

Hashing 2 :

Feb 18 2021

"Two-level hashing"

First level : table size $O(N)$

At each location : hash table
(Collision-free)

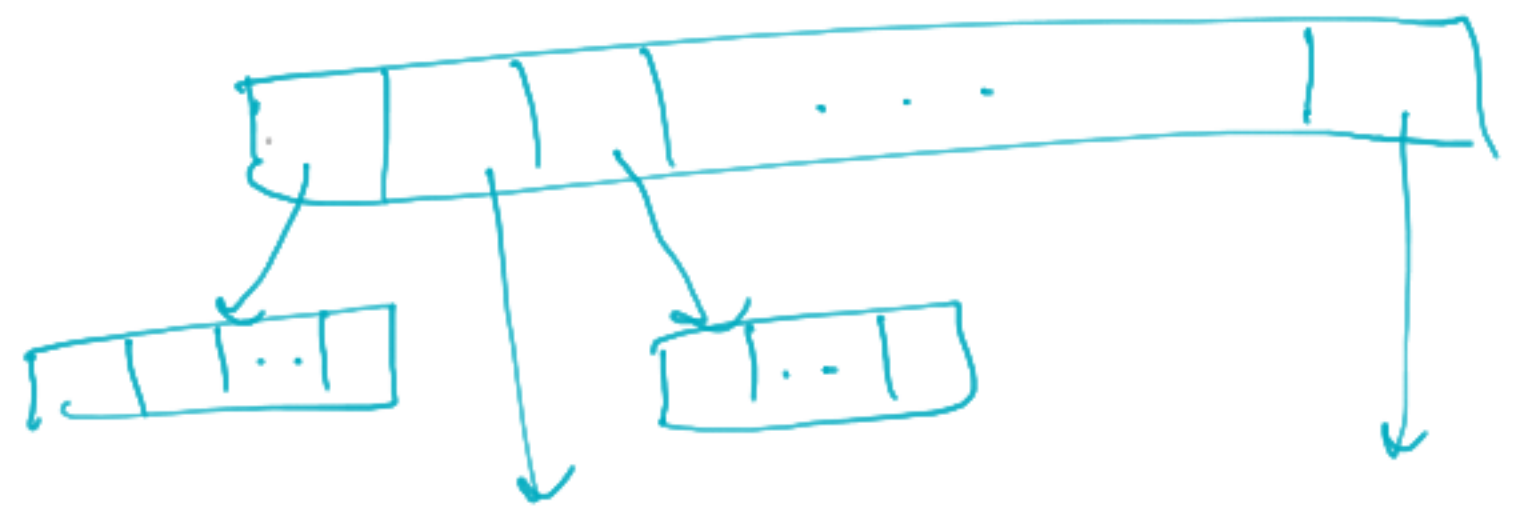
$M =$ table size

$N = |S|$

Let

$i =$ location index in the first level table.

$C(i) =$ num. of elements mapped to location i in the first level table



Q: What should be the size of table at i for collision free?
 $C(i)^2$

Total table size:

$$\sum_{i=1}^M C(i)^2$$

Recall

$$E[C] = \binom{N}{2} \cdot \frac{1}{M}$$

$$\Rightarrow E \left[\sum_{i=1}^M \binom{C(i)}{2} \right] = \binom{N}{2} \frac{1}{M}$$

$$E \left[\sum_{i=1}^M C(i)^2 - \sum_{i=1}^M C(i) \right] = O(N)$$

$$E \left[\sum_{i=1}^M C(i)^2 \right] = O(N)$$

$$\therefore E \left[\sum_{i=1}^M C(i) \right] = O(N)$$

C = total num. of collisions

$$C = \sum_{i=1}^M \binom{C(i)}{2}$$

since $M = O(N)$

Total table size = $O(N)$

Collision-free in $O(N)$ space

"Perfect hashing"

k-wise independent hash functions:

Defn: A family \mathcal{H} of hash functions $U \rightarrow [M]$ is

k-wise independent if for any k distinct keys x_1, \dots, x_k
and any k values d_1, \dots, d_k we have

$$P(h(x_1) = d_1 \wedge h(x_2) = d_2 \wedge \dots \wedge h(x_k) = d_k) = \frac{1}{M^k}$$

For $k=2$, "pairwise independent".

Suppose H is k -wise indep. ($k \geq 2$)

① $(k-1)$ wise indep

② for any $x \in U$ and value $a \in [M]$

$$P [h(x) = a] \leq \frac{1}{M}$$

③ Universal

Q: Which is stronger : pairwise indep. or universal?

E.g. Construction from last class A

$$h(x) = Ax$$

$$h(0) = 0$$

$$P [h(0) = 0] = 1$$

Construction 1: 2-wise independent

$$h(x) = Ax + b$$

$(m \times u)$ $(m \times 1)$

\rightarrow m -length random binary vector.

(modulo 2)

Claim: ... is 2-wise independent

Q: num. of hash fns? 2^{um+m}

Q: bits to store? $O(um)$

Can we do this with fewer bits?

Construction 2 (using fewer bits)

- fill first row & column with uniform random binary entries

- $A_{i,j} = A_{i-1,j-1}$

- $h(x) = Ax + b \pmod{2}$

Claim: ... 2-wise indep.

