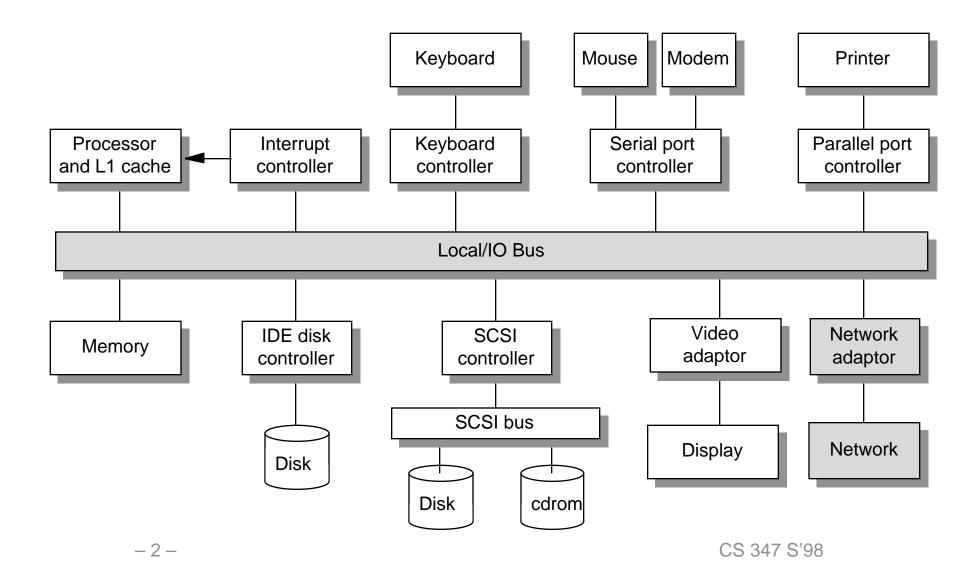
Networks

Randal E. Bryant CS347 Lecture 24 April 16, 1997

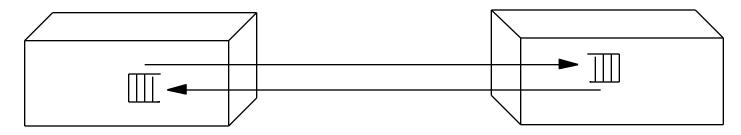
Topics

- Overview
- Telephone system
- Ethernet
- ATM

Computer system



Simple example



Starting Point: Want to send bits between 2 computers

- FIFO queue on each end
- Can send both ways ("full duplex")
- Rules for communication ("protocol")
- Name for standard group of bits sent: "packet"

Simple request/response protocol and packet format:

0/1 data/address	
------------------	--

- 0: please send the data word at "address"
- 1: here is the data word you asked for.

Questions about simple example

What if more than 2 computers want to communicate?

• Need computer address field in packet?

How do multiple machines share the interconnect?

• multiple paths? arbitration? congestion control?

What if a packet is garbled in transit?

• Add error detection field in packet?

What if a packet is lost?

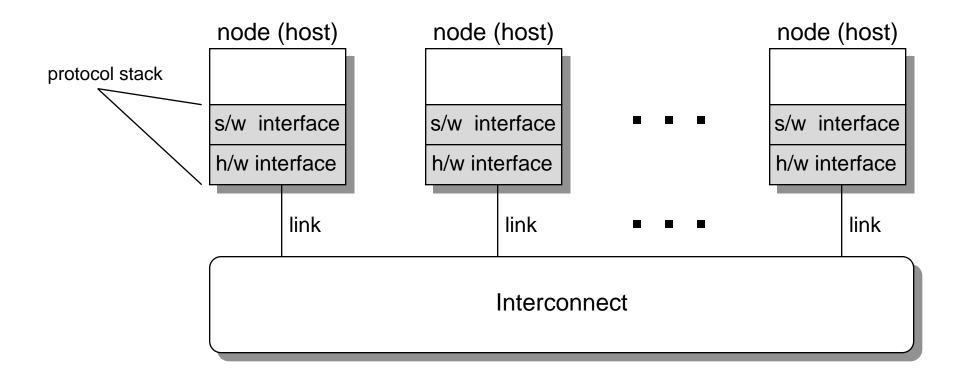
• More elaborate protocols to detect loss?

What if multiple processes/machine?

• one queue per process? separate field in packet to identify process?

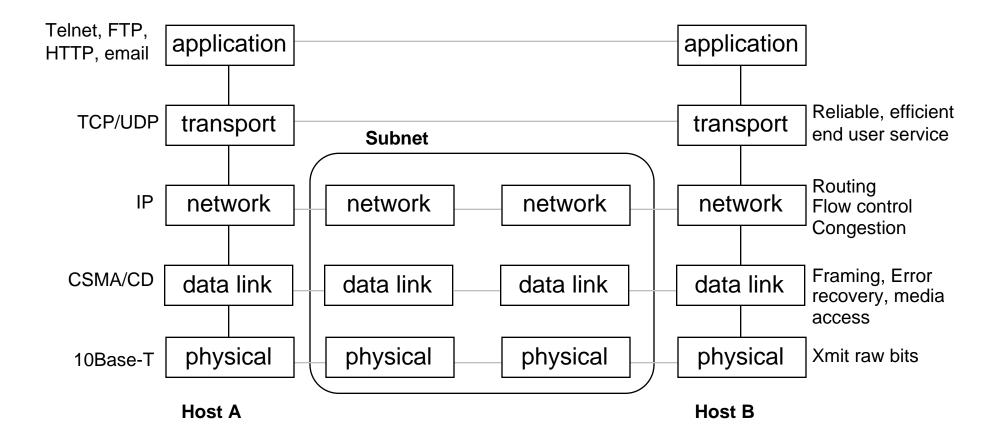
Questions such as these lead to more complex interconnects, protocols and packet formats

Generic network



Warning: You are entering a buzzword-rich environment!

Protocol stacks



Basic network types

System area network (SAN)

- same room (meters)
- 150 MB/s Cray T3D

Local area network (LAN)

- same bldg or campus (kilometers)
- 10 Mb/sEthernet
- 100 Mb/s Ethernet
- 100 Mb/s FDDI
- 150 Mb/s OC-3 ATM
- 622 Mb/s OC-12 ATM

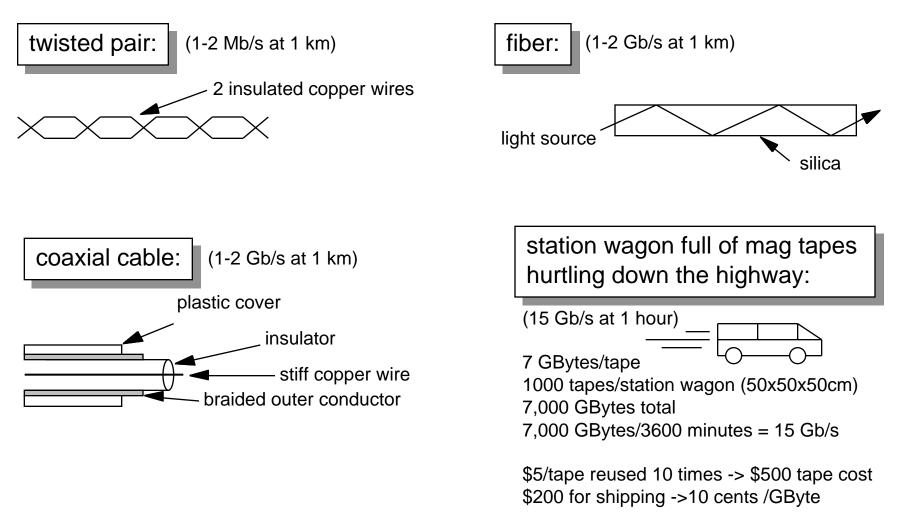
Metropolitan area network (MAN)

- same city (10's of kilometers)
- 800 Mb/s Gigabit Nectar

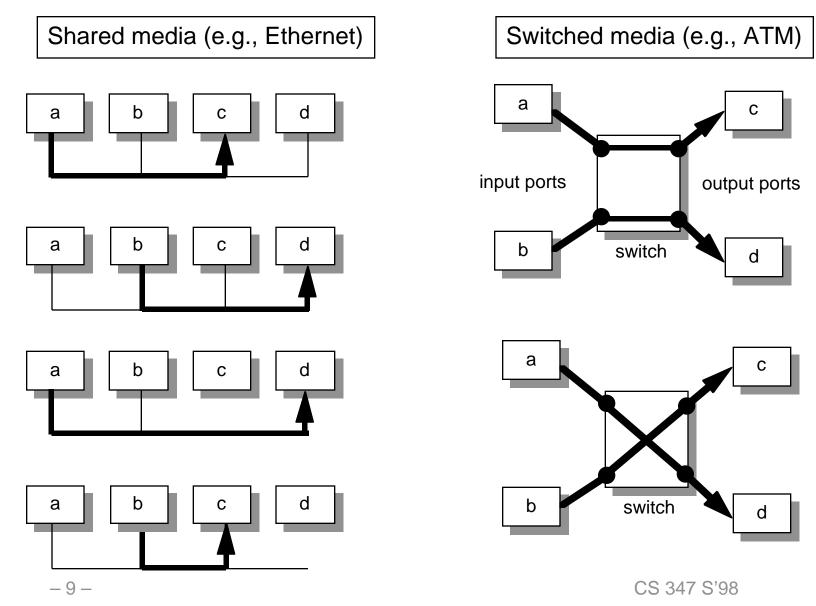
Wide area network (WAN)

- nationwide or worldwide (1000's of kilometers)
- telephone system
- 1.544 Mb/s T1 carrier
- 44.736 Mb/s T3 carrier

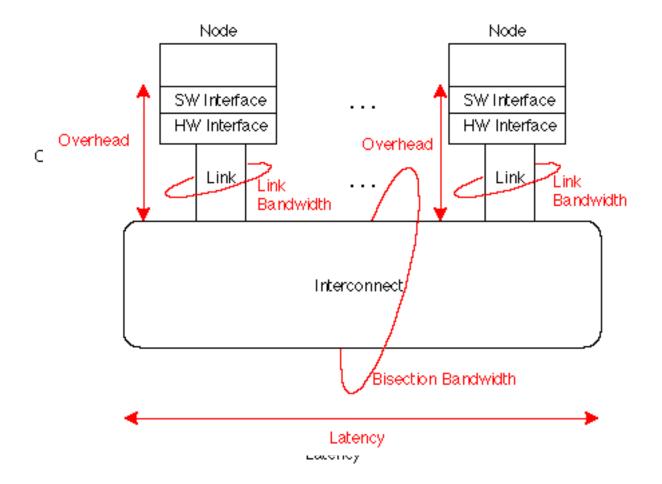
Transmission media



Shared vs switched media



Network performance measures

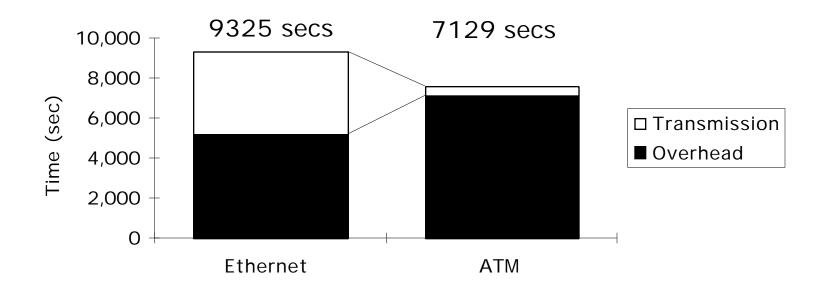


Example performance measures

SAN	LAN	WAN
CM-5	Ethernet	ATM
N x 5MB/s	1.125 MB/s	N x 10 MB/s
20 MB/s	1.125 MB/s	10 MB/s
5 µsec	15 µsec	50 to 10,000 µs
0.5/0.5 µs	6/6 µs	6/6 µs
1.6/12.4 µs	200/241 µs <i>(TCP/IP on L</i>	207/360 µs AN/WAN)
	CM-5 N x 5MB/s 20 MB/s 5 µsec 0.5/0.5 µs	CM-5EthernetN x 5MB/s1.125 MB/s20 MB/s1.125 MB/s5 μsec15 μsec0.5/0.5 μs6/6 μs

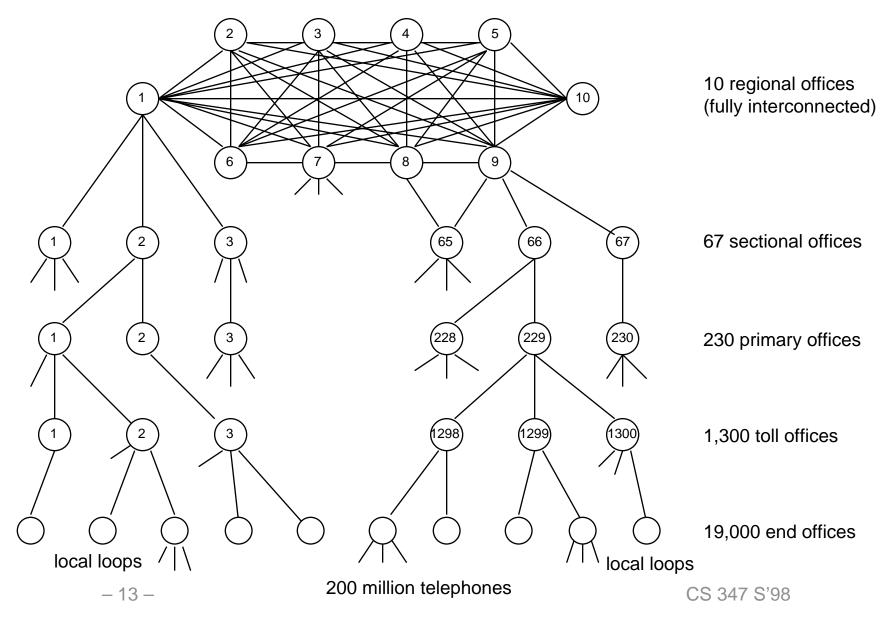
Importance of Overhead (+ Latency)

Ethernet / SS10:9 Mb/s BW,900 µsecs ovhdATM Synoptics:78 Mb/s BW,1,250 µsecs ovhd.NFS trace over 1 week:95% msgs < 200 bytes</td>

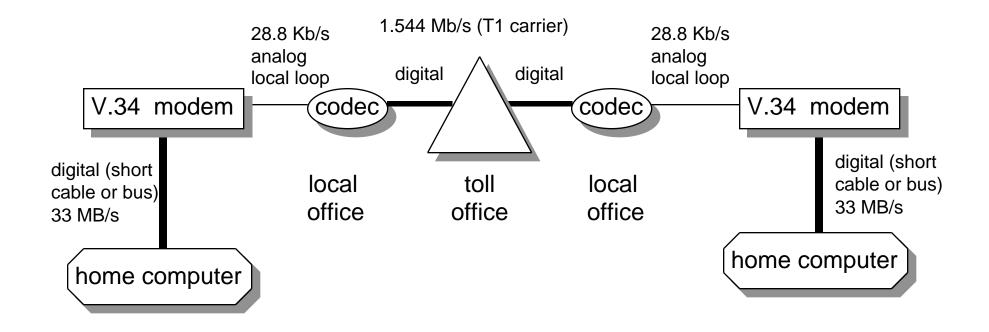


• Link bandwidth is as misleading as MIPS

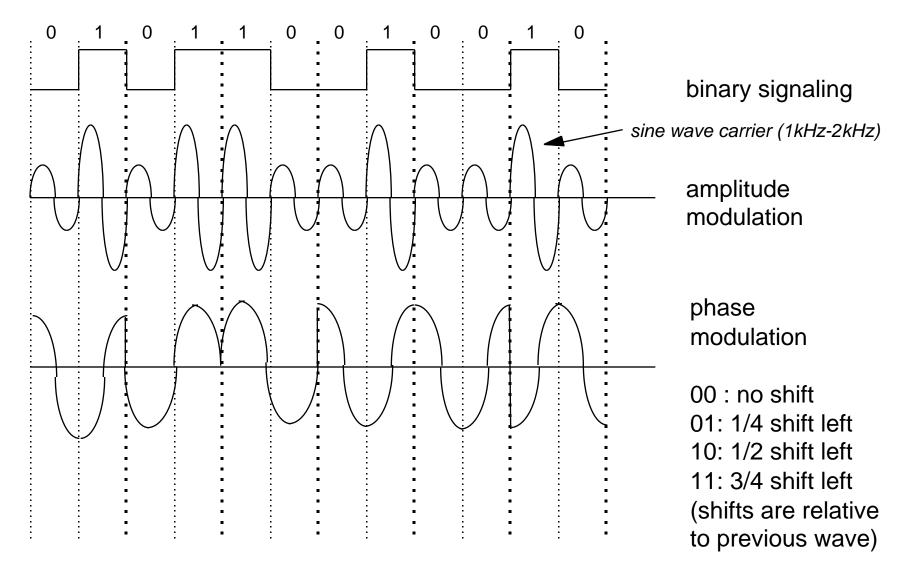
AT&T Telephone Hierarchy



Computer-to-computer calls



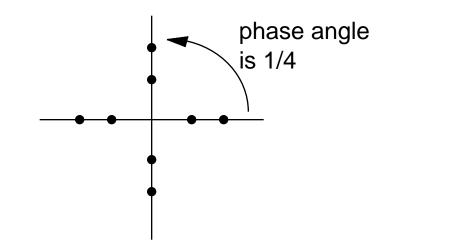
Modulating digital signals



CS 347 S'98

Quadrature amplitude modulation (QAM)

Modern modems use a combination of of amplitude and phase modulation to encode multiple bits per "symbol", i.e. amplitude/phase pair.



3 bits/symbol QAM modulation (8 symbols)

4 bits/symbol QAM modulation (16 symbols)

Conventional Modems

MOdulate: convert from digital to analog DEModulate: convert from analog to digital

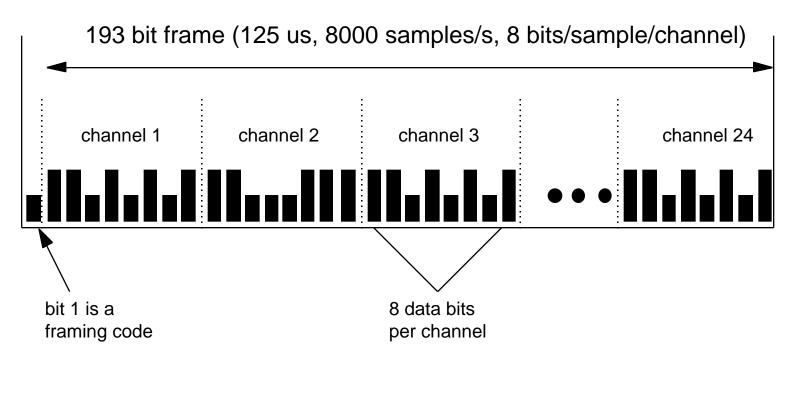
modem standards:					
type	symbols/sec	bits/symbol	Kb/s		
v.32	2400	4	9.6		
v.32.bis	2400	6	14.4		
v.34	3200	9	28.8		

Theoretical limit for modulated signals is approx 35 Kb/s:

Shannon's law: max bits/s = H log2(1 +S/N), where H is bdwdth and S/N is signal to noise ratio. For phone network, H~3,600 and S/N is 30 dB. Thus max rate is ~35 Kb/s.

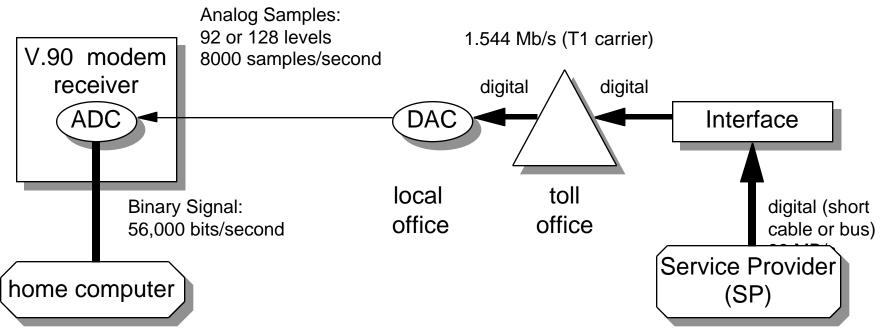
T1 carrier (1.544 Mb/s)

Digital part of phone system based on the T1 carrier:



Each channel has a data rate of 8000 samples/s * 8 bits/channel = 64 Kb/s

56KB "Modems"



• Asymmetric: home to SP uses conventional v.34 modem

• SP has digital connection into phone system

- Channel sending 8000 samples / second, up to 8-bits/sample

- DAC encodes each sample with 92 or 128 voltage levels
 - -Not enough precision on analog side to handle finer resolution
- Receiver converts samples back to digital values
 - -Must match frequency & phase of senders DAC
 - -Establish using "training" signals from sender

- 19 -

Ethernet

History

- 1976- proposed by Metcalfe and Boggs at Xerox PARC
- 1978 standardized by Xerox, Intel, DEC

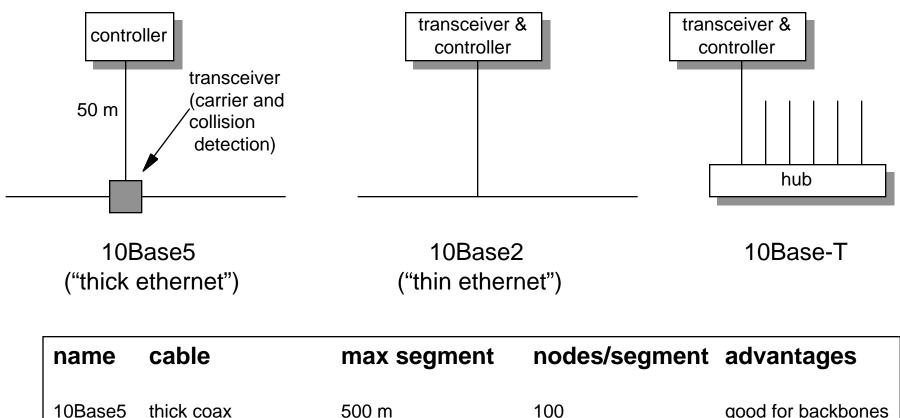
Bandwidth

• 10 Mbits/sec (old), 100 Mbits/sec (new)

Key features

- broadcast over shared bus (the ether)
 - -no centralized bus arbiter
 - -each adapter sees all bits
- each adapter has a unique (for all time!) 48-bit address
- variable length frames (packets) (64 1518 bytes)

Ethernet cabling



pest
pesi
maintenance
between bldgs

Ethernet frame format

64 - 1518 bytes

Preamble	Dest addr	Src addr	Frame type	Payload	CRC
64 bits	48 bits	48 bits	16 bits	368-12000 bits	32 bits

```
Preamble: 101010101 (synch)
```

dest and src addr: unique ethernet addresses

payload: data

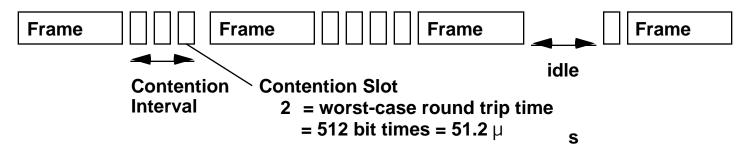
CRC: cyclic redundancy check (error detection)

Ethernet bus arbitration

10 Mbit/s centralized bus with no centralized control.

Arbitration scheme: Carrier sense multiple access with collision detection (CSMA/CD):

- Listen. If nobody is talking, go ahead and talk.
- If you hear someone else talking, yell SORRY and try again later

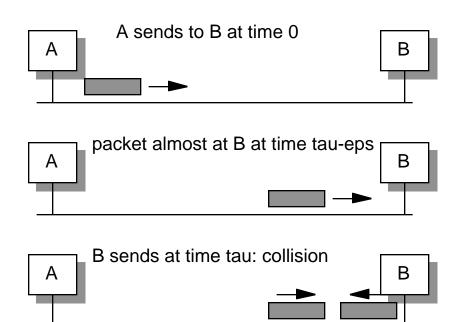


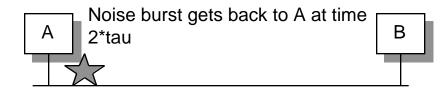
Binary exponential backoff:

- after 1st collision, wait 0 or 1 slots, at random.
- after 2nd collision, wait 0, 1, 2, 3 slots at random.
- etc up to 1023 slots.
- after 16 collisions, exception.

Why the 64 byte minimum packet size?

Assume propogation delay from A to B is tau.





Conclusions: Senders must take more than 2*tau seconds to send their packets.

For ethernet, tau is specified by IEEE 803.3 (2500 m cable w/ 4 repeaters) to be 51.2 usecs, which at 10 Mb/s is 512 bit times, or 64 bytes.

Rough estimate: propogation through copper is about 20 cm/ns. With a 2500 m cable, tau is 12.5 us and 2*tau is 25 us.

As speeds increase there are two possibilities:

- 1. increase packet sizes
- 2. decrease maximum cable length

Neither is particularly appealing.

Ethernet pros and cons

Pros:

- simple
- robust
- cheap (\$50/adapter in 1998)

Cons:

- no quality of service guarantees
 - -OK for data
 - -not OK for real-time bit streams like video or voice
- fixed bit rate
 - not keeping up with faster processors
 - -workstation can produce data at 10-50 MBytes/sec

prone to congestion

- -processors getting faster
- -bridged Ethernets can help some

Asynchronous transfer mode (ATM)

History

- 1988- proposed by international ATM forum
- telecomunications and computer vendors

Goal:

- mechanism for integrated transport of bit streams with different performance and reliability requirements (quality of service)
 - -video: 1.5 Mbits/sec, latency and variance sensitive, some bit loss OK
 - -voice: 64 Kbits/sec, latency and variance sensitive, some bit loss OK
 - -data: high data rates, latency and variance insensitive, bit loss not OK

ATM overview (cont)

Bandwidths

- OC-1: 51.84 Mbits/sec
- OC-3: 155.52 Mbits/sec (current LAN rate)
- OC-12: 622.08 Mbits/sec (current LAN rate)
- OC-24: 1244.16 (Gigabit network)

Key features:

- virtual connections (VC's)
 - -allow bandwidth reservation
- fixed cell (frame) size of 53 bytes
 - -simplifies high-speed switching
- small cell size
 - -allows fine-grained allocation of network bandwidth

ATM cell format

53 bytes (fixed)

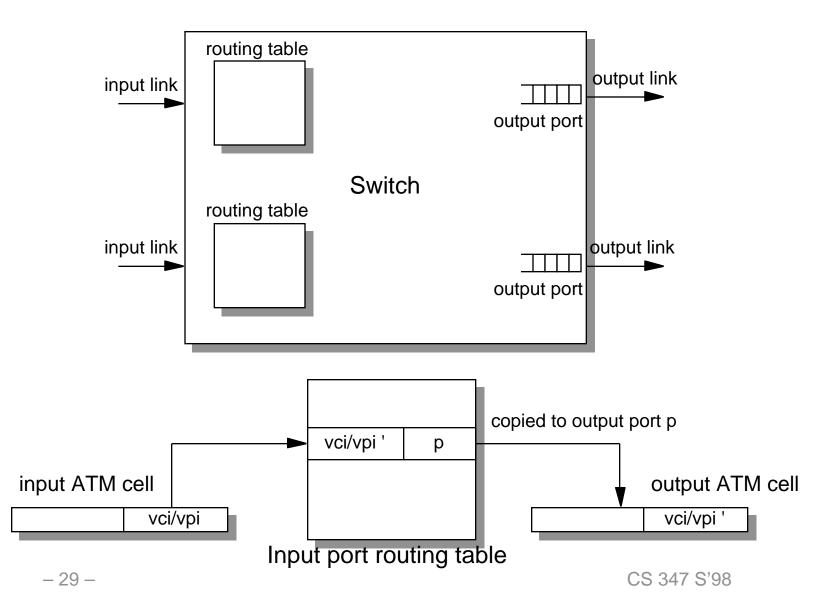
Generic flow ctl	VCI/ VPI	Payload type	Priority	Header checksum	Payload
4 bits	24 bits	2 bits	2 bits	8 bits	48 bytes

VCI: virtual connection (channel, circuit) identifier

VPI: virtual path identifier

payload: data

ATM cell routing



ATM pros and cons

Pros:

- bandwidth can be reserved (connections)
- scalable aggragate bandwidth (wide range of supported bit rates)
- support for network traffic with different quality of service requirements (small, fixed, easily multiplexed cells)
- potential for high speed switching (small fixed-size cells)

Cons:

- maximum user bandwidth still limited by link bandwidth
- connections make broadcast and multicast more difficult
- quality of service is still a research issue