Thesis Oral

On the Communication Complexity of Classical Correlation Distillation and

Quantum Entanglement Distillation

Ke Yang

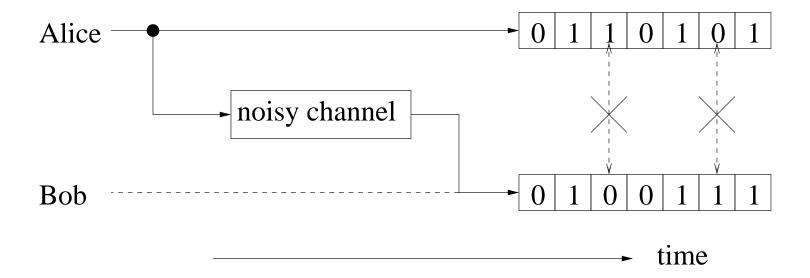
Thesis Committee
Steven Rudich (chair)
Avrim Blum
Robert Griffiths
Andris Ambainis (IAS)

On Repairing Corrupted Correlation

Recurring Theme in Information Theory

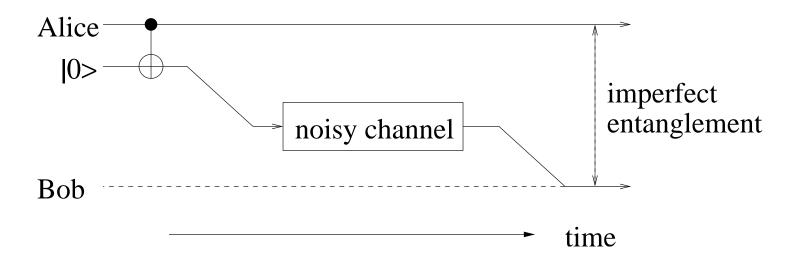
- Correlation Corruption
 Alice and Bob share imperfectly correlated information
- Correlation Recovery
 Alice and Bob take action to recover perfect correlation

Classical Noisy Channel



- Alice sends bits to Bob
- Correlation corruption by the noisy channel

Quantum Noisy Channel



- Alice sends qubits to Bob
- Entanglement corruption by the noisy channel

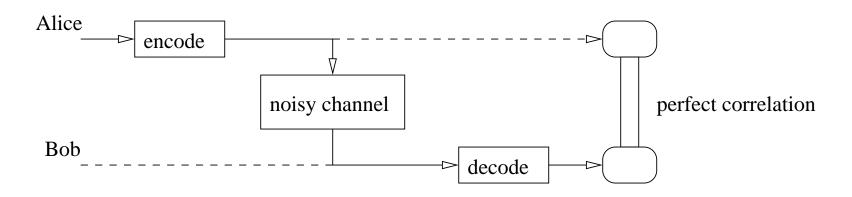
"Correlation" Overloading

- classical::correlation = correlation
- quantum::correlation = entanglement

Strategies for Correlation Recovery

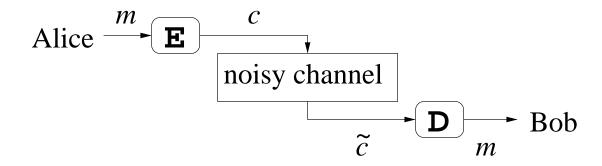
- Preventive Strategy
 Adding redundancy before the corruption
- Reparative Strategy
 Recovering correlation only after corruption

Preventive Strategy



- Information encoded before the corruption
- Error Correcting Codes (ECCs)
- Quantum Error Correcting Codes (QECCs)

Error Correcting Codes



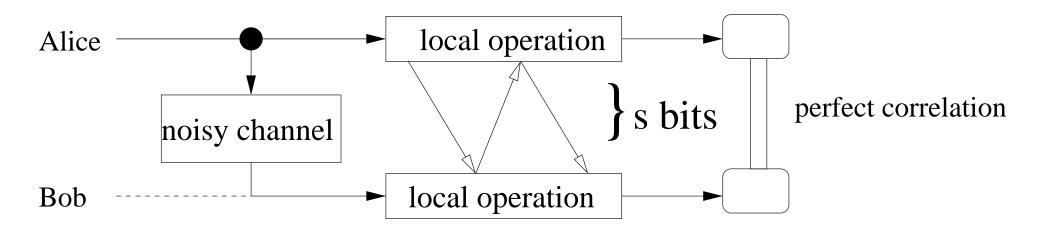
• (n, k, d)-ECC: $\{0, 1\}^k \mapsto \{0, 1\}^n$, such that

$$\mathsf{DIST}(E(m_1), E(m_2)) \geq d$$

- Code Overhead: (n-k) bits
- Noise Tolerance: $\leq (d-1)/2$ bit flips

(encoding/decoding complexity not our focus)

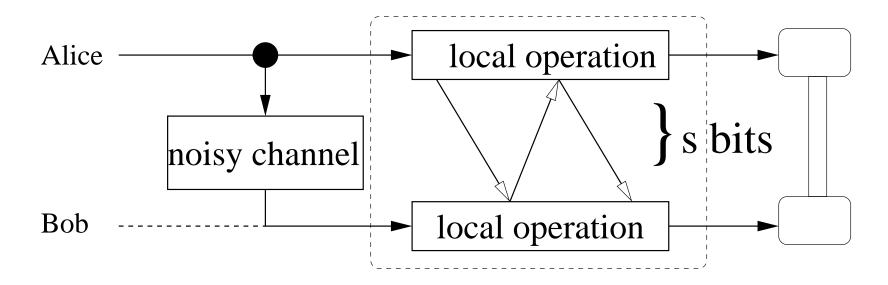
Reparative Strategy



- Correlation repaired after the corruption
- Alice and Bob exchange s bits to recover the correlation
 - ASSUMPTION: noiseless classical communication
 - GOAL: minimize s

(computational complexity not our focus)

Correlation Distillation

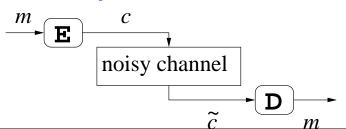


- Classical Correlation Distillation Protocol (CDP)
- Quantum Entanglement Distillation Protocol (EDP)

Information Transmission

Alice wishes to transmit m to Bob, noiselessly

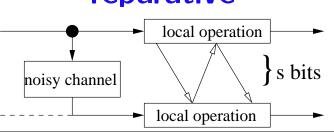




- 1. Encoding: c = E(m)
- 2. Transmission: $c \rightarrow \tilde{c}$
- 3. Decoding: $m = D(\tilde{c})$

Overhead = |c| - |m|

reparative

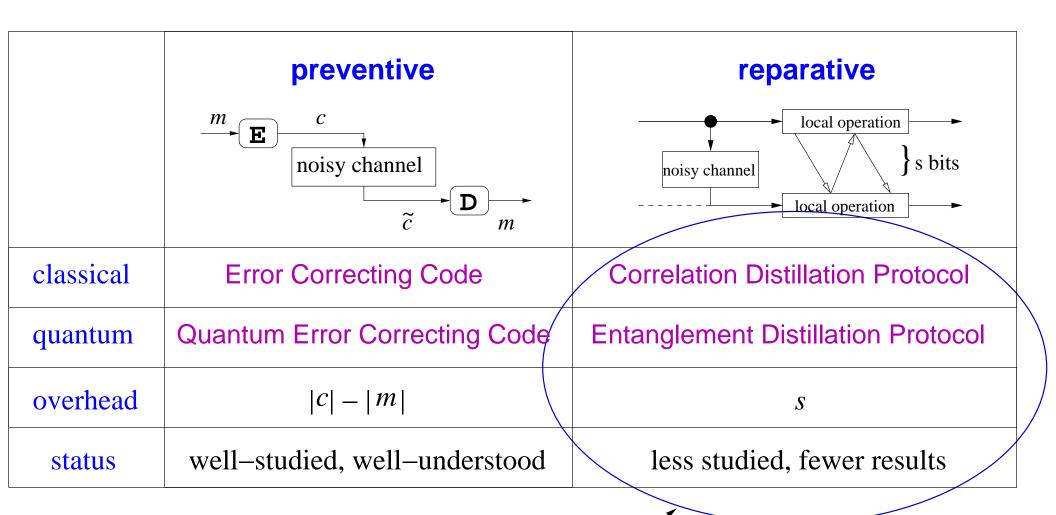


- 1. Transmission: $m \to \tilde{m}$
- 2. Distillation:

$$(m, \tilde{m}) \xrightarrow{\mathcal{P}} (m, m)$$

Overhead = s

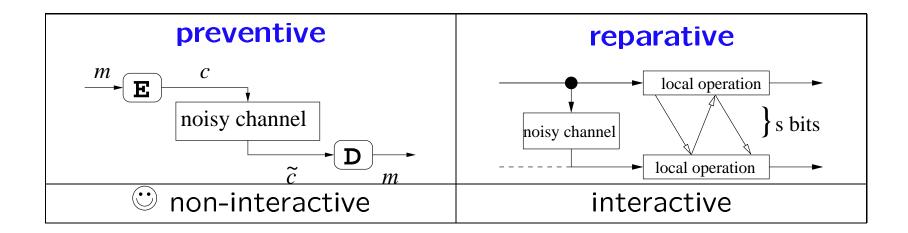
	preventive	reparative
	$ \begin{array}{c c} & m & c \\ \hline & noisy channel \\ \hline & \tilde{c} & m \end{array} $	local operation noisy channel local operation local operation
classical	Error Correcting Code	Correlation Distillation Protocol
quantum	Quantum Error Correcting Code	Entanglement Distillation Protocol
overhead	c - m	S
status	well-studied, well-understood	less studied, fewer results



My thesis

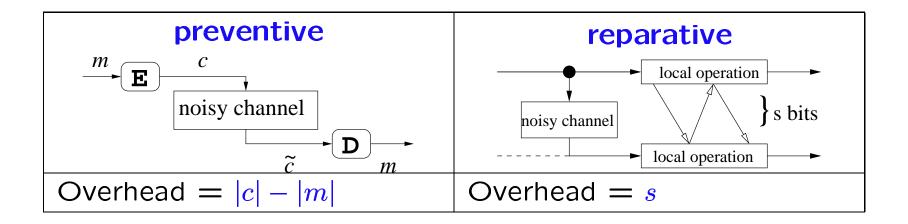
why?

Error Correction is Great!



"An ounce of prevention is worth a pound of cure."

"An Ounce of Prevention is Worth a Pound of Cure."



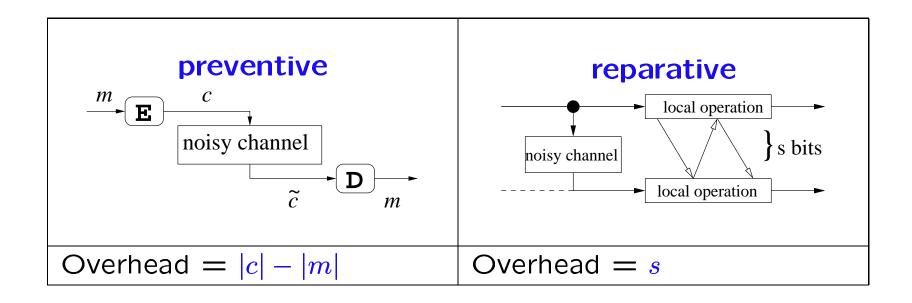
same level of corruption, $16\times$ more efficient?

Not Necessarily

Correlation distillation is ...

- 1. as efficient as error correction
- 2. applicable to a wider range of applications

Information Transmission

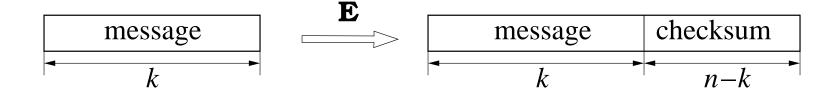


THM (n, k, d)-linear ECC \Rightarrow CDP of overhead s = (n - k)

THM (n, k, d)-stabilizer QECC \Rightarrow EDP of overhead s = (n - k)

Proof

THM (n, k, d)-linear ECC \Rightarrow (n - k)-bit CDP



PROOF

- 1. Alice sends the (n-k)-bit check-sum
- 2. Bob decodes

"An ounce of prevention is worth a pound of cure."

an ounce

Correlation Distillation Beats ECCs

THM Correlation distillation is provably more powerful than ECCs

∃ noisy channel, s.t.

- No ECC can achieve a non-trivial rate.
- But Correlation Distillation Protocols can

Entanglement Distillation Beats QECCs

[Bennett, Di Vincenzo, Smolin, Wootters 1996]

Entanglement Distillation is provably more powerful than QECCs

∃ noisy channel, s.t.

- No QECC can work
- But Entanglement Distillation Protocols can

"An ounce of prevention is worth a pound of cure."

"In a corrupted world, prevention is useless, yet there is cure."

Correlation Distillation has More Applications

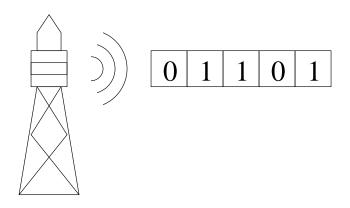
Assumptions made by error correction —

Preventive encoding must precede the noise "What if encoding is impossible?"

Noise model identical independent noise, known noise rate "What if the noise model is different?"

Have to guess an upper bound on noise rate

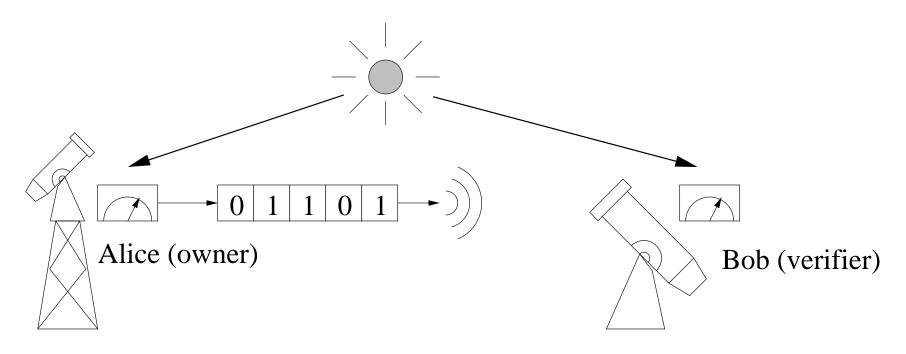
Random Beacon



A real-time, verifiable random source

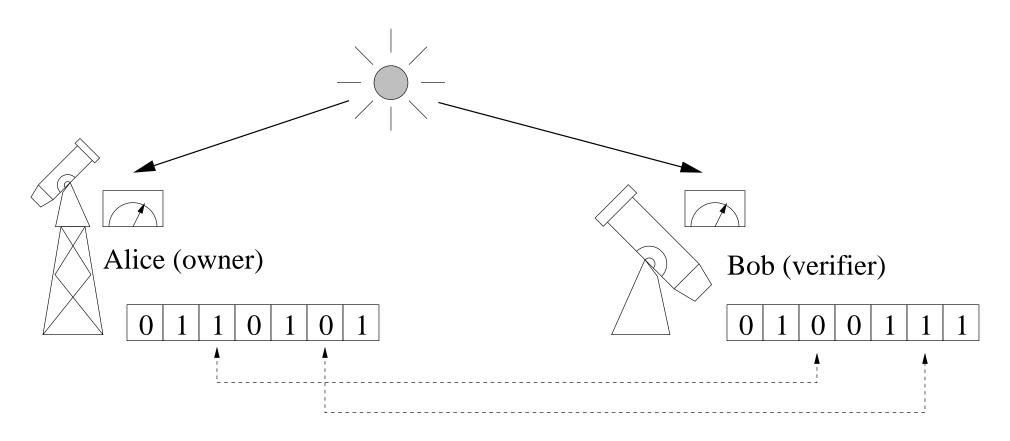
- verifiable lottery
- information-theoretically secure cryptography key-exchange, encryption... (assuming bounded storage)

How to Build a Random Beacon



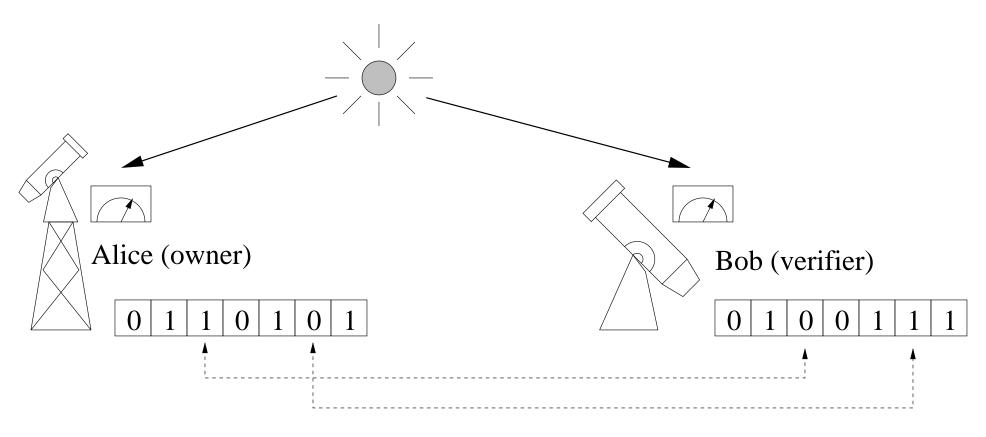
- Point a telescope to a pulsar
- Measure the signal, convert to random bits
- Real-time verifiable: (almost) everyone can see the pulsar

Noisy Measurement



Measurement errors — corrupted correlation

Correlation Recovery for Random Beacons

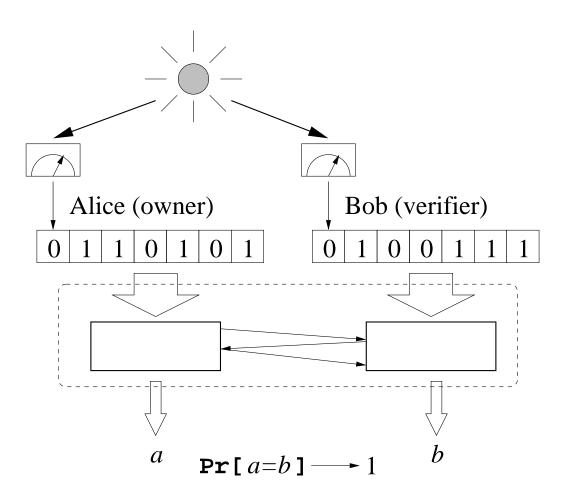


GOAL = to achieve (almost) perfect correlation

Error Correction on a Pulsar ?!

- Both Alice and Bob have corrupted information
- Preventive strategy doesn't work
- Okay to produce "fresh" random bits

Correlation Distillation for Random Beacon

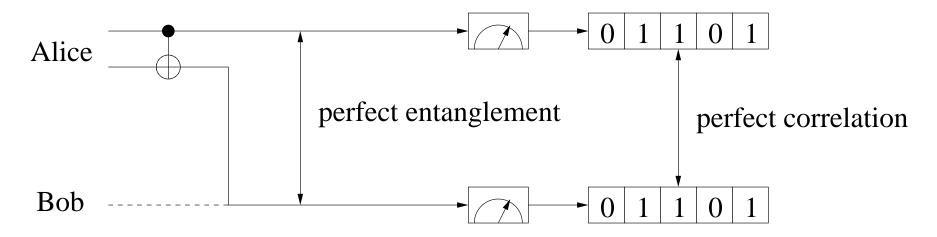


Random Beacon: error correction doesn't apply

Storing EPR Pairs

- EPR pairs are useful quantum objects, but hard to store
- Constantly decaying varying noise rate
- QECC has to guess an upper bound of noise rate

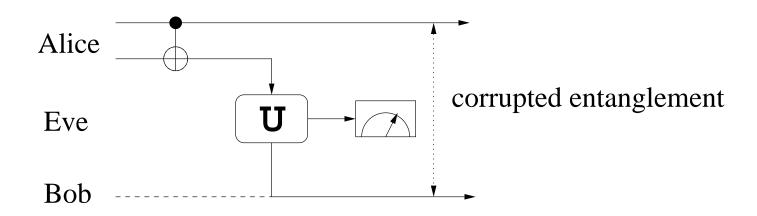
Quantum Key Distribution (Ideal)



[Bennett-Brassard 84, Bennett 92] (modified)

- Alice sends random qubits to Bob and keeps a copy herself
- (Ideally) perfectly entangled qubits
- Both measure ⇒ (Ideally) perfectly correlated bits

Quantum Key Distribution (Real life)



- Eve intercepts some qubits and distorts them
- ◆ corrupted entanglement ⇒ corrupted correlation

Error Correction for Eve?

QECC assumes identical independent noise but...

Eve is adversarial

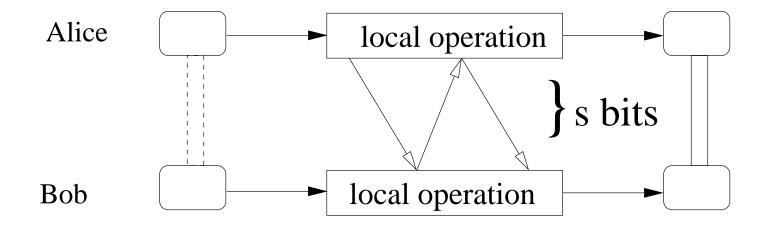
Quantum Key Distribution: error correction uses a different model Carnegie Mellon 37

Why Reparative?

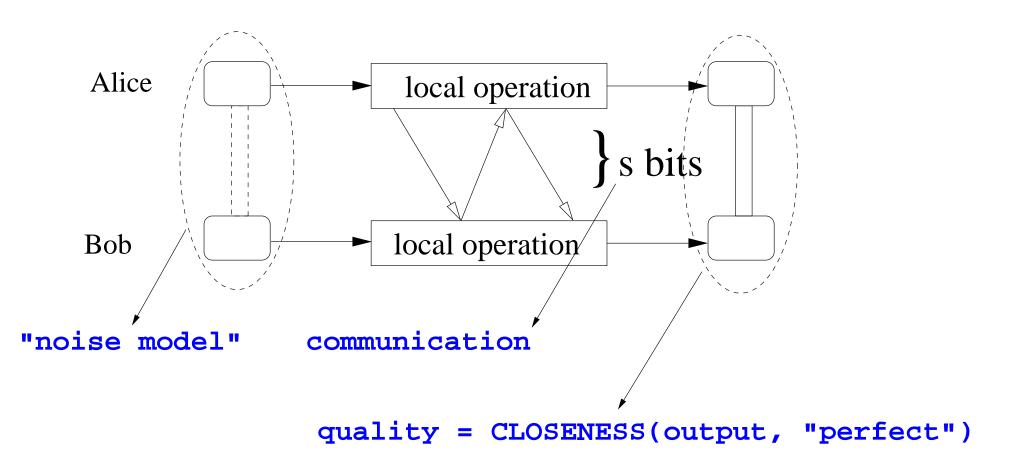
Scenario	Reason
Information Transmission	Correlation distillation is as efficient as error correction (and can be more useful)
Random Beacon	ECCs don't apply (can't error correct a pulsar)
Storing EPR pairs	QECCs are inefficient (varying noise rate)
Quantum Key Distribution	QECCs don't apply (different noise models)

What's known?

Quantifying Distillation Protocols



Fix Noise Model, Study Communication vs. Quality



communication

noise model	0	1	many	quality
bounded corruption				C
binary symmetric				SO CO
binary erasure				្ត ក- ន
tensor product				a L
bounded corruption				
bounded measurement				nr.
depolarization				quantum
entanglement				B
fidelity				

communication

noise model

bounded corruption binary symmetric binary erasure

tensor product

bounded corruption

bounded measurement

depolarization

entanglement

fidelity

0	1	many
		L
·· u	© L	L
⊕ U		L
υ υ		L
⊕ U		L
∵ ∪		L
· u	⊕ u	U 🙂
⊕ L U	⊕ L U	⊕ L U

L = lower bound

U = upper bound

○ = my orignal result

○ = independent result

quantum

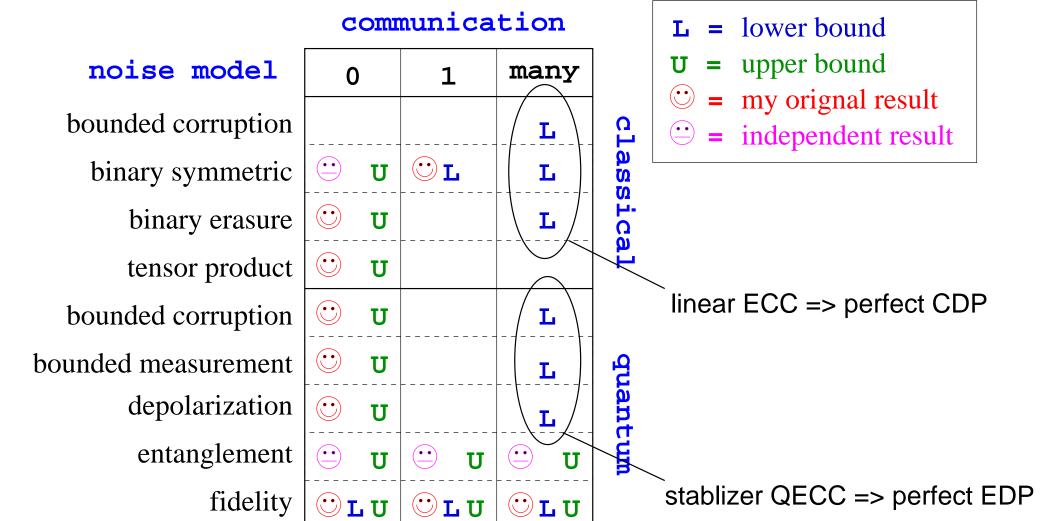
classical

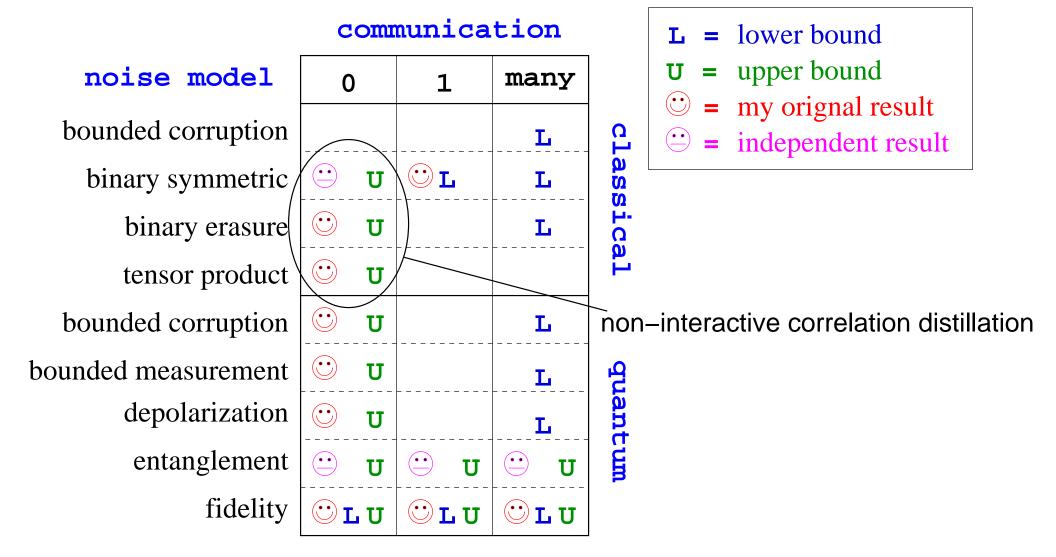
Carnegie Mellon

43

Related Publications

- [Ambainis, Smith, Yang 2002] "Extracting Quantum Entanglement (General Entanglement Purification Protocols)", IEEE Conference on Computational Complexity 2002.
- [Yang 2004] "On the (Im)possibility of Non-interactive Correlation Distillation", Latin American Theoretical Informatics (LATIN 2004).
- [Ambainis, Yang 2004] "Towards the Classical Communication Complexity of Entanglement Distillation Protocols with Incomplete Information", *IEEE Conference of Computational Complexity (CCC 2004)*.





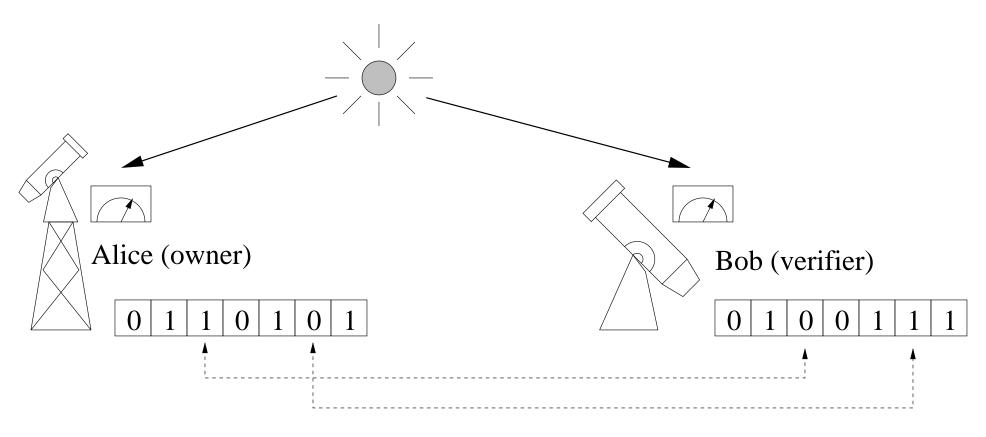
Carnegie Mellon

46

Non-interactive Correlation Distillation

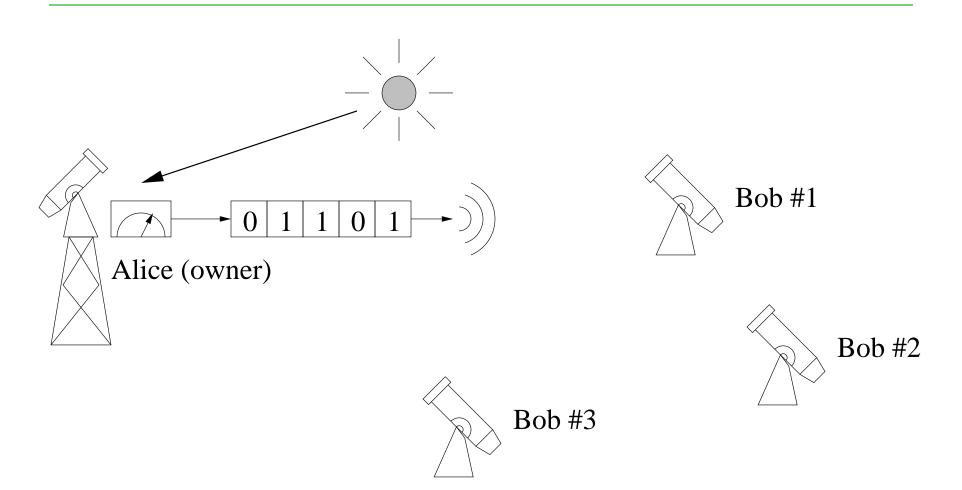
Alice and Bob distill correlation without communicating

Correlation Recovery for Random Beacons

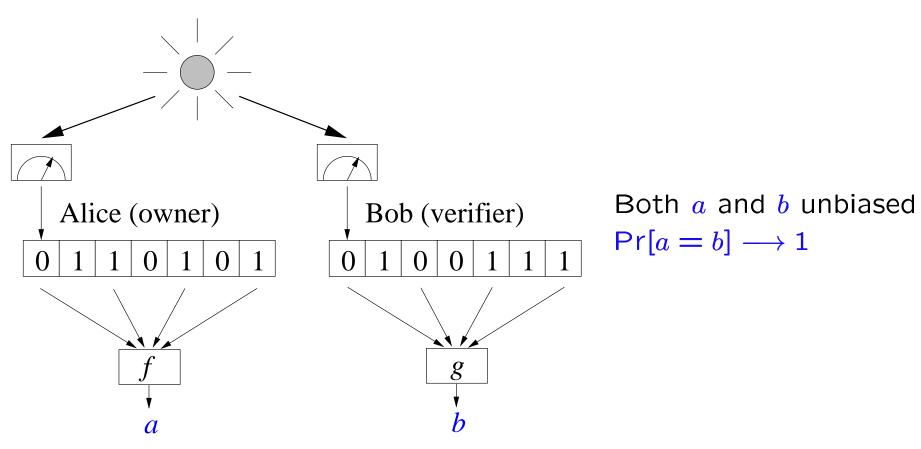


GOAL = to achieve (almost) perfect correlation

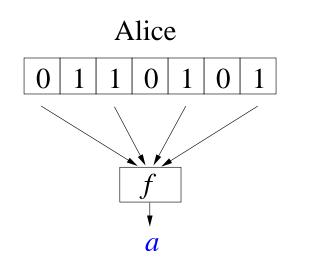
One Alice, Many Bobs

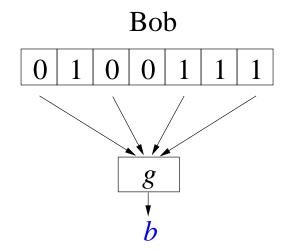


Non-Interactive Correlation Distillation for Random Beacon



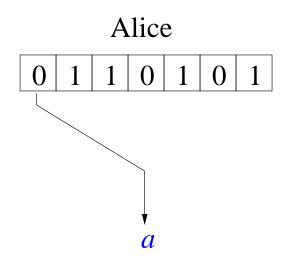
Correlation Extraction, Mathematically

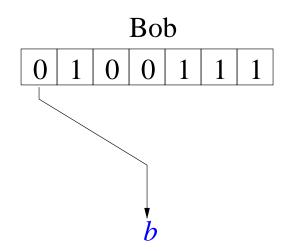




- Alice $x_1, x_2, ..., x_n$, Bob $y_1, y_2, ..., y_n$, s.t. $\Pr[x_k = y_k] = 1 p$
- Alice $a = f(x_1, x_2, ..., x_n)$; Bob $b = g(y_1, y_2, ..., y_n)$
- Unbiased bits Pr[a = 0] = 1/2, Pr[b = 0] = 1/2
- Maximize Pr[a = b]

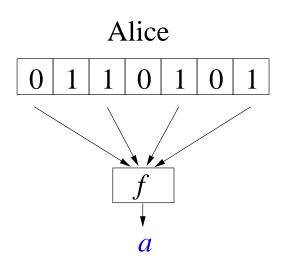
Naïve Strategy

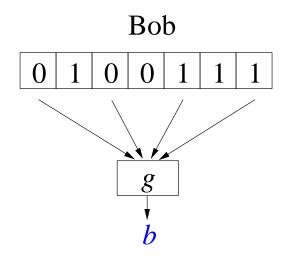




- Both output the first bit
- Pr[a = b] = 1 p

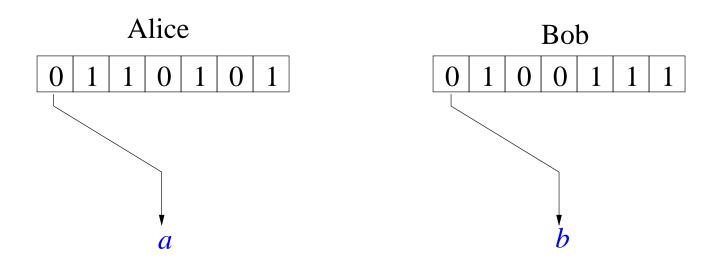
Can We do Better?





- Alice $x_1, x_2, ..., x_7$, Bob $y_1, y_2, ..., y_7$, $\Pr[x_k = y_k] = 0.9$
- Can $\Pr[a = b] \ge 0.91$? (mutual information = 3.72)

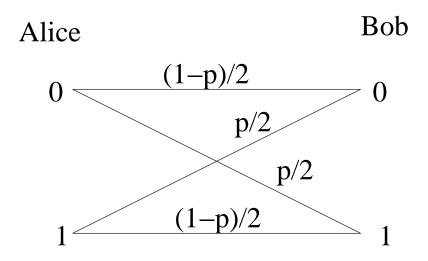
No



[Alon, Maurer, Wigderson], [Mossel, O'Donnell], [Yang 2004]

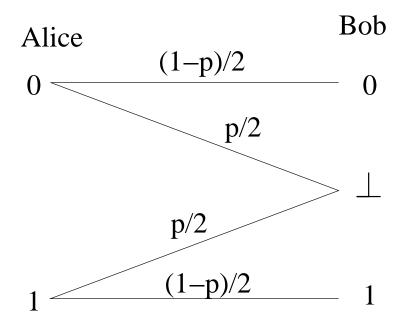
- The naïve strategy is optimal
- All optimal strategies are naïve

Binary Symmetric Model

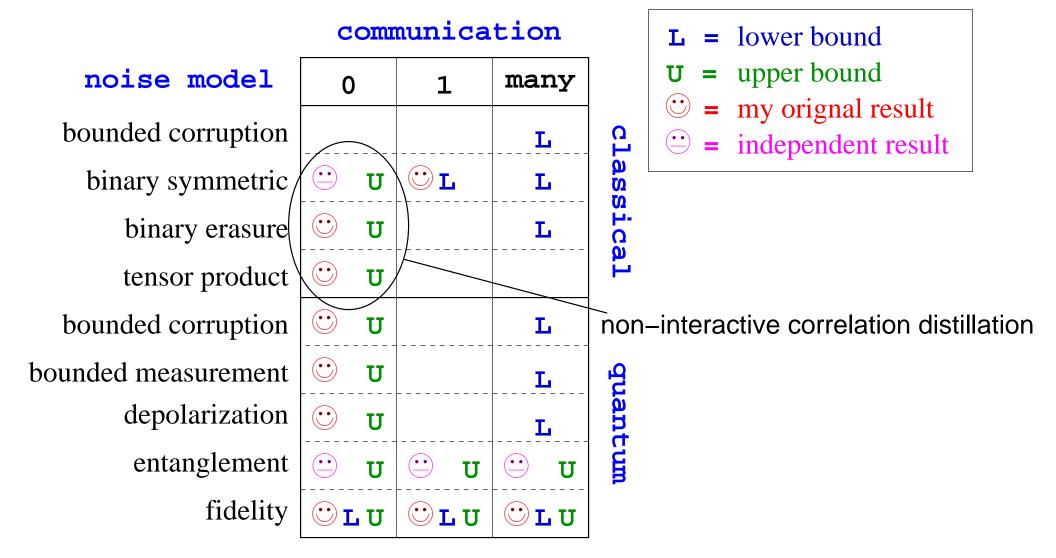


[Yang 2004] generalization to Tensor Product Model (large alphabet, more general noise)

Binary Erasure Model



[Yang 2004] The naïve strategy is asymptotically optimal



Carnegie Mellon

57

communication

noise model	0	1	many
bounded corruption			L
binary symmetric	∵ ʊ	□ L	L
binary erasure	υ σ		L
tensor product	υ σ		
bounded corruption	υ σ		L
bounded measurement	υ σ		L
depolarization	U U		L
entanglement	U :	\odot \mathbf{v}	U U
fidelity	⊕ L U	Ů L U	⊕ L U

L = lower bound

U = upper bound

assical

○ = my orignal result

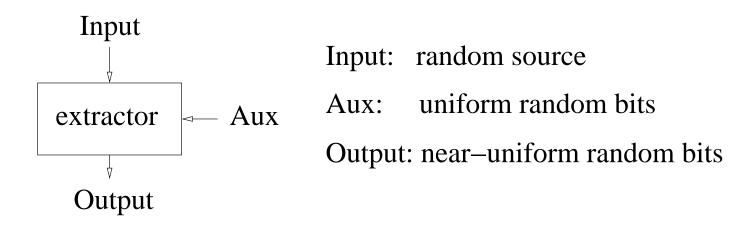
impossibility for general EPR extraction

Carnegie Mellon

58

Motivation: classical randomness extraction

Randomness Extractors



produce near-uniform random bits from arbitrary random sources

Facts About Extractors

Very useful, works with very general input

- input = arbitrary random source.
- output ← min-entropy(input)
- $|auxiliary input| = \Theta(log(|input|))$
- [Ta-Shma, Umans, Zuckerman 2001] Near-optimal constructions exist

"General Entanglement Distillation?"

classical	quantum		
uniform bits	EPR pairs		
randomness in purest form	entanglement in purest form		
extractor	entanglement distillation		
low-quality randomness	low-quality entanglement		
↓	↓		
high-quality randomness	high-quality entanglement		
input	input		
arbitrary random bits	arbitrary entangled state?		

No

THM General entanglement distillation is impossible (no protocol extracts EPR pairs from arbitrary entangled states)

Proof Sketch

classical unique distribution of max entropy

quantum infinitely many maximally entangled states

The 4 Bell states:

$$\Phi^{+} = \frac{1}{\sqrt{2}} (|0\rangle^{A}|0\rangle^{B} + |1\rangle^{A}|1\rangle^{B})$$

$$\Phi^{-} = \frac{1}{\sqrt{2}} (|0\rangle^{A}|0\rangle^{B} - |1\rangle^{A}|1\rangle^{B})$$

$$\Psi^{+} = \frac{1}{\sqrt{2}} (|0\rangle^{A}|1\rangle^{B} + |1\rangle^{A}|0\rangle^{B})$$

$$\Psi^{-} = \frac{1}{\sqrt{2}} (|0\rangle^{A}|1\rangle^{B} - |1\rangle^{A}|0\rangle^{B})$$

Proof Sketch, cont'd

Suppose there exists such a protocol \mathcal{P} , s.t.,

$$\mathcal{P}(\Phi^+) \rightarrow \Phi^+, \ \mathcal{P}(\Phi^-) \rightarrow \Phi^+, \ \mathcal{P}(\Psi^+) \rightarrow \Phi^+, \ \mathcal{P}(\Psi^-) \rightarrow \Phi^+$$

Let ρ be a mixed state:

$$\rho = \frac{1}{4} \left(|\Phi^{+}\rangle\!\langle \Phi^{+}| + |\Phi^{-}\rangle\!\langle \Phi^{-}| + |\Psi^{+}\rangle\!\langle \Psi^{+}| + |\Psi^{-}\rangle\!\langle \Psi^{-}| \right)$$

We should also have:

$$\mathcal{P}(\rho) \to \Phi^+$$

Change of Basis

$$\rho = \frac{1}{4} \left(|\Phi^{+}\rangle\langle \Phi^{+}| + |\Phi^{-}\rangle\langle \Phi^{-}| + |\Psi^{+}\rangle\langle \Psi^{+}| + |\Psi^{-}\rangle\langle \Psi^{-}| \right)$$

By changing of basis:

$$\rho = \frac{1}{4} (|00\rangle\langle00| + |01\rangle\langle01| + |10\rangle\langle10| + |11\rangle\langle11|)$$

 ρ is disentangled \Rightarrow impossible to produce EPR pairs $\Rightarrow \Leftarrow$

communication

noise model	0	1	many
bounded corruption			L
binary symmetric	⊕ u	© L	L
binary erasure	U U		L
tensor product	· u		
bounded corruption	U U		L
bounded measurement	U U		L
depolarization	U U		L
entanglement	⊕ U	⊕ u	⊕ U
fidelity	⊕ L U	⊕ L U	⊕ L U

L = lower bound

U = upper bound

classical

○ = my orignal result

impossibility for general EPR extraction

Carnegie Mellon

67

Why General Entanglement Extraction Fails?

- No protocol can do well on average
- Useful protocol only if input is "close" to some state

The Fidelity Noise Model

[Ambainis, Smith, Yang 2002]

fidelity(input, "perfect") $\geq 1 - \epsilon$

[Lo, Chau 1999], [Shor, Preskill 2000] used it in proof of security of [BB84] key distribution protocol

communication

noise model many 0 1 bounded corruption \mathbf{L} ··· L binary symmetric U \mathbf{L} binary erasure U L tensor product U bounded corruption U L bounded measurement U L depolarization U L entanglement U U U fidelity ⊕ L U ⊕ L U ⊕ L U L = lower bound

u = upper bound

36

sical

quantum

○ = my orignal result

= independent result

matching lower/upper bounds

Carnegie Mellon

70

Lower Bound: a Construction

[Ambainis, Smith, Yang 2002]

 $\forall n, s, \exists s$ -bit protocol, on n qubit pairs of fidelity $1 - \epsilon$, either:

- \bullet fails with probability ϵ (nothing is output), or
- outputs (n-s) pairs of qubits of fidelity $1-\frac{2^{-s}}{(1-\epsilon)}$

(output fidelity = output quality)

- + Can increase the fidelity as close to 1 as possible, sacrificing logarithmic number of qubit pairs and using logarithmic bit of communication
- **–** Fails with probability ϵ .

Failure is Unavoidable

[Ambainis, Smith, Yang 2002]

 \exists n qubit pairs in state ρ of fidelity $1-\epsilon$, s.t. any protocol taking ρ as input and outputting m qubit pairs, has average fidelity at most $1-\frac{1-2^{-m}}{1-2^{-n}}\epsilon \approx 1-\epsilon$.

Cannot increase the overall fidelity

Optimality of Our Construction

[Ambainis, Smith, Yang 2002]

 $\forall n, s, \exists s$ -bit protocol, on n qubit pairs of fidelity $1 - \epsilon$, either:

- \bullet fails with probability ϵ (nothing is output), or
- outputs (n-s) pairs of qubits of fidelity $1-\frac{2^{-s}}{(1-\epsilon)}$

Optimal...

- ullet Failure Probability Must fail with probability ϵ in order to achieve close-to-one "lucky fidelity"
- Yield (n-s) qubit pairs, asymptotically optimal

More Optimality

[Ambainis, Smith, Yang 2002]

 $\forall n, s, \exists s$ -bit protocol, on n qubit pairs of fidelity $1 - \epsilon$, either:

- \bullet fails with probability ϵ (nothing is output), or
- outputs (n-s) pairs of qubits of fidelity $1-\frac{2^{-s}}{(1-\epsilon)}$

[Ambainis, Yang 2004] ♥

Communication complexity optimal up to an additive constant

A Bit More Technically...

Analysis of general two-party protocols prior to [Ambainis, Yang 2004]

[Nielsen 1999] "Simulation-based Reduction"

- For pure state input, Alice can "simulate" Bob's actions
- Arbitrary protocol → single-message protocol

(Alice measures; Alice sends message to Bob; Bob measures)

Simulation-based Reduction

"reducing any protocol to a single-message protocols"

- Does not work for protocols with mixed states as input
- [Bennett, Di Vincenzo, Smolin, Wootters 1996]
 Two-way protocols more powerful than one-way protocols
- Reduction doesn't work!
- Other techniques do not seem to work with mixed states either (e.g [Hayden, Winter 2002])

Our Contribution

[Ambainis, Yang 2004]

Novel technique for mixed states and two-way protocols

- Keep track of the local density matrices of Alice and Bob
- Communication causes a density matrix to "split"
- Maintain an invariant with communication history

communication

SSB

sical

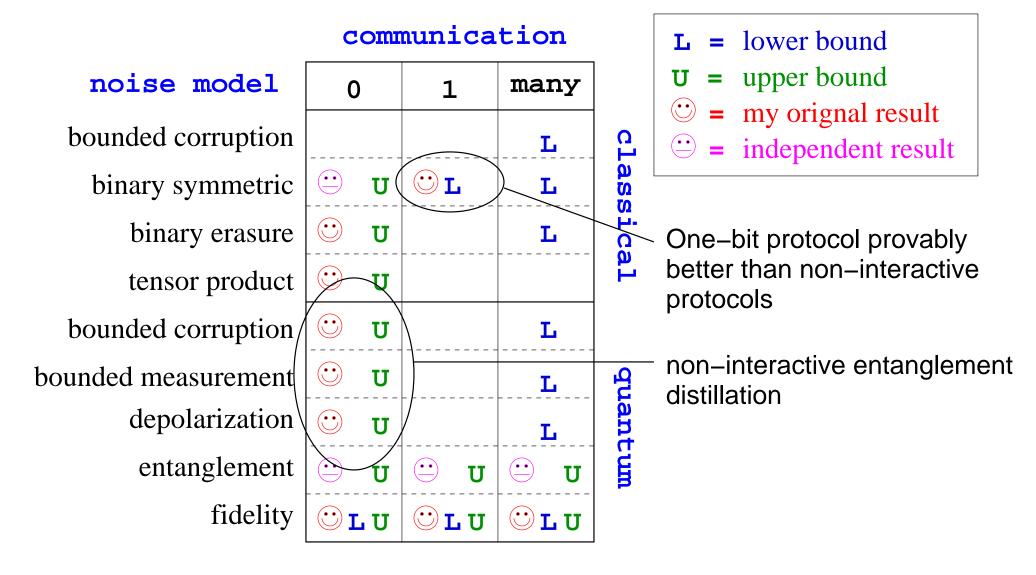
quantum

noise model many 0 1 bounded corruption L binary symmetric U L binary erasure U L tensor product U bounded corruption U L U bounded measurement L depolarization U L entanglement U U U fidelity **∵L** υ \odot L σ ⊕ L U independent result

matching lower/upper bounds

Carnegie Mellon

78



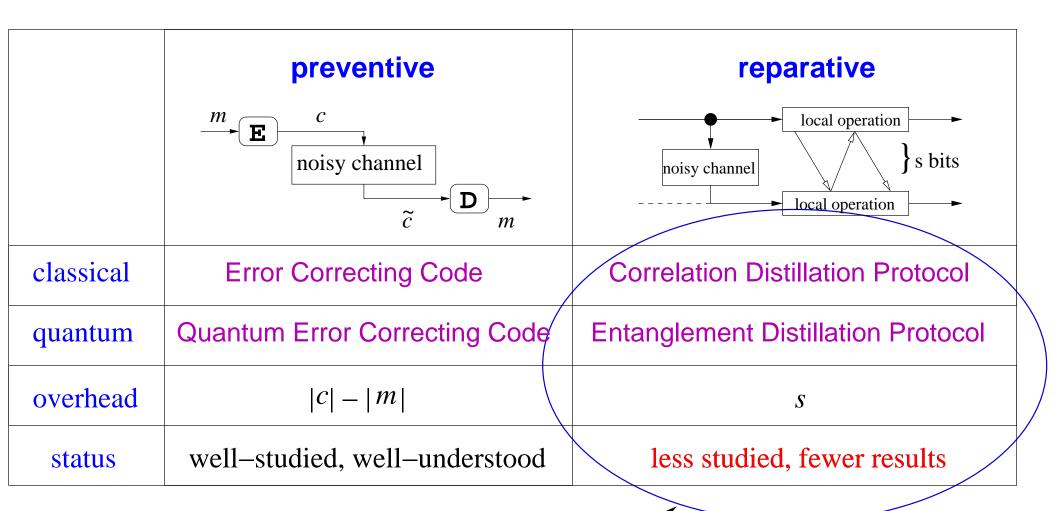
Summary

- Reparative: Correlation/Entanglement Distillation Protocols
- CDP/EDPs as efficient and ECC/QECC, maybe more
- Wider applications
- Results:
 - Impossibility of NICD/NIED
 - Impossibility of general EPR extraction
 - Optimal protocl for fidelity model
 - One-bit protocol for binary symmetric model

Thanks!

Questions?

What's next?



My thesis

communication

noise model

bounded corruption
binary symmetric
binary erasure
tensor product
bounded corruption
bounded measurement
depolarization
entanglement

0	1	many
		L
·· u	CL	L
· u		L
· u		
· u		L
· u		L
· u		L
··· u	U 🙂	∵ u
⊕ L U	⊕ L U	⊕ L U

L = lower bound

U = upper bound

: w = my orignal result

○ = independent result

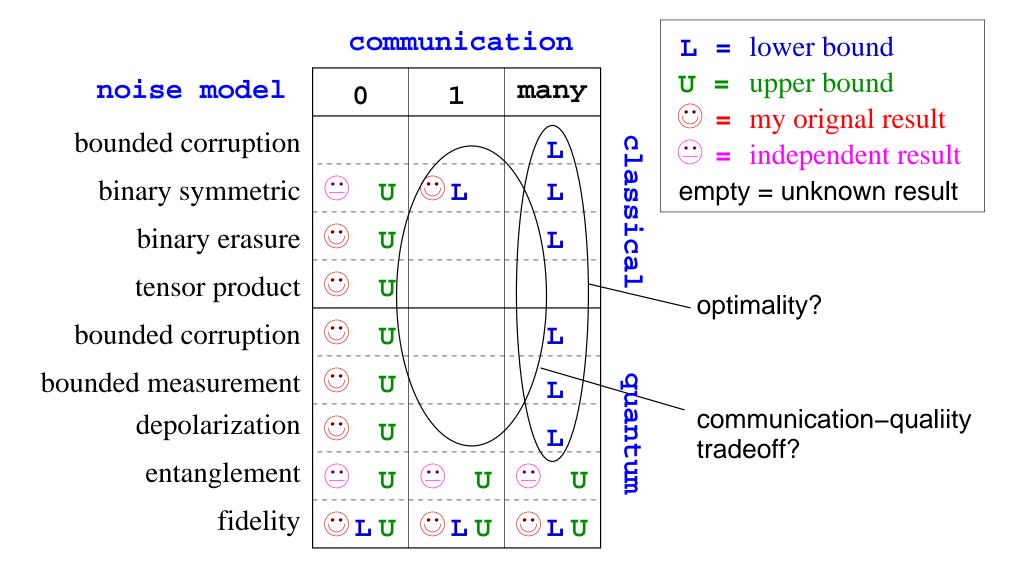
lantum

classical

Carnegie Mellon

fidelity

84



Big Questions

Optimality of constructions

"Linear ECC \Rightarrow CDP, Stabilizer QECC \Rightarrow EDP, are they optimal?"

More Trade-off on interactive correlation distillation

"What's the optimal quality Alice and Bob can get with s bits of communication?"

Unified results

"Are there noise models more general than, say, the fidelity model?"

"Can we merge the results to make the table smaller?"

More Immediate Questions: one-bit Protocols

- Can we upper bound the quality of one-bit CDP/EDPs?
- Is the protocol with the binary symmetric model optimal?

Time-line

[2003/3 — 2004/3] Continue research

[2004/4 — 2004/9] Write thesis