15-213

"The course that gives CMU its Zip!"

Linking
Oct. 9, 2003

Topics

- Static linking
- Dynamic linking
- Case study: Library interpositioning

Example C Program

main.c

```
int buf[2] = (1,2)

int main()
{
  swap();
  return 0;
}
```

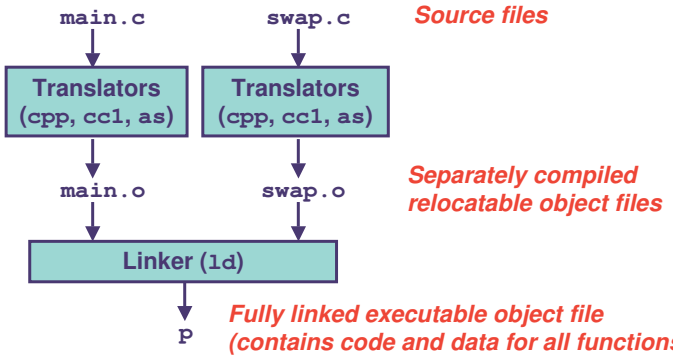
swap.c

```
extern int buf[];

int *bufp0 = &buf[0];
int *bufp1;

void swap()
{
  int temp;

  bufp1 = &buf[1];
  temp = *bufp0;
  *bufp0 = *bufp1;
  *bufp1 = temp;
}
```


Static Linking

Programs are translated and linked using a *compiler driver*:

- `unix> gcc -O2 -g -o p main.c swap.c`
- `unix> ./p`



Why Linkers?

Modularity

- Program can be written as a collection of smaller source files, rather than one monolithic mass.
- Can build libraries of common functions (more on this later)
  - e.g., Math library, standard C library

Efficiency

- Time:
  - Change one source file, compile, and then relink.
  - No need to recompile other source files.
- Space:
  - Libraries of common functions can be aggregated into a single file...
  - Yet executable files and running memory images contain only code for the functions they actually use.

What Does a Linker Do?

Symbol resolution

- Programs define and reference *symbols* (variables and functions):
  - code: `a();  /* reference to symbol a */`
  - data: `int *xp=&x;  /* reference to symbol x */`
- Symbol definitions are stored (by compiler) in *symbol table*.
  - Each entry includes name, size, and location of object.
- Linker associates each symbol reference with exactly one symbol definition.

Relocation

- Merges separate code and data sections into single sections
- Relocates symbols from their relative locations in the `.o` files to new absolute positions in the executable.
- Updates all references to these symbols to reflect their new positions.

Object Files (Modules)

Relocatable object file (`.o` file)

- Contains code and data in a form that can be combined with other relocatable object files at compile time to form an executable.

Executable object file

- Contains code and data in a form that can be copied directly into memory and then executed.

Shared object file (`.so` file)

- Special type of relocatable object file that can be loaded into memory and linked dynamically, at either load time or run-time.
- Called *Dynamic Link Libraries* (DLLs) by Windows

Executable and Linkable Format (ELF)

Standard binary format for object files

Derives from AT&T System V Unix

- Later adopted by BSD Unix variants and Linux

One unified format for

- Relocatable object files (`.o`),
- Executable object files
- Shared object files (`.so`)

Generic name: ELF binaries

ELF Object File Format

Elf header

- Magic number, type (.o, exec, .so), machine, byte ordering, etc.

Program header table

- Page size, virtual addresses memory segments (sections), segment sizes.

`.text` section

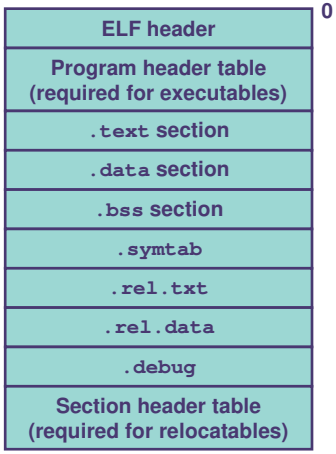
- Code

`.data` section

- Initialized (static) data

`.bss` section

- Uninitialized (static) data
- "Block Started by Symbol"
- "Better Save Space"
- Has section header but occupies no space



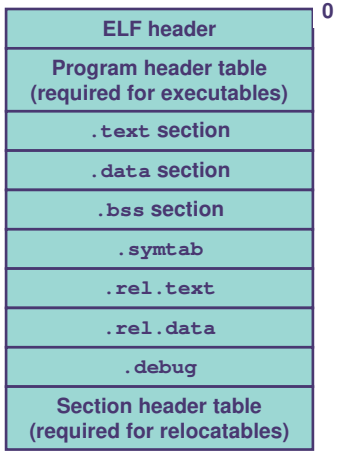
ELF Object File Format (cont)

.symtab section
- Symbol table
- Procedure and static variable names
- Section names and locations

.rel.text section
- Relocation info for .text section
- Addresses of instructions that will need to be modified in the executable
- Instructions for modifying.

.rel.data section
- Relocation info for .data section
- Addresses of pointer data that will need to be modified in the merged executable

.debug section
- Info for symbolic debugging (gcc -g)



Linker Symbols

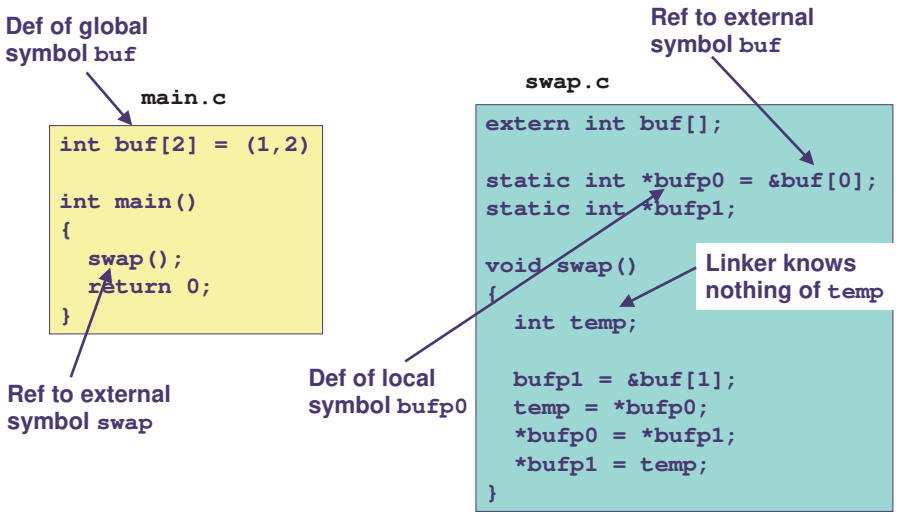
Global symbols
- Symbols defined by module *m* that can be referenced by other modules.
- Ex: non-static C functions and non-static global variables.

External symbols
- Global symbols that are referenced by module *m* but defined by some other module.

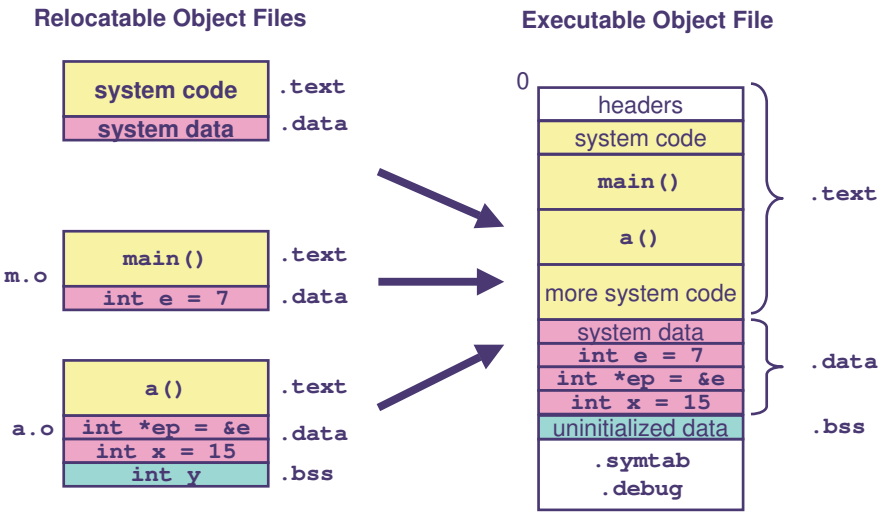
Local symbols
- Symbols that are defined and referenced exclusively by module *m*.
- Ex: C functions and variables that are defined with the static attribute.

Key Point: Local linker symbols are *not* local program variables

Resolving Symbols



Relocating Symbol Definitions and References



main.o Relocation Info

```
int buf[2] = (1,2)

int main()
{
  swap();
  return 0;
}
```

```
0000000 <main>:
  0:   55              push   %ebp
  1:   89 e5           mov    %esp,%ebp
  3:   83 ec 08        sub    $0x8,%esp
  6:   e8 fc ff ff ff  call   7 <main+0x7>
                       7: R_386_PC32 swap
  b:   31 c0           xor    %eax,%eax
  d:   89 ec           mov    %ebp,%esp
  f:   5d              pop    %ebp
 10:   c3              ret
```

```
Disassembly of section .data:

 00000000 <buf>:
   0:   01 00 00 00 02 00 00 00
```

Source: objdump

swap.o Relocation Info (.text)

```
extern int buf[];

static int *bufp0 =
           &buf[0];
static int *bufp1;

void swap()
{
  int temp;

  bufp1 = &buf[1];
  temp = *bufp0;
  *bufp0 = *bufp1;
  *bufp1 = temp;
}
```

```
Disassembly of section .text:

00000000 <swap>:
 0: 55                    push   %ebp
 1: 8b 15 00 00 00 00     mov    0x0,%edx
                          3: R_386_32 bufp0
 7: a1 0 00 00 00         mov    0x4,%eax
                          8: R_386_32 buf
 c: 89 e5                 mov    %esp,%ebp
 e: c7 05 00 00 00 00 04  movl   $0x4,0x0
15: 00 00 00
                          10: R_386_32 bufp1
                          14: R_386_32 buf
18: 89 ec                mov    %ebp,%esp
1a: 8b 0a                mov    (%edx),%ecx
1c: 89 02                mov    %eax,(%edx)
1e: a1 00 00 00 00       mov    0x0,%eax
                          1f: R_386_32 bufp1
23: 89 08                mov    %ecx,(%eax)
25: 5d                   pop    %ebp
26: c3                   ret
```


a.o Relocation Info (.data)

```
extern int buf[];

static int *bufp0 =
           &buf[0];
static int *bufp1;

void swap()
{
  int temp;

  bufp1 = &buf[1];
  temp = *bufp0;
  *bufp0 = *bufp1;
  *bufp1 = temp;
}
```

```
Disassembly of section .data:

00000000 <bufp0>:
   0:   00 00 00 00

        0: R_386_32 buf
```


Executable After Relocation (.text)

```
080483b4 <main>:
 80483b4:       55                      push   %ebp
 80483b5:       89 e5                   mov    %esp,%ebp
 80483b7:       83 ec 08                sub    $0x8,%esp
 80483ba:       e8 09 00 00 00          call   80483c8 <swap>
 80483bf:       31 c0                   xor    %eax,%eax
 80483c1:       89 ec                   mov    %ebp,%esp
 80483c3:       5d                      pop    %ebp
 80483c4:       c3                      ret
080483c8 <swap>:
 80483c8:       55                      push   %ebp
 80483c9:       8b 15 5c 94 04 08       mov    0x804945c,%edx
 80483cf:       a1 58 94 04 08          mov    0x8049458,%eax
 80483d4:       89 e5                   mov    %esp,%ebp
 80483d6:       c7 05 48 95 04 08 58    movl   $0x8049458,0x8049548
 80483dd:       94 04 08
 80483e0:       89 ec                   mov    %ebp,%esp
 80483e2:       8b 0a                   mov    (%edx),%ecx
 80483e4:       89 02                   mov    %eax,(%edx)
 80483e6:       a1 48 95 04 08          mov    0x8049548,%eax
 80483eb:       89 08                   mov    %ecx,(%eax)
 80483ed:       5d                      pop    %ebp
 80483ee:       c3                      ret
```

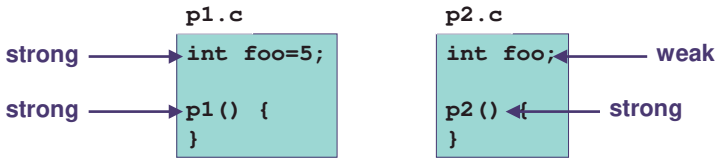
Executable After Relocation (.data)

```
Disassembly of section .data:

08049454 <buf>:
 8049454:       01 00 00 00 02 00 00 00

0804945c <bufp0>:
 804945c:       54 94 04 08
```


Strong and Weak Symbols

Program symbols are either strong or weak

- *strong*: procedures and initialized globals
- *weak*: uninitialized globals



Linker's Symbol Rules

Rule 1. A strong symbol can only appear once.

Rule 2. A weak symbol can be overridden by a strong symbol of the same name.

- references to the weak symbol resolve to the strong symbol.

Rule 3. If there are multiple weak symbols, the linker can pick an arbitrary one.

Linker Puzzles

`int x; p1() {}`	`p1() {}`	Link time error: two strong symbols (`p1`)
`int x; p1() {}`	`int x; p2() {}`	References to `x` will refer to the same uninitialized int. Is this what you really want?
`int x; int y; p1() {}`	`double x; p2() {}`	Writes to `x` in `p2` might overwrite `y`! Evil!
`int x=7; int y=5; p1() {}`	`double x; p2() {}`	Writes to `x` in `p2` will overwrite `y`! Nasty!
`int x=7; p1() {}`	`int x; p2() {}`	References to `x` will refer to the same initialized variable.

Nightmare scenario: two identical weak structs, compiled by different compilers with different alignment rules.

Packaging Commonly Functions

How to package functions commonly used by programmers?

- Math, I/O, memory management, string manipulation, etc.

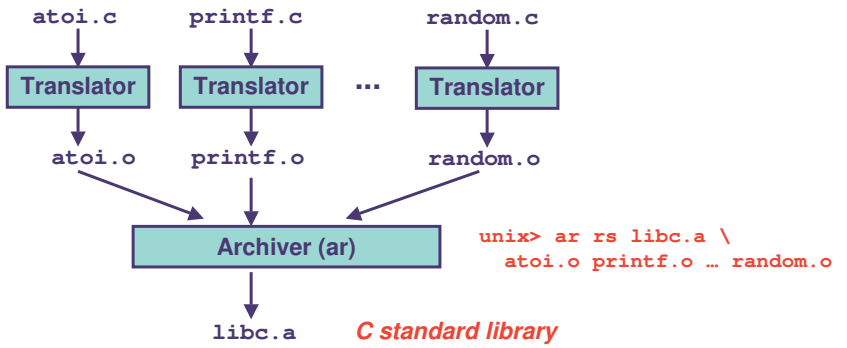
Awkward, given the linker framework so far:

- Option 1: Put all functions in a single source file
  - Programmers link big object file into their programs
  - Space and time inefficient
- Option 2: Put each function in a separate source file
  - Programmers explicitly link appropriate binaries into their programs
  - More efficient, but burdensome on the programmer

Solution: *static libraries* (`.a` archive files)

- Concatenate related relocatable object files into a single file with an index (called an archive).
- Enhance linker so that it tries to resolve unresolved external references by looking for the symbols in one or more archives.
- If an archive member file resolves reference, link into executable.

Creating Static Libraries



Archiver allows incremental updates:

• Recompile function that changes and replace .o file in archive.

Commonly Used Libraries

`libc.a` (the C standard library)

- 8 MB archive of 900 object files.
- I/O, memory allocation, signal handling, string handling, data and time, random numbers, integer math

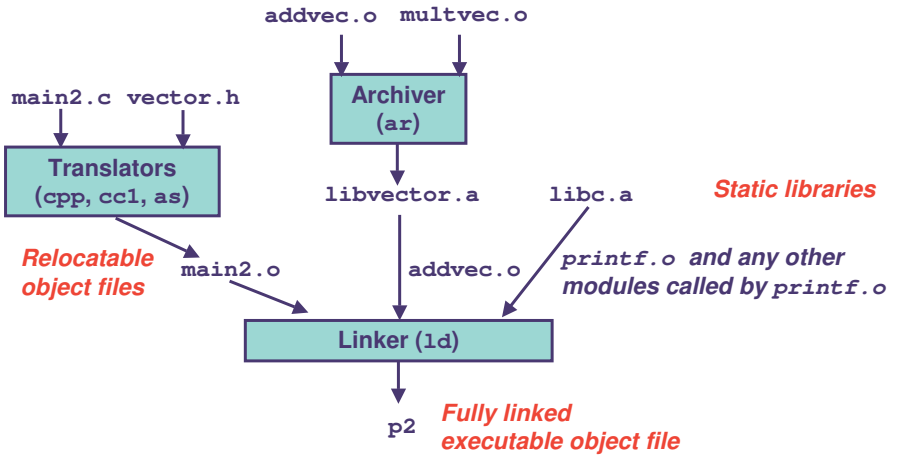
`libm.a` (the C math library)

- 1 MB archive of 226 object files.
- floating point math (sin, cos, tan, log, exp, sqrt, …)

```
% ar -t /usr/lib/libc.a | sort
…
fork.o
…
fprintf.o
fpu_control.o
fputc.o
freopen.o
fscanf.o
fseek.o
fstab.o
…
```

```
% ar -t /usr/lib/libm.a | sort
…
e_acos.o
e_acosf.o
e_acosh.o
e_acoshf.o
e_acoshl.o
e_acosl.o
e_asin.o
e_asinf.o
e_asinl.o
…
```


Linking with Static Libraries



Using Static Libraries

Linker's algorithm for resolving external references:

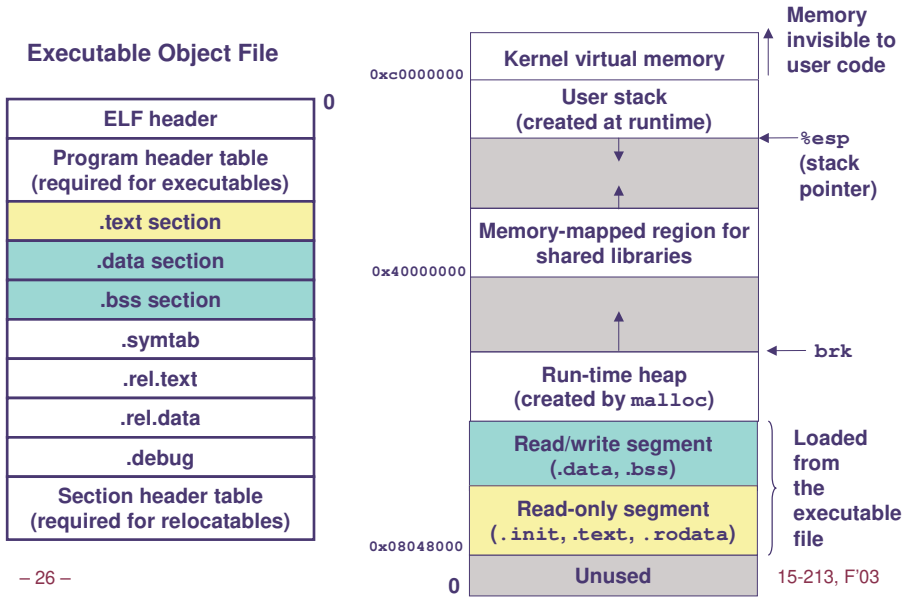
- Scan `.o` files and `.a` files in the command line order.
- During the scan, keep a list of the current unresolved references.
- As each new `.o` or `.a` file, *obj*, is encountered, try to resolve each unresolved reference in the list against the symbols defined in *obj*.
- If any entries in the unresolved list at end of scan, then error.

Problem:

- Command line order matters!
- Moral: put libraries at the end of the command line.

```
bass> gcc -L. libtest.o -lmine
bass> gcc -L. -lmine libtest.o
libtest.o: In function `main':
libtest.o(.text+0x4): undefined reference to `libfun'
```


Loading Executable Object Files



Shared Libraries

Static libraries have the following disadvantages:

- Potential for duplicating lots of common code in the executable files on a filesystem.
  - e.g., every C program needs the standard C library
- Potential for duplicating lots of code in the virtual memory space of many processes.
- Minor bug fixes of system libraries require each application to explicitly relink

Modern Solution: Shared Libraries

- Object files that contain code and data that are loaded and linked into an application *dynamically*, at either *load-time* or *run-time*
- Dynamic link libraries, DLLs, .so files

Shared Libraries (cont)

Dynamic linking can occur when executable is first loaded and run (load-time linking).

- Common case for Linux, handled automatically by the dynamic linker (`ld-linux.so`).
- Standard C library (`libc.so`) usually dynamically linked.

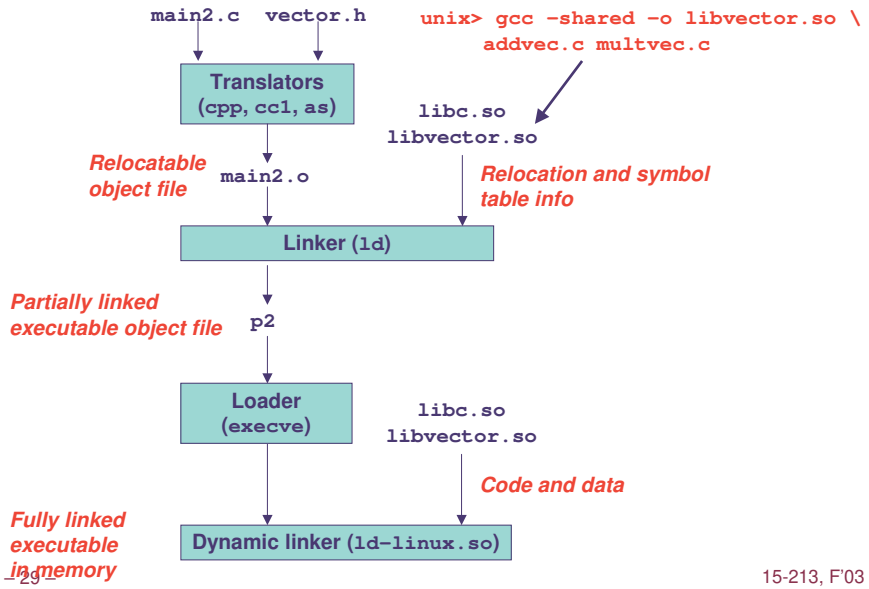
Dynamic linking can also occur after program has begun (run-time linking).

- In Unix, this is done by calls to the `dlopen()` interface.
  - High-performance web wervers.
  - Runtime library interpositioning

Shared library routines can be shared by multiple processes.

- More on this when we learn about virtual memory.

Dynamic Linking at Load-time



Dynamic Linking at Run-time

```
#include <stdio.h>
#include <dlfcn.h>

int x[2] = {1, 2};
int y[2] = {3, 4};
int z[2];

int main()
{
    void *handle;
    void (*addvec)(int *, int *, int *, int);
    char *error;

    /* dynamically load the shared lib that contains addvec() */
    handle = dlopen("./libvector.so", RTLD_LAZY);
    if (!handle) {
        fprintf(stderr, "%s\n", dlerror());
        exit(1);
    }
```


Dynamic Linking at Run-time

```
    ...

    /* get a pointer to the addvec() function we just loaded */
    addvec = dlsym(handle, "addvec");
    if ((error = dlerror()) != NULL) {
        fprintf(stderr, "%s\n", error);
        exit(1);
    }

    /* Now we can call addvec() it just like any other function */
    addvec(x, y, z, 2);
    printf("z = [%d %d]\n", z[0], z[1]);

    /* unload the shared library */
    if (dlclose(handle) < 0) {
        fprintf(stderr, "%s\n", dlerror());
        exit(1);
    }
    return 0;
}
```


Some Interpositioning Applications

Security

- Confinement (sandboxing)
  - Autolab interposes calls to libc functions.
  - Running `/bin/xterm` from student code on server a no-no.
- Behind the scenes encryption
  - Automatically encrypt otherwise unencrypted network connections.

Monitoring and Profiling

- Count number of calls to functions
- Characterize call sites and arguments to functions
- Malloc tracing
  - Detecting memory leaks
  - Generating malloc traces

Case Study: Library Interpositioning

Library interpositioning is a powerful linking technique that allows programmers to intercept calls to arbitrary functions

Interpositioning can occur at:

- compile time
  - When the source code is compiled
- link time
  - When the relocatable object files are linked to form an executable object file
- load/run time
  - When an executable object file is loaded into memory, dynamically linked, and then executed.