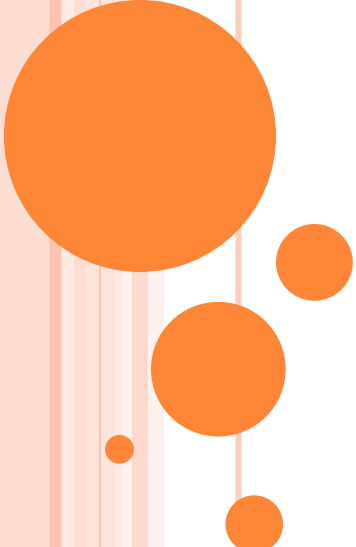




ANITA'S SUPER AWESOME RECITATION SLIDES

**15/18-213: Introduction to Computer Systems
Memory and Caches, 18 Feb 2013**

Anita Zhang, Section M

UP TO SPEED YET?

○ Buflab

- Due tomorrow, midnight

○ Cachelab

- Out tomorrow, midnight
- Due Thursday, February 28, 2013, midnight
 - Labs will be going back to regular Thursday due date
- 10 days to get your C skills back up!



THIS AND THAT AND WHAT

- Alignment
- Memory Organization
- Caching
 - Buzzword: locality
 - Cache organization
- Cachelab
 - Part A – Implement a (hardware) cache simulator
 - Part B – Efficient matrix transpose
 - “Bro, do you even C?” – helpful C stuff

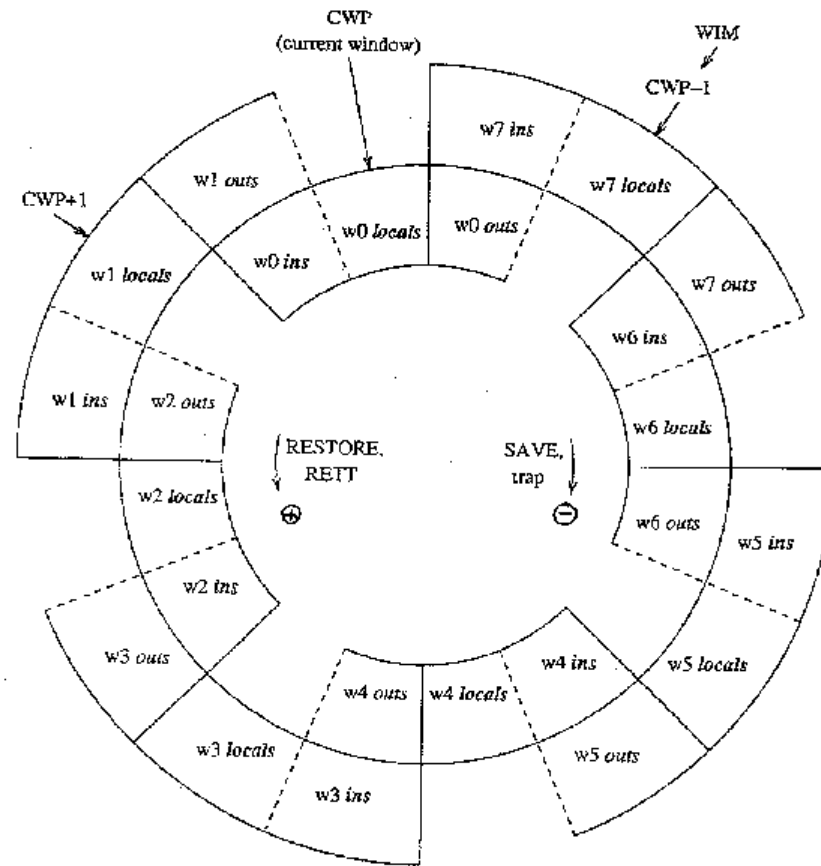


BEFORE WE BEGIN...

- Encouraging Female Reverse Engineers
 - Contest to reverse engineer malicious software
 - And document it with utmost understanding
- Prize
 - Ticket to Symposium on Security for Asia Network (SysScan)
- Details
 - <http://addxorrol.blogspot.de/2013/01/encouraging-female-reverse-engineers.html>
 - Only women can submit (sorry guys)
 - One of the judges went to CMU!
 - Deadline: 24th of March 2013, 23:59 GMT+1



BAM! CIRCULAR STACK!



SPARC (scalable processor architecture) Architecture



SUPER BRIEF ON ALIGNMENT

○ Structs

- Align according to the largest alignment requirement
 - Must be multiple of K (largest alignment requirement)
 - System dependent alignments requirements
 - Compilers enforce this
- Overall structure length a multiple of K
 - Optimize length by declaring largest elements first

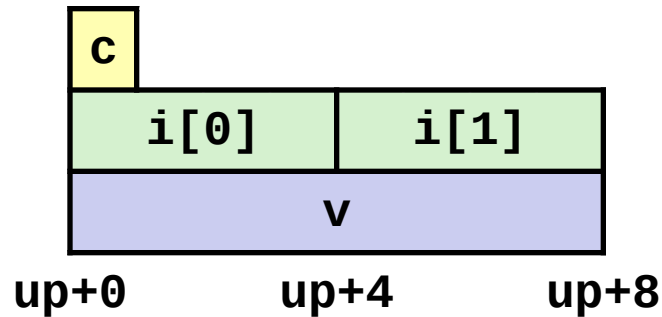
○ Unions

- Size allocated according to largest element
- Only one field used at a time

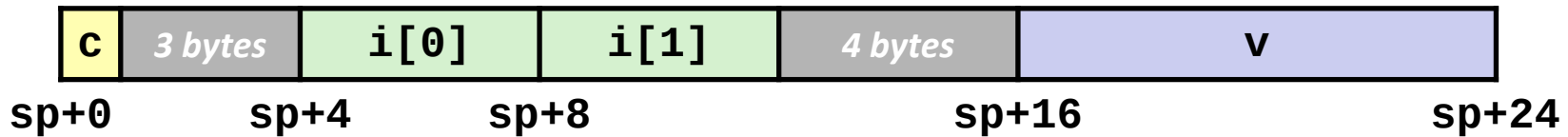


QUICK EXAMPLE FROM LECTURE

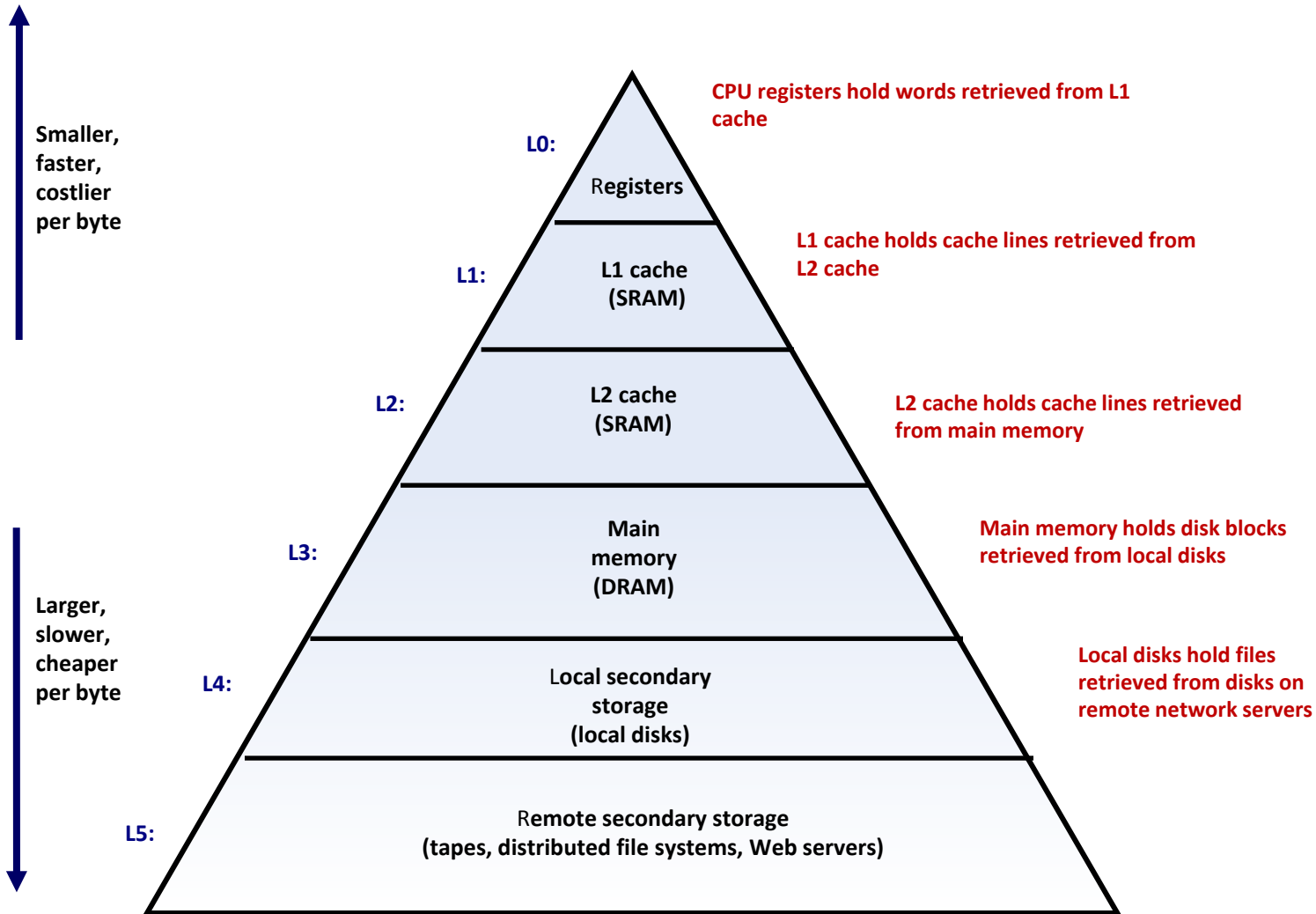
```
union U1 {  
    char c;  
    int i[2];  
    double v;  
} *up;
```



```
struct S1 {  
    char c;  
    int i[2];  
    double v;  
} *sp;
```



MEMORY HIERARCHY (FROM LECTURE)



SRAM vs DRAM

○ SRAM

- Faster (L1 Cache: 1 CPU cycle)
- Smaller (L1 in kilobytes; L2 in megabytes)
- More expensive and “energy-hungry”

○ DRAM (Main memory)

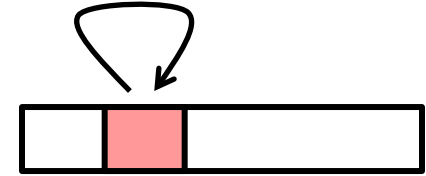
- Relatively slower (hundreds of CPU cycles)
- Larger (Gigabytes)
- Cheaper



LOCALITY

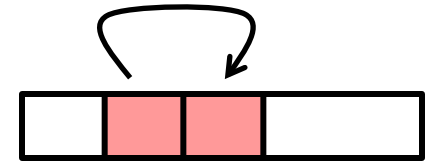
○ Temporal locality

- Recently referenced items are likely to be referenced again in the near future
- After accessing address X in memory, save the bytes in cache for future access

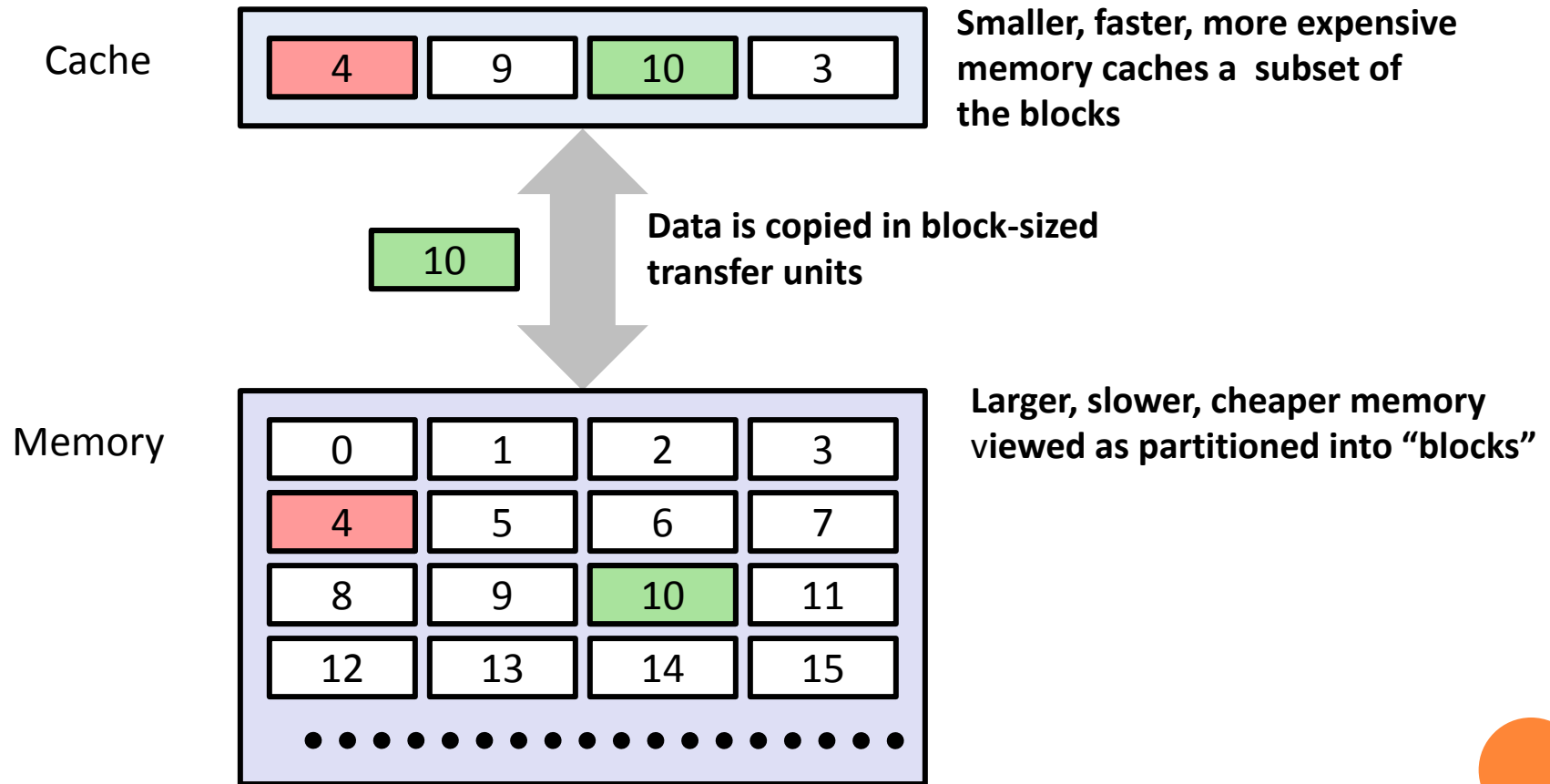


○ Spatial locality

- Items with nearby addresses tend to be referenced close together in time
- After accessing address X, save the block of memory around X in cache for future access



GENERAL CACHING (FROM LECTURE)



ADDRESS DIVISION IN CACHES

- On the Sharks, addresses are 64-bits
- Dividing a memory address
 - Block offset: b bits
 - Set index: s bits

memory address



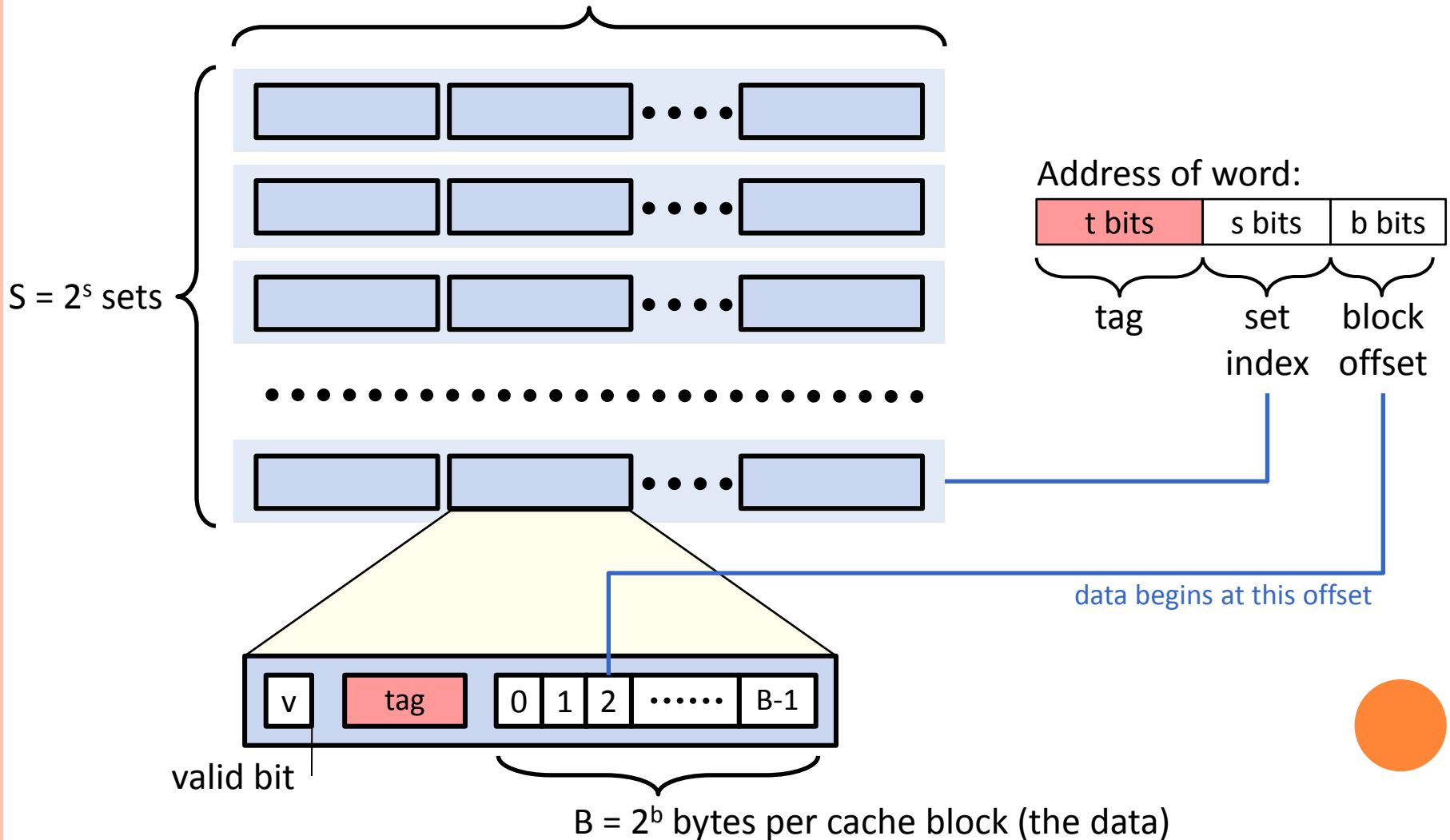
CACHES

- A cache is a set of 2^s **cache sets**
- A **cache set** is a set of E **cache lines**
 - E is called associativity
 - If $E=1$, it is called “direct-mapped”
- Each **cache line** stores a block
 - Each block has 2^b bytes



VISUAL CACHE TERMINOLOGY

$E = 2^e$ lines per set



CACHE LAB PART A

○ Cache Simulator

- Implement for variable s, b, and E values
 - Values read in from a trace file (at runtime)
- Least Recently Used (LRU) Policy

○ Cache Simulator != Cache

- This simulator does NOT store memory contents
- We do NOT care about block offsets here
- Your goal: **implement the policy and count the number of hits, misses, and evictions**



GENERAL SIMULATOR DESIGN HINTS

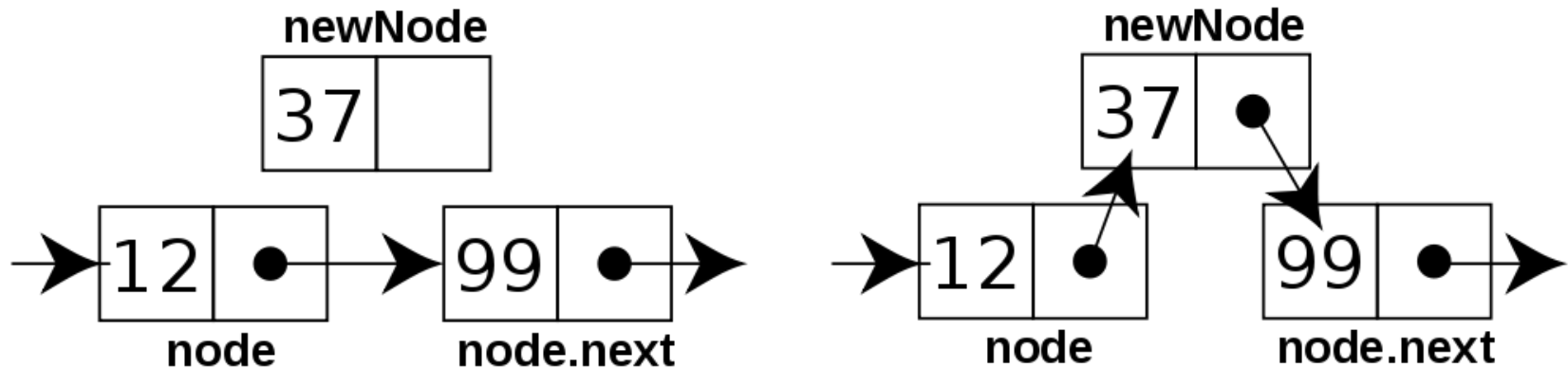
- A cache is just 2D array of cache lines:
 - `struct cache_line cache[S][E];`
 - $S = 2^s$ is the number of sets
 - E is associativity
- Each `cache_line` has:
 - Valid bit
 - Tag
 - LRU counter



ANITA'S FAVORITE DATA STRUCTURE

○ Linked lists

- “The only data structure you will ever need”
- (Heavily) used in cache and malloc lab
- A lesson on linked list in the credits page



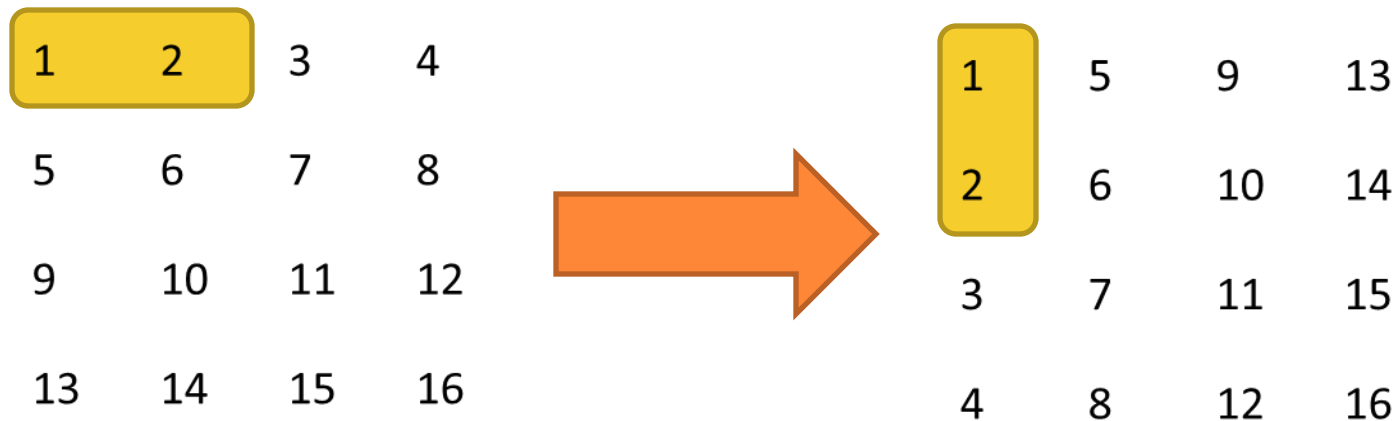
FOOD FOR THOUGHT/ OTHER DESIGNS

- How necessary is the LRU counter?
 - We have the power to insert nodes wherever we want
 - So why use a counter?
- As a C programmer, implementing a linked list should be second nature
 - 5-10 minutes tops
 - The same deal every time
 - Pointers to each node
 - Traversal helper functions
 - Making the right checks



CACHELAB PART B

- Efficient matrix transpose
 - Goal: Increasing locality via blocking



CACHELAB PART B

○ Cache:

- 1 kilobytes of cache
- Directly mapped ($E=1$)
- Block size is 32 bytes ($b=5$)
- $S = 32$ sets ($s=5$)

○ Test Matrices:

- 32 x 32, 64 x 64, 61 x 67
- You only need to optimize for these sizes



“BRO, DO YOU EVEN C?”

- In this section:
 - Warnings are errors
 - Headers
 - Useful C functions



WARNINGS ARE ERRORS

- Strict compilation flags
 - Avoid potential errors that are hard to debug
 - Learn good habits from the beginning
- Add “-Werror” to your compilation flags
- DO NOT ignore the compiler errors



WHAT ABOUT HEADERS?

- Remember to include files that we will be using functions from
- If function declaration is missing
 - Find corresponding header files
 - `unix> man function-name`
 - Skim the man pages, they'll tell you what you need to know



FUNCTION 1: GETOPT

- `getopt` automates parsing elements on the unix command line
 - Typically called in a loop to retrieve arguments
 - Use a switch statement to handle options
 - Returns -1 when there are no more arguments
- Must include the header file `unistd.h`



FUNCTION 1: GETOPT

- Switch statement used on the (local) variable holding the return value from getopt
 - Each command line input can be handled separately
 - optarg – Points to the value of the option argument
 - This is set by the getopt function
- Food for thought
 - How do we handle invalid inputs?



FUNCTION 1: GETOPT EXAMPLE

- Suppose we had an executable called “foo”
 - Example call from shell: `unix> ./foo -x 1`
- Next slide: Parsing the argument to the x option
 - Notice: We passed in an int which is read as a char *
 - We use `atoi` to convert the string to an int



FUNCTION 1: GETOPT EXAMPLE

```
int main(int argc, char** argv){
    int opt, x;

    /* looping over arguments */
    while(-1 != (opt = getopt(argc, argv, "x:"))){

        /* determine which argument it's processing */
        switch(opt) {
            case 'x':
                x = atoi(optarg);
                break;
            default:
                printf("wrong argument\n");
                break;
        }
    }
}
```



FUNCTION 2: FSCANF

- The `fscanf` function is just like `scanf`
 - But it can specify a stream to read from
 - `scanf` always reads from `stdin`
- Parameters:
 - File pointer
 - Format string with information on how to read file
 - Variable number of pointers to with locations for storing data from file
- Typically use in a loop until it hits the end of file
- `fscanf` will be useful in reading from the trace files



FUNCTION 2: FSCANF EXAMPLE

```
FILE * pFile; // pointer to FILE object

/* open file for reading */
pFile = fopen ("myfile.txt", "r");

int x, y;
char c;

/* read two ints and a char from file */
while(fscanf(pFile, "%d %d %c", &x, &y, &c) > 0){
    // Do stuff
}

fclose(pFile); // remember to close file when done
```



FUNCTION 3 AND 4: MALLOC/FREE

- Use `malloc` to allocate memory on the heap
 - Returns a pointer to location in memory
- Always free what you `malloc`
 - Or you'll suffer from memory leaks
- Example usage:
 - `int *pointer = malloc( sizeof(int) );`
 - `free(pointer);`
- **DO NOT free memory you didn't allocate**
 - This includes double free-ing



STYLE AND TIPS FOR LIFE

- Read the style guideline
 - “But I already read it!”
 - Good, read it again.
- Check for failures and errors ALWAYS
 - Functions don’t always succeed
 - What happens when a system call fails?
- Common cases of failure:
 - Not checking the return of `malloc`
 - Not handling invalid inputs
 - Generally, not checking returns of functions



I STOLE FROM THESE PLACES

- [Understanding the SPARC Architecture](#)
- [Wikipedia: Linked Lists](#)
- [C Linked List Example](#)
- [getopt from GNU](#)
- [fscanf from CPlusPlus.com](#)

