

Recitation 9: Tshlab + VM

Instructor: TAs

29 October 2018

Outline

- Labs
- Signals
- IO
- Virtual Memory

tshlab and malloclab

- **tshlab due Tuesday**
- **malloclab is released immediately after**
 - Start early
 - Do the checkpoint first, don't immediately go for the final
 - Expect a recitation next week
 - Working for several hours will improve the value significantly

Signals

- **Parent process sends SIGINT to a child process.
What is the behavior of the child?**
- **What is the default?**
- **What else could the child do?**

More Signals

- **Parent process sends SIGKILL to a child process.
What is the behavior of the child?**
- **What is the default?**
- **What else could the child do?**

Sending Signals

- Parent sends SIGKILL to a child process.

...

```
pid_t pid = ...; // child pid
kill(pid, SIGKILL);
// At this point, what could have
// happened to the child process?
```

Blocking Signals

- The shell is currently running its handler for SIGCHLD.
- What signals can it receive?
- What signals can it not receive (i.e., blocked)?

Errno

- Included from `<errno.h>`
- Global integer variable – usually 0
- When a system call fails (usually indicated by returning -1), it also will set `errno` to a value describing what went wrong

- Example: let's assume there is no "foo.txt" in our path

```
int fd = open("foo.txt", O_RDONLY);  
if (fd < 0) printf("%d\n", errno);
```

- The code above will print 2 – in the man pages, we can see that 2 is ENOENT "No such file or directory"
- In shell lab, your signal handlers must preserve `errno`

IO functions

Needed for tshlab

- `int open(const char *pathname, int flags);`
 - Some important flags:
 - `O_CREAT` – creates file if needed, opens for read/write
 - `O_RDWR` – opens for read/write
 - `O_RDONLY` – opens for read only
 - Various permission modes
- `int close(int fd);`
- `int dup2(int oldfd, int newfd);`

Needed for life

- `ssize_t read(int fd, void *buf, size_t count);`
- `ssize_t write(int fd, const void *buf, size_t count);`
- `off_t lseek(int fd, off_t offset, int whence);`

More on open

- `int open(const char *pathname, int flags, mode_t mode);`
- For *flags*, you can pass a bitwise-OR of one or more flags
- **Three kinds of flags (we only discuss the important ones)**
 - Access modes (one of them must be included):
 - `O_RDONLY`, `O_WRONLY`, `O_RDWR`
 - File creation flags:
 - `O_CREAT`, `O_TRUNC`, etc.
 - File status flags

Access mode flags and file creation flags

■ O_RDONLY

- Open the file read-only.

■ O_WRONLY

- Open the file write-only.

■ O_RDWR

- Open the file read/write.

■ O_CREAT

- If the provided *pathname* does not exist, create it as a regular file.

■ O_TRUNC

- If the file already exists and if the access mode allows writing (i.e. is **O_RDWR** or **O_WRONLY**), then the file will be truncated to length 0.

More on open

- `int open(const char *pathname,
 int flags, mode_t mode);`
- For *mode*, you can pass a bitwise-OR of one or more constants
- Specifies, when creating a file, what permission the file will be created with
- Only useful when *flags* contain `O_CREAT` (or `O_TMPFILE`)

Linux permissions

- **Every file and directory has permission information**
- **You've seen it before**
 - `ls -l` prints the permissions for each file/directory like:
-rw-r--r-- ... drwxr-xr-x ...
 - `chmod` changes the permissions for files/directories
 - `$ chmod -R 777 /`
- **There are read (R), write (W) and executable (X) permissions for user (USR), group (GRP) and other (OTH)**

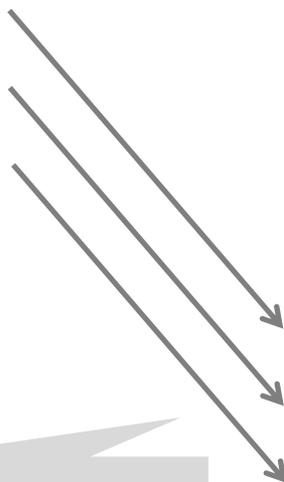
Specify permissions in open()

	Read (R)	Write (W)	Executable (X)	All (RWX)
User (USR)	S_IRUSR	S_IWUSR	S_IXUSR	S_IRWXU
Group (GRP)	S_IRGRP	S_IWGRP	S_IXGRP	S_IRWXG
Other (OTH)	S_IROTH	S_IWOTH	S_IXOTH	S_IRWXO

- These constants can be bitwise-OR'd and passed to the third argument of open()
- What does `S_IRWXG | S_IXUSR | S_IXOTH` mean?
- How to create a file which everyone can read from but only the user can write to it or execute it?

File descriptors

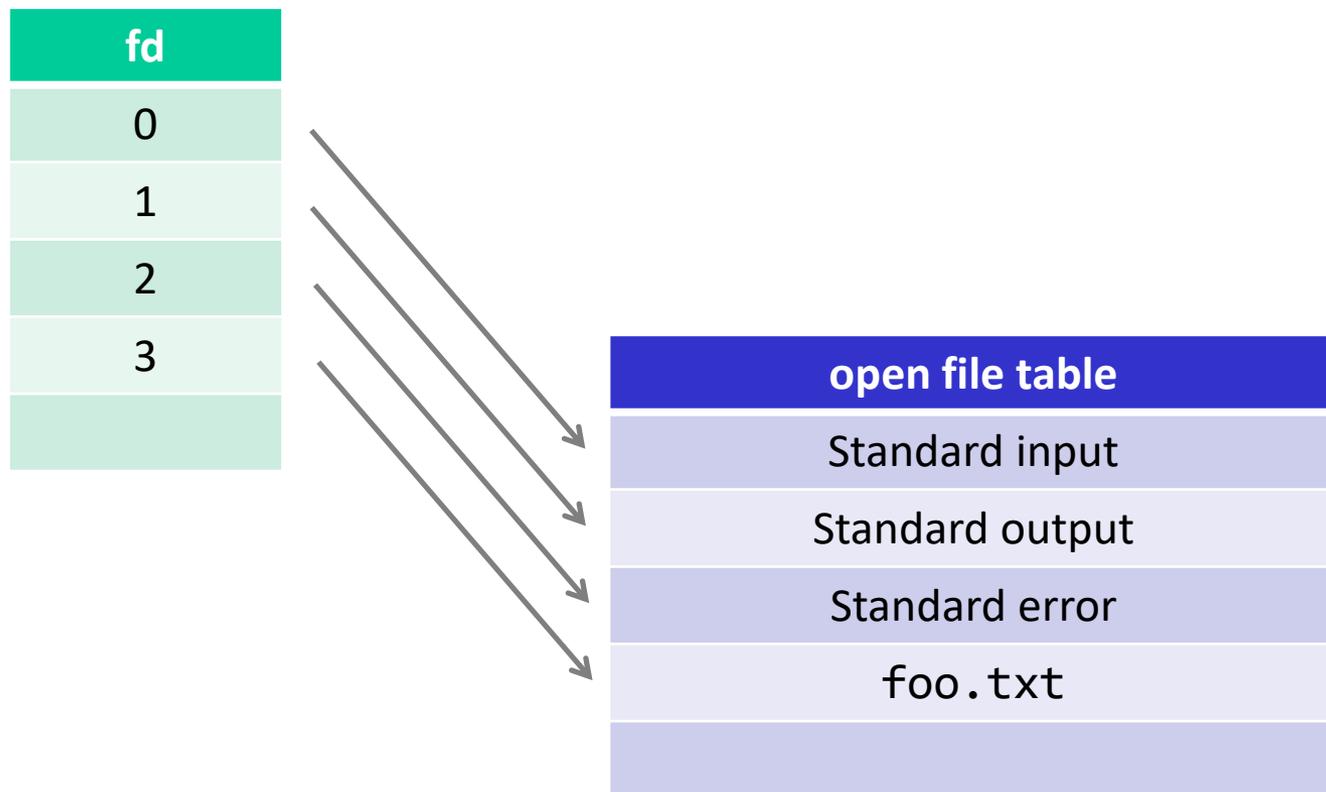
fd
0
1
2



open file table
Standard input
Standard output
Standard error

stdin, stdout, stderr are opened automatically and closed by normal termination or exit()

open("foo.txt")



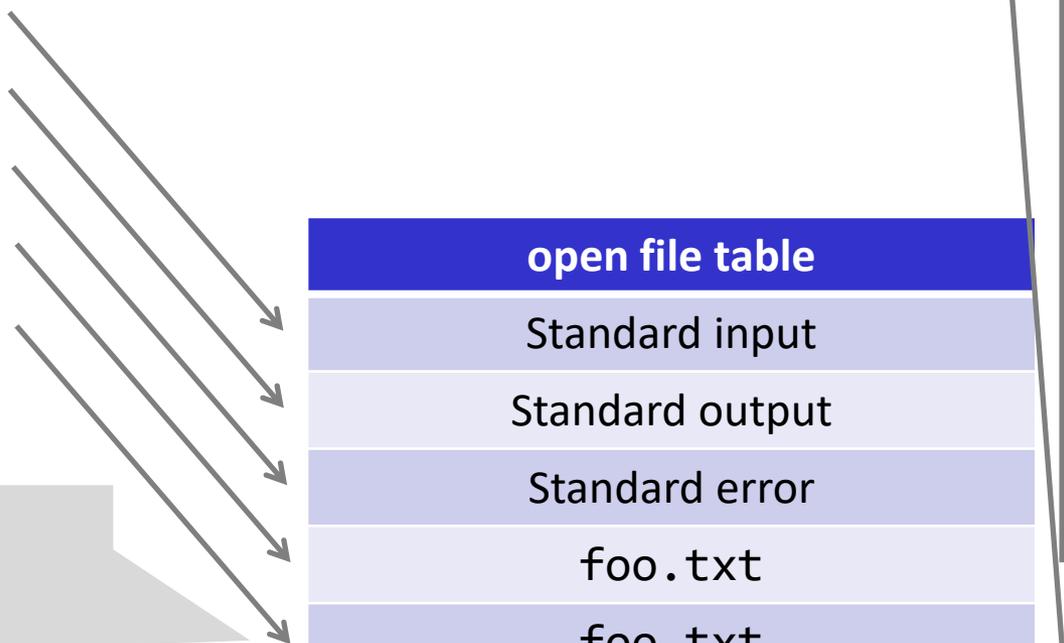
open("foo.txt")

fd
0
1
2
3
4

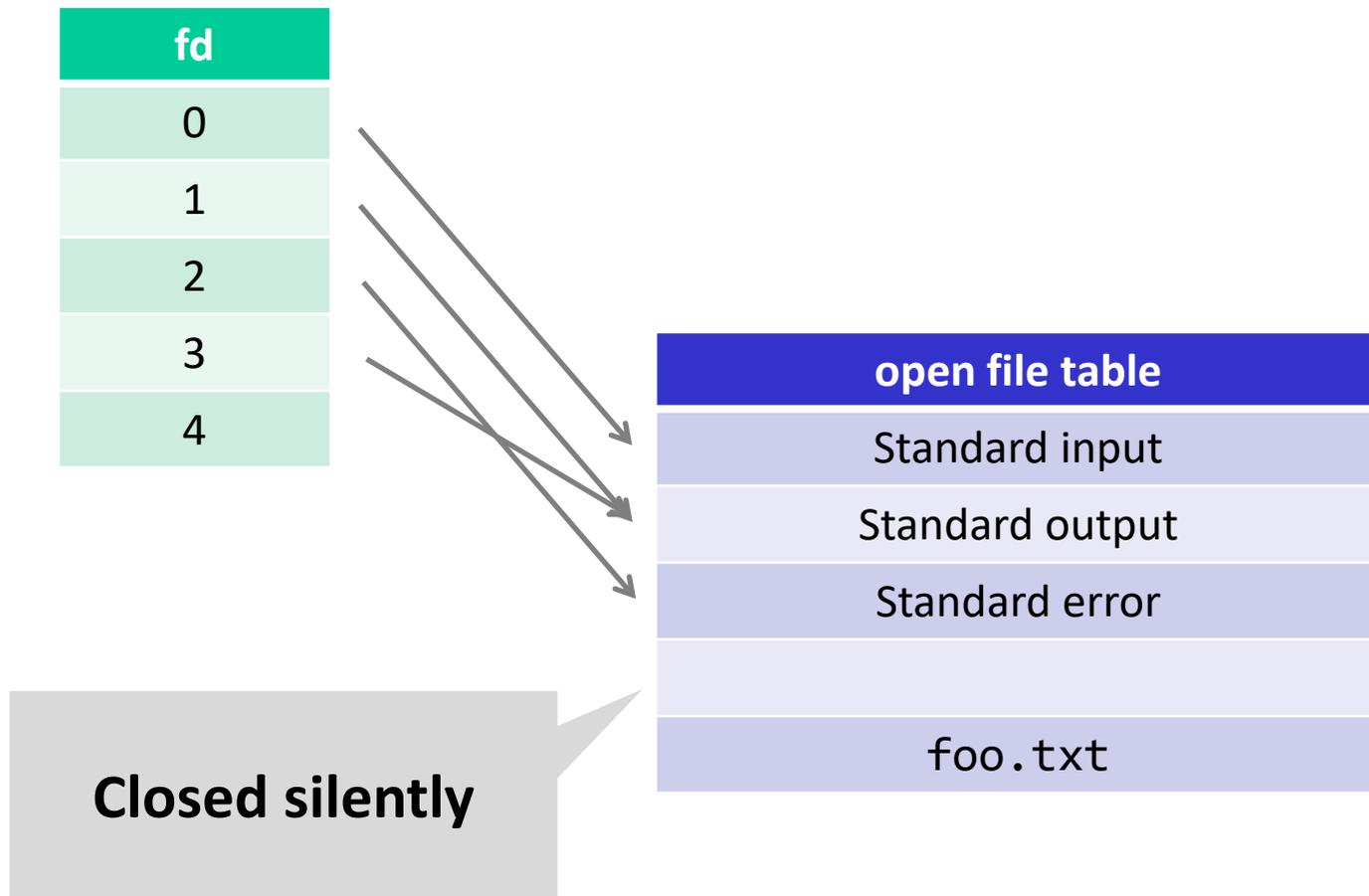
inode table
foo.txt

open file table
Standard input
Standard output
Standard error
foo.txt
foo.txt

Each call to `open()`
creates a new open file
description



dup2(STDOUT_FILENO, 3)



IO and Fork()

- File descriptor management can be tricky.
- How many file descriptors are open in the parent process at the indicated point?
- How many does each child have open at the call to `execve`?

```
int main(int argc, char** argv)
{
    int i;
    for (i = 0; i < 4; i++)
    {
        int fd = open("foo", O_RDONLY);
        pid_t pid = fork();
        if (pid == 0)
        {
            int ofd = open("bar", O_RDONLY);
            execve(...);
        }
    }
    // How many file descriptors are open in the parent?
```

Redirecting IO

- File descriptors can be directed to identify different open files.

```
int main(int argc, char** argv) {
    int i;
    for (i = 0; i < 4; i++)
    {
        int fd = open("foo", O_RDONLY);
        pid_t pid = fork();
        if (pid == 0)
        {
            int ofd = open("bar", O_WRONLY);
            dup2(fd, STDIN_FILENO);
            dup2(ofd, STDOUT_FILENO);
            execve(...);
        }
    }
    // How many file descriptors are open in the parent?
}
```

Redirecting IO

- At the two points (A and B) in main, how many file descriptors are open?

```
int main(int argc, char** argv)
{
    int i, fd;
    fd = open("foo", O_WRONLY);
    dup2(fd, STDOUT_FILENO);
    // Point A
    close(fd);
    // Point B
    ...
}
```

Memory Access

- **The processor tries to write to a memory address.**
- **List different steps that are required to complete this operation.**

Memory Access

- The processor tries to write to a memory address.
- List different steps that are required to complete this operation. (non exhaustive list)
- Virtual to physical address conversion (TLB lookup)
- TLB miss
- Page fault, page loaded from disk
- TLB updated, check permissions
- L1 Cache miss (and L2 ... and)
- Request sent to memory
- Memory sends data to processor
- Cache updated

Address Translation with TLB

- Translate 0x15213, given the contents of the TLB and the first 32 entries of the page table below.
- 1MB Virtual Memory
256KB Physical Memory
4KB page size

2-way
set
associative

Index	Tag	PP N	Valid
0	05	13	1
	3F	15	1
1	10	0F	1
	0F	1E	0
2	1F	01	1
	11	1F	0
3	03	2B	1
	1D	23	0

VPN	PPN	Valid	VPN	PPN	Valid
00	17	1	10	26	0
01	28	1	11	17	0
02	14	1	12	0E	1
03	0B	0	13	10	1
04	26	0	14	13	1
05	13	0	15	18	1
06	0F	1	16	31	1
07	10	1	17	12	0
08	1C	0	18	23	1
09	25	1	19	04	0
0A	31	0	1A	0C	1
0B	16	1	1B	2B	0
0C	01	0	1C	1E	0
0D	15	0	1D	3E	1
0E	0C	0	1E	27	1
0F	2B	1	1F	15	1

If you get stuck on tshlab

- **Read the writeup!**
- **Do manual unit testing before `runtrace` and `sdriver`!**
- **Post private questions on piazza!**

- **Read the man pages on the syscalls.**
 - Especially the error conditions
 - What errors should terminate the shell?
 - What errors should be reported?

man wait

Taken from <http://man7.org/linux/man-pages/man2/wait.2.html>

WAIT(2)

Linux Programmer's Manual

WAIT(2)

NAME

wait, waitpid, waitid - wait for process to change state

SYNOPSIS

```
#include <sys/types.h>
#include <sys/wait.h>

pid_t wait(int *wstatus);

pid_t waitpid(pid_t pid, int *wstatus, int options);

int waitid(idtype_t idtype, id_t id, siginfo_t *infop, int options);
/* This is the glibc and POSIX interface; see
   NOTES for information on the raw system call. */
```

man pages (probably) cover all you need

- **What arguments does the function take?**
 - read SYNOPSIS
- **What does the function do?**
 - read DESCRIPTION
- **What does the function return?**
 - read RETURN VALUE
- **What errors can the function fail with?**
 - read ERRORS
- **Is there anything I should watch out for?**
 - read NOTES
- **Different categories for man page entries with the same name**
- **Looking up man pages online is not an academic integrity violation**

Function arguments

- Should I do `dup2(old, new)` or `dup2(new, old)`?
- Read the man page:

\$ man dup2

SYNOPSIS

```
#include <unistd.h>

int dup(int oldfd);
int dup2(int oldfd, int newfd);
```

Function behavior

- How should I write my format string when I need to print a long double in octals with precision 5 and zero-padded?
- Read the man page:

\$ man printf

DESCRIPTION

Flag characters

The character % is followed by zero or more of the following flags:

- # The value should be converted...
- 0 The value should be zero padded...
- The converted value is to be left adjusted...
- ' ' (a space) A blank should be left before...
- + A sign (+ or -) should always ...

Function return

- What does `waitpid()` return with and without `WNOHANG`?
- Read the man page:

\$ man waitpid

RETURN VALUE

`waitpid()`: on success, returns the process ID of the child whose state has changed; if `WNOHANG` was specified and one or more child(ren) specified by `pid` exist, but have not yet changed state, then `0` is returned. On error, `-1` is returned.

Each of these calls sets `errno` to an appropriate value in the case of an error.

Potential errors

- How should I check `waitpid` for errors?
- Read the man page:

```
$ man waitpid
```

ERRORS

ECHILD (for `waitpid()` or `waitid()`) The process specified by *pid* (`waitpid()`) or *idtype* and *id* (`waitid()`) does not exist or is not a child of the calling process. (This can happen for one's own child if the action for **SIGCHLD** is set to **SIG_IGN**. See also the Linux Notes section about threads.)

EINTR **WNOHANG** was not set and an unblocked signal or a **SIGCHLD** was caught; see `signal(7)`.

EINVAL The *options* argument was invalid.

Get advice from the developers

- I sprintf from a string into itself, is this okay?
- Read the man page:

\$ man sprintf

NOTES

Some programs imprudently rely on code such as the following

```
    sprintf(buf, "%s some further text", buf);
```

to append text to *buf*. However, the standards explicitly note that the results are undefined if source and destination buffers overlap when calling `sprintf()`, `snprintf()`, `vsprintf()`, and `vsnprintf()`. Depending on the version of `gcc(1)` used, and the compiler options employed, calls such as the above will **not** produce the expected results.

The `glibc` implementation of the functions `snprintf()` and `vsnprintf()` conforms to the C99 standard, that is, behaves as described above, since `glibc` version 2.1. Until `glibc` 2.0.6, they would return `-1` when the output was truncated.