

dyf d(0,0) = 0

d(0,0) = 2x dy = 1  
d(0,0) = 1 dy = 1  
d(0,0) = 1 dy = 1

# Lecture 3 Lect III

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(1) Recap. We saw some important convex functions and some ways of constructing them.

(2) Thm: Let  $h(x,y)$  be convex in  $x$ , for every  $y \in Y$ . Then,

$$f(x) := \sup_{y \in Y} h(x,y) \text{ is convex.}$$

(3) Special case: Fenchel-Conjugate:

$$f^*(x) := \sup_{y \in \mathbb{R}^n} \langle x, y \rangle - f(y)$$

Observe  $f^*$  is convex even if  $f$  is not.

(4) Basic theorem: Let  $f$  be a closed convex function

(i.e. a function whose epi is a closed set, i.e. a lower semi-cont. w.r.t.  $\text{epi}$ )

then  $f^{**} = f$ ;

(5) Automatic characterization of the Legendre-Fenchel transform

- Why study  $f^*$ : intimately related to duality
- Ultimately grounded in the simple idea.



#  $f$  being convex is globally underestimated by linear functions. (apart out)

# The highest affine underestimator is

$$\lambda(x) := \inf_y f(y) - \langle x, y \rangle$$

i.e. at a point  $x$  the value of highest linear underestimator.

\* encoding epi via its supporting hyperplanes.



~~Theorem: (Artstein-Avidan, Milman)~~ Let  $\Gamma_0(\mathbb{R}^n)$  denote the class of closed convex fns on  $\mathbb{R}^n$ . The Legendre-Fenchel transform of a fn.  $f \in \Gamma_0$ , is defined by

(2)

$$L: f \mapsto \sup_y \langle x, y \rangle - f(y)$$

i.e.  $(Lf)(x) = \sup_y \langle x, y \rangle - f(y)$

Theorem! Let  $T$  be any transform that maps  $\Gamma_0 \rightarrow \Gamma_0$  and satisfies

- 1)  $T(Tf) = f$  (involution)
- 2)  $f \leq g \Rightarrow Tf \geq Tg$  (order-reversing)

Then  $T$  must be "essentially" the Legendre-Fenchel transform.

That is,  $\exists C \in \mathbb{R}, U \in \mathbb{R}^n, B \in GL_n$  st.

$$(Tf)(x) = (Lf)(Bx + U) + \langle z, U \rangle + C$$

Remarks: gives axiomatic support to great success of LF-transform.

Playful: Let  $\mathcal{S}$  be a distinguished class of functions in  $\mathbb{R}^n$ .

- we call  $T$  a duality transform on  $\mathcal{S}$  if

- for any  $f \in \mathcal{S}$ ,  $T(Tf) = f$
- for any two fns  $f, g \in \mathcal{S}$ , st  $f \leq g$ ,  $Tf \geq Tg$ .

Eg. Say  $\mathcal{S}$  is class of log-concave fns

then  $(Tf)(x) := \inf_{y \in \mathbb{R}^n} \frac{e^{-\langle x, y \rangle}}{f(y)}$  is such a transform.

(In fact, essentially unique).

# Try your own class!! Eg try  $(\mathbb{R}_+, |\log x - \log 1|)$

# Concept of subdifferentials, intimately related to ~~convex~~ Fenchel-Conjugates

(3)

Let  $f: \mathbb{R}^n \rightarrow \mathbb{R}$  be any function.

Define:  $\partial f(x) := \{ g \in \mathbb{R}^n \mid f(y) \geq f(x) + \langle g, y-x \rangle \forall y \in \text{dom} f \}$

Notice: if  $x \notin \text{dom} f$ , then  $\partial f(x) = \emptyset$

Theorem: Let  $f$  be a proper (i.e.  $\neq +\infty$ ) convex function.

Then  $g \in \partial f(x)$  iff  $f^*(g) + f(x) = \langle g, x \rangle$

Proof:  $g \in \partial f(x) \Leftrightarrow f(y) \geq f(x) + \langle g, y-x \rangle$   
 $\Leftrightarrow f(y) - \langle g, y \rangle \geq f(x) - \langle g, x \rangle$   
 $\Leftrightarrow \langle g, y \rangle - f(y) \leq \langle g, x \rangle - f(x) \quad \forall y \in \text{dom} f$

i.e.  $f^*(g) = \sup_y \langle g, y \rangle - f(y) = \langle g, x \rangle - f(x)$

But in view of def  $f^*$

$$f^*(g) \leq \langle g, x \rangle - f(x)$$

But by defn of  $f^*$  this is an equality.

# Some basic theorems you should be aware of.

(a) Let  $f \in \Gamma_0(\mathbb{R}^n)$ ; (i)  $\partial f(x) \neq \emptyset$  whenever  $x \in \text{ri}(\text{dom} f)$ .

(ii) Let  $f$  be a convex function.

(b) If  $f \in \Gamma_0(\mathbb{R}^n)$ , then

$$g \in \partial f(x) \Leftrightarrow x \in \partial f^*(g)$$

Thm: (Rockafellar) ~~more~~ let  $f_1, \dots, f_m$  be proper convex functions on  $\mathbb{R}^n$  and let  $f = f_1 + \dots + f_m$

(9)

Then

$$\partial f_1 + \dots + \partial f_m \subset \partial f(x), \forall x$$

If the convex sets  $\bigcap_{i=1}^m \text{ri}(\text{dom } f_i) \neq \emptyset$

then, actually

$$\partial(f_1 + \dots + f_m) = \partial f_1 + \dots + \partial f_m$$

Moreover, if say  $f_1, \dots, f_m$  are polyhedral, then

it is enough that

$$\left[ \bigcap_{i=1}^m \text{dom } f_i \right] \cap \left( \bigcap_{i=1}^m \text{ri}(\text{dom } f_i) \right) \neq \emptyset$$

Proof:  $\subset$  part is easy.

Assume  $\bigcap_{i=1}^m \text{ri}(\text{dom } f_i) \neq \emptyset$ , we have

$$(f_1 + \dots + f_m)^*(g) = \inf \{ f_1^*(g_1) + \dots + f_m^*(g_m) \mid g_1 + \dots + g_m = g \}$$

when for each  $i$  the inf is attained.

(I'll keep the form in the notes)

Say  $\exists g = g_1 + \dots + g_m$  where  $g_i \in \partial f_i(x)$ .

So  $g \in \partial f(x)$

@ Thm, ~~Key~~ }  $f(z) = f_1(z) + \dots + f_m(z) \geq f_1(x) + \dots + f_m(x) + \langle g_1, z-x \rangle + \dots + \langle g_m, z-x \rangle = f(x) + \langle g_1 + \dots + g_m, z-x \rangle = f(x) + \langle g, z-x \rangle$

Example: Define  $f_1, f_2$  by

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$$f_1(x) := \begin{cases} -2\sqrt{x} & \text{if } x \geq 0 \\ +\infty & \text{if } x < 0 \end{cases}$$

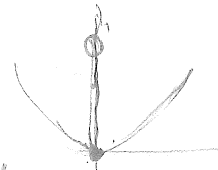
$$f_2(x) := \begin{cases} +\infty & \text{if } x > 0 \\ -2\sqrt{-x} & \text{if } x \leq 0 \end{cases}$$

$\sqrt{x}$  is concave on  $[0, \infty)$  so  $-2\sqrt{x}$  is convex;  $\sqrt{-x}$  is concave.

$$f \equiv \underline{f_1(x) + f_2(x)} \equiv \begin{cases} 0, & x = 0 \\ +\infty, & \text{otherwise} \end{cases} \equiv \delta_{\{0\}}(x)$$

so  $\underline{\partial f(0) = \mathbb{R}}$

But  $\underline{\partial f_1(0) = \emptyset}$   
 $\underline{\partial f_2(0) = \emptyset}$  ( Exercise  
 $\sqrt{x}$  is non diff at 0  
in fact  $\rightarrow \infty$  )



## # Subgradient calculus

- Reading assignment: ( Link posted on class website )

# • Determinis  $\underline{\partial f(x)}$ : Usually very hard

• Usually we need just one  $\underline{g_x \in \partial f}$ .

• However testing optimality may req. knowing  $\partial f$ .

Thm. Let  $f(x) := \sup \{ f_i(x) \mid i \in I \} < +\infty$   
 $\forall x \in \mathbb{R}^n$ .

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$I$  is an arbitrary index set.

Let  $I(x) := \{ i \in I \mid f_i(x) = f(x) \}$  admit that

then

$$\partial f(x) = \text{conv} \bigcup_{i \in I(x)} \partial f_i(x)$$

Question: what happens when  $f(x)$  can be  $+\infty$  for some values?

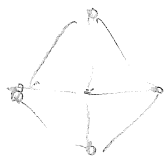
Example:  ~~$f(x) = \|x\|$~~  = what is  $\partial \|x\|_\infty$ ?

clearly:  ~~$\|x\|_\infty = \sup \{ \langle x, u \rangle \mid \|u\|_1 \leq 1 \}$~~

$$\|x\|_\infty = \max_{1 \leq i \leq n} |x_i|$$

$$= \max_{1 \leq i \leq n} |e_i^T x|$$

$$\Rightarrow \partial \|x\|_\infty(0) = \text{conv} \{ \pm e_1, \dots, \pm e_n \}$$



# For other examples and details please refer to notes

PROPT: (If  $g \in \partial f(x) \Rightarrow x \in \text{dom } f \Rightarrow g \in \text{dom } f^*$   
 But dom  $f^*$  can be larger than  $\partial f(\mathbb{R}^n)$ ! )

# Optimization Problems:

(7)

①  $\min_{x \in \mathbb{R}^n} f(x)$

→ say  $f$  is differentiable we know that

$\nabla f(x^*) = 0$  is necc. for  $x^*$  to be a local min.

→ If  $f \in C^2$ ,  $\nabla^2 f(x^*) \succ 0$  suff for local min.

~~Nonconv world~~

→ For conv  $f$ ,  $\nabla f(x^*) = 0$  also sufficient.

→ Nonconv world: No known "easy" necc + suff conditions for optimality.

②  ~~$\min_{x \in \mathbb{R}^n} f(x)$  s.t.  $x \in X$~~

~~write as unconstrained problem~~

Theorem: (Fermat's law): Let  $f: \mathbb{R}^n \rightarrow \mathbb{R} \cup \{+\infty\}$ . Then

$$\text{Argmin } f = \text{Zer}(\partial f) := \{x \in \mathbb{R}^n \mid 0 \in \partial f(x)\}$$

Proof:  $x \in \text{Argmin } f \Rightarrow f(x) \leq f(y) \quad \forall y \in \mathbb{R}^n$ .

So  $f(y) \geq f(x) + \langle 0, y-x \rangle \quad \forall y \Rightarrow 0 \in \partial f(x)$ .

(convex is obvious!)

Example:  $\min_{x \in X} f(x)$

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$$\Leftrightarrow \min_{x \in \mathbb{R}^n} f(x) + \delta_X(x)$$

So  $x^* \in \text{Argmin}$  iff

$$0 \in \partial (f + \delta_X)(x^*)$$

Recall Rockafellar's theorem: if  $0 \in \text{ri}(\text{dom} f) \cap \text{ri}(\text{dom} \delta_X) \neq \emptyset$

then  $\partial (f + \delta_X) = \partial f + \delta_X$

(So assuming the above "constraint qualification")

$$0 \in \partial f(x^*) + \partial \delta_X(x^*)$$

What is  $\partial \delta_X$ ?

$$g \in \partial \delta_X \Leftrightarrow \delta_X(y) \geq \delta_X(x) + \langle g, y-x \rangle \quad \forall y \in \mathbb{R}^n$$

$x \in X$  so  $\delta_X(x) = 0$

$$\Rightarrow 0 \geq \langle g, y-x \rangle \quad \forall y \in X$$

Def: (Normal cone)

$$\mathcal{N}_X(x) := \{ g \in \mathbb{R}^n \mid 0 \geq \langle g, y-x \rangle, \forall y \in X \}$$

$$\Rightarrow 0 \in \partial f(x^*) + \mathcal{N}_X(x^*)$$



If  $f$  is differentiable,  $\partial f(x^*) = \nabla f(x^*)$

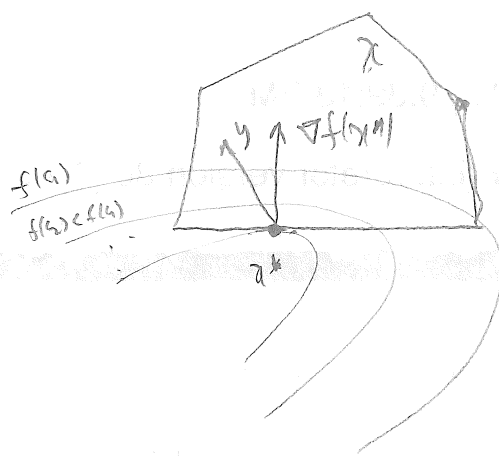
(9)

So above condition translates into:

$$\langle \nabla f(x^*), y - x^* \rangle \geq 0 \quad \forall y \in \mathcal{X}$$

(So-called Kolmogorov condition in optimality)

Picture.



$\nabla f(x^*)$

defines

supporting

hyperplane at

$x^*$

Example: Projection Lemma:

$$\min_{x \in \mathcal{X}} \frac{1}{2} \|x - x_0\|^2$$

iff.  $\langle x^* - x_0, y - x^* \rangle \geq 0 \quad \forall y \in \mathcal{X}$

(Usual result)

Question: What about optimality conditions in normed convex, quadratic metric spaces ( $\mathcal{X}, d$ )

HW: Show that if  $f$  is q.c. on a q.c. set  $\mathcal{X}$

then local optimality  $\implies$  global optimality

# Specific Examples of OPT problems

① Linear program: 
$$\begin{array}{l} \min \quad C^T x \\ \text{st.} \quad Ax = b \\ \quad \quad x \geq 0 \end{array} \quad \text{Standard form}$$

Wk!

② Quadratic: 
$$\begin{array}{l} \min \quad x^T A x + b^T x \\ \text{st.} \quad Ax \leq b \end{array}$$

③ ~~SOCP~~ SOCP, and SDPs:

Ex: write:  $\min f(x) = \left( \text{min}_{|K| \leq m} a^T x + b \right)$  as an LIP

Q: Does QP always have a solution?

## Semidefinite Programming

$$\begin{array}{l} \min \quad C^T x \\ \text{st.} \quad x \in \mathbb{R}^n \\ \quad \quad \mathcal{A}(x) := A_0 + x_1 A_1 + \dots + x_n A_n \succeq 0 \end{array}$$

$A_0, \dots, A_n \in S^{n \times n}$

Exercise:  $\text{SDP} \supset \text{SOCP} \supset \text{QP} \supset \text{LP}$

Equivalent / Std. form.

$$\begin{array}{l} \min \quad \langle C, X \rangle \\ \quad \quad \langle A_i, X \rangle = b_i \\ \quad \quad X \succeq 0 \end{array}$$

Using duals  
or conveys  
between  
the two. (to be later)

Where does the idea come from?



LP:  $\min c^T x$ ,  
 $Ax = b$   
 $x \geq 0$

Simplex, Successive (pivot,  
 Dantzig, Smallest, CR, etc. etc.)

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How to generalize:  $c^T x$ :  $\Rightarrow$  replace by nonlinear fn.

$Ax = b$  replace A by a nonlinear mat?  $\parallel$  new  $x_i = 1$   
 $\cup$   $x_i = 1!!$

Both quickly become non-tractable.

# Generalize structural constraint.  $\mathbb{R}_+^n$

$\mathbb{R}_+^n$  is a cone (polyhedral convex)

\* Replace by general convex cone!

\* Importantly: replace  $x \geq 0$  by  $x \in K$

\* Nesterov-Nemirovskii - 80's Convex Opt. Theory

Seeking to gen.  $x \geq 0$

$$K := \{x \in \mathbb{R}^n \mid x \geq 0\}$$

$$x \geq y \Leftrightarrow x - y \geq 0 \Leftrightarrow x - y \in K$$

Lemmas (Nec. & suff. for a set  $K \subset \mathbb{R}^n$  to define a

useful vector ineq. is that  $K$  should be a nonempty  
convex cone: i.e.

- $K \neq \emptyset$
- $x, y \in K \Rightarrow x + y \in K$
- $\alpha x \in K, \alpha \geq 0$
- $x, -x \in K \Rightarrow x = 0$

Two other important properties

Recall for  $\mathbb{R}_+^n$ : Closed with nonempty interior (12)

- (  $\{x_u\} \subset \mathbb{R}_+^n \rightarrow x \Rightarrow x \in \mathbb{R}_+^n$  )
- (  $n \mathbb{R}_+^n$  : contains Ball with the radius 1 )

ED Conic problem : 
$$\begin{aligned} \min & c^T x \\ \text{A}x &= b \\ x &\succeq_K 0 \end{aligned}$$

- LP:  $K = \mathbb{R}_+^n$ ;
- SOCP:  $K = \mathbb{Q}^n := \{(y,t) \in \mathbb{R}^n \mid \|y\|_2 \leq t\}$
- SDP:  $K = S_+^n := \{X = X^T \succeq 0, X \in S^n\}$
- OR Cartesian products of these

Other cones could be used, but there none are "conist"  $\rightarrow$  Belav  
project could investigate Some other norm cones.

(Not all cones are nice: Recall Chn)

SDP (conic program): 
$$\min c^T x, Ax = b, x \in K$$

Think of  $x$  as a matrix  
 # so  $c^T x = \text{tr}(C^* x) = \underline{C(\cdot)^T x(\cdot)}$  matrix notation  
 # wlog assume  $K = S_+^n$

#  $Ax = b \Leftrightarrow a_i^T x = b_i \Leftrightarrow \langle A_i, X \rangle = b_i$   
 $x \in K \Leftrightarrow X \succeq 0$

$$\begin{aligned} \min & \langle C, X \rangle \\ & \langle A_i, X \rangle = b_i \\ & X \succeq 0 \end{aligned}$$

MW: try solving a few SDPs in CUX

(Don't forget)

Important Topic: Which course (and number) 13

Optimization Problems can be formulated as SDPs?

↳ Related to major open problems in convex geometry  
 (We'll come to that later in the course)

Ex: 
$$\begin{aligned} \min \quad & x^T A x \\ \text{s.t.} \quad & x_i^2 = 1 \end{aligned}$$

Notice  $A \succeq 0$   
 But constraint is nonlinear

$x^T A x = \text{tr}(A x x^T)$

$$= \begin{cases} \min \langle A, X \rangle \\ \text{Diag}(X) = \mathbb{1} \\ X \succeq 0 \end{cases}$$

$\text{rank}(X) = 1$

$\Rightarrow X = x x^T$

$\left( \begin{matrix} X \\ x x^T \end{matrix} \right) \text{ exists!}$

$X \succeq x x^T$

$$\begin{bmatrix} X & x \\ x^T & 1 \end{bmatrix} \succeq 0$$

