Linking

15-213/15-503: Introduction to Computer Systems 9th Lecture, May 28, 2025

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Disclaimer

- Linkers continue to improve their functionality to help avoid programmer mistakes
 - This lecture follows the textbook, although some examples may be out of date with linker defaults / terminology

Today

- Linking
 - Motivation
 - What it does
 - How it works

Example C Program

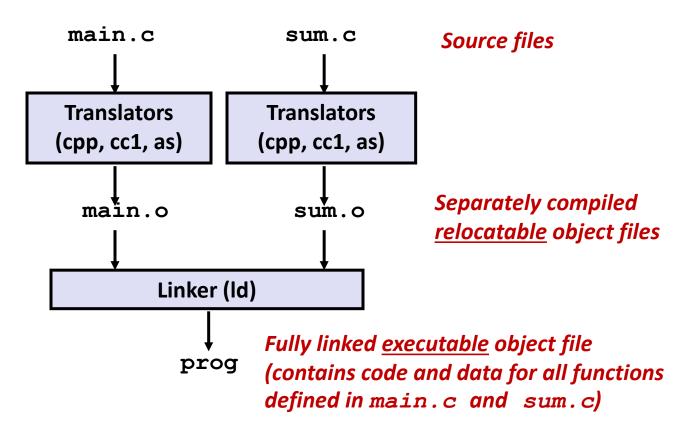
```
int sum(int *a, int n);
int array[2] = {1, 2};
int main(int argc, char** argv)
{
   int val = sum(array, 2);
   return val;
}
```

```
int sum(int *a, int n)
{
   int i, s = 0;

   for (i = 0; i < n; i++) {
       s += a[i];
   }
   return s;
}</pre>
```

Linking

- Programs are translated and linked using a compiler driver:
 - linux> gcc -Og -o prog main.c sum.c
 - linux> ./prog



Why Linkers?

- Reason 1: Modularity
 - Program can be written as a collection of smaller source files, rather than one monolithic mass.
 - Can build libraries of common functions
 - e.g., Math library, standard C library
 - Header files in C declare types that are defined in libraries

Why Linkers? (cont)

Reason 2: Efficiency

- Time: Separate compilation
 - Change one source file, compile, and then relink.
 - No need to recompile other source files.
 - Can compile multiple files concurrently.
- Space: Libraries
 - Common functions can be aggregated into a single file...
 - Option 1: Static Linking
 - Executable files and running memory images contain only the library code they actually use
 - Option 2: Dynamic linking
 - Executable files contain no library code
 - During execution, single copy of library code can be shared across all executing processes

What Do Linkers Do?

Step 1: Symbol resolution

Programs define and reference symbols (global variables and functions):

```
void swap() {...} /* define symbol swap */
swap(); /* reference symbol swap */
int *xp = &x; /* define symbol xp, reference x */
```

- Symbol definitions are stored in object file (by assembler) in symbol table.
 - Symbol table is an array of entries
 - Each entry includes name, size, and location of symbol.
- During symbol resolution step, the linker associates each symbol reference with exactly one symbol definition.

Symbols in Example C Program

Definitions

```
int sum(int *a, int n),
int array[2] = {1, 2};
int main(int argc, char** argv)
{
   int val = sum(array, 2);
   return val;
}
```

```
int sum(int *a, int n)
{
   int i, s = 0;

   for (i = 0; i < n; i++) {
       s += a[i];
   }
   return s;
}</pre>
```

Reference

What Do Linkers Do? (cont'd)

- Step 2: Relocation
 - Merges separate code and data sections into single sections
 - Relocates symbols from their relative locations in the .o files to their final absolute memory locations in the executable.
 - Updates all references to these symbols to reflect their new positions.

Let's look at these two steps in more detail....

Three Kinds of Object Files (Modules)

Relocatable object file (. o file)

- Contains code and data in a form that can be combined with other relocatable object files to form executable object file.
 - Each .o file is produced from exactly one source (.c) file

Executable object file (a.out file)

 Contains code and data in a form that can be copied directly into memory and then executed.

Shared object file (.so file)

- Special type of relocatable object file that can be loaded into memory and linked dynamically, at either load time or run-time.
- Called Dynamic Link Libraries (DLLs) by Windows

Executable and Linkable Format (ELF)

- Standard binary format for object files
- One unified format for
 - Relocatable object files (.o),
 - Executable object files (a.out)
 - Shared object files (.so)
- Generic name: ELF binaries

ELF Object File Format

Elf header

 Word size, byte ordering, file type (.o, exec, .so), machine type, etc.

Segment header table

 Page size, virtual address memory segments (sections), segment sizes.

. text section

Code

.rodata section

Read only data: jump tables, string constants, ...

. data section

Initialized global variables

.bss section

- Uninitialized global variables
- "Block Started by Symbol"
- "Better Save Space"
- Has section header but occupies no space

ELF header
Segment header table (required for executables)
. text section
.rodata section
. data section
.bss section
.symtab section
.rel.txt section
.rel.data section
.debug section
Section header table

ELF Object File Format (cont.)

. symtab section

- Symbol table
- Procedure and static variable names
- Section names and locations

.rel.text section

- Relocation info for .text section
- Addresses of instructions that will need to be modified in the executable
- Instructions for modifying

.rel.data section

- Relocation info for .data section
- Addresses of pointer data that will need to be modified in the merged executable

debug section

■ Info for symbolic debugging (gcc -g)

Section header table

Offsets and sizes of each section

ELF header
Segment header table (required for executables)
. text section
.rodata section
. data section
. bss section
.symtab section
.rel.txt section
.rel.data section
. debug section
Section header table

Linker Symbols

Global symbols

- Symbols defined by module m that can be referenced by other modules.
- e.g., non-static C functions and non-static global variables.

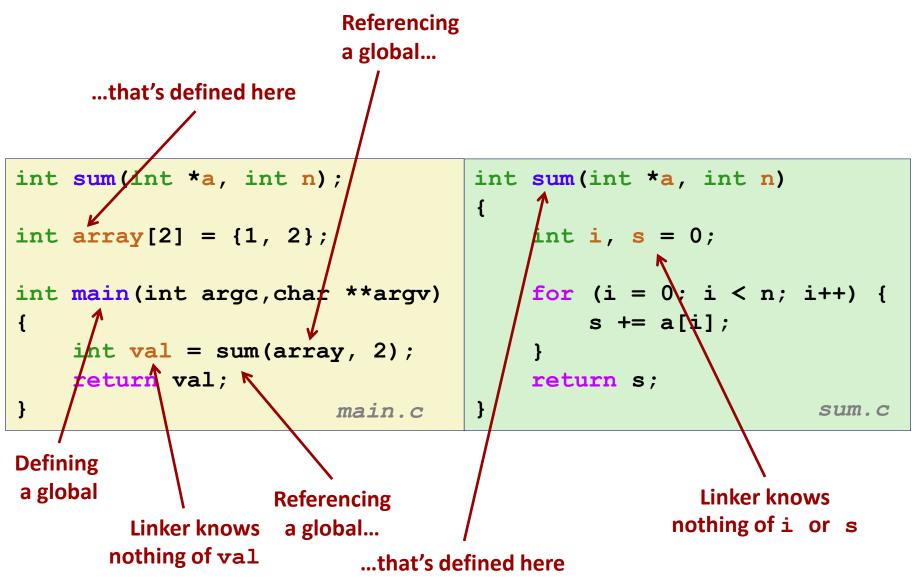
External symbols

 Global symbols that are referenced by module m but defined by some other module.

Local symbols

- Symbols that are defined and referenced exclusively by module m.
- e.g, C functions and global variables defined with the **static** attribute.
- Local linker symbols are not local program variables

Step 1: Symbol Resolution



Symbol Identification

Which of the following names will be in the symbol table of symbols.o?

symbols.c:

```
int incr = 1;
static int foo(int a) {
  int b = a + incr;
  return b;
int main (int argc,
         char* argv[]) {
 printf("%d\n", foo(5));
  return 0;
```

Names:

- incr
- foo
- argc
- argv
- main
- printf
- "%d\n"

Can find this with readelf: linux> readelf -s symbols.o

Local Symbols

Local non-static C variables vs. local static C variables

- Local non-static C variables: stored on the stack
- Local static C variables: stored in either .bss or .data

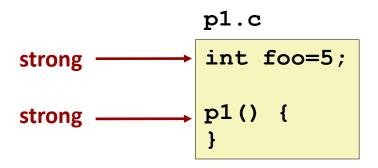
```
static int x = 15;
int f() {
    static int x = 17;
    return x++;
int q() {
    static int x = 19;
    return x += 14;
int h() {
    return x += 27;
        static-local.c
```

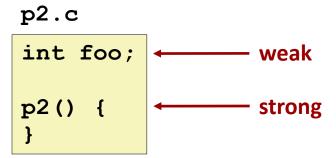
Compiler allocates space in .data for each definition of x

Creates local symbols in the symbol table with unique names, e.g., x, x.1721 and x.1724.

How Linker Resolves Duplicate Symbol Definitions

- Program symbols are either strong or weak
 - Strong: procedures and initialized globals
 - Weak: uninitialized globals
 - Or ones declared with specifier extern





Linker's Symbol Rules

- Rule 1: Multiple strong symbols are not allowed
 - Each item can be defined only once
 - Otherwise: Linker error
- Rule 2: Given a strong symbol and multiple weak symbols, choose the strong symbol
 - References to the weak symbol resolve to the strong symbol
- Rule 3: If there are multiple weak symbols, pick an arbitrary one
 - Can override this with gcc -fno-common
- Puzzles on the next slide

Linker Puzzles

```
int x;
p1() {}
```

Link time error: two strong symbols (p1)

```
int x;
p1() {}
```

References to **x** will refer to the same uninitialized int. Is this what you really want?

```
int x;
int y;
p1() {}
```

Writes to **x** in **p2** might overwrite **y**! Evil!

```
int x=7;
int y=5;
p1() {}
```

Writes to **x** in **p2** might overwrite **y**! Nasty!

References to **x** will refer to the same initialized variable.

Important: Linker does not do type checking.

Type Mismatch Example

- Compiles without any errors or warnings
- What gets printed?

```
-bash-4.2$ ./mismatch
4614253070214989087
```

Global Variables

Avoid if you can

Otherwise

- Use static if you can
- Initialize if you define a global variable
- Use extern if you reference an external global variable
 - Treated as weak symbol
 - But also causes linker error if not defined in some file

Use of extern in .h Files (#1)

c1.c

```
#include "global.h"
int f() {
  return g+1;
}
```

global.h

```
extern int g;
int f();
```

c2.c

```
#include <stdio.h>
#include "global.h"

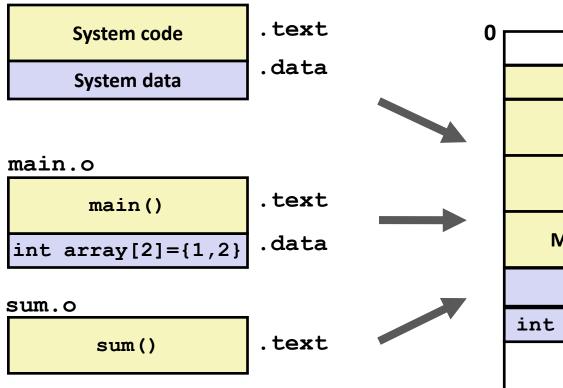
int g = 0;

int main(int argc, char argv[]) {
   int t = f();
   printf("Calling f yields %d\n", t);
   return 0;
}
```

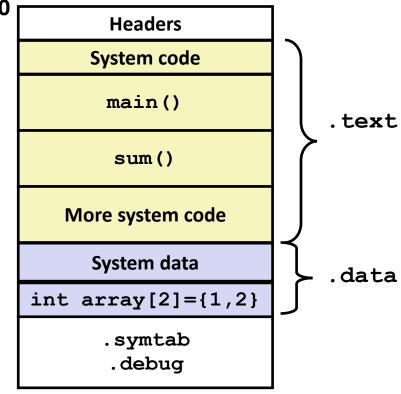
Linking Example

Step 2: Relocation

Relocatable Object Files



Executable Object File



Relocation Entries

```
int array[2] = {1, 2};
int main(int argc, char**
argv)
{
   int val = sum(array, 2);
   return val;
}
```

```
0000000000000000 <main>:
  0: 48 83 ec 08
                                    $0x8,%rsp
                              sub
  4: be 02 00 00 00
                                    $0x2,%esi
                             mov
                                    $0x0, %edi  # %edi = &array
  9: bf 00 00 00 00
                             mov
                      a: R X86 64 32 array
                                                  # Relocation entry
       e8 00 00 00 00
                              callq 13 < main + 0x13 > \# sum()
  e:
                      f: R X86 64 PC32 sum-0x4 # Relocation entry
 13: 48 83 c4 08
                              add
                                    $0x8,%rsp
 17:
    c3
                              retq
                                                              main.o
```

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Relocated .text section

```
00000000004004d0 <main>:
 4004d0:
                48 83 ec 08
                                         $0x8,%rsp
                                  sub
 4004d4:
                be 02 00 00 00
                                         $0x2,%esi
                                  mov
                                         $0x601018, %edi # %edi = &array
 4004d9:
               bf 18 10 60 00
                                  mov
 4004de:
                e8 05 00 00 00
                                         4004e8 <sum>
                                                          # sum()
                                  callq
 4004e3:
               48 83 c4 08
                                         $0x8,%rsp
                                  add
 4004e7:
                c3
                                  reta
00000000004004e8 <sum>:
 4004e8:
                b8 00 00 00 00
                                                $0x0, %eax
                                        mov
               ba 00 00 00 00
                                                $0x0,%edx
 4004ed:
                                        mov
                                                4004fd < sum + 0x15 >
 4004f2:
                eb 09
                                        jmp
 4004f4:
               48 63 ca
                                        movslq %edx,%rcx
                03 04 8f
 4004f7:
                                        add
                                               (%rdi,%rcx,4),%eax
 4004fa:
               83 c2 01
                                        add
                                               $0x1, %edx
 4004fd:
                39 £2
                                               %esi,%edx
                                        cmp
 4004ff:
                7c f3
                                               4004f4 < sum + 0xc >
                                        il
 400501:
                f3 c3
                                        repz retq
```

callq instruction uses PC-relative addressing for sum():

0x4004e8 = 0x4004e3 + 0x5

Loading Executable Object Files

Executable Object File

ELF header	
Program header table (required for executables)	
.init section	
.text section	
.rodata section	
.data section	
.bss section	
.symtab	
.debug	
.line	
.strtab	
Section header table (required for relocatables)	

Memory invisible to **Kernel virtual memory** user code User stack (created at runtime) %rsp (stack pointer) Memory-mapped region for shared libraries brk **Run-time heap** (created by malloc) Loaded Read/write data segment from (.data, .bss) the Read-only code segment executable (.init,.text,.rodata) file Unused

0x400000

Quiz

https://canvas.cmu.edu/courses/47415/quizzes/143244

Packaging Commonly Used Functions

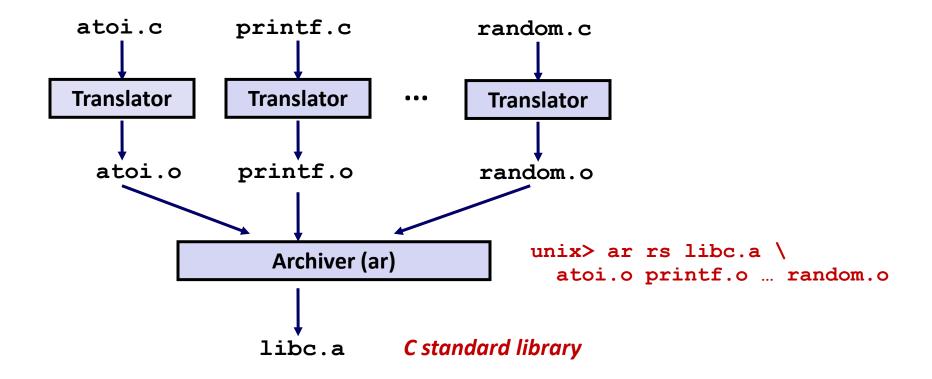
- How to package functions commonly used by programmers?
 - Math, I/O, memory management, string manipulation, etc.
- Awkward, given the linker framework so far:
 - Option 1: Put all functions into a single source file
 - Programmers link big object file into their programs
 - Space and time inefficient
 - Option 2: Put each function in a separate source file
 - Programmers explicitly link appropriate binaries into their programs
 - More efficient, but burdensome on the programmer

Old-fashioned Solution: Static Libraries

Static libraries (.a archive files)

- Concatenate related relocatable object files into a single file with an index (called an archive).
- Enhance linker so that it tries to resolve unresolved external references by looking for the symbols in one or more archives.
- If an archive member file resolves reference, link it into the executable.

Creating Static Libraries



- Archiver allows incremental updates
- Recompile function that changes and replace .o file in archive.

Commonly Used Libraries

libc.a (the C standard library)

- 4.6 MB archive of 1496 object files.
- I/O, memory allocation, signal handling, string handling, data and time, random numbers, integer math

libm. a (the C math library)

- 2 MB archive of 444 object files.
- floating point math (sin, cos, tan, log, exp, sqrt, ...)

```
% ar -t libc.a | sort
...
fork.o
...
fprintf.o
fpu_control.o
fputc.o
freopen.o
fscanf.o
fseek.o
fstab.o
...
```

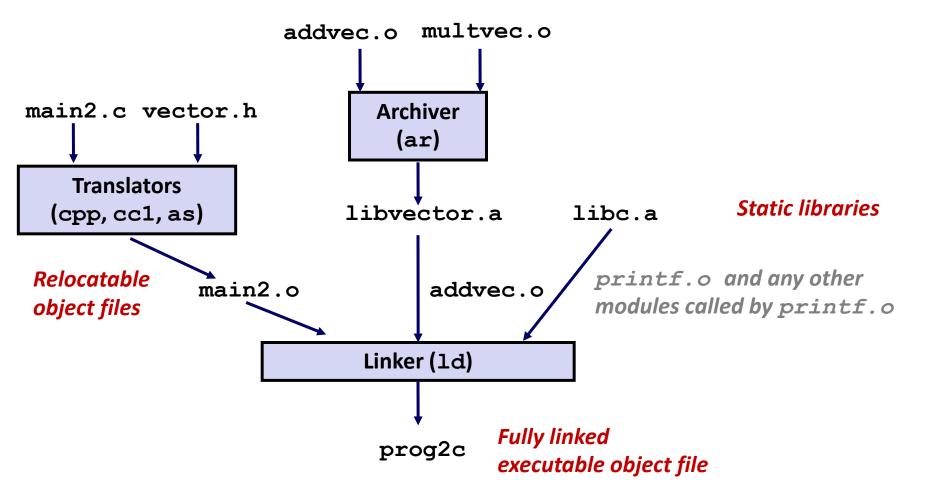
```
% ar -t libm.a | sort
...
e_acos.o
e_acosf.o
e_acosh.o
e_acoshf.o
e_acoshl.o
e_acosl.o
e_asin.o
e_asinf.o
e_asinf.o
e_asinf.o
...
```

Linking with Static Libraries

```
#include <stdio.h>
#include "vector.h"
int x[2] = \{1, 2\};
int y[2] = \{3, 4\};
int z[2];
int main(int argc, char**
argv)
{
    addvec(x, y, z, 2);
    printf("z = [%d %d] \n",
           z[0], z[1]);
    return 0;
                    main2.c
```

libvector.a

Linking with Static Libraries



"c" for "compile-time"

Using Static Libraries

Linker's algorithm for resolving external references:

- Scan .o files and .a files in the command line order.
- During the scan, keep a list of the current unresolved references.
- As each new .o or .a file, obj, is encountered, try to resolve each unresolved reference in the list against the symbols defined in obj.
- If any entries in the unresolved list at end of scan, then error.

Problem:

- Command line order matters!
- Moral: put libraries at the end of the command line.

```
unix> gcc -L. libtest.o -lmine
unix> gcc -L. -lmine libtest.o
libtest.o: In function 'main':
libtest.o(.text+0x4): undefined reference to 'libfun'
```

Modern Solution: Shared Libraries

Static libraries have the following disadvantages:

- Duplication in the stored executables (every function needs libc)
- Duplication in the running executables
- Minor bug fixes of system libraries require each application to explicitly relink
 - Rebuild everything with glibc?
 - https://security.googleblog.com/2016/02/cve-2015-7547-glibc-getaddrinfo-stack.html

Modern solution: Shared Libraries

- Object files that contain code and data that are loaded and linked into an application dynamically, at either load-time or run-time
- Also called: dynamic link libraries, DLLs, .so files

Shared Libraries (cont.)

- Dynamic linking can occur when executable is first loaded and run (load-time linking).
 - Common case for Linux, handled automatically by the dynamic linker (ld-linux.so).
 - Standard C library (libc.so) usually dynamically linked.
- Dynamic linking can also occur after program has begun (run-time linking).
 - In Linux, this is done by calls to the dlopen() interface.
 - Distributing software.
 - High-performance web servers.
 - Runtime library interpositioning.
- Shared library routines can be shared by multiple processes.
 - More on this when we learn about virtual memory

What dynamic libraries are required?

.interp section

Specifies the dynamic linker to use (i.e., ld-linux.so)

.dynamic section

- Specifies the names, etc of the dynamic libraries to use
- Follow an example of csim-ref from cachelab

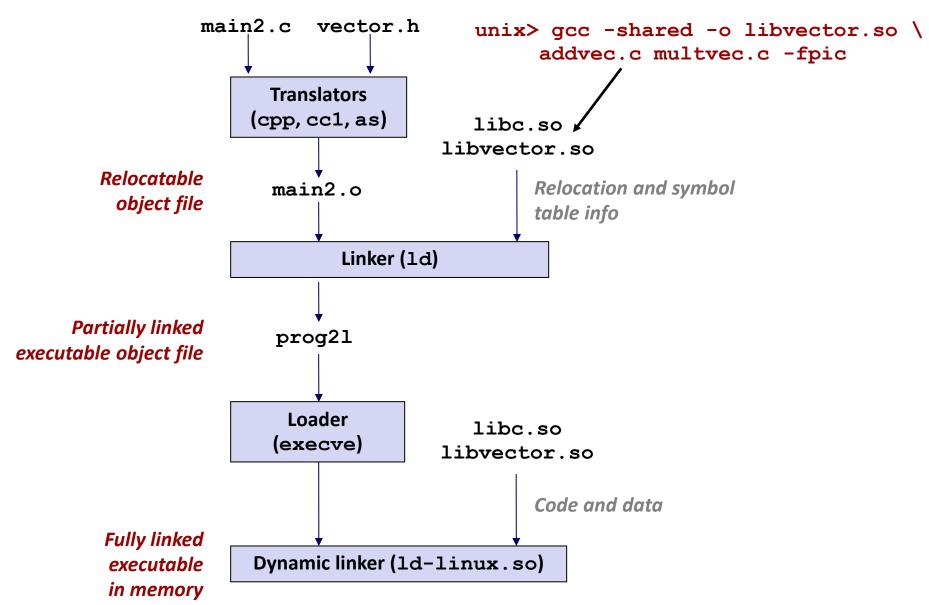
```
(NEEDED) Shared library: [libm.so.6]
```

Where are the libraries found?

Use "1dd" to find out:

```
unix> ldd csim-ref
  linux-vdso.so.1 => (0x00007ffc195f5000)
  libc.so.6 => /lib/x86_64-linux-gnu/libc.so.6 (0x00007f345eda6000)
  /lib64/ld-linux-x86-64.so.2 (0x00007f345f181000)
```

Dynamic Linking at Load-time



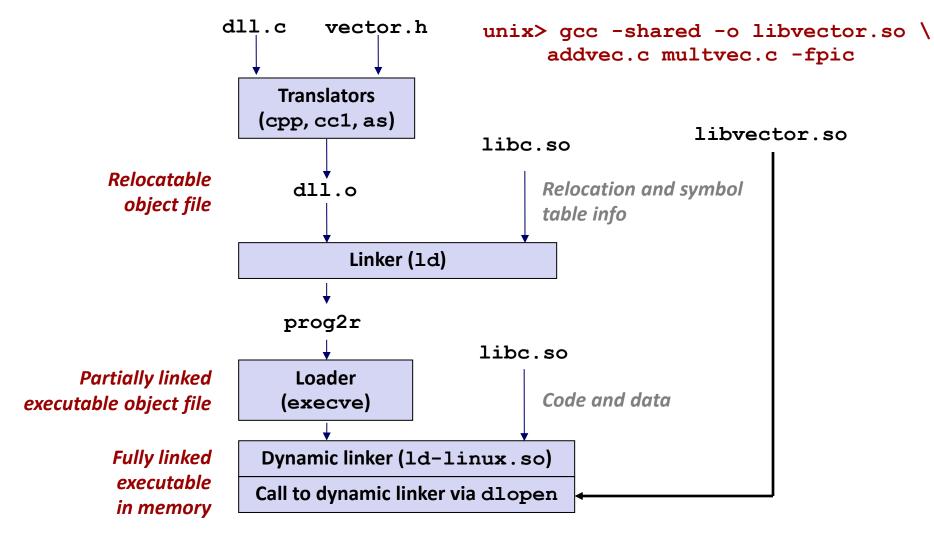
Dynamic Linking at Run-time

```
#include <stdio.h>
#include <stdlib.h>
#include <dlfcn.h>
int x[2] = \{1, 2\};
int y[2] = \{3, 4\};
int z[2];
int main(int argc, char** argv)
{
   void *handle;
   void (*addvec)(int *, int *, int *, int);
    char *error;
    /* Dynamically load the shared library that contains addvec() */
    handle = dlopen("./libvector.so", RTLD LAZY);
    if (!handle) {
        fprintf(stderr, "%s\n", dlerror());
       exit(1);
                                                                 d11.c
```

Dynamic Linking at Run-time (cont)

```
/* Get a pointer to the addvec() function we just loaded */
addvec = dlsym(handle, "addvec");
if ((error = dlerror()) != NULL) {
    fprintf(stderr, "%s\n", error);
    exit(1);
/* Now we can call addvec() just like any other function */
addvec(x, y, z, 2);
printf("z = [%d %d] \n", z[0], z[1]);
/* Unload the shared library */
if (dlclose(handle) < 0) {</pre>
    fprintf(stderr, "%s\n", dlerror());
    exit(1);
return 0;
                                                        d11.c
```

Dynamic Linking at Run-time



Linking Recap

- Usually: Just happens, no big deal
- Sometimes: Strange errors

Supplemental

Dynamic linking can also be