

Trees

6B

Heaps & Other Trees



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Heap



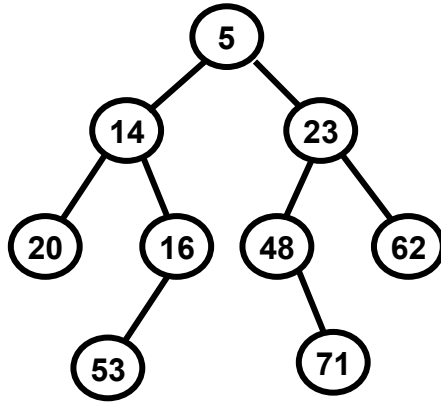
- A min-heap is a binary tree such that
 - the data contained in each node is less than (or equal to) the data in that node's children.
 - the binary tree is complete

- A max-heap is a binary tree such that
 - the data contained in each node is greater than (or equal to) the data in that node's children.
 - the binary tree is complete

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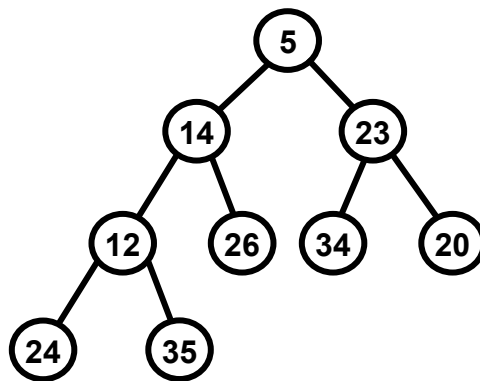
Is it a min-heap?



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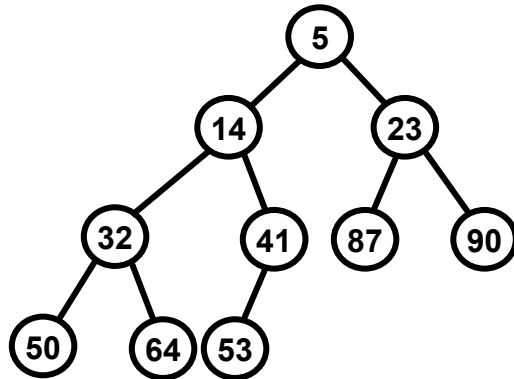
Is it a min-heap?



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Is it a min-heap?



Using heaps



What are min-heaps good for?
(What operation is extremely fast when using a min-heap?)

The difference in level between any two leaves in a heap is at most what?

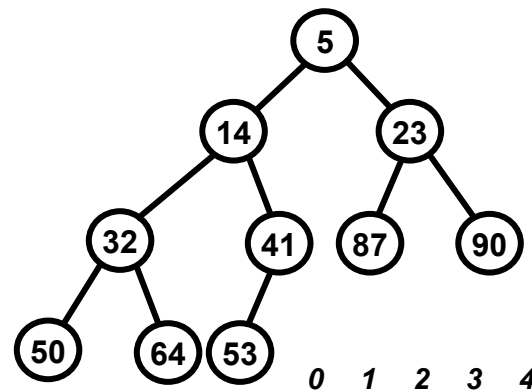


Storage of a heap

- Use an array to hold the data.
- Store the root in position 1.
 - We won't use index 0 for this implementation.
- For any node in position i ,
 - its left child (if any) is in position $2i$
 - its right child (if any) is in position $2i + 1$
 - its parent (if any) is in position $i/2$
(use integer division)



Storage of a heap



For node at i :
Left child is at $2i$
Right child is at $2i+1$
Parent is at $\lfloor i/2 \rfloor$

0	1	2	3	4	5	6	7	8	9	10
	5	14	23	32	41	87	90	50	64	53

Inserting into a min-heap



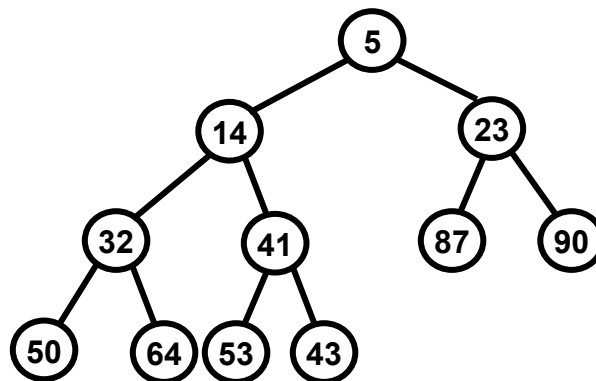
- Place the new element in the next available position in the array.
- Compare the new element with its parent. If the new element is smaller, than swap it with its parent.
- Continue this process until either
 - the new element's parent is smaller than or equal to the new element, or
 - the new element reaches the root (index 0 of the array)

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Inserting into a min-heap

Insert 43

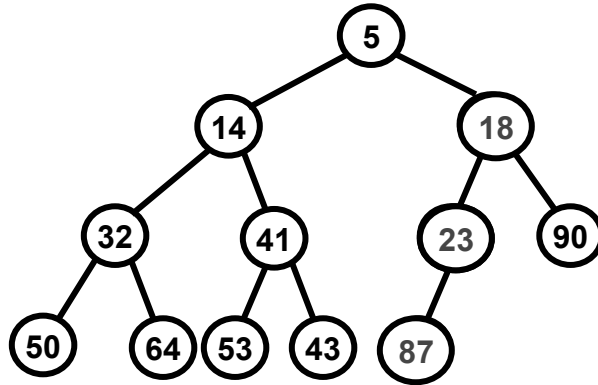


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Inserting into a min-heap

Insert 18

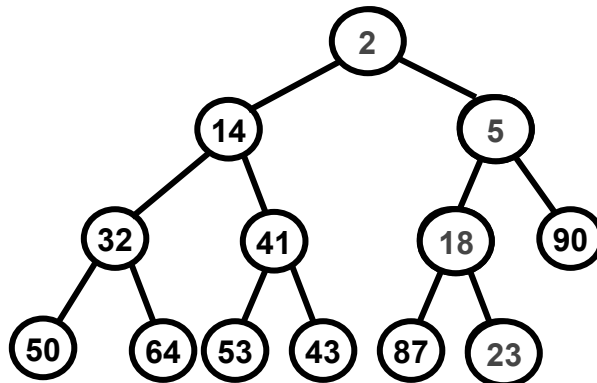


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Inserting into a min-heap

Insert 2



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Removing from a heap



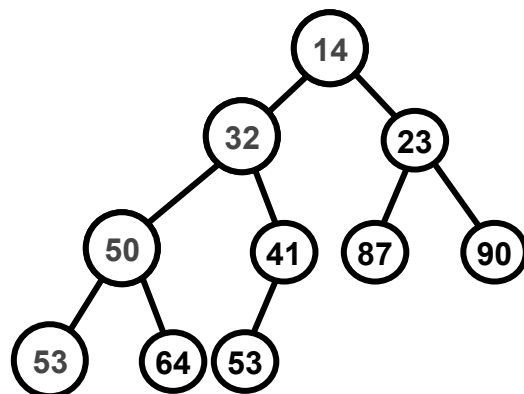
- Place the root element in a variable to return later.
- Remove the last element in the deepest level and move it to the root.
- While the moved element has a value greater than at least one of its children, swap this value with the smaller-valued child.
- Return the original root that was saved.

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Removing from a min-heap

Remove min



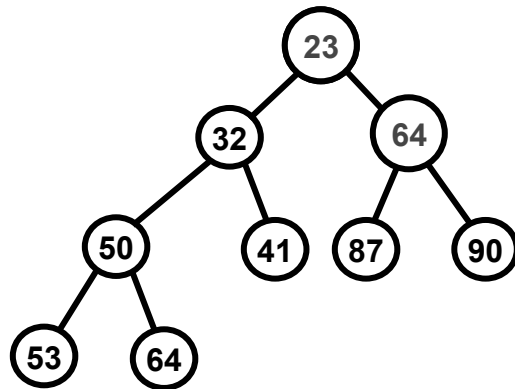
returnValue 5

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Removing from a min-heap

Remove min



returnValue 14

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Efficiency of heaps



Assume the heap has N nodes.

Then the heap has $\lceil \log_2(N+1) \rceil$ levels.

- Insert
Since the insert swaps at most once per level, the order of complexity of insert is $O(\log N)$
- Remove
Since the remove swaps at most once per level, the order of complexity of remove is also $O(\log N)$

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Priority Queues



- A priority queue PQ is like an ordinary queue except that we can only remove the “maximum” element at any given time (not the “front” element necessarily).
- If we use an array to implement a PQ, enqueue is $O(\text{_____})$ dequeue is $O(\text{_____})$
- If we use a sorted array to implement a PQ enqueue is $O(\text{_____})$ dequeue is $O(\text{_____})$
- If we use a max-heap to implement a PQ enqueue is $O(\text{_____})$ dequeue is $O(\text{_____})$

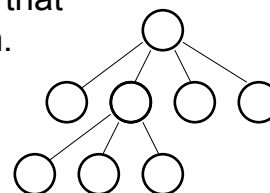
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General Trees

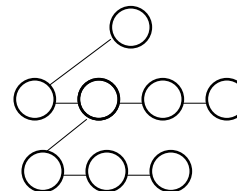


- A general tree consists of nodes that can have any number of children.



- Implementation using a binary tree:

Each node has 2 fields:
firstChild, nextSibling



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Balanced Trees



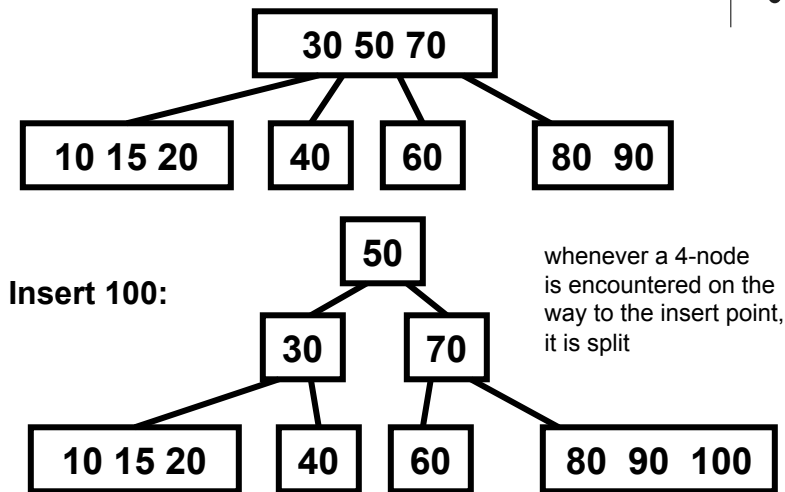
- Binary search trees can become quite unbalanced, with some branches being much longer than others.
 - Searches can become $O(n)$ operations
- These variants allow for searching while keeping the tree (nearly) balanced:
 - 2-3-4 trees
 - Red-black trees

2-3-4-trees



- A 2-3-4 Tree is a tree in which each internal node (nonleaf) has two, three, or four children, and all leaves are at the same depth.
 - A node with 2 children is called a "2-node".
 - A node with 3 children is called a "3-node".
 - A node with 4 children is called a "4-node".

Sample 2-3-4-tree



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Red-Black Trees

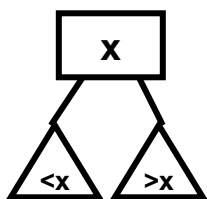


- A red-black tree has the advantages of a 2-3-4 tree but requires less storage.
- Red-black tree rules:
 - Every node is colored either red or black.
 - The root is black.
 - If a node is red, its children must be black.
 - Every path from a node to a null link must contain the same number of black nodes.

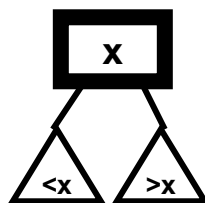
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2-3-4 Trees vs. Red-Black Trees

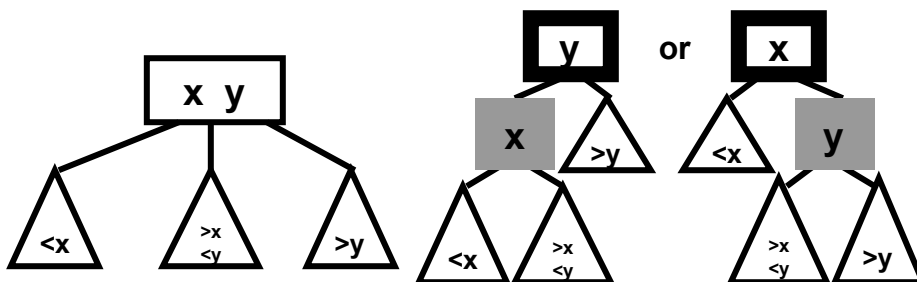


"2-node" in a 2-3-4 tree



equivalent red-black tree configuration

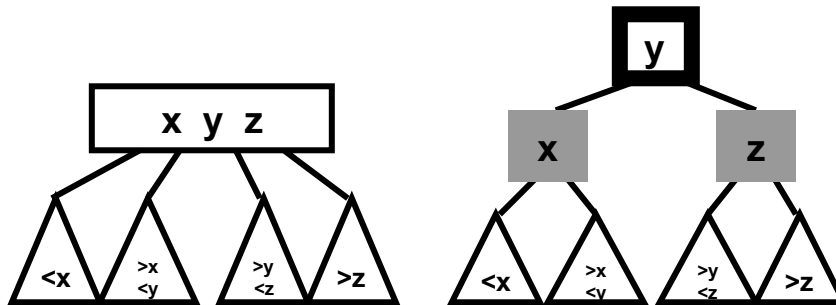
2-3-4 Trees vs. Red-Black Trees



"3-node" in a 2-3-4 tree

equivalent red-black tree configurations

2-3-4 Trees vs. Red-Black Trees



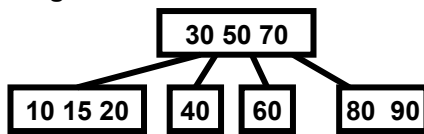
"4-node" in a 2-3-4 tree

equivalent red-black tree configuration

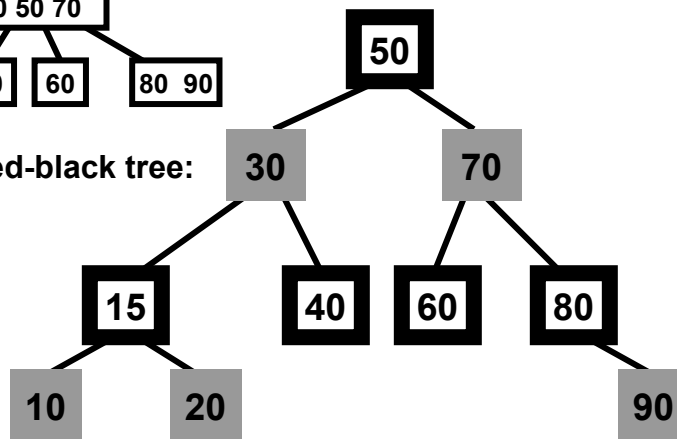
Sample Red-Black Tree



Original 2-3-4 tree:

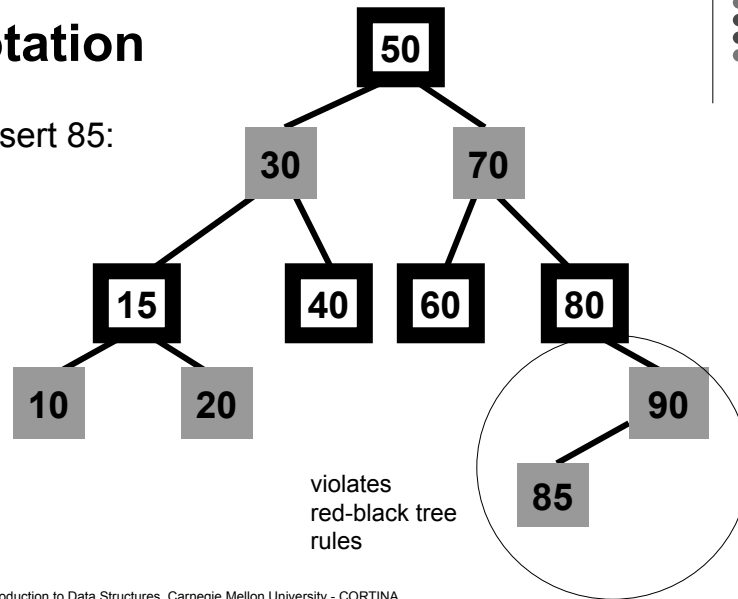


Equivalent red-black tree:



Rotation

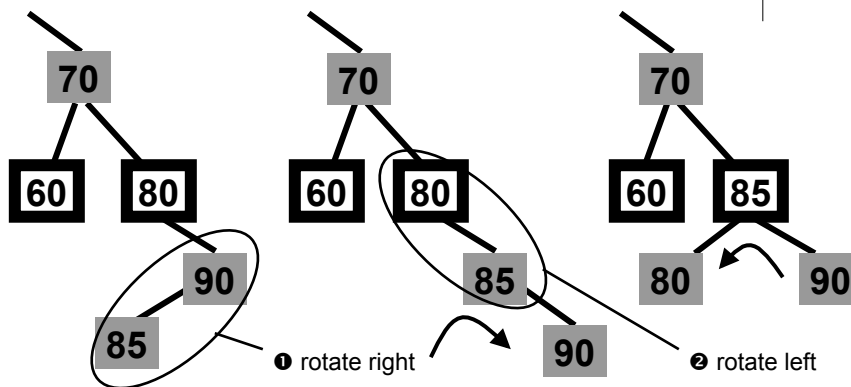
- Insert 85:



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Rotation (cont'd)



See textbook for additional cases where rotation is required.

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Additional Self-Balancing Trees



- AVL Trees
- 2-3 Trees
- B-Trees
- Splay Trees
 - (co-invented by Prof. Danny Sleator at CMU)