# **Advanced Optimization**

(10-801: CMU)

Lecture 19
Parallel proximal; Incremental gradient

26 Mar, 2014

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Suvrit Sra

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$$\begin{array}{rcl} \operatorname{prox}_f + \operatorname{prox}_{f^*} & = & I \\ & 2\operatorname{prox}_f & = & 2I - 2\operatorname{prox}_{f^*} \end{array}$$

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$$2 \operatorname{prox}_f - I = I - 2 \operatorname{prox}_{f^*}$$

$$R_f = -R_{f^*}$$

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$$z \leftarrow \frac{1}{2}(I + R_f R_g R_h)z$$

$$\min f(x) + g(x) + h(x)$$

$$\begin{array}{rcl} 0 & \in & \partial f(x) + \partial g(x) + \partial h(x) \\ 3x & \in & (I+\partial f)(x) + (I+\partial g)(x) + (I+\partial h)(x) \\ 3x & \in & (I+\partial f)(x) + z + w \\ & \text{now what?} \end{array}$$

$$\min f(x) + g(x) + h(x)$$

#### Partial solution (Borwein, Tam (2013))

$$T_{hf} := \frac{1}{2}(I + R_f R_h)$$
$$T_{[fgh]} := T_{hf} T_{gh} T_{fg}$$
$$z \leftarrow T_{[fgh]} z$$

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- Works for more than 3 functions too!
- $\circ~$  For two functions  $T_{[fg]} = T_{gf}T_{fg}$
- o Does not coincide with usual DR.
- o Finding "correct" generalization an open problem

### Parallel proximal methods

#### Optimizing separable objective functions

$$f(x) := \frac{1}{2} ||x - y||_2^2 + \sum_i f_i(x)$$
  
 $f(x) := \sum_i f_i(x)$ 

### Parallel proximal methods

#### Optimizing separable objective functions

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Let us consider

$$\min \quad f(x) = \sum_{i=1}^{m} f_i(x), \qquad x \in \mathbb{R}^n.$$

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- ▶ Suppose we have  $\sum_{i=1}^{m} f_i(x)$

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- ▶ Now problem is over domain  $\mathcal{H}^m := \mathcal{H} \times \mathcal{H} \times \cdots \times \mathcal{H}$  (*m*-times)
- $\blacktriangleright$  New constraint:  $x_1 = x_2 = \ldots = x_m$

$$\min_{(x_1,\dots,x_m)} \quad \sum_i f_i(x_i)$$

s.t. 
$$x_1 = x_2 = \dots = x_m$$
.

Technique due to: G. Pierra (1976)

$$\min_{\boldsymbol{x}} f(\boldsymbol{x}) + \mathbb{I}_{\mathcal{B}}(\boldsymbol{x})$$

where 
$$\boldsymbol{x} \in \mathcal{H}^m$$
 and  $\mathcal{B} = \{ \boldsymbol{z} \in \mathcal{H}^m \mid \boldsymbol{z} = (x, x, \dots, x) \}$ 

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$$\blacktriangleright$$
 Let  $\boldsymbol{y}=(y_1,\ldots,y_m)$ 

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#### Two block problem

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- $\blacktriangleright$  Let  $\boldsymbol{y}=(y_1,\ldots,y_m)$
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$$\min_{\boldsymbol{z} \in \mathcal{B}} \quad \frac{1}{2} \|\boldsymbol{z} - \boldsymbol{y}\|_{2}^{2} 
\min_{\boldsymbol{x} \in \mathcal{H}} \quad \sum_{i} \frac{1}{2} \|\boldsymbol{x} - y_{i}\|_{2}^{2} 
\implies \quad \boldsymbol{x} = \frac{1}{m} \sum_{i} y_{i}$$

#### Two block problem

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$$\implies \quad x = \frac{1}{m} \sum_{i} y_{i}$$

**Exercise:** Work out the details of DR using the product space idea

This technique commonly exploited in ADMM too

$$\min_{x} \ \frac{1}{2} ||x - y||_{2}^{2} + f(x) + h(x)$$

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#### Proximal-Dykstra method

- 1 Let  $x_0 = y$ ;  $u_0 = 0$ ,  $z_0 = 0$
- **2** k-th iteration  $(k \ge 0)$

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Why does it work?

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#### Why does it work?

**Exercise:** Use the product-space technique to extend this to a parallel prox-Dykstra method for  $m \geq 3$  functions.

Combettes, Pesquet (2010); Bauschke, Combettes (2012)

# Proximal-Dykstra – some insight

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$$g(\nu,\mu) \quad := \quad \inf_{x,z,w} L(x,z,\nu,\mu)$$
 
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#### **Equivalent dual problem**

$$\min \quad G(\nu, \mu) := \frac{1}{2} \|\nu + \mu - y\|_2^2 + f^*(\nu) + h^*(\mu).$$

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- ▶  $0 \in \nu_{k+1} + \mu_k y + \partial f^*(\nu_{k+1})$
- $\blacktriangleright 0 \in \nu_{k+1} + \mu_{k+1} y + \partial h^*(\mu_{k+1}).$

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$$0 \in \nu_{k+1} + \mu_k - y + \partial f^*(\nu_{k+1}) \implies y - \mu_k \in \nu_{k+1} + \partial f^*(\nu_{k+1})$$

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$$0 \in \nu_{k+1} + \mu_k - y + \partial f^*(\nu_{k+1}) \Longrightarrow y - \mu_k \in \nu_{k+1} + \partial f^*(\nu_{k+1})$$
  
$$\Longrightarrow \nu_{k+1} = \operatorname{prox}_{f^*}(y - \mu_k) \Longrightarrow \nu_{k+1} = y - \mu_k - \operatorname{prox}_f(y - \mu_k)$$

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$$0 \in \nu_{k+1} + \mu_k - y + \partial f^*(\nu_{k+1}) \implies y - \mu_k \in \nu_{k+1} + \partial f^*(\nu_{k+1}) \\ \implies \nu_{k+1} = \operatorname{prox}_{f^*}(y - \mu_k) \implies \nu_{k+1} = y - \mu_k - \operatorname{prox}_f(y - \mu_k)$$
 Similarly,  $\mu_{k+1} = y - \nu_{k+1} - \operatorname{prox}_h(y - \nu_{k+1})$ 

▶ 
$$0 \in \nu_{k+1} + \mu_k - y + \partial f^*(\nu_{k+1})$$

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▶ 
$$0 \in \nu_{k+1} + \mu_k - y + \partial f^*(\nu_{k+1})$$

$$0 \in \nu_{k+1} + \mu_{k+1} - y + \partial h^*(\mu_{k+1}).$$

$$\nu_{k+1} = y - \mu_k - \text{prox}_f(y - \mu_k)$$
  

$$\mu_{k+1} = y - \nu_{k+1} - \text{prox}_h(y - \nu_{k+1})$$

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## Now use Lagrangian stationarity condition

$$x = y - \nu - \mu \implies y - \mu = x + \nu$$

to rewrite BCD using primal and dual variables.

▶ 
$$0 \in \nu_{k+1} + \mu_k - y + \partial f^*(\nu_{k+1})$$

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#### **BCD**

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$$\mu_{k+1} = \operatorname{argmin}_{\mu} \ G(\nu_{k+1}, \mu).$$

▶ 
$$0 \in \nu_{k+1} + \mu_k - y + \partial f^*(\nu_{k+1})$$

$$0 \in \nu_{k+1} + \mu_{k+1} - y + \partial h^*(\mu_{k+1}).$$

$$\nu_{k+1} = y - \mu_k - \text{prox}_f(y - \mu_k)$$

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## **Prox-Dykstra**

$$w_k \leftarrow \operatorname{prox}_f(x_k + \nu_k)$$
$$\nu_{k+1} \leftarrow x_k + \nu_k - w_k$$
$$x_{k+1} \leftarrow \operatorname{prox}_h(w_k + \mu_k)$$
$$\mu_{k+1} \leftarrow \mu_k + w_k - x_{k+1}$$

# **Example practical use**

## **Anisotropic 2D-TV Proximity operator**

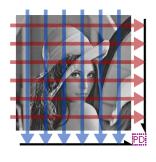
$$\min_{X} \quad \tfrac{1}{2} \|X - Y\|_{\mathsf{F}}^2 + \sum\nolimits_{ij} w_{ij}^c |x_{i, {\color{blue} j+1}} - x_{ij}| + \sum\nolimits_{ij} w_{ij}^r |x_{{\color{blue} i+1}, j} - x_{ij}|$$



## **Example practical use**

#### **Anisotropic 2D-TV Proximity operator**

$$\min_{X} \quad \tfrac{1}{2} \|X - Y\|_{\mathsf{F}}^2 + \sum\nolimits_{ij} w_{ij}^c |x_{i, {\color{blue} j+1}} - x_{ij}| + \sum\nolimits_{ij} w_{ij}^r |x_{{\color{blue} i+1}, j} - x_{ij}|$$



- Amenable to prox-Dykstra
- Used in (Barbero, Sra, ICML 2011).
- The subproblem:  $\min \frac{1}{2} ||a - b||_2^2 + \sum_i w_i |a_i - a_{i+1}|$ itself has been subject of over 15 papers!
- I still use it now and then



# Incremental first-order methods

# **Separable objectives**

$$\min \quad f(x) = \sum_{i=1}^{m} f_i(x) + \lambda r(x)$$

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## **Gradient / subgradient methods**

$$x_{k+1} = x_k - \alpha_k \nabla f(x_k) \qquad \lambda = 0,$$
  

$$x_{k+1} = x_k - \alpha_k g(x_k), \qquad g(x_k) \in \partial f(x_k) + \lambda \partial r(x_k)$$
  

$$x_{k+1} = \operatorname{prox}_{\alpha_k r}(x_k - \alpha_k \nabla f(x_k))$$

#### **Product-space based methods**

$$\min F(x_1, \dots, x_m) + \mathbb{I}_{\mathcal{B}}(x_1, \dots, x_m)$$
$$(x_{1,k+1}, \dots, x_{m,k+1}) \leftarrow \operatorname{prox}_F(y_{1,k}, \dots, y_{m,k})$$

# **Separable objectives**

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$$x_{k+1} = \operatorname{prox}_{\alpha_k r}(x_k - \alpha_k \nabla f(x_k))$$

#### Product-space based methods

$$\min F(x_1, \dots, x_m) + \mathbb{I}_{\mathcal{B}}(x_1, \dots, x_m)$$
$$(x_{1,k+1}, \dots, x_{m,k+1}) \leftarrow \operatorname{prox}_F(y_{1,k}, \dots, y_{m,k})$$

How much computation does one iteration take?

What if at iteration k, we randomly pick an integer  $i(k) \in \{1, 2, \dots, m\}$ ?

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And instead just perform the update?

$$x_{k+1} = x_k - \alpha_k \nabla f_{i(k)}(x_k)$$

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And instead just perform the update?

$$x_{k+1} = x_k - \alpha_k \nabla f_{i(k)}(x_k)$$

- lacktriangle The update requires only gradient for  $f_{i(k)}$
- ▶ One iteration now m times faster than with  $\nabla f(x)$

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- ▶ One iteration now m times faster than with  $\nabla f(x)$



But does this make sense?

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- ♠ Usually randomization greatly simplifies convergence analysis

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(Use: 
$$\sum_i a_i b_i = \sum_i a_i^2 (b_i/a_i)$$
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- $\blacktriangleright$  But once inside region R, no guarantee that incremental method will make progress towards optimum.

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Convergence rate analysis?

Fermat-Weber problem (historically the first facility-location problem)

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- ▶ Assuming  $\|\cdot\| = \|\cdot\|_2$
- $\blacktriangleright$  Also assume no  $a_i$  is an optimum
- ▶ [Weiszfeld; '37] Let  $T:=x\mapsto \left(\sum_i \frac{w_i a_i}{\|x-a_i\|}\right)/\left(\sum_i \frac{w_i}{\|x-a_i\|}\right)$
- ▶ Assuming T is well-defined,  $T^k(x_0) \to \operatorname{argmin}$
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- $\blacktriangleright \text{ What if } \|\cdot\| = \|\cdot\|_p?$
- ▶ 100s of papers discuss the Fermat-Weber problem

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**Exercise:** Obtain closed form solution to  $x_{k+1}$ 

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Rate of convergence? Most likely, sublinear? Can we somehow get linear convergence?

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Moreover, analysis easier if we go through the  $f_i$  randomly (so-called stochastic)

### Incremental methods: deterministic

$$\min \quad (f(x) = \sum_{i} f_i(x)) + r(x)$$

#### Gradient with error

$$\nabla f_{i(k)}(x) = \nabla f(x) + \frac{e}{e}$$
$$x_{k+1} = \operatorname{prox}_{\alpha r} [x_k - \alpha_k (\nabla f(x_k) + \frac{e_k}{e})]$$

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#### **Gradient with error**

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So if in the limit error  $\alpha_k e_k$  disappears, we should be ok!

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### Some stepsize choices

- $\spadesuit$   $\alpha_k = c$ , a small enough constant
- $\spadesuit$   $\alpha_k \to 0$ ,  $\sum_k \alpha_k = \infty$  (diminishing scalar)
- Constant for some iterations, diminish, again constant, repeat

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- ♠ Some extend to parallel and distributed computation

### References

- ♠ EE227A slides, S. Sra
- ♠ Introductory Lectures on Convex Optimization, Yu. Nesterov
- Proximal splitting methods, Combettes & Pesquet