

15-213
 "The course that gives CMU its Zip!"
**Structured Data I:
 Homogenous Data**
 February 8, 2001

Topics

- Arrays
 - Single
 - Nested
- Pointers
 - Multilevel Arrays
- Optimized Array Code

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Basic Data Types

Integral

- Stored & operated on in general registers
- Signed vs. unsigned depends on instructions used

Intel	GAS	Bytes	C
byte	b	1	[unsigned] char
word	w	2	[unsigned] short
double word	l	4	[unsigned] int

Floating Point

- Stored & operated on in floating point registers

Intel	GAS	Bytes	C
Single	s	4	float
Double	l	8	double
Extended	t	10/12	long double

Pointer

- Stored & operated on in general registers

Intel	GAS	Bytes	C
Double	l	8	char *

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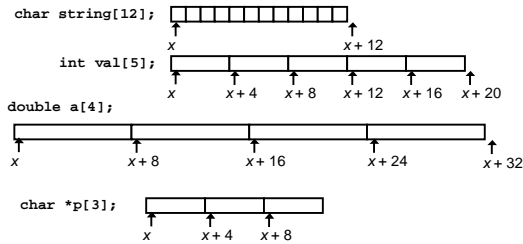
- 2 -

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Array Allocation

Basic Principle

- $T A[L];$
- Array named A of data type T with L elements
 - Contiguously allocated region of $L * \text{sizeof}(T)$ bytes



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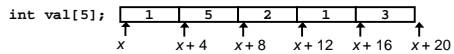
- 3 -

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Array Access

Basic Principle

- $T A[L];$
- Array named A of data type T and length L
 - Identifier A can be used as a pointer to starting element of the array



Reference	Type	Value
$\text{val}[4]$	int	3
val	int *	x
$\text{val}+1$	int *	$x+4$
$\&\text{val}[2]$	int *	$x+8$
$\text{val}[5]$	int	??
$\text{*(val}+1)$	int	5
$\text{val} + i$	int *	$x+4i$

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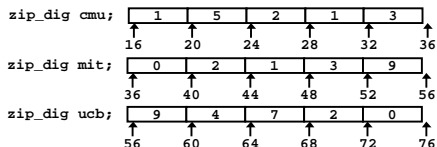
- 4 -

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Array Example

```
typedef int zip_dig[5];

zip_dig cmu = { 1, 5, 2, 1, 3 };
zip_dig mit = { 0, 2, 1, 3, 9 };
zip_dig ucb = { 9, 4, 7, 2, 0 };
```



Notes

- Declaration "zip_dig cmu" equivalent to "int cmu[5]"
- Example arrays were allocated in successive 20 byte blocks
–Not guaranteed to happen in general

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-5-

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Array Accessing Example

Computation

- Register %edx contains starting address of array
- Register %eax contains array index
- Desired digit at $4 * \text{eax} + \text{edx}$
- Use memory reference (`%edx, %eax, 4`)

```
int get_digit
(zip_dig z, int dig)
{
    return z[dig];
}
```

Memory Reference Code

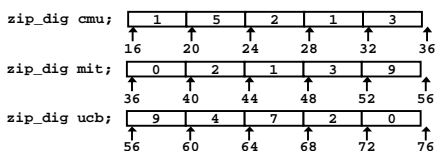
```
# %edx = z
# %eax = dig
movl (%edx,%eax,4),%eax # z[dig]
```

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-6-

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Referencing Examples



Code Does Not Do Any Bounds Checking!

Reference	Address	Value	Guaranteed?
mit[3]	$36 + 4 * 3 = 48$	3	Yes
mit[5]	$36 + 4 * 5 = 56$	9	No
mit[-1]	$36 + 4 * -1 = 32$	3	No
cmu[15]	$16 + 4 * 15 = 76$??	No

- Out of range behavior implementation-dependent
–No guaranteed relative allocation of different arrays

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-7-

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Array Loop Example

Original Source

```
int zd2int(zip_dig z)
{
    int i;
    int zi = 0;
    for (i = 0; i < 5; i++) {
        zi = 10 * zi + z[i];
    }
    return zi;
}
```

Transformed Version

- Eliminate loop variable `i`
- Convert array code to pointer code
- Express in do-while form
–No need to test at entrance

```
int zd2int(zip_dig z)
{
    int zi = 0;
    int *zend = z + 4;
    do {
        zi = 10 * zi + *z;
        z++;
    } while(z <= zend);
    return zi;
}
```

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-8-

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Array Loop Implementation

Registers

```
%ecx z
%eax zi
%ebx zend
```

Computations

- $10 * zi + *z$ implemented as $*z + 2 * (zi + 4 * zi)$
- $z++$ increments by 4

```
int zd2int(zip_dig z)
{
    int zi = 0;
    int *zend = z + 4;
    do {
        zi = 10 * zi + *z;
        z++;
    } while(z <= zend);
    return zi;
}
```

```
# %ecx = z
xorl %eax,%eax # zi = 0
leal 16(%ecx),%ebx # zend = z+4
.L59:
leal (%eax,%eax,4),%edx # 5*zi
movl (%ecx),%eax # *z
addl $4,%ecx # z++
leal (%eax,%edx,2),%eax # zi = *z + 2*(5*zi)
cmpl %ebx,%ecx # z : zend
jle .L59 # if <= goto loop
```

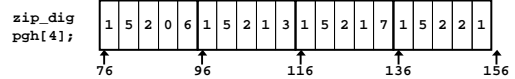
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- 9 -

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Nested Array Example

```
#define PCOUNT 4
zip_dig pgh[PCOUNT] =
{{1, 5, 2, 0, 6},
 {1, 5, 2, 1, 3},
 {1, 5, 2, 1, 7},
 {1, 5, 2, 2, 1}};
```



- Declaration "`zip_dig pgh[4]`" equivalent to "`int pgh[4][5]`"
 - Variable `pgh` denotes array of 4 elements
 - » Allocated contiguously
 - Each element is an array of 5 int's
 - » Allocated contiguously
- "Row-Major" ordering of all elements guaranteed

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- 10 -

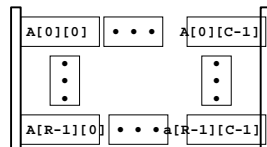
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Nested Array Allocation

Declaration

```
T A[R][C];
```

- Array of data type T
- R rows
- C columns
- Type T element requires K bytes



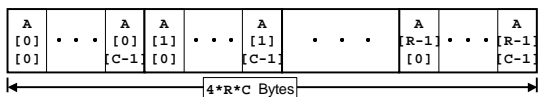
Array Size

- $R * C * K$ bytes

Arrangement

- Row-Major Ordering

```
int A[R][C];
```



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- 11 -

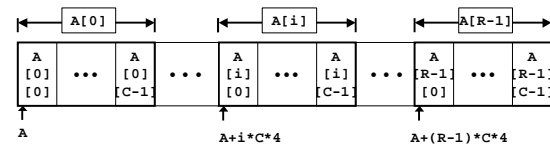
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Nested Array Row Access

Row Vectors

- $A[i]$ is array of C elements
- Each element of type T
- Starting address $A + i * C * K$

```
int A[R][C];
```



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- 12 -

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Nested Array Row Access Code

```
int *get_pgh_zip(int index)
{
    return pgh[index];
}
```

Row Vector

- `pgh[index]` is array of 5 int's
- Starting address `pgh+20*index`

Code

- Computes and returns address
- Compute as `pgh + 4*(index+4*index)`

```
# %eax = index
leal (%eax,%eax,4),%eax # 5 * index
leal pgh(,%eax,4),%eax # pgh + (20 * index)
```

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- 13 -

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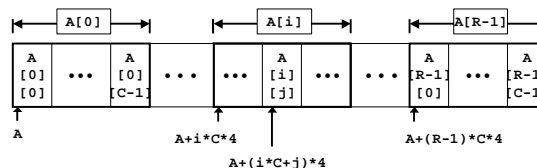
Nested Array Element Access

Array Elements

- `A[i][j]` is element of type `T`
- Address `A + (i * C + j) * K`

```
A
[i]
[j]
```

```
int A[R][C];
```



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- 14 -

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Nested Array Element Access Code

Array Elements

- `pgh[index][dig]` is int
- Address: `pgh + 20*index + 4*dig`

```
int get_pgh_digit
(int index, int dig)
{
    return pgh[index][dig];
}
```

Code

- Computes address `pgh + 4*dig + 4*(index+4*index)`
- `movl` performs memory reference

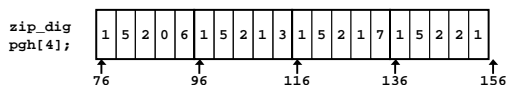
```
# %ecx = dig
# %eax = index
leal 0(,%ecx,4),%edx # 4*dig
leal (%eax,%eax,4),%eax # 5*index
movl pgh(%edx,%eax,4),%eax # *(pgh + 4*dig + 20*index)
```

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- 15 -

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Strange Referencing Examples



Reference	Address	Value	Guaranteed?
<code>pgh[3][3]</code>	$76+20*3+4*3 = 148$	2	Yes
<code>pgh[2][5]</code>	$76+20*2+4*5 = 136$	1	Yes
<code>pgh[2][-1]</code>	$76+20*2+4*-1 = 112$	3	Yes
<code>pgh[4][-1]</code>	$76+20*4+4*-1 = 152$	1	Yes
<code>pgh[0][19]</code>	$76+20*0+4*19 = 152$	1	Yes
<code>pgh[0][-1]</code>	$76+20*0+4*-1 = 72$??	No

- Code does not do any bounds checking
- Ordering of elements within array guaranteed

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- 16 -

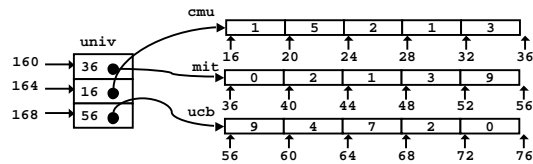
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Multi-Level Array Example

- Variable `univ` denotes array of 3 elements
- Each element is a pointer
-4 bytes
- Each pointer points to array of int's

```
zip_dig cmu = { 1, 5, 2, 1, 3 };
zip_dig mit = { 0, 2, 1, 3, 9 };
zip_dig ucb = { 9, 4, 7, 2, 0 };

#define UCOUNT 3
int *univ[UCOUNT] = {mit, cmu, ucb};
```



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- 17 -

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Referencing "Row" in Multi-Level Array

Row Vector

- `univ[index]` is pointer to array of int's
- Starting address `Mem[univ+4*index]`

```
int* get_univ_zip(int index)
{
    return univ[index];
}
```

Code

- Computes address within `univ`
- Reads pointer from memory and returns it

```
# %edx = index
leal 0(, %edx, 4), %eax # 4*index
movl univ(%eax), %eax # *(univ+4*index)
```

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- 18 -

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Accessing Element in Multi-Level Array

Computation

- Element access
`Mem[Mem[univ+4*index]+4*dig]`
- Must do two memory reads
 - First get pointer to row array
 - Then access element within array

```
int get_univ_digit
(int index, int dig)
{
    return univ[index][dig];
}
```

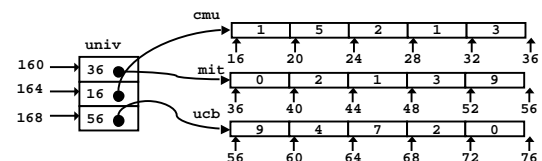
```
# %ecx = index
# %eax = dig
leal 0(, %ecx, 4), %edx # 4*index
movl univ(%edx), %edx # Mem[univ+4*index]
movl (%edx, %eax, 4), %eax # Mem[univ+4*index+4*dig]
```

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- 19 -

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Strange Referencing Examples



Reference	Address	Value	Guaranteed?
<code>univ[2][3]</code>	$56+4*3 = 68$	2	Yes
<code>univ[1][5]</code>	$16+4*5 = 36$	0	No
<code>univ[2][-1]</code>	$56+4*-1 = 52$	9	No
<code>univ[3][-1]</code>	??	??	No
<code>univ[1][12]</code>	$16+4*12 = 64$	7	No

- Code does not do any bounds checking
- Ordering of elements in different arrays not guaranteed

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- 20 -

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Using Nested Arrays

Strengths

- C compiler handles doubly subscripted arrays
- Generates very efficient code
 - Avoids multiply in index computation

Limitation

- Only works if have fixed array size

```
#define N 16
typedef int fix_matrix[N][N];

/* Compute element i,k of
fixed matrix product */
int fix_prod_ele
(fix_matrix a, fix_matrix b,
int i, int k)
{
    int j;
    int result = 0;
    for (j = 0; j < N; j++)
        result += a[i][j]*b[j][k];
    return result;
}
```

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- 21 -
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Dynamic Nested Arrays

Strength

- Can create matrix of arbitrary size

Programming

- Must do index computation explicitly

Performance

- Accessing single element costly
- Must do multiplication

```
int * new_var_matrix(int n)
{
    return (int *)
        calloc(sizeof(int), n*n);
}

int var_ele
(int *a, int i,
int j, int n)
{
    return a[i*n+j];
}

movl 12(%ebp),%eax # i
movl 8(%ebp),%edx # a
imull 20(%ebp),%eax # n*i
addl 16(%ebp),%eax # n*i+j
movl (%edx,%eax,4),%eax # Mem[a+4*(i*n+j)]
```

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- 22 -
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Dynamic Array Multiplication

Without Optimizations

- Multiplies**
 - 2 for subscripts
 - 1 for data
- Adds**
 - 4 for array indexing
 - 1 for loop index
 - 1 for data

```
/* Compute element i,k of
variable matrix product */
int var_prod_ele
(int *a, int *b,
int i, int k, int n)
{
    int j;
    int result = 0;
    for (j = 0; j < n; j++)
        result +=
            a[i*n+j] * b[j*n+k];
    return result;
}
```

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- 23 -
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Optimizing Dynamic Array Multiplication

Optimizations

- Performed when set optimization level to -O2

Code Motion

- Expression $i*n$ can be computed outside loop

Strength Reduction

- Incrementing j has effect of incrementing $j*n+k$ by n

Performance

- Compiler can optimize regular access patterns

```
{
    int j;
    int result = 0;
    for (j = 0; j < n; j++)
        result +=
            a[i*n+j] * b[j*n+k];
    return result;
}

{
    int j;
    int result = 0;
    int iTn = i*n;
    int jTnPk = k;
    for (j = 0; j < n; j++) {
        result +=
            a[iTn+j] * b[jTnPk];
        jTnPk += n;
    }
    return result;
}
```

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- 24 -
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```

{
int j;
int result = 0;
int iTn = i*n;
int jTnPk = k;
for (j = 0; j < n; j++) {
    result += a[iTn+j] * b[jTnPk];
    jTnPk += n;
}
return result;
}

```

Dynamic Array Multiplication

%ecx	result
%edx	j
%esi	n
%ebx	jTnPk
Mem[-4(%ebp)]	iTn

```

.L44:                # loop
movl -4(%ebp),%eax  # iTn
movl 8(%ebp),%edi   # a
addl %edx,%eax     # iTn+j
movl (%edi,%eax,4),%eax # a[iTn+j]
movl 12(%ebp),%edi # b
incl %edx          # j++
imull (%edi,%ebx,4),%eax # b[jTnPk]*a[iTn+j]
addl %eax,%ecx    # result += ..
addl %esi,%ebx   # jTnPk += n
cmpl %esi,%edx  # j : n
jle .L44        # if < goto loop

```

Inner Loop

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Summary

Arrays in C

- Contiguous allocation of memory
- Pointer to first element
- No bounds checking

Compiler Optimizations

- Compiler often turns array code into pointer code
zdzint
- Uses addressing modes to scale array indices
– code motion
- Lots of tricks to improve array indexing in loops
– reduction in strength

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- 26 -
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