

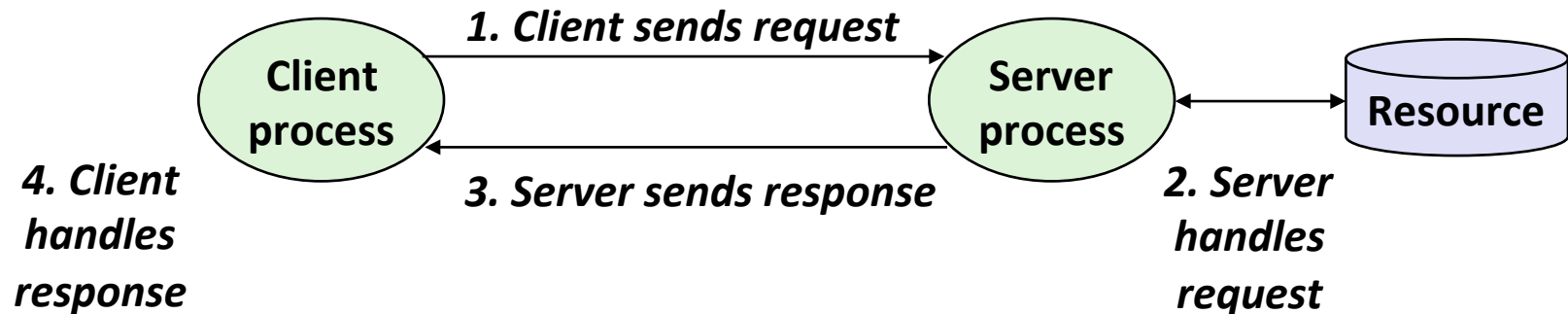
# Internetworking

15-213/18-213: Introduction to Computer Systems  
20<sup>th</sup> Lecture, April. 2, 2013

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# A Client-Server Transaction

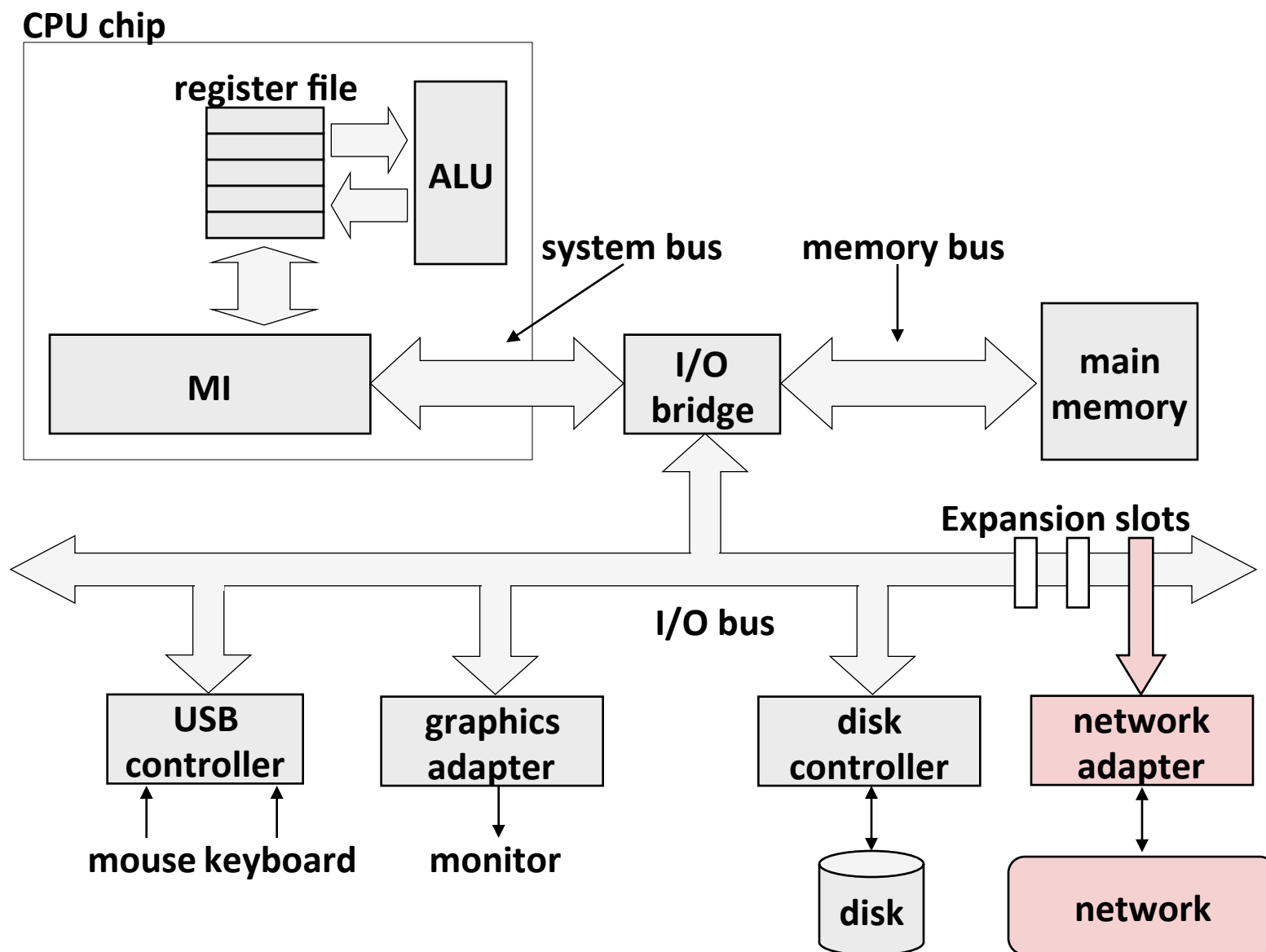


*Note: clients and servers are processes running on hosts  
(can be the same or different hosts)*

## ■ Most network applications are based on the client-server model:

- A **server** process and one or more **client** processes
- Server manages some **resource**
- Server provides **service** by manipulating resource for clients
- Server activated by request from client (vending machine analogy)

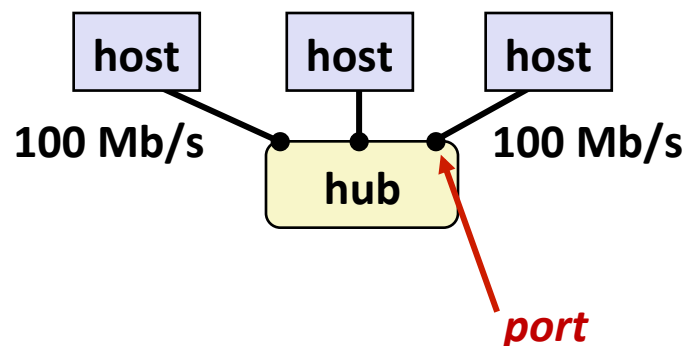
# Hardware Organization of a Network Host



# Computer Networks

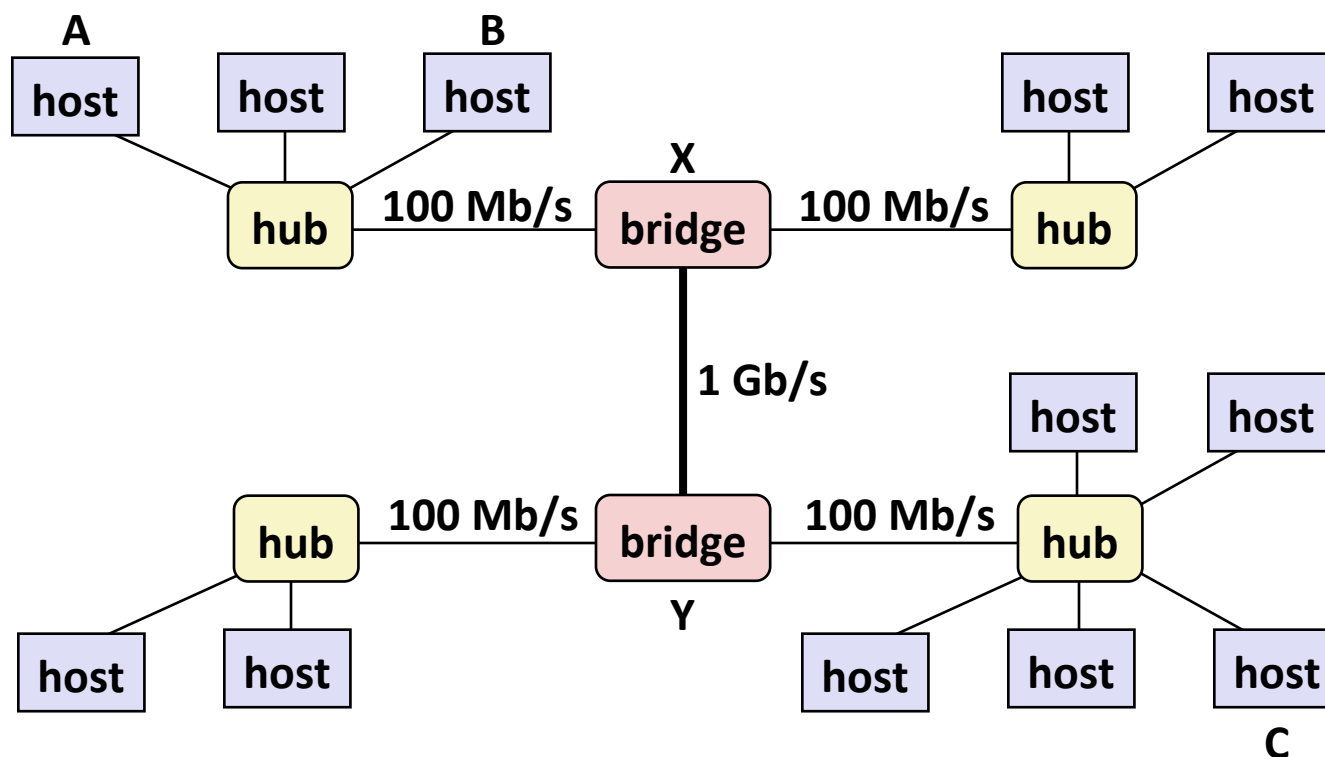
- A ***network*** is a hierarchical system of boxes and wires (or not) organized by geographical proximity
  - SAN (System Area Network) spans cluster or machine room
    - Switched Ethernet, Quadrics QSW, ...
  - LAN (Local Area Network) spans a building or campus
    - Ethernet is most prominent example
  - WAN (Wide Area Network) spans country or world
    - Typically high-speed point-to-point phone lines
- An ***internetwork*** (*internet*) is an interconnected set of networks
  - The Global IP Internet (uppercase “I”) is the most famous example of an internet (lowercase “i”)
- Let’s see how an internet is built from the ground up

# Lowest Level: Ethernet Segment



- Ethernet segment consists of a collection of *hosts* connected by wires (twisted pairs) to a *hub*
- Spans room or floor in a building
- Operation
  - Each Ethernet adapter has a unique 48-bit address (MAC address)
    - E.g., 00:16:ea:e3:54:e6
  - Hosts send bits to any other host in chunks called *frames*
  - Hub slavishly copies each bit from each port to every other port
    - Every host sees every bit
    - Note: Hubs are on their way out. Bridges (switches, routers) became cheap enough to replace them (means no more broadcasting)

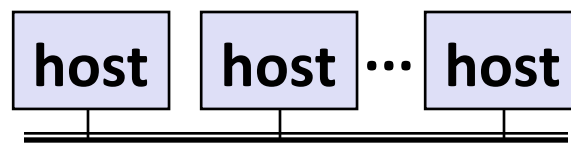
# Next Level: Bridged Ethernet Segment



- Spans building or campus
- Bridges cleverly learn which hosts are reachable from which ports and then selectively copy frames from port to port

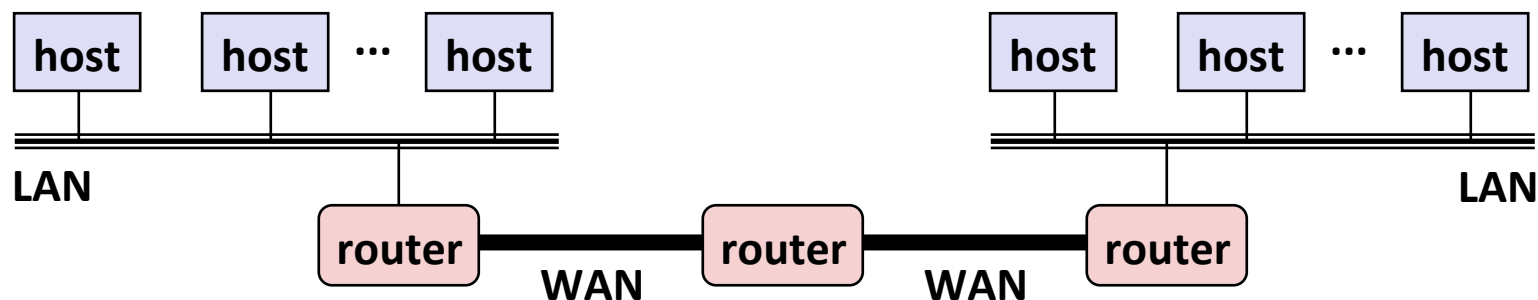
# Conceptual View of LANs

- For simplicity, hubs, bridges, and wires are often shown as a collection of hosts attached to a single wire:



# Next Level: internets

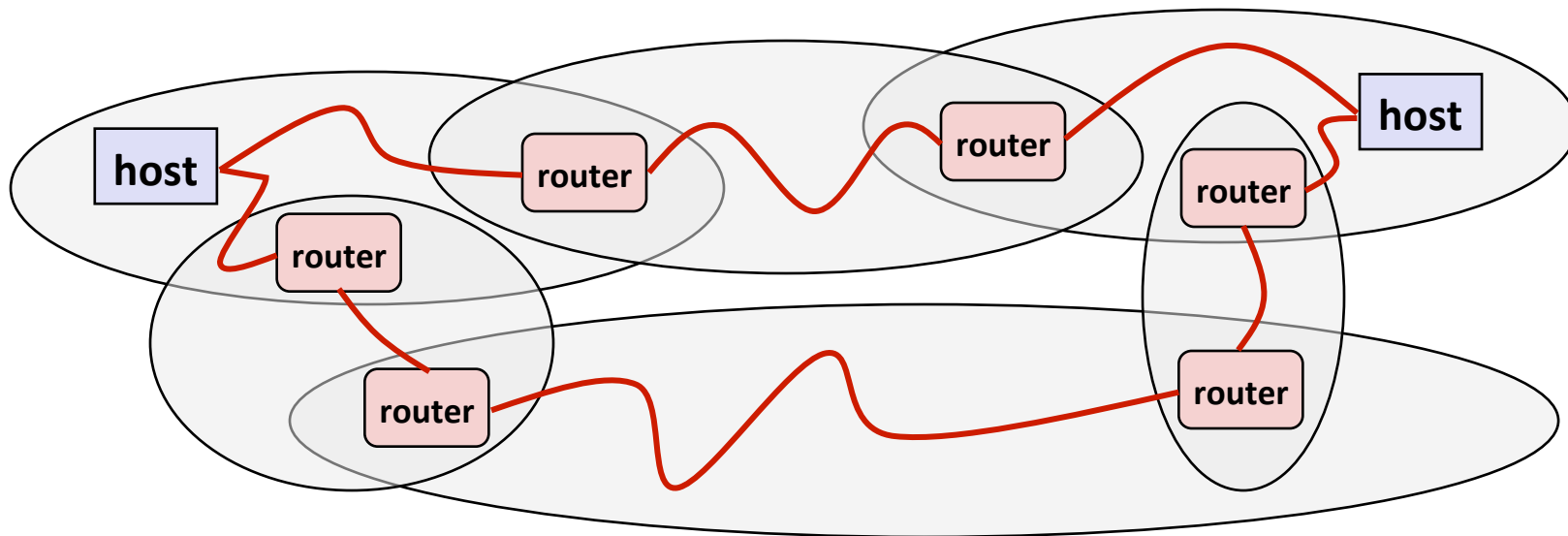
- Multiple incompatible LANs can be physically connected by specialized computers called **routers**
- The connected networks are called an **internet**



*LAN 1 and LAN 2 might be completely different, totally incompatible (e.g., Ethernet and Wifi, 802.11\*, T1-links, DSL, ...)*



# Logical Structure of an internet



- **Ad hoc interconnection of networks**
  - No particular topology
  - Vastly different router & link capacities
- **Send packets from source to destination by hopping through networks**
  - Router forms bridge from one network to another
  - Different packets may take different routes

# The Notion of an internet Protocol

- **How is it possible to send bits across incompatible LANs and WANs?**
- **Solution:**
  - protocol software running on each host and router
  - smooths out the differences between the different networks
- **Implements an internet protocol (i.e., set of rules)**
  - governs how hosts and routers should cooperate when they transfer data from network to network
  - TCP/IP is the protocol for the global IP Internet

# What Does an internet Protocol Do?

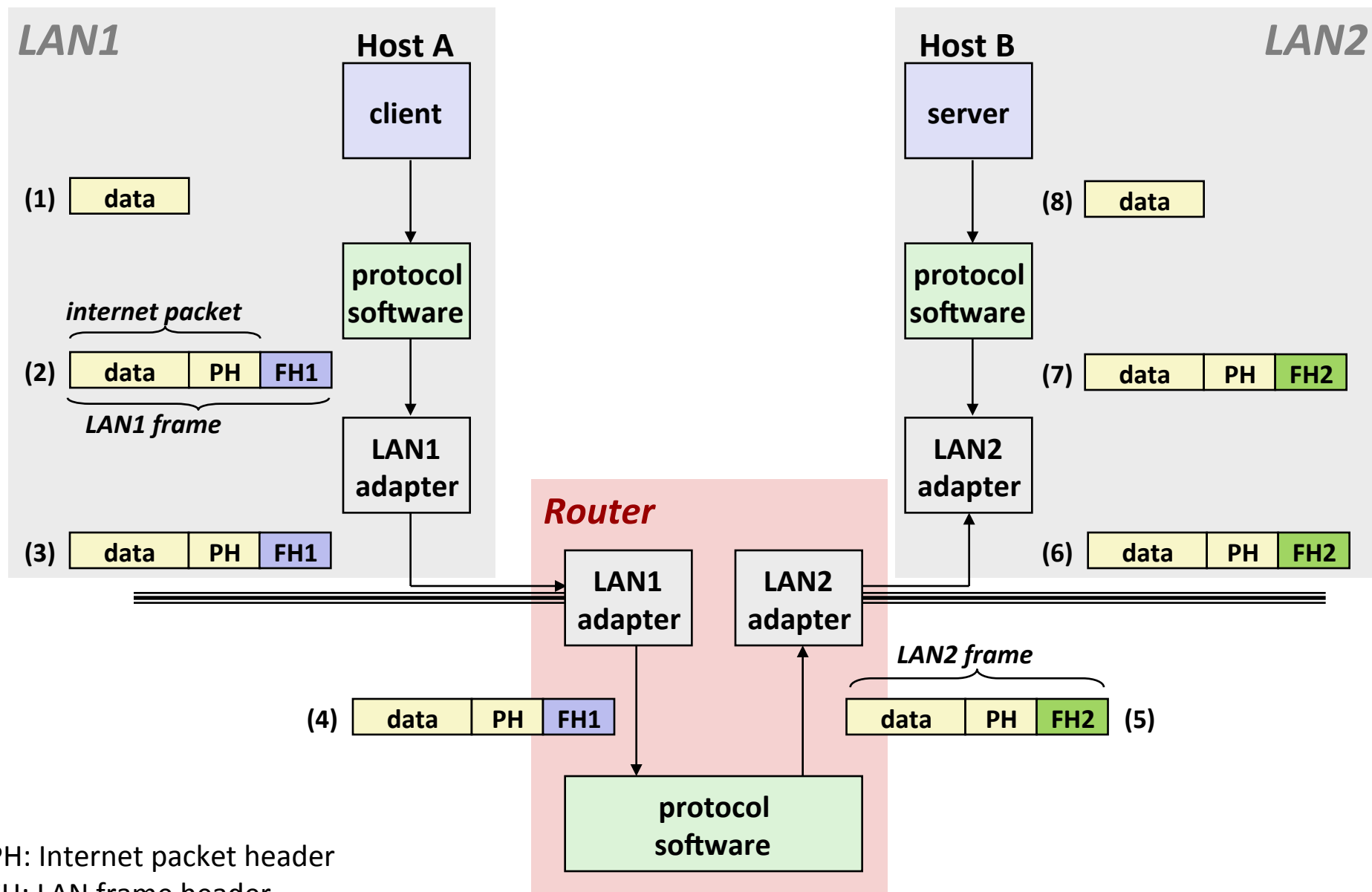
## ■ Provides a naming scheme

- An internet protocol defines a uniform format for *host addresses*
- Each host (and router) is assigned at least one of these internet addresses that uniquely identifies it

## ■ Provides a delivery mechanism

- An internet protocol defines a standard transfer unit (*packet*)
- Packet consists of *header* and *payload*
  - Header: contains info such as packet size, source and destination addresses
  - Payload: contains data bits sent from source host

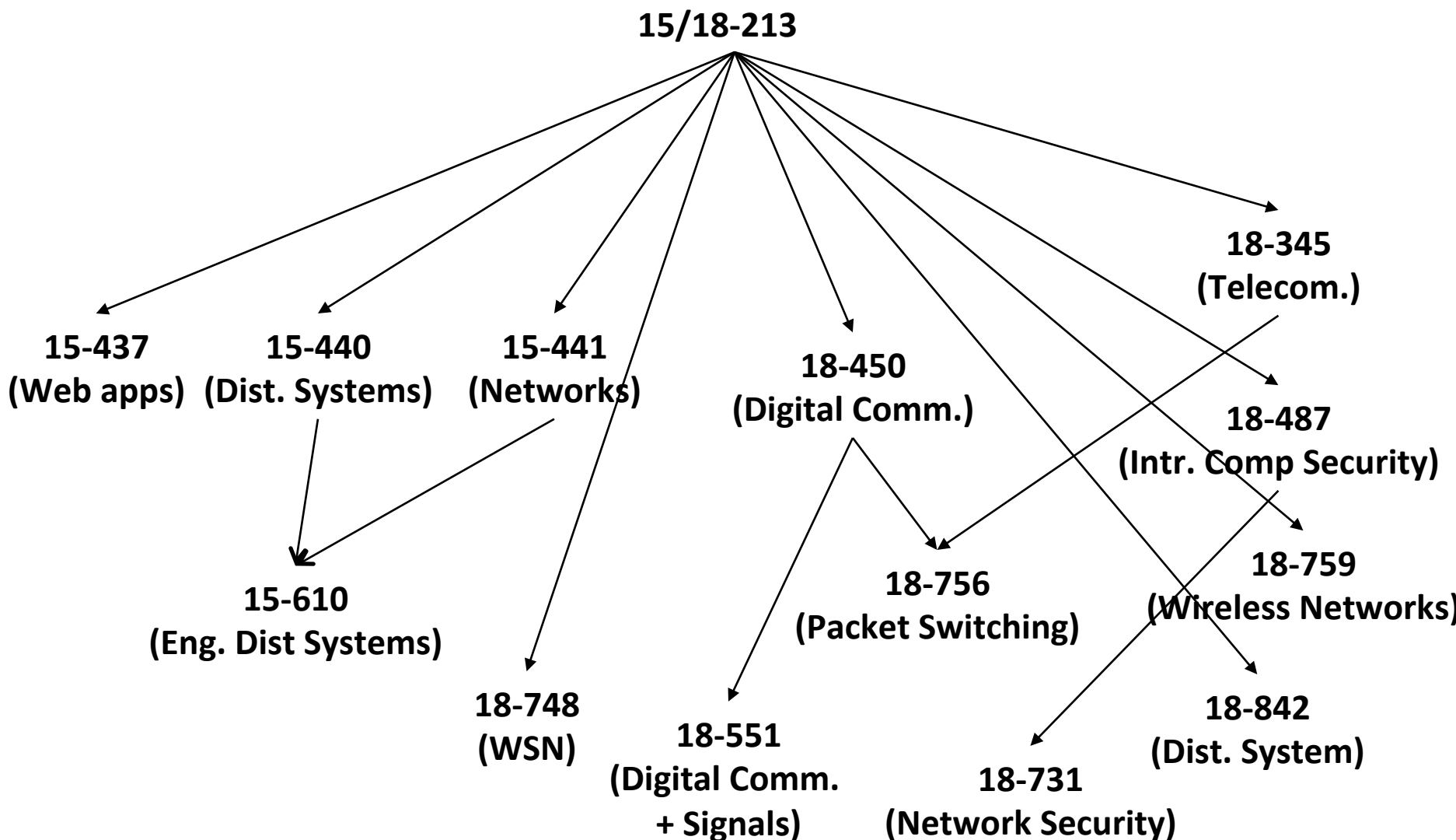
# Transferring Data Over an internet



# Other Issues

- **We are glossing over a number of important questions:**
  - What if different networks have different maximum frame sizes? (segmentation)
  - How do routers know where to forward frames?
  - How are routers informed when the network topology changes?
  - What if packets get lost?
  
- **These (and other) questions are addressed by the area of systems known as *computer networking***

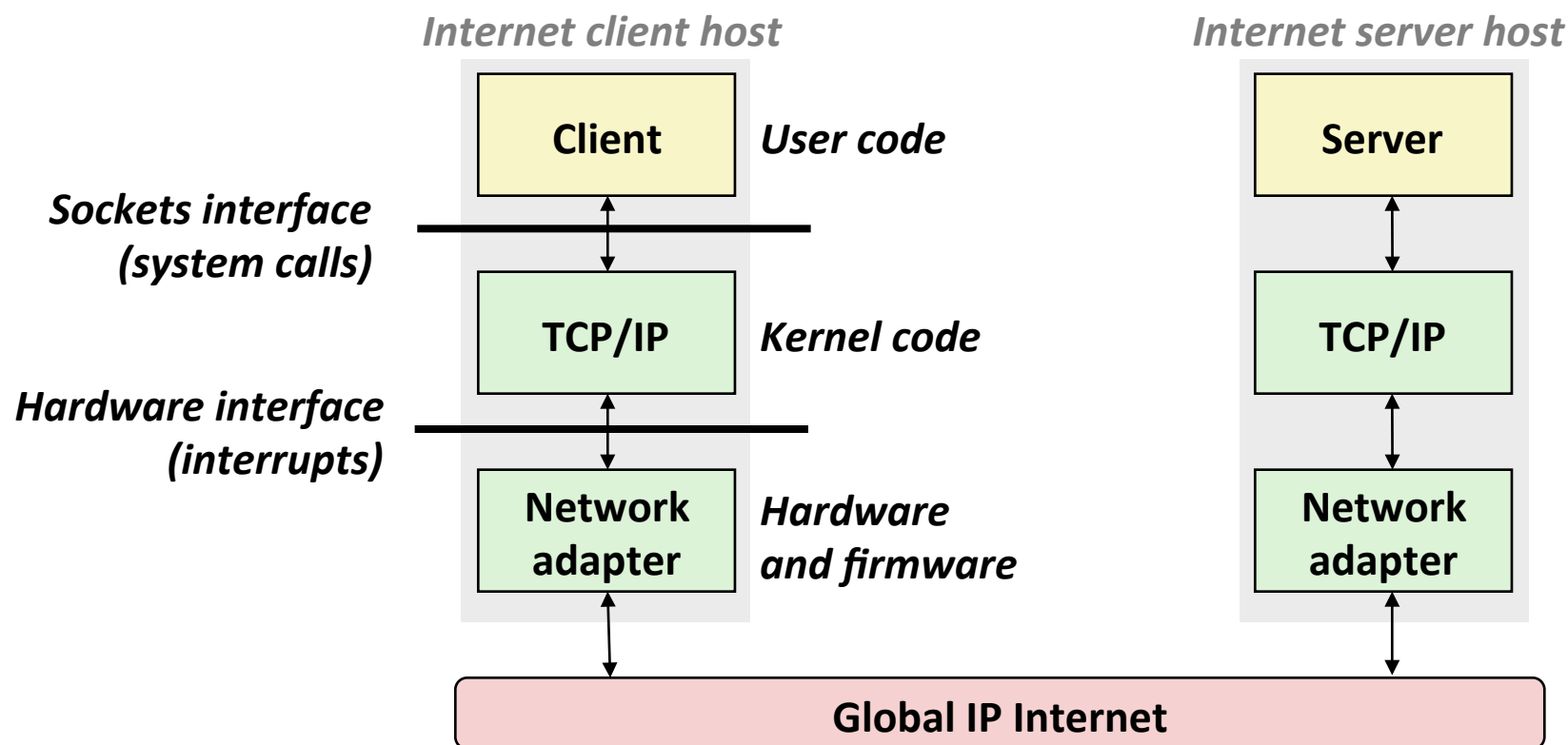
# If you are interested in networking...



# Global IP Internet

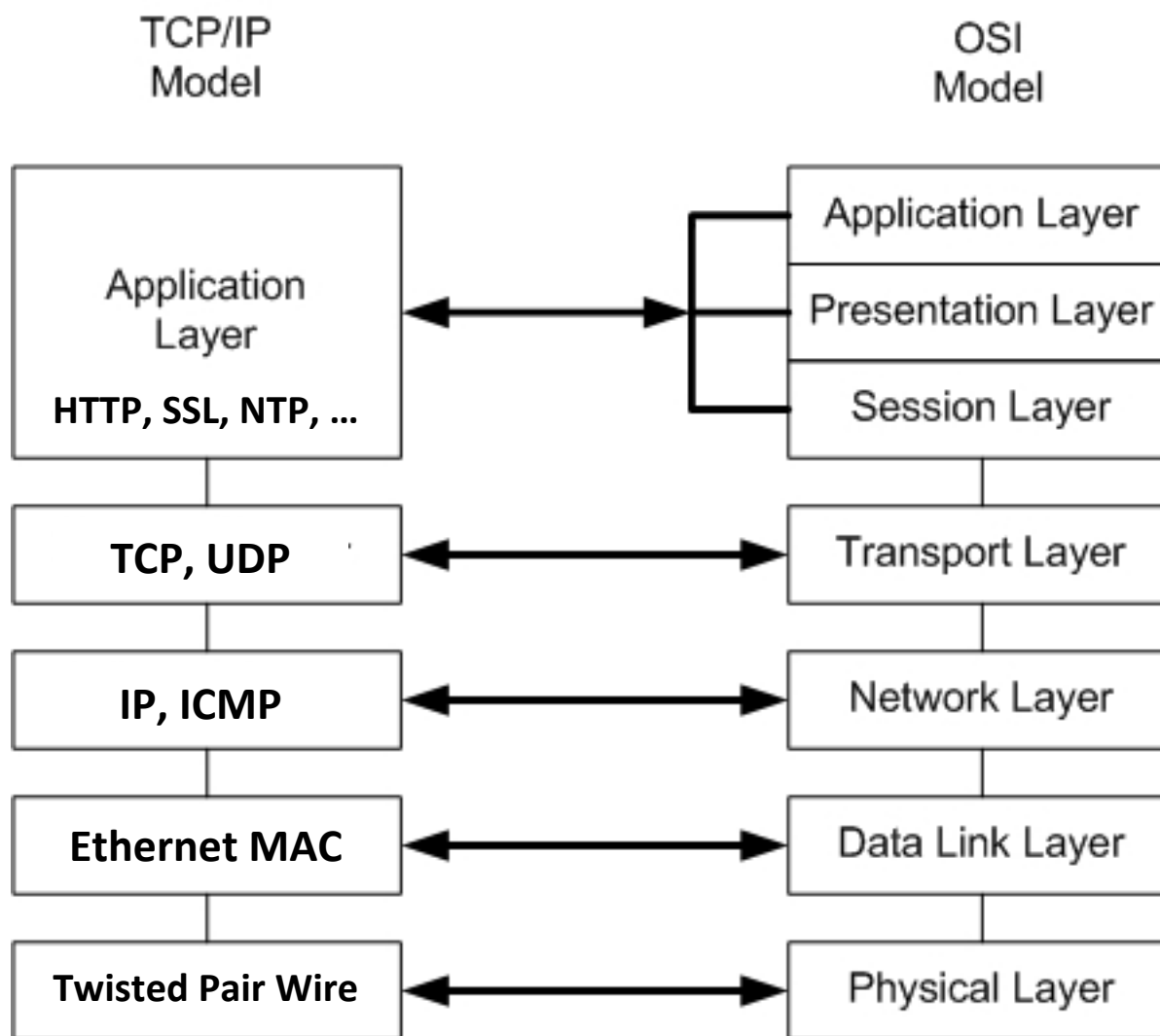
- Most famous example of an internet
- Based on the TCP/IP protocol family
  - IP (Internet protocol) :
    - Provides *basic naming scheme* and unreliable *delivery capability* of packets (datagrams) from host-to-host
  - UDP (Unreliable Datagram Protocol)
    - Uses IP to provide unreliable datagram delivery from *process-to-process*
  - TCP (Transmission Control Protocol)
    - Uses IP to provide *reliable* byte streams from process-to-process over connections
- Accessed via a mix of Unix file I/O and functions from the *sockets interface*

# Hardware and Software Organization of an Internet Application





# Protocol Layers



# Basic Internet Components

## ■ Internet backbone:

- collection of routers (nationwide or worldwide) connected by high-speed point-to-point networks

## ■ Network Access Point (NAP):

- router that connects multiple backbones (often referred to as peers)

## ■ Regional networks:

- smaller backbones that cover smaller geographical areas (e.g., cities or states)

## ■ Point of presence (POP):

- machine that is connected to the Internet

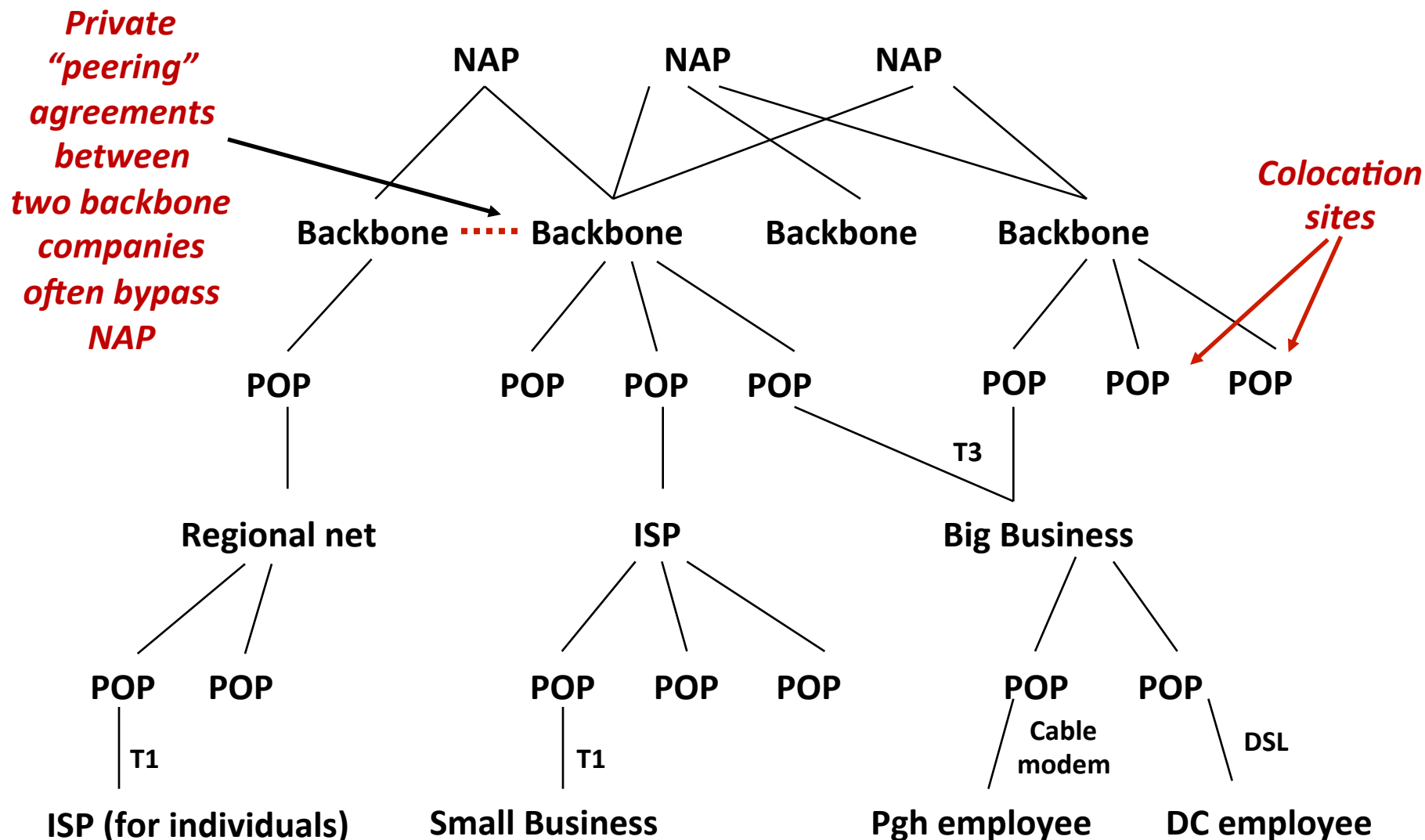
## ■ Internet Service Providers (ISPs):

- provide dial-up or direct access to POPs

# NAP-Based Internet Architecture

- **NAPs link together commercial backbones provided by companies such as AT&T and Worldcom**
- **Currently in the US there are about 50 commercial backbones connected by ~12 NAPs (peering points)**
- **Similar architecture worldwide connects national networks to the Internet**

# Internet Connection Hierarchy



# Naming and Communicating on the Internet

## ■ Original Idea

- Every node on Internet would have unique IP address
  - Everyone would be able to talk directly to everyone
- No secrecy or authentication
  - Messages visible to routers and hosts on same LAN
  - Possible to forge source field in packet header

## ■ Shortcomings

- There aren't enough IP addresses available
- Don't want everyone to have access or knowledge of all other hosts
- Security issues mandate secrecy & authentication

# Evolution of Internet: Naming

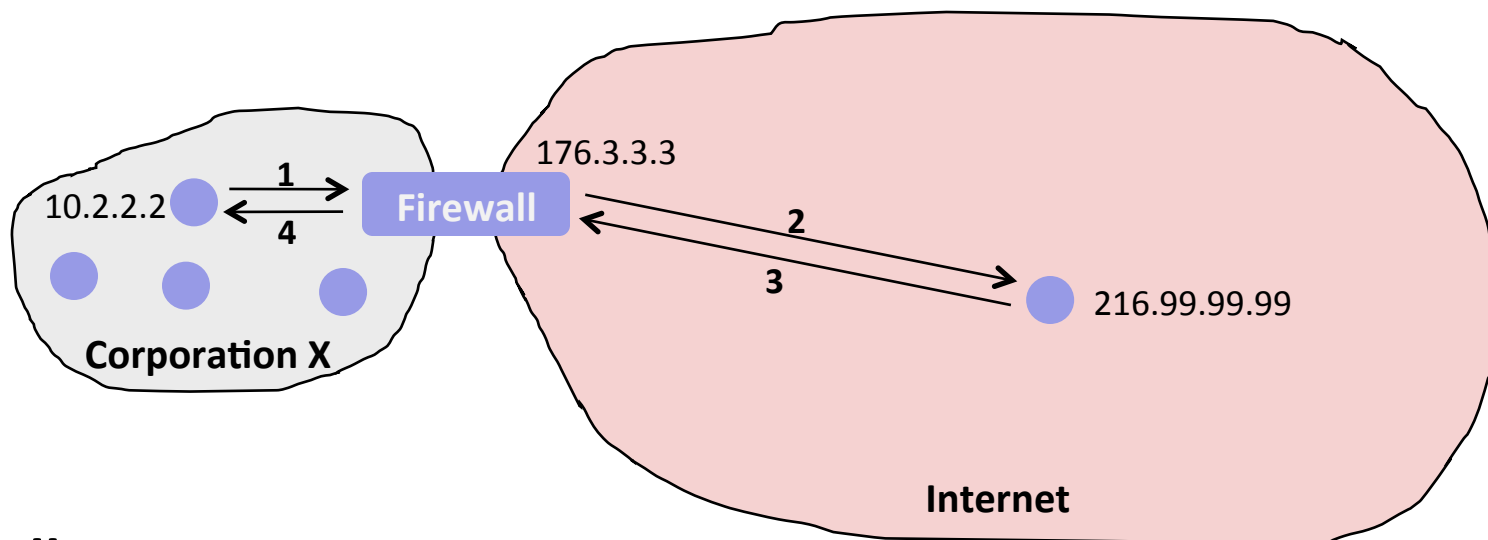
## ■ Dynamic address assignment

- Most hosts don't need to have known address
  - Only those functioning as servers
- DHCP (Dynamic Host Configuration Protocol)
  - Local ISP assigns address for temporary use

## ■ Example:

- My laptop at CMU (wired connection)
  - IP address 128.237.207.3 (**corvina.ece.cmu.edu**)
  - Assigned statically
- My laptop at home
  - IP address 192.168.1.5
  - Only valid within home network

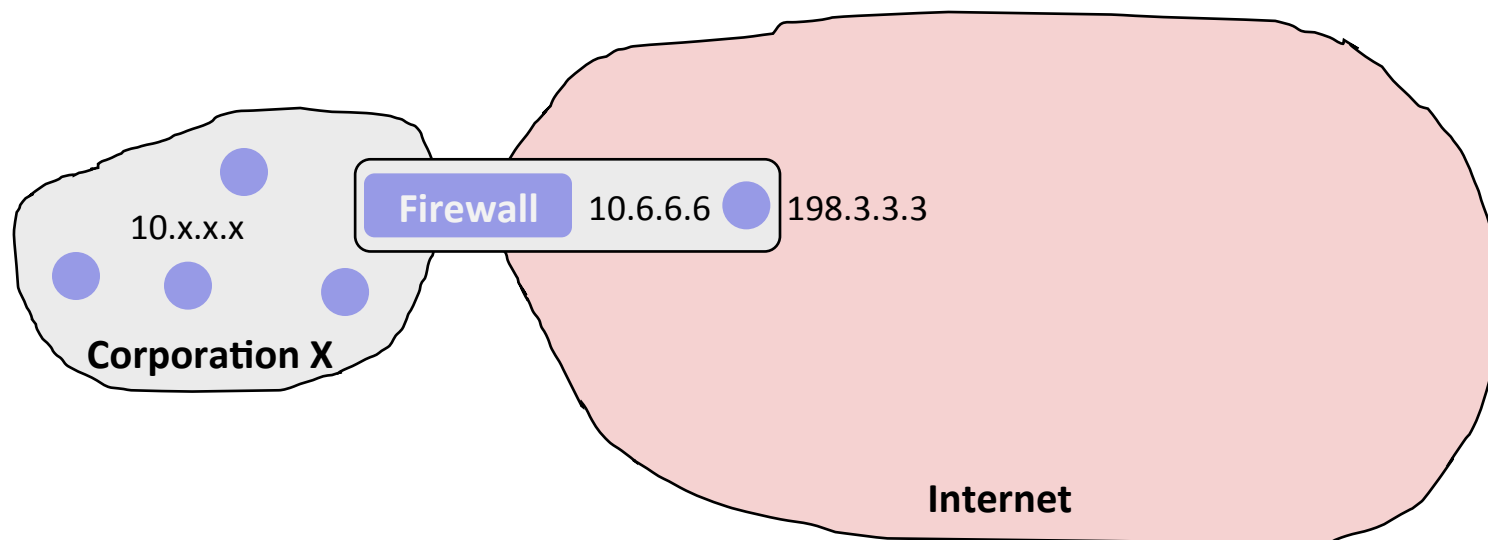
# Evolution of Internet: Firewalls



## ■ Firewalls

- Hides organizations nodes from rest of Internet
- Use local IP addresses within organization
- For external service, provides proxy service
  1. Client request: src=10.2.2.2, dest=216.99.99.99
  2. Firewall forwards: src=176.3.3.3, dest=216.99.99.99
  3. Server responds: src=216.99.99.99, dest=176.3.3.3
  4. Firewall forwards response: src=216.99.99.99, dest=10.2.2.2

# Virtual Private Networks



## ■ Supporting road warrior

- Employee working remotely with assigned IP address 198.3.3.3
- Wants to appear to rest of corporation as if working internally
  - From address 10.6.6.6
  - Gives access to internal services (e.g., ability to send mail)

## ■ Virtual Private Network (VPN)

- Overlays private network on top of regular Internet



# A Programmer's View of the Internet

- Hosts are mapped to a set of 32-bit *IP addresses*
  - 128.2.203.179
- The set of IP addresses is mapped to a set of identifiers called Internet *domain names*
  - 128.2.203.179 is mapped to `www.cs.cmu.edu`
- A process on one Internet host can communicate with a process on another Internet host over a *connection*

# IP Addresses

- **32-bit IP addresses are stored in an *IP address struct***
  - IP addresses are always stored in memory in network byte order (big-endian byte order)
  - True in general for any integer transferred in a packet header from one machine to another.
    - E.g., the port number used to identify an Internet connection.

```
/* Internet address structure */  
struct in_addr {  
    unsigned int s_addr; /* network byte order (big-endian) */  
};
```

**Useful network byte-order conversion functions (“l” = 32 bits, “s” = 16 bits)**

**htonl**: convert uint32\_t from host to network byte order

**htons**: convert uint16\_t from host to network byte order

**ntohl**: convert uint32\_t from network to host byte order

**ntohs**: convert uint16\_t from network to host byte order

# Dotted Decimal Notation

- By convention, each byte in a 32-bit IP address is represented by its decimal value and separated by a period
  - IP address: `0x8002C2F2` = `128.2.194.242`
- Functions for converting between binary IP addresses and dotted decimal strings:
  - `inet_aton`: dotted decimal string → IP address in network byte order
  - `inet_ntoa`: IP address in network byte order → dotted decimal string
  - “n” denotes network representation
  - “a” denotes application representation

# IP Address Structure

## ■ IP (V4) Address space divided into classes:

|         | 0 | 1      | 2      | 3      | 8                        | 16      | 24      | 31 |
|---------|---|--------|--------|--------|--------------------------|---------|---------|----|
| Class A | 0 | Net ID |        |        | Host ID                  |         |         |    |
| Class B | 1 | 0      | Net ID |        |                          | Host ID |         |    |
| Class C | 1 | 1      | 0      | Net ID |                          |         | Host ID |    |
| Class D | 1 | 1      | 1      | 0      | Multicast address        |         |         |    |
| Class E | 1 | 1      | 1      | 1      | Reserved for experiments |         |         |    |

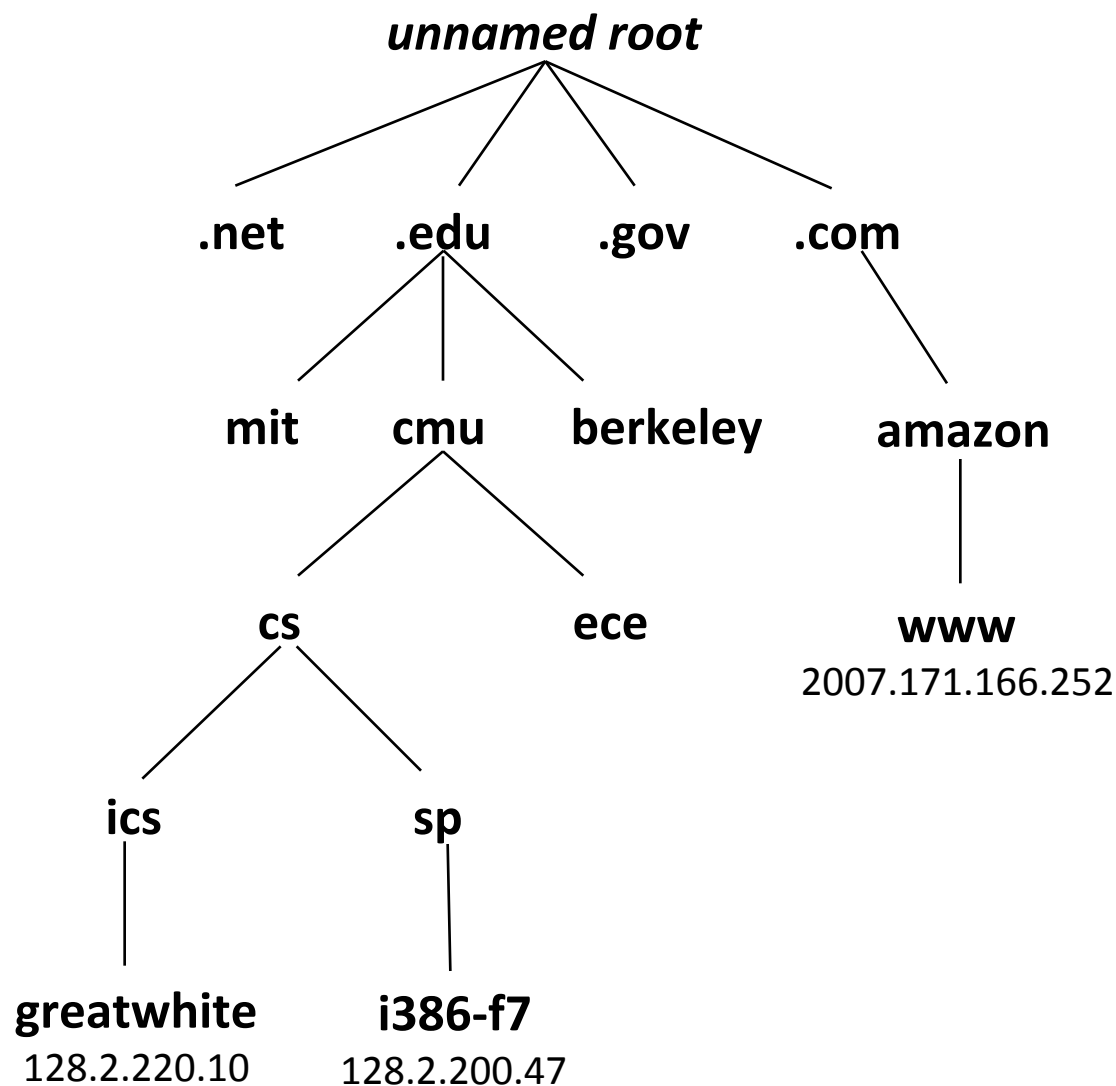
## ■ Network ID Written in form w.x.y.z/n

- n = number of bits in host address
- E.g., CMU written as 128.2.0.0/16
  - Class B address

## ■ Unrouted (private) IP addresses:

10.0.0.0/8   172.16.0.0/12   192.168.0.0/16

# Internet Domain Names



# Domain Naming System (DNS)

- The Internet maintains a mapping between IP addresses and domain names in a huge worldwide distributed database called **DNS**

- Conceptually, programmers can view the DNS database as a collection of millions of *host entry structures*:

```
/* DNS host entry structure */
struct hostent {
    char    *h_name;           /* official domain name of host */
    char    **h_aliases;       /* null-terminated array of domain names */
    int     h_addrtype;        /* host address type (AF_INET) */
    int     h_length;          /* length of an address, in bytes */
    char    **h_addr_list;     /* null-terminated array of in_addr structs
*/
};
```

- Functions for retrieving host entries from DNS:

- **gethostbyname**: query key is a DNS domain name.
- **gethostbyaddr**: query key is an IP address.

# Properties of DNS Host Entries

- Each host entry is an equivalence class of domain names and IP addresses
- Each host has a locally defined domain name `localhost` which always maps to the *loopback address* `127.0.0.1`
- Different kinds of mappings are possible:
  - Simple case: one-to-one mapping between domain name and IP address:
    - `greatwhile.ics.cs.cmu.edu` maps to `128.2.220.10`
  - Multiple domain names mapped to the same IP address:
    - `eeecs.mit.edu` and `cs.mit.edu` both map to `18.62.1.6`
  - Multiple domain names mapped to multiple IP addresses:
    - `google.com` maps to multiple IP addresses
  - Some valid domain names don't map to any IP address:
    - for example: `ics.cs.cmu.edu`

# A Program That Queries DNS

```
int main(int argc, char **argv) { /* argv[1] is a domain name */
    char **pp;                    /* or dotted decimal IP addr */
    struct in_addr addr;
    struct hostent *hostp;

    if (inet_aton(argv[1], &addr) != 0)
        hostp = Gethostbyaddr((const char *)&addr, sizeof(addr),
                               AF_INET);
    else
        hostp = Gethostbyname(argv[1]);
    printf("official hostname: %s\n", hostp->h_name);

    for (pp = hostp->h_aliases; *pp != NULL; pp++)
        printf("alias: %s\n", *pp);

    for (pp = hostp->h_addr_list; *pp != NULL; pp++) {
        addr.s_addr = ((struct in_addr *)*pp)->s_addr;
        printf("address: %s\n", inet_ntoa(addr));
    }
}
```



# Using DNS Program

```
linux> ./dns greatwhite.ics.cs.cmu.edu  
official hostname: greatwhite.ics.cs.cmu.edu  
address 128.2.220.10
```

```
linux> ./dns 128.2.220.11  
official hostname: ANGELSHARK.ICS.CS.CMU.EDU  
address: 128.2.220.11
```

```
linux> ./dns www.google.com  
official hostname: www.l.google.com  
alias: www.google.com  
address: 72.14.204.99  
address: 72.14.204.103  
address: 72.14.204.104  
address: 72.14.204.147  
linux> dig +short -x 72.14.204.103  
iad04s01-in-f103.1e100.net.
```

# Querying DIG

- Domain Information Groper (dig) provides a scriptable command line interface to DNS

```
linux> dig +short greatwhite.ics.cs.cmu.edu
128.2.220.10
linux> dig +short -x 128.2.220.11
ANGELSHARK.ICS.CS.CMU.EDU.
linux> dig +short google.com
72.14.204.104
72.14.204.147
72.14.204.99
72.14.204.103
linux> dig +short -x 72.14.204.103
iad04s01-in-f103.1e100.net.
```

# More Exotic Features of DIG

- Provides more information than you would ever want about DNS

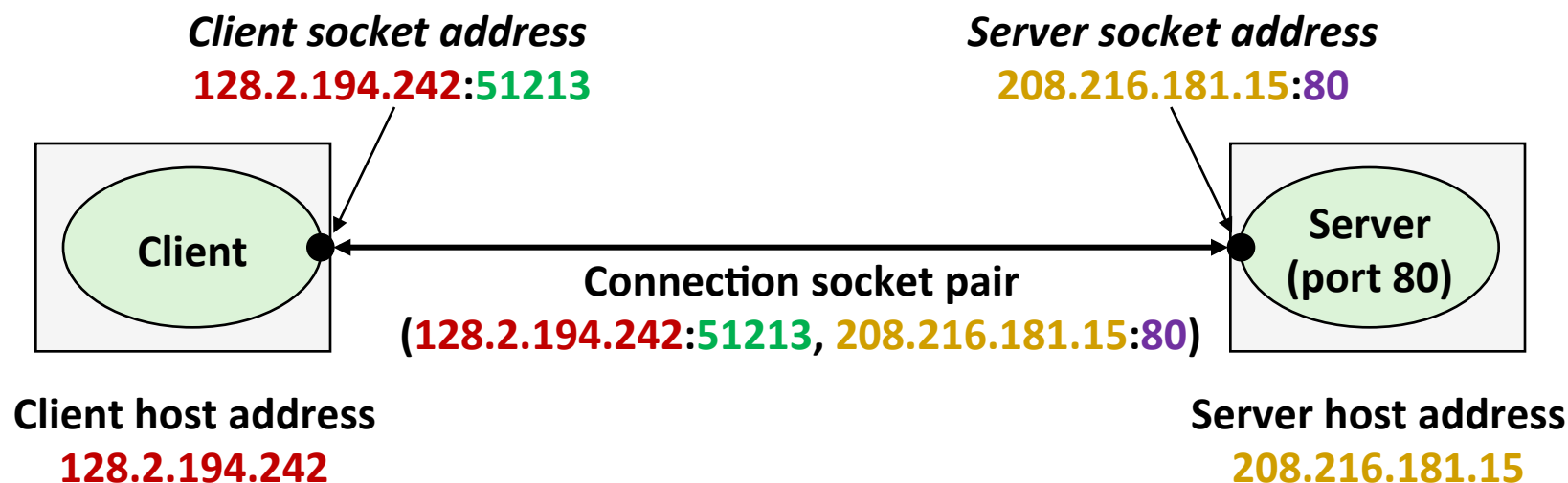
```
linux> dig www.phys.msu.ru a +trace  
128.2.220.10
```

```
linux> dig www.google.com a +trace
```

# Internet Connections

- Clients and servers communicate by sending streams of bytes over **connections**:
  - Point-to-point, full-duplex (2-way communication), and reliable.
- A **socket** is an endpoint of a connection
  - Socket address is an `IPAddress:port` pair
- A **port** is a 16-bit integer that identifies a process:
  - **Ephemeral port**: Assigned automatically on client when client makes a connection request
  - **Well-known port**: Associated with some service provided by a server (e.g., port 80 is associated with Web servers)
- A connection is uniquely identified by the socket addresses of its endpoints (**socket pair**)
  - `(cliaddr:cliport, servaddr:servport)`

# Putting it all Together: Anatomy of an Internet Connection



# Next Time

- **How to use the sockets interface to establish Internet connections between clients and servers**
- **How to use Unix I/O to copy data from one host to another over an Internet connection**