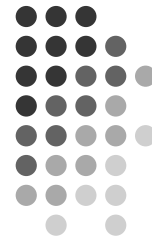


# Algorithmic Methods

## Tricks of the Trade

# 5A

### Recursion



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## Recursion



- A recursive operation is an operation that is defined in terms of itself.



<http://fusionanomaly.net/recursion.jpg>



Sierpinski's  
Gasket

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## Recursion



- Every recursive definition includes two parts:
  - Base case (non-recursive)  
A simple case that can be done without solving the same problem again.
  - Recursive case(s)  
One or more cases that are “simpler” versions of the original problem.
    - By “simpler”, we sometimes mean “smaller” or “shorter” or “closer to the base case”.

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## Factorial



- Definition:  $n! = n(n-1)(n-2)\dots(2)(1)$
- Since  $(n-1)(n-2)\dots(2)(1) = (n-1)!$ 
  - $n! = n(n-1)!$ , for  $n > 0$
  - $n! = 1$  for  $n = 0$  (base case)
- Example:

$4! = 4(3!)$   
 $3! = 3(2!)$   
 $2! = 2(1!)$   
 $1! = 1(0!) = 1(1) = 1$

$= 4(6) = 24$   
 $= 3(2) = 6$   
 $= 2(1) = 2$

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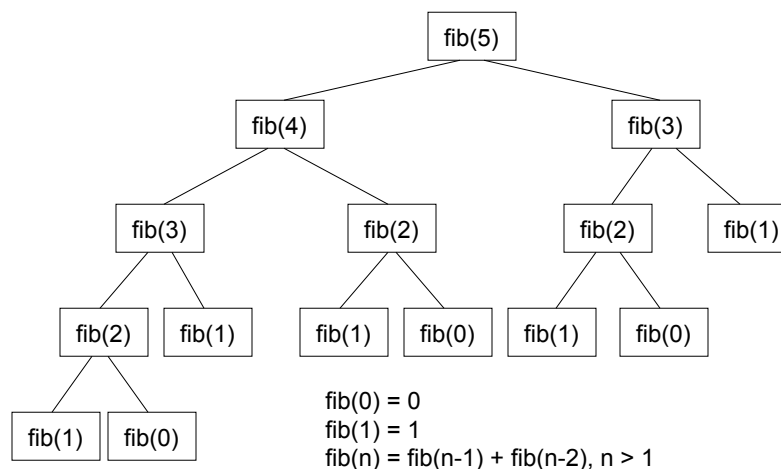
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# Fibonacci Numbers

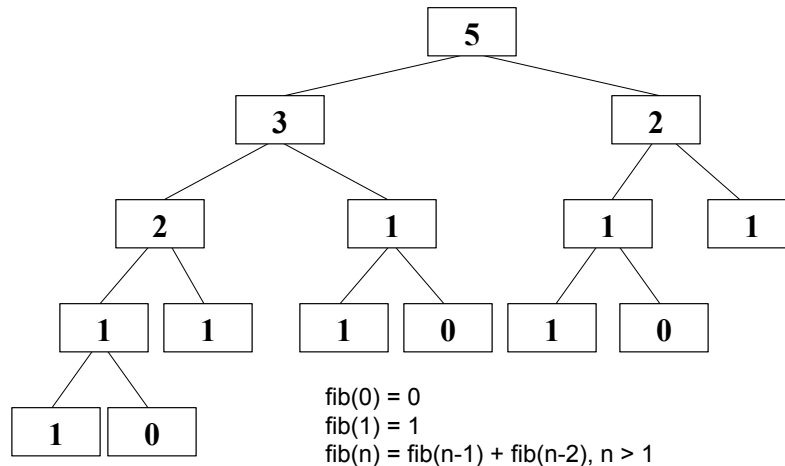


- A sequence of numbers each number is the sum of the previous two numbers in the sequence, starting the sequence with 0 and 1.
- 0, 1, 1, 2, 3, 5, 8, 13, 21, 34, 55, 89, etc.
- Let  $\text{fib}(n)$  = the  $n^{\text{th}}$  Fibonacci number,  $n \geq 0$ 
  - $\text{fib}(0) = 0$  (base case)
  - $\text{fib}(1) = 1$  (base case)
  - $\text{fib}(n) = \text{fib}(n-1) + \text{fib}(n-2)$ ,  $n > 1$

# Fibonacci Numbers



# Fibonacci Numbers

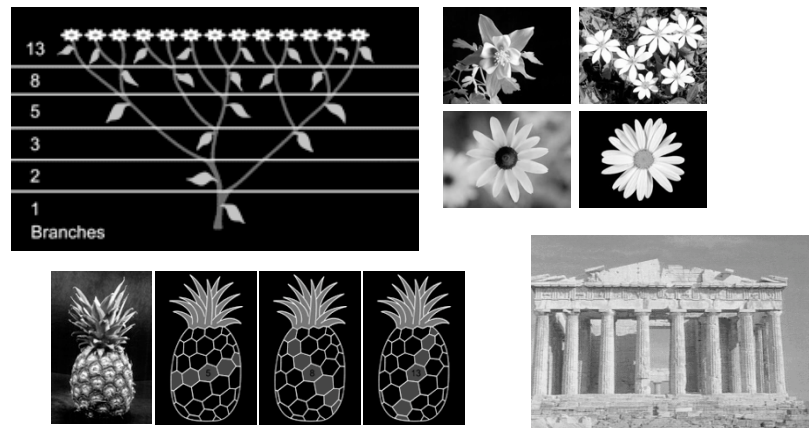


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# Fibonacci Numbers in Nature

<http://britton.disted.camosun.bc.ca/fibslide/jbfibslide.htm>  
<http://www.geom.uiuc.edu/~demo5337/s97b/art.htm>



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## Recursive Sum (in Scheme)



- Computing the sum of a list of numbers.

Use: `(sum (list 30 28 45 12))`

- Recursive Definition:

```
(define (sum numlist)
  (if (null? numlist) 0
      (+ (first numlist)
          (sum (rest numlist)))
  )
)
```

Annotations:

- ← is numlist empty?
- ← if yes, result is 0
- } if no, result is the first number + the sum of the rest of the numbers

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## Recursive Sum (in Scheme)



```
(sum (list 30 28 45 12))
(+ 30 (sum (list 28 45 12)))
(+ 30 (+ 28 (sum (list 45 12))))
(+ 30 (+ 28 (+ 45 (sum (list 12)))))
(+ 30 (+ 28 (+ 45 (+ 12 (sum (list ))))))
(+ 30 (+ 28 (+ 45 (+ 12 0 ))))
(+ 30 (+ 28 (+ 45 12 )))
(+ 30 (+ 28 57 ))
(+ 30 85 )
115
```

Annotation: empty list (pointing to the empty list in the recursive call)

```
(if (null? numlist)
    0
    (+ (first numlist)
        (sum (rest numlist))
    )
)
```

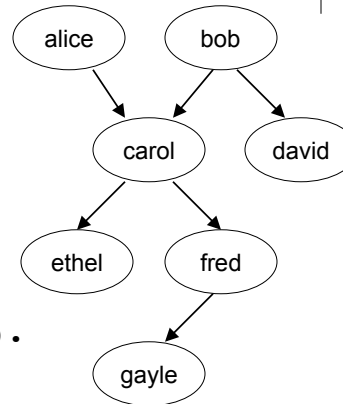
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## Recursion in Prolog

```
ancestor(X, Z) :-  
    parent(X, Z).  
ancestor(X, Z) :-  
    parent(X, Y),  
    ancestor(Y, Z).
```

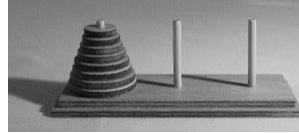
```
?- ancestor(alice, gayle).  
yes
```



## Recursion in Prolog

```
ancestor(alice, gayle) :-  
    parent(alice, Y), ancestor(Y, gayle).  
Try Y = carol.  
    parent(alice, carol). YES  
    ancestor(carol, gayle) :-  
        parent(carol, gayle). NO  
    ancestor(carol, gayle) :-  
        parent(carol, Y'), ancestor(Y', gayle).  
Try Y' = fred.  
    parent(carol, fred). YES  
    ancestor(fred, gayle) :-  
        parent(fred, gayle). YES
```

## Towers of Hanoi



Towers of Hanoi with 8 discs.

- A puzzle invented by French mathematician Edouard Lucas in 1883.
- At a temple far away, priests were led to a courtyard with three pegs and 64 discs stacked on one peg in size order.
  - Priests are only allowed to move one disc at a time from one peg to another.
  - Priests may not put a larger disc on top of a smaller disc at any time.
- The goal of the priests was to move all 64 discs from the leftmost peg to the rightmost peg.
- According to the legend, the world would end when the priests finished their work.

## Towers of Hanoi

**Move  $N$  discs from peg  $X$  to peg  $Y$   
(Let  $Z$  represent the other peg.)**

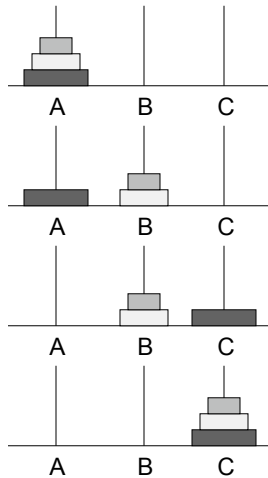
a. Move  $N-1$  discs from peg  $X$  to peg  $Z$  (if  $N > 1$ ).

b. Move 1 disc from peg  $X$  to peg  $Y$ .

c. Move  $N-1$  discs from peg  $Z$  to peg  $Y$  (if  $N > 1$ ).

extra peg  $Z$

## Towers of Hanoi (N=3)



**Move 3 discs from peg A to peg C.  
(extra peg is B)**

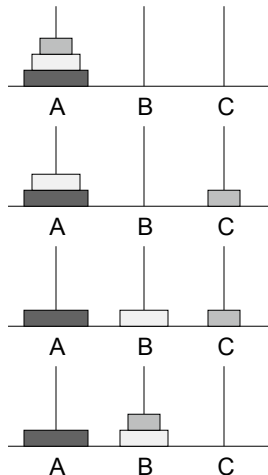
- Move 2 discs from peg A to peg B.  
(RECURSIVE... see next slide)
- Move 1 disc from peg A to peg C.
- Move 2 discs from peg B to peg C.  
(RECURSIVE... see two slides ahead)

extra  
peg

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## Towers of Hanoi (N=2)



**Move 2 discs from peg A to peg B.  
(extra peg is C)**

- Move 1 disc from peg A to peg C.
- Move 1 disc from peg A to peg B.
- Move 1 disc from peg C to peg B.

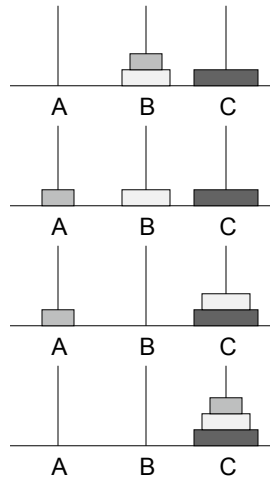
extra  
peg

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## Towers of Hanoi (N=2)



**Move 2 discs from peg B to peg C.  
(extra peg is A)**

a. Move 1 disc from peg B to peg A.

b. Move 1 disc from peg B to peg C.

c. Move 1 disc from peg A to peg C.

extra  
peg

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## Towers of Hanoi



Discs

Moves

1

\_\_\_\_\_

2

\_\_\_\_\_

3

\_\_\_\_\_

4

\_\_\_\_\_

5

\_\_\_\_\_

...

n

\_\_\_\_\_

What is the  
fewest number of  
disc moves needed  
for a problem  
with n discs?

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## Towers of Hanoi



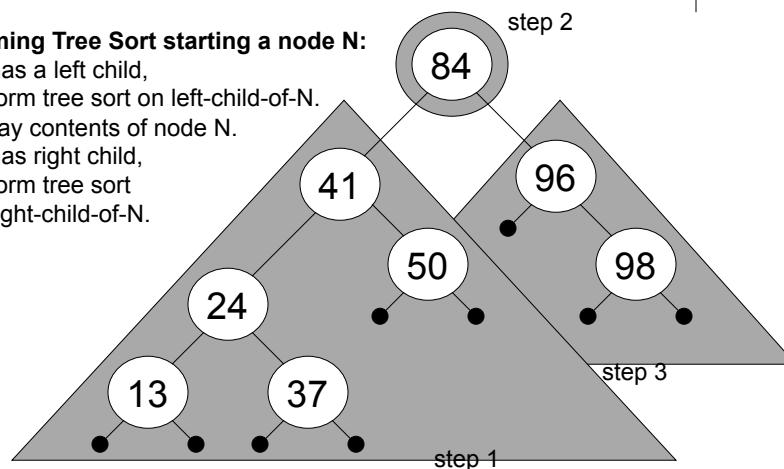
- If the priests moved a disc at a rate of 1 per second using the fewest number of disc moves, it would take the priests roughly 585 billion years to complete this puzzle!
  - The universe is currently about 13.7 billion years old.

## Tree Sort Algorithm revisited



### Performing Tree Sort starting a node N:

1. If N has a left child, perform tree sort on left-child-of-N.
2. Display contents of node N.
3. If N has right child, perform tree sort on right-child-of-N.

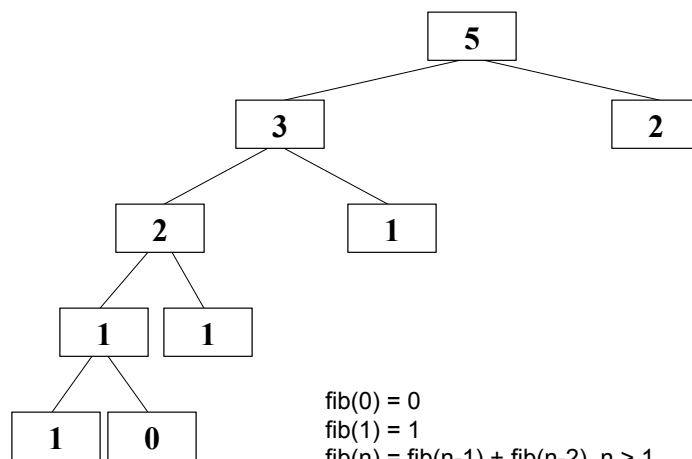


# Dynamic Programming



- A technique that...
  - The problem is broken into sub-problems, and these sub-problems are solved and the solutions remembered, in case they need to be solved again.
  - All sub-problems that might be needed are solved in advance and then used to build up solutions to larger problems.

## Fibonacci Numbers revisited



n	fib(n)
0	0
1	1
2	1
3	2
4	3
5	5

memoization