)$

 }

Until $W = \emptyset$