## **Recitation 14: Proxy Lab Part 2**

Instructor: TA(s)

### **Outline**

- Proxylab
- Makefiles
- Threading
- Threads and Synchronization

#### **Makefiles**

- Make a separate file for your cache!
  - Need to update the makefile
  - Push your new makefile to GitHub!
- Makefile: tells program how to compile and link files

```
# List of all header files (for fake cache.c file)
DEPS = csapp.h transpose.h

# Rules for building cache
cache: cache.o transpose.o csapp.o

transpose.o: transpose.c $(DEPS)

cache.o: cache.c $(DEPS)
csapp.o: csapp.c csapp.h
```

## **ProxyLab**

#### ProxyLab is due next Thursday. Checkpoint is due tomorrow.

- One grace day for each
- Make sure to submit well in advance of the deadline in case there are errors in your submission.
- Build errors are a common source of failure

#### A proxy is a server process

- It is expected to be long-lived
- To not leak resources
- To be robust against user input

#### Note on CSAPP

- Most CSAPP functions have been removed
- Error check all system calls and exit only on critical failure

#### **Proxies and Threads**

#### Network connections can be handled concurrently

- Three approaches were discussed in lecture for doing so
- Your proxy should (eventually) use threads
- Threaded echo server is a good example of how to do this

#### Multi-threaded cache design

- Be careful how you use mutexes. Do not hold locks over network / file operations (read, write, etc)
- Using semaphores is not permitted
- Be careful how you maintain your object age

#### Tools

 Use Firefox's Network Monitor (Developer > Network) to see if all requests have been fulfilled

# Join / Detach

Does the following code terminate? Why or why not?

```
int main(int argc, char** argv)
{
    pthread create(&tid, NULL, work, NULL);
    if (pthread join(tid, NULL) != 0) printf("Done.\n");
void* work(void* a)
    pthread detatch(pthread self());
    while(1);
```

# Join / Detach cont.

Does the following code terminate now? Why or why not?

```
int main(int argc, char** argv)
    pthread create(&tid, NULL, work, NULL); sleep(1);
    if (pthread_join(tid, NULL) != 0) printf("Done.\n");
void* work(void* a)
    pthread detach(pthread self());
   while(1);
```

#### When should threads detach?

- In general, pthreads will wait to be reaped via pthread\_join.
- When should this behavior be overridden?
- When termination status does not matter.
  - pthread\_join provides a return value
- When result of thread is not needed.
  - When other threads do not depend on this thread having completed

#### **Threads**

- What is the range of value(s) that main will print?
- A programmer proposes removing j from thread and just directly accessing count. Does the answer change?

## **Synchronization**

- Is not cheap
  - 100s of cycles just to acquire without waiting
- Is also not that expensive
  - Recall your malloc target of 15000kops => ~100 cycles
- May be necessary
  - Correctness is always more important than performance

## Which synchronization should I use?

- Counting a shared resource, such as shared buffers
  - Semaphore
- Exclusive access to one or more variables
  - Mutex
- Most operations are reading, rarely writing / modifying
  - RWLock

For proxy it's sufficient to just use mutexes!

### **Threads Revisited**

- Which lock type should be used?
- Where should it be acquired / released?

## **Associating locks with data**

- Given the following key-value store
  - Key and value have separate RWLocks: klock and vlock
  - When an entry is replaced, both locks are acquired.
- Describe why the printf may not be accurate.

```
typedef struct _data_t {
  int key;
  size_t value;
} data_t;

#define SIZE 10
data_t space[SIZE];
int search(int k)
{
  for(int j = 0; j < SIZE; j++)
    if (space[j].key == k) return j;
  return -1;
}</pre>
```

```
pthread_rwlock_rdlock(klock);
match = search(k);
pthread_rwlock_unlock(klock);

if (match != -1)
{
   pthread_rwlock_rdlock(vlock);
   printf("%zd\n", space[match]);
   pthread_rwlock_unlock(vlock);
}
```

## Locks gone wrong

- 1. RWLocks are particularly susceptible to which issue:
- a. Starvation b. Livelock c. Deadlock
- If some code acquires rwlocks as readers: LockA then LockB, while other readers go LockB then LockA. What, if any, order can a writer acquire both LockA and LockB? No order is possible without a potential deadlock.
- Design an approach to acquiring two semaphores that avoids deadlock and livelock, while allowing progress to other threads needing only one semaphore.

### **Client-to-Client Communication**

- Clients don't have to fetch content from servers
  - Clients can communicate with each other
  - In a chat system, a server acts as a facilitator between clients
  - Clients could also send messages directly to each other, but this is more complicated (peer-to-peer networking)
- Running the chat server
  - ./chatserver <port>
- Running the client
  - telnet <hostname> <port>
- What race conditions could arise from having communication between multiple clients?

## **Proxylab Reminders**

- Plan out your implementation
  - "Weeks of programming can save you hours of planning"
    - Anonymous
  - Arbitrarily using mutexes will not fix race conditions
- Read the writeup
- Submit your code (days) early
  - Test that the submission will build and run on Autolab
- Final exam is only a few weeks away!

## **Appendix**

- Calling exit() will terminate all threads
- Calling pthread\_join on a detached thread is technically undefined behavior. Was defined as returning an error.