Exceptional Control Flow: Signals and Nonlocal Jumps

15-213 / 18-213: Introduction to Computer Systems 14th Lecture, Feb 28, 2013

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ECF Exists at All Levels of a System

Exceptions

Hardware and operating system kernel software

Process Context Switch

Hardware timer and kernel software

Signals

Kernel software and application software

Nonlocal jumps

Application code

Previous Lecture

This Lecture

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Today

- Multitasking, shells
- Signals
- Nonlocal jumps

The World of Multitasking

- System runs many processes concurrently
- Process: executing program
 - State includes memory image + register values + program counter
- Regularly switches from one process to another
 - Suspend process when it needs I/O resource or timer event occurs
 - Resume process when I/O available or given scheduling priority
- Appears to user(s) as if all processes executing simultaneously
 - Even though most systems can only execute one process at a time
 - Except possibly with lower performance than if running alone

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Programmer's Model of Multitasking

Basic functions

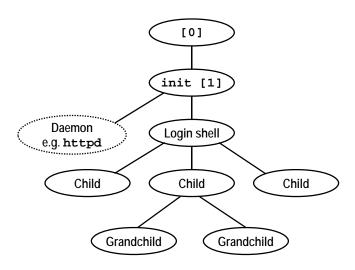
- fork spawns new process
 - Called once, returns twice
- exit terminates own process
 - Called once, never returns
 - Puts it into "zombie" status
- wait and waitpid wait for and reap terminated children
- execve runs new program in existing process
 - Called once, (normally) never returns

Programming challenge

- Understanding the nonstandard semantics of the functions
- Avoiding improper use of system resources
 - E.g. "Fork bombs" can disable a system

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Unix Process Hierarchy



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Shell Programs

- A shell is an application program that runs programs on behalf of the user.
 - sh Original Unix shell (Stephen Bourne, AT&T Bell Labs, 1977)
 - csh BSD Unix C shell (tcsh: enhanced csh at CMU and elsewhere)
 - bash "Bourne-Again" Shell

```
int main() {
    char cmdline[MAXLINE];

while (1) {
    /* read */
    printf("> ");
    Fgets(cmdline, MAXLINE, stdin);
    if (feof(stdin))
        exit(0);

    /* evaluate */
    eval(cmdline);
    }
}
```

Execution is a sequence of read/evaluate steps

Simple Shell eval Function

```
void eval(char *cmdline) {
    char *argv[MAXARGS]; /* argv for execve() */
   int bg;
                        /* should the job run in bg or fg? */
   pid_t pid;
                         /* process id */
   bg = parseline(cmdline, argv);
   if (!builtin command(argv)) {
                                   /* child runs user job */
       if ((pid = Fork()) == 0) {
           if (execve(argv[0], argv, environ) < 0) {</pre>
               printf("%s: Command not found.\n", argv[0]);
               exit(0);
       if (!bg) { /* parent waits for fg job to terminate */
           int status;
           if (waitpid(pid, &status, 0) < 0)</pre>
               unix_error("waitfg: waitpid error");
                    /* otherwise, don't wait for bg job */
           printf("%d %s", pid, cmdline);
```

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Simple Shell eval Function

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void eval(char *cmdline) {
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           printf("%d %s", pid, cmdline);
```

What Is a "Background Job"?

- Users generally run one command at a time
 - Type command, read output, type another command
- Some programs run "for a long time"
 - Example: "delete this file in two hours"

```
unix> sleep 7200; rm /tmp/junk # shell stuck for 2 hours
```

A "background" job is a process we don't want to wait for

```
unix> (sleep 7200 ; rm /tmp/junk) & [1] 907
unix> # ready for next command
```

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Problem with Simple Shell Example

- Our example shell correctly waits for and reaps foreground jobs
- But what about background jobs?
 - Will become zombies when they terminate
 - Will never be reaped because shell (typically) will not terminate
 - Will create a memory leak that could run the kernel out of memory
 - Modern Unix: once you exceed your process quota, your shell can't run any new commands for you: fork() returns -1

```
unix> limit maxproc # csh syntax
maxproc 202752
unix> ulimit -u # bash syntax
202752
```

ECF to the Rescue!

- Problem
 - The shell doesn't know when a background job will finish
 - By nature, it could happen at any time
 - The shell's regular control flow can't reap exited background processes in a timely fashion
 - Regular control flow is "wait until running job completes, then reap it"
- Solution: Exceptional control flow
 - The kernel will interrupt regular processing to alert us when a background process completes
 - In Unix, the alert mechanism is called a signal

Today

- Multitasking, shells
- Signals
- Nonlocal jumps

Signals

- A signal is a small message that notifies a process that an event of some type has occurred in the system
 - akin to exceptions and interrupts
 - sent from the kernel (sometimes at the request of another process) to a process
 - signal type is identified by small integer ID's (1-30)
 - only information in a signal is its ID and the fact that it arrived

ID	Name	Default Action	Corresponding Event
2	SIGINT	Terminate	Interrupt (e.g., ctl-c from keyboard)
9	SIGKILL	Terminate	Kill program (cannot override or ignore)
11	SIGSEGV	Terminate & Dump	Segmentation violation
14	SIGALRM	Terminate	Timer signal
17	SIGCHLD	Ignore	Child stopped or terminated

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Sending a Signal

- Kernel *sends* (delivers) a signal to a *destination process* by updating some state in the context of the destination process
- Kernel sends a signal for one of the following reasons:
 - Kernel has detected a system event such as divide-by-zero (SIGFPE) or the termination of a child process (SIGCHLD)
 - Another process has invoked the kill system call to explicitly request the kernel to send a signal to the destination process

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Receiving a Signal

- A destination process receives a signal when it is forced by the kernel to react in some way to the delivery of the signal
- Three possible ways to react:
 - Ignore the signal (do nothing)
 - *Terminate* the process (with optional core dump)
 - Catch the signal by executing a user-level function called signal handler
 - Akin to a hardware exception handler being called in response to an asynchronous interrupt

Pending and Blocked Signals

- A signal is pending if sent but not yet received
 - There can be at most one pending signal of any particular type
 - Important: Signals are not queued
 - If a process has a pending signal of type k, then subsequent signals of type k that are sent to that process are discarded
- A process can *block* the receipt of certain signals
 - Blocked signals can be delivered, but will not be received until the signal is unblocked
- A pending signal is received at most once

No Counting!!!

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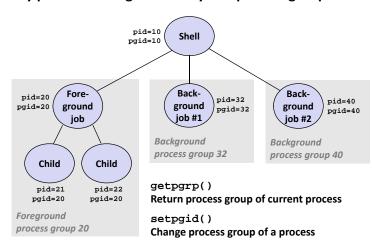
Signal Concepts

- Kernel maintains pending and blocked bit vectors in the context of each process
 - pending: represents the set of pending signals
 - Kernel sets bit k in **pending** when a signal of type k is delivered
 - Kernel clears bit k in **pending** when a signal of type k is received
 - **blocked**: represents the set of blocked signals
 - Can be set and cleared by using the **sigprocmask** function

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Process Groups

Every process belongs to exactly one process group



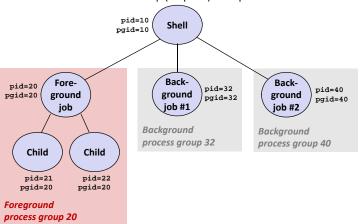
Sending Signals with /bin/kill Program

- /bin/kill program sends arbitrary signal to a process or process group
- Examples
 - /bin/kill -9 24818 Send SIGKILL to process 24818
 - /bin/kill -9 -24817
 Send SIGKILL to every process in process group 24817

linux> ./forks 16 Child1: pid=24818 pgrp=24817 Child2: pid=24819 pgrp=24817				
linux> ps				
PID TTY TIME CMD				
24788 pts/2 00:00:00 tcsh				
24818 pts/2 00:00:02 forks				
24819 pts/2 00:00:02 forks				
24820 pts/2 00:00:00 ps				
linux> /bin/kill -9 -24817				
linux> ps				
PID TTY TIME CMD				
24788 pts/2 00:00:00 tcsh				
24823 pts/2 00:00:00 ps				
linux>				

Sending Signals from the Keyboard

- Typing ctrl-c (ctrl-z) sends a SIGINT (SIGTSTP) to every job in the foreground process group.
 - SIGINT default action is to terminate each process
 - SIGTSTP default action is to stop (suspend) each process



Example of ctrl-c and ctrl-z

```
bluefish> ./forks 17
Child: pid=28108 pgrp=28107
                                               First letter:
Parent: pid=28107 pgrp=28107
<types ctrl-z>
                                               S: sleeping
Suspended
bluefish> ps w
 PID TTY
               STAT
                       TIME COMMAND
                       0:00 -tcsh
27699 pts/8
               Ss
28107 pts/8
                       0:01 ./forks 17
28108 pts/8
                       0:01 ./forks 17
28109 pts/8
                       0:00 ps w
bluefish> fg
./forks 17
<types ctrl-c>
                                               details
bluefish> ps w
 PID TTY
               STAT
                       TIME COMMAND
27699 pts/8
               Ss
                       0:00 -tcsh
28110 pts/8
                       0:00 ps w
```

STAT (process state) Legend:

T: stopped

R: running

Second letter:

s: session leader

+: foreground proc group

See "man ps" for more

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Sending Signals with kill Function

```
void fork12()
   pid_t pid[N];
   int i, child_status;
   for (i = 0; i < N; i++)
        if ((pid[i] = fork()) == 0)
            while(1); /* Child infinite loop */
   /* Parent terminates the child processes */
   for (i = 0; i < N; i++) {
        printf("Killing process %d\n", pid[i]);
        kill(pid[i], SIGINT);
   /* Parent reaps terminated children */
   for (i = 0; i < N; i++) {
        pid_t wpid = wait(&child_status);
        if (WIFEXITED(child status))
            printf("Child %d terminated with exit status %d\n",
                    wpid, WEXITSTATUS(child status));
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```

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Sending Signals with kill Function

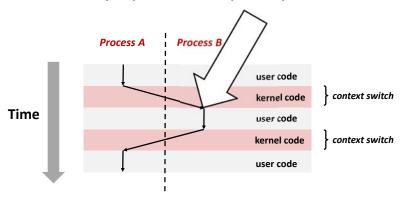
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Receiving Signals

 Suppose kernel is returning from an exception handler and is ready to pass control to process p



Important: All context switches are initiated by calling some exceptional hander.

Receiving Signals

 Suppose kernel is returning from an exception handler and is ready to pass control to process p

■ Kernel computes pnb = pending & ~blocked

The set of pending nonblocked signals for process p

 $\blacksquare \text{ If (pnb } == 0)$

Pass control to next instruction in the logical flow for p

Else

 Choose least nonzero bit k in pnb and force process p to receive signal k

• The receipt of the signal triggers some *action* by *p*

Repeat for all nonzero k in pnb

Pass control to next instruction in logical flow for p

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Default Actions

Each signal type has a predefined default action, which is one of:

The process terminates

The process terminates and dumps core

The process stops until restarted by a SIGCONT signal

The process ignores the signal

Installing Signal Handlers

The signal function modifies the default action associated with the receipt of signal signum:

handler t *signal(int signum, handler t *handler)

■ Different values for handler:

SIG IGN: ignore signals of type signum

SIG DFL: revert to the default action on receipt of signals of type signum

Otherwise, handler is the address of a signal handler

Called when process receives signal of type signum

Referred to as "installing" the handler

Executing handler is called "catching" or "handling" the signal

 When the handler executes its return statement, control passes back to instruction in the control flow of the process that was interrupted by receipt of the signal

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Signal Handling Example

```
void int_handler(int sig) {
    safe_printf("Process %d received signal %d\n", getpid(), sig);
    exit(0);
void fork13() {
    pid_t pid[N];
    int i, child_status;
    signal(SIGINT, int handler);
                                   linux> ./forks 13
   for (i = 0; i < N; i++)
       if ((pid[i] = fork()) == 0 Killing process 25417
           while(1); /* child inf Killing process 25418
                                   Killing process 25419
                                   Killing process 25420
   for (i = 0; i < N; i++) {
       printf("Killing process %d Killing process 25421
                                   Process 25417 received signal 2
        kill(pid[i], SIGINT);
                                   Process 25418 received signal 2
                                   Process 25420 received signal 2
   for (i = 0; i < N; i++) {
       pid_t wpid = wait(&child_s Process 25421 received signal 2
       if (WIFEXITED(child_status Process 25419 received signal 2
            printf("Child %d termi Child 25417 terminated with exit status 0
                   wpid, WEXITSTAT Child 25418 terminated with exit status 0
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        else
                                   Child 25419 terminated with exit status 0
            printf("Child %d termi:
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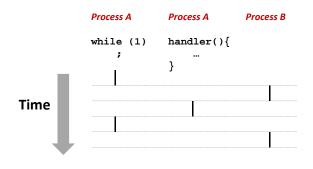
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        kill(pid[i], SIGINT);
                                   Process 25418 received signal 2
                                   Process 25420 received signal 2
    for (i = 0; i < N; i++) {
        pid_t wpid = wait(&child_s Process 25421 received signal 2
        if (WIFEXITED(child_status Process 25419 received signal 2
            printf("Child %d termi Child 25417 terminated with exit status 0
                   wpid, WEXITSTAT Child 25418 terminated with exit status 0
                                   Child 25420 terminated with exit status 0
            printf("Child %d termi Child 25419 terminated with exit status 0
                                   Child 25421 terminated with exit status 0
                                   linux>
```

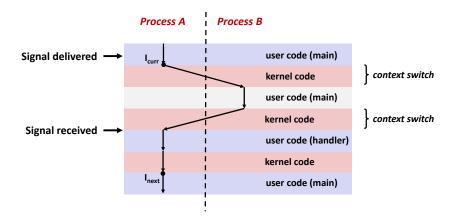
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Signals Handlers as Concurrent Flows

- A signal handler is a separate logical flow (not process) that runs concurrently with the main program
 - "concurrently" in the "not sequential" sense



Another View of Signal Handlers as Concurrent Flows



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Signal Handler Funkiness

```
Pending signals are not
int ccount = 0;
void child_handler(int sig)
                                                         queued
                                                         ■ For each signal type, just
    int child_status;
    pid_t pid = wait(&child_status);
                                                          have single bit indicating
    ccount--;
                                                          whether or not signal is
    safe_printf(
                                                          pending
           "Received signal %d from process %d\n",
          sig, pid);
                                                         Even if multiple processes
void fork14()
                                                          have sent this signal
   pid_t pid[N];
   int i, child_status;
    ccount = N:
    signal(SIGCHLD, child_handler);
   for (i = 0; i < N; i+ linux> ./forks 14
        if ((pid[i] = for Received SIGCHLD signal 17 for process 21344
            sleep(1); /* Received SIGCHLD signal 17 for process 21345
    while (ccount > 0)
        pause(); /* Suspend until signal occurs */
```

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Living With Nonqueuing Signals

- Must check for all terminated jobs
 - Typically loop with waitpid

```
void child handler2(int sig)
   int child status;
   while ((pid = waitpid(-1, &child_status, WNOHANG)) > 0) {
       safe_printf("Received signal %d from process %d\n",
                     sig, pid);
               greatwhite> forks 15
void fork15()
               Received signal 17 from process 27476
               Received signal 17 from process 27477
               Received signal 17 from process 27478
   signal(SIGC Received signal 17 from process 27479
               Received signal 17 from process 27480
               greatwhite>
```

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More Signal Handler Funkiness

- Signal arrival during long system calls (say a read)
- Signal handler interrupts read call
 - Linux: upon return from signal handler, the read call is restarted automatically
 - Some other flavors of Unix can cause the read call to fail with an **EINTR** error number (**errno**) in this case, the application program can restart the slow system call
- Subtle differences like these complicate the writing of portable code that uses signals
 - Consult your textbook for details

A Program That Reacts to Internally Generated Events

```
#include <stdlib.h>
#include <stdio.h>
#include <signal.h>
void handler(int sig) {
 safe_printf("You think hitting ctrl-c will stop the bomb?\n");
 safe_printf("Well...");
                                linux> ./external
 sleep(1);
                                 <ctrl-c>
 printf("OK\n");
                                You think hitting ctrl-c will stop
 exit(0);
                                 the bomb?
                                 Well...OK
                                linux>
main() {
```

```
#include <stdio.h>
#include <signal.h>

int beeps = 0;

/* SIGALRM handler */
void handler(int sig) {
    safe_printf("BEEP\n");

    if (++beeps < 5)
        alarm(1);
    else {
        safe_printf("BOOM!\n");
        exit(0);
    }
}
internal.c</pre>
```

```
linux> ./internal
BEEP
BEEP
BEEP
BEEP
BEEP
BOOM!
bass>
```

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Async-Signal-Safety

while(1) {

external.c

- Function is async-signal-safe if either reentrant (all variables stored on stack frame, CS:APP2e 12.7.2) or non-interruptible by signals.
- Posix guarantees 117 functions to be async-signal-safe
 - write is on the list, printf is not

A Program That Reacts to

Externally Generated Events (Ctrl-c)

signal(SIGINT, handler); /* installs ctl-c handler */

One solution: async-signal-safe wrapper for printf:

Today

- Multitasking, shells
- Signals
- Nonlocal jumps

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Nonlocal Jumps: setjmp/longjmp

- Powerful (but dangerous) user-level mechanism for transferring control to an arbitrary location
 - Controlled to way to break the procedure call / return discipline
 - Useful for error recovery and signal handling
- int setjmp(jmp_buf j)
 - Must be called before longimp
 - Identifies a return site for a subsequent longimp
 - Called once, returns one or more times
- Implementation:
 - Remember where you are by storing the current register context, stack pointer, and PC value in jmp_buf
 - Return 0

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Limitations of Nonlocal Jumps

■ Works within stack discipline

 Can only long jump to environment of function that has been called but not yet completed

env

```
if (setjmp(env)) {
    /* Long Jump to here */
} else {
    P2();
}
}
```

longjmp(env, 1);

P3()

P2
P2
P2
P3

Before longjmp

After longjmp

P1

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setjmp/longjmp (cont)

- void longjmp(jmp_buf j, int i)
 - Meaning:
 - return from the set jmp remembered by jump buffer j again ...
 - ... this time returning instead of 0
 - Called after setjmp
 - Called once, but never returns

■ longjmp Implementation:

- Restore register context (stack pointer, base pointer, PC value) from jump buffer j
- Set %eax (the return value) to i
- Jump to the location indicated by the PC stored in jump buf j

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setjmp/longjmp Example

```
#include <setjmp.h>
jmp_buf buf;

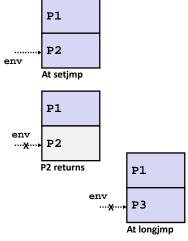
main() {
    if (setjmp(buf) != 0) {
        printf("back in main due to an error\n");
    else
        printf("first time through\n");
    pl(); /* pl calls p2, which calls p3 */
}
...
p3() {
    <error checking code>
    if (error)
        longjmp(buf, 1)
}
```

Limitations of Long Jumps (cont.)

Works within stack discipline

 Can only long jump to environment of function that has been called but not yet completed

```
jmp_buf env;
P1()
{
     P2(); P3();
}
P2()
{
     if (setjmp(env)) {
        /* Long Jump to here */
     }
}
P3()
{
     longjmp(env, 1);
}
```



Putting It All Together: A Program That Restarts Itself When ctrl-c'd

```
#include <stdio.h>
#include <signal.h>
#include <setjmp.h>

sigjmp_buf buf;

void handler(int sig) {
    siglongjmp(buf, 1);
}

main() {
    signal(SIGINT, handler);
    if (!sigsetjmp(buf, 1))
        printf("starting\n");
    else
        printf("restarting\n");
    while(1) {
        sleep(1);
        printf("processing...\n");
    }
}
```

```
greatwhite> ./restart
starting
processing...
processing...
restarting
processing...
processing...
ctrl-c
processing...
restarting
processing...
restarting
processing...
processing...
processing...
```

restart.c

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Summary

- Signals provide process-level exception handling
 - Can generate from user programs
 - Can define effect by declaring signal handler
- Some caveats
 - Very high overhead
 - >10,000 clock cycles
 - Only use for exceptional conditions
 - Don't have queues
 - Just one bit for each pending signal type
- Nonlocal jumps provide exceptional control flow within process
 - Within constraints of stack discipline

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Midterm Exam info

- Midterm on 3/5 at 6:30 for everyone
- Lastname starts with
 - A-J -> Wh 7500
 - K-Z -> Rashid
- Covers all material through today
- Review on Sunday at 3pm at Dh2315 and Dh2210
- Recitation will also do midterm review on Monday