

# Computational Learning Theory

Reading:

- Mitchell chapter 7

Suggested exercises:

- 7.1, 7.2, 7.5, 7.7

Machine Learning 10-601

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(originally, from slides by Prof. Tom Mitchell)

# Computational Learning Theory

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What general laws constrain inductive learning?

We seek theory to relate:

- Probability of successful learning
- Number of training examples
- Complexity of hypothesis space
- Accuracy to which target function is approximated
- Manner in which training examples presented

# Sample Complexity: What it means

[Haussler, 1988]: probability that the version space is not  $\epsilon$ -exhausted after  $m$  training examples is at most  $|H|e^{-\epsilon m}$

$$\Pr[(\exists h \in H) \text{ s.t. } (error_{train}(h) = 0) \wedge (error_{true}(h) > \epsilon)] \leq |H|e^{-\epsilon m}$$

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↑

Suppose we want this probability to be at most  $\delta$

1. How many training examples suffice?

$$m \geq \frac{1}{\epsilon}(\ln |H| + \ln(1/\delta))$$

2. If  $error_{train}(h) = 0$  then with probability at least  $(1-\delta)$ :

$$error_{true}(h) \leq \frac{1}{m}(\ln |H| + \ln(1/\delta))$$

# Agnostic Learning

**Result we proved:** probability, after  $m$  training examples, that  $H$  contains a hypothesis  $h$  with zero training error, but true error greater than  $\epsilon$  is bounded

$$\Pr[(\exists h \in H) s.t. (error_{train}(h) = 0) \wedge (error_{true}(h) > \epsilon)] \leq |H|e^{-\epsilon m}$$

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probabilistic argument

**Agnostic case:** don't know whether  $H$  contains a perfect hypothesis

$$\Pr[(\exists h \in H) s.t. (error_{true}(h) > \epsilon + error_{train}(h))] \leq |H|e^{-2\epsilon^2 m}$$

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Hoeffding bound

# General Hoeffding Bounds

- When estimating the mean  $\theta$  inside  $[a,b]$  from  $m$  examples

$$P(|\hat{\theta} - E[\hat{\theta}]| > \epsilon) \leq 2e^{\frac{-2m\epsilon^2}{(b-a)^2}}$$

- When estimating a probability  $\theta$  is inside  $[0,1]$ , so

$$P(|\hat{\theta} - E[\hat{\theta}]| > \epsilon) \leq 2e^{-2m\epsilon^2}$$

- And if we're interested in only one-sided error, then

$$P((E[\hat{\theta}] - \hat{\theta}) > \epsilon) \leq e^{-2m\epsilon^2}$$

# PAC Learning

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Consider a class  $C$  of possible target concepts defined over a set of instances  $X$  of length  $n$ , and a learner  $L$  using hypothesis space  $H$ .

*Definition:*  $C$  is **PAC-learnable** by  $L$  using  $H$  if for all  $c \in C$ , distributions  $\mathcal{D}$  over  $X$ ,  $\epsilon$  such that  $0 < \epsilon < 1/2$ , and  $\delta$  such that  $0 < \delta < 1/2$ ,

learner  $L$  will with probability at least  $(1 - \delta)$  output a hypothesis  $h \in H$  such that  $error_{\mathcal{D}}(h) \leq \epsilon$ , in time that is polynomial in  $1/\epsilon$ ,  $1/\delta$ ,  $n$  and  $size(c)$ .

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Sufficient condition:

Holds if  $L$  requires only a polynomial number of training examples, and processing per example is polynomial

## What if $H$ is not finite?

- Can't use our sample complexity results for infinite  $H$
- Need some other measure of complexity for  $H$ 
  - Vapnik-Chervonenkis (VC) dimension!



# Sample Complexity based on VC dimension

How many randomly drawn examples suffice to  $\epsilon$ -exhaust  $VS_{H,D}$  with probability at least  $(1-\delta)$ ?

ie., to guarantee that any hypothesis that perfectly fits the training data is probably  $(1-\delta)$  approximately  $(\epsilon)$  correct

$$m \geq \frac{1}{\epsilon} (4 \log_2(2/\delta) + 8VC(H) \log_2(13/\epsilon))$$

Compare to our earlier results based on  $|H|$ :

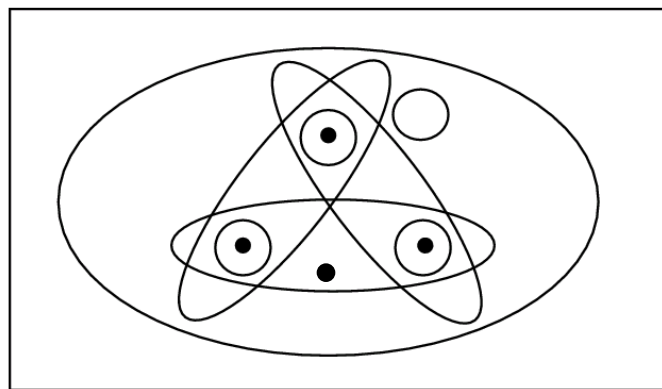
$$m \geq \frac{1}{\epsilon} (\ln(1/\delta) + \ln |H|)$$

# The Vapnik-Chervonenkis Dimension

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*Definition:* The **Vapnik-Chervonenkis dimension**,  $VC(H)$ , of hypothesis space  $H$  defined over instance space  $X$  is the size of the largest finite subset of  $X$  shattered by  $H$ . If arbitrarily large finite sets of  $X$  can be shattered by  $H$ , then  $VC(H) \equiv \infty$ .

Instance space  $X$



$$VC(H)=3$$

# VC dimension: examples

What is VC dimension of lines in a plane?

- $H_2 = \{ ((w_0 + w_1x_1 + w_2x_2) > 0 \rightarrow y=1) \}$



# VC dimension: examples

What is VC dimension of

- $H_2 = \{ ((w_0 + w_1x_1 + w_2x_2) > 0 \rightarrow y=1) \}$ 
  - $VC(H_2)=3$
- For  $H_n$  = linear separating hyperplanes in  $n$  dimensions,  
 $VC(H_n)=n+1$



Can you give an upper bound on  $VC(H)$  in terms of  $|H|$ , for any hypothesis space  $H$ ?  
(hint: yes)

# More VC Dimension Examples to Think About

- Logistic regression over  $n$  continuous features
  - Over  $n$  boolean features?
- Linear SVM over  $n$  continuous features
- Decision trees defined over  $n$  boolean features
$$F: \langle X_1, \dots, X_n \rangle \rightarrow Y$$
- Decision trees of depth 2 defined over  $n$  features
- How about 1-nearest neighbor?
- is there a hypothesis class with infinite VC dimension?

# Tightness of Bounds on Sample Complexity

How many examples  $m$  suffice to assure that any hypothesis that fits the training data perfectly is probably  $(1-\delta)$  approximately  $(\epsilon)$  correct?

$$m \geq \frac{1}{\epsilon} (4 \log_2(2/\delta) + 8VC(H) \log_2(13/\epsilon))$$

How tight is this bound?

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How tight is this bound?

**Lower bound on sample complexity** (Ehrenfeucht et al., 1989):

Consider any class  $C$  of concepts such that  $VC(C) > 1$ , any learner  $L$ , any  $0 < \epsilon < 1/8$ , and any  $0 < \delta < 0.01$ . Then there exists a distribution  $\mathcal{D}$  and a target concept in  $C$ , such that if  $L$  observes fewer examples than

$$\max \left[ \frac{1}{\epsilon} \log(1/\delta), \frac{VC(C) - 1}{32\epsilon} \right]$$

Then with probability at least  $\delta$ ,  $L$  outputs a hypothesis with  $error_{\mathcal{D}}(h) > \epsilon$

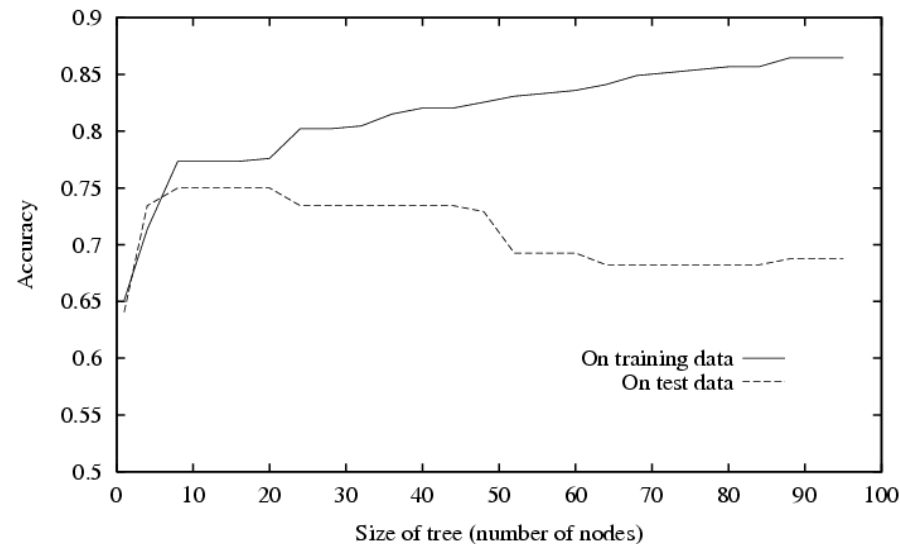


# Agnostic Learning: VC Bounds

[Schölkopf and Smola, 2002]

With probability at least  $(1-\delta)$  every  $h \in H$  satisfies

$$error_{true}(h) < error_{train}(h) + \sqrt{\frac{VC(H)(\ln \frac{2m}{VC(H)} + 1) + \ln \frac{4}{\delta}}{m}}$$

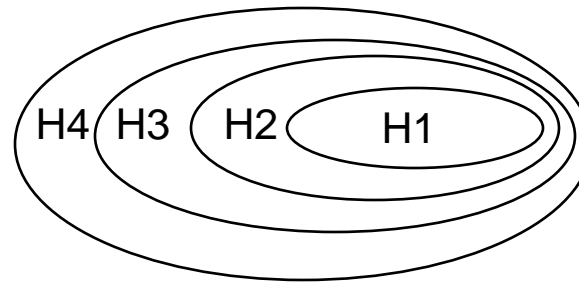


# Structural Risk Minimization

[Vapnik]

Which hypothesis space should we choose?

- Bias / variance tradeoff



SRM: choose  $H$  to minimize bound on true error!

$$error_{true}(h) < error_{train}(h) + \sqrt{\frac{VC(H)(\ln \frac{2m}{VC(H)} + 1) + \ln \frac{4}{\delta}}{m}}$$

\* unfortunately a somewhat loose bound...

# Mistake Bounds

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So far: how many examples needed to learn?

What about: how many mistakes before convergence?

Let's consider similar setting to PAC learning:

- Instances drawn at random from  $X$  according to distribution  $\mathcal{D}$
- Learner must classify each instance before receiving correct classification from teacher
- Can we bound the number of mistakes learner makes before converging?

# Mistake Bounds: Find-S

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Consider Find-S when  $H$  = conjunction of boolean literals

FIND-S:

- Initialize  $h$  to the most specific hypothesis  
 $l_1 \wedge \neg l_1 \wedge l_2 \wedge \neg l_2 \dots l_n \wedge \neg l_n$
- For each positive training instance  $x$ 
  - Remove from  $h$  any literal that is not satisfied by  $x$
- Output hypothesis  $h$ .

How many mistakes before converging to correct  $h$ ?

# Mistake Bounds: Halving Algorithm

Consider the Halving Algorithm:

- Learn concept using version space CANDIDATE-ELIMINATION algorithm
- Classify new instances by majority vote of version space members

1. Initialize  $VS \leftarrow H$
2. For each training example,
  - remove from VS every hypothesis that misclassifies this example

How many mistakes before converging to correct  $h$ ?

- ... in worst case?
- ... in best case?

# Optimal Mistake Bounds

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Let  $M_A(C)$  be the max number of mistakes made by algorithm  $A$  to learn concepts in  $C$ . (maximum over all possible  $c \in C$ , and all possible training sequences)

$$M_A(C) \equiv \max_{c \in C} M_A(c)$$

*Definition:* Let  $C$  be an arbitrary non-empty concept class. The **optimal mistake bound** for  $C$ , denoted  $Opt(C)$ , is the minimum over all possible learning algorithms  $A$  of  $M_A(C)$ .

$$Opt(C) \equiv \min_{A \in \text{learning algorithms}} M_A(C)$$

$$VC(C) \leq Opt(C) \leq M_{Halving}(C) \leq \log_2(|C|).$$

# Weighted Majority Algorithm

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$a_i$  denotes the  $i^{th}$  prediction algorithm in the pool  $A$  of algorithms.  $w_i$  denotes the weight associated with  $a_i$ .

- For all  $i$  initialize  $w_i \leftarrow 1$
- For each training example  $\langle x, c(x) \rangle$ 
  - \* Initialize  $q_0$  and  $q_1$  to 0
  - \* For each prediction algorithm  $a_i$ 
    - If  $a_i(x) = 0$  then  $q_0 \leftarrow q_0 + w_i$
    - If  $a_i(x) = 1$  then  $q_1 \leftarrow q_1 + w_i$
  - \* If  $q_1 > q_0$  then predict  $c(x) = 1$
  - If  $q_0 > q_1$  then predict  $c(x) = 0$
  - If  $q_1 = q_0$  then predict 0 or 1 at random for  $c(x)$
  - \* For each prediction algorithm  $a_i$  in  $A$  do
    - If  $a_i(x) \neq c(x)$  then  $w_i \leftarrow \beta w_i$

when  $\beta=0$ ,  
equivalent to  
the Halving  
algorithm...

# Weighted Majority

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Even algorithms  
that learn or  
change over time...

[Relative mistake bound for  
WEIGHTED-MAJORITY] Let  $D$  be any sequence of  
training examples, let  $A$  be any set of  $n$  prediction  
algorithms, and let  $k$  be the minimum number of  
mistakes made by any algorithm in  $A$  for the  
training sequence  $D$ . Then the number of mistakes  
over  $D$  made by the WEIGHTED-MAJORITY  
algorithm using  $\beta = \frac{1}{2}$  is at most

$$2.4(k + \log_2 n)$$



# What You Should Know

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- Sample complexity varies with the learning setting
  - Learner actively queries trainer
  - Examples provided at random
- Within the PAC learning setting, we can bound the probability that learner will output hypothesis with given error
  - For ANY consistent learner (case where  $c \in H$ )
  - For ANY “best fit” hypothesis (agnostic learning, where perhaps  $c$  not in  $H$ )
- VC dimension as measure of complexity of  $H$
- Mistake bounds
- Conference on Learning Theory: <http://www.learningtheory.org>