

# 15-744: Computer Networking

## L-1 Intro to Computer Networks



## Short Diagnostic Quiz



- Problem
  - Different undergraduate backgrounds
- Objective
  - Gauge relative skill level
- Time
  - 15 minutes

## Outline



- **Administrivia**
- Whirlwind tour of networking

## Who's Who?



- Professor: Srinivasan Seshan
  - <http://www.cs.cmu.edu/~srini>
  - [srini@cmu.edu](mailto:srini@cmu.edu)
  - Office hours: Wed 11:00-12:00 (for now...)
- TA: Bianca Schroeder
  - [bianca@cs.cmu.edu](mailto:bianca@cs.cmu.edu)
  - Office hours: Mon 12:00-1:00
- Course info
  - <http://www.cs.cmu.edu/~srini/15-744/S02/>

## Objectives



- Understand the state-of-the-art in network protocols, architectures and applications
- Understand how networking research is done
- How is class different from undergraduate networking (15-441)
  - Training network programmers vs. training network researchers

## Web Page



- Check regularly!!
- Course schedule
- Reading list
- Lecture notes
- Announcements
- Assignments
- Project ideas
- Exams
- Student list

## Course Materials



- Research papers
  - Links to ps or pdf on Web page
  - Combination of classic and recent work
  - ~40 papers
  - Optional readings
- Recommended textbook
  - For students not familiar with networking
  - Peterson & Davie 2<sup>nd</sup> edition
  - 2 copies on reserve

## Grading



- Homework assignments (20%)
  - Problem sets
  - Hands-on assignments
- Class participation (10%)
  - Pop quizzes
- 2 person project (30%)
- Midterm exam (20%)
- Final exam (not cumulative) (20%)

## Outline



- Administrivia
- Whirlwind tour of networking

## What is the Objective of Networking?



- Communication between applications on different computers
- Must understand application needs/demands
  - Traffic data rate
  - Traffic pattern (bursty or constant bit rate)
  - Traffic target (multipoint or single destination, mobile or fixed)
  - Delay sensitivity
  - Loss sensitivity

## Four Steps to Networking



- Communicating across a link
- Connecting together multiple links (internetworking)
- Finding and routing data to nodes on internetwork
- Matching application requirements

## A First Step



- Creating a link between nodes
- Link: path followed by bits
  - Wired or wireless
  - Broadcast or point-to-point (or both)
- Node: any device connected to a link

## Types of Links



### Point-to-Point



### Multiple Access



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## Packet Transmission Modes



- Unicast
  - Transmission to single specific receiver
- Broadcast
  - Transmission to all network nodes
- Multicast
  - Transmission to specific subset of nodes
- Anycast
  - Transmission to one of a specific subset of nodes

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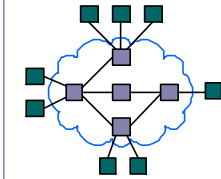
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## What are Switched Networks?



- Switch: moves bits between links
  - Packet switching
  - Circuit switching

### Switched Network

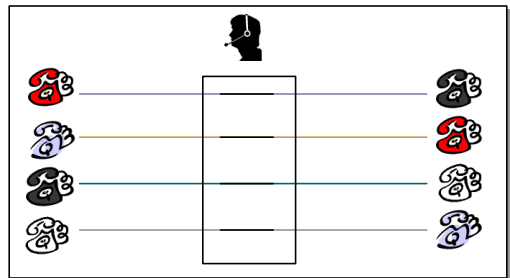


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## Back in the Old Days...



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## Then Came TDM...



- Synchronous time division multiplexing

### Multiplex (mux)



### Demultiplex (demux)

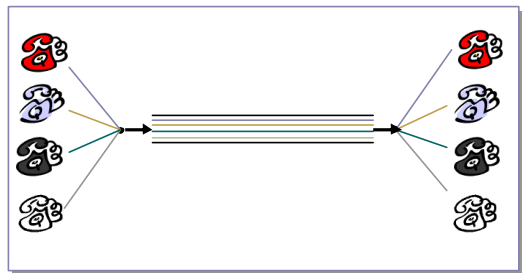


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## TDM Logical Network View

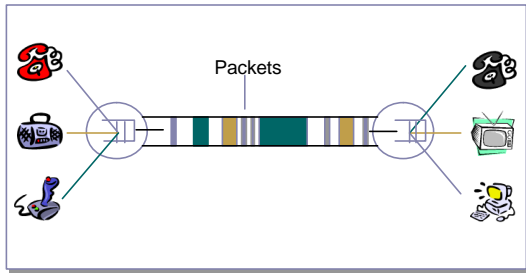


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## Packet Switching (Internet)



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## Packet Switching

- Interleave packets from different sources
- Efficient: resources used on demand
  - Statistical multiplexing
- General
  - Multiple types of applications
- Accommodates bursty traffic
  - Addition of queues

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## Statistical Multiplexing Gain

- 1 Mbps link; users require 0.1 Mbps when transmitting; users active only 10% of the time
- Circuit switching: can support 10 users
- Packet switching: with 35 users, probability that  $\geq 10$  are transmitting at the same time  $< 0.0017$

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## Characteristics of Packet Switching

- Store and forward
  - Packets are self contained units
  - Can use alternate paths - reordering
- Contention
  - Congestion
  - Delay

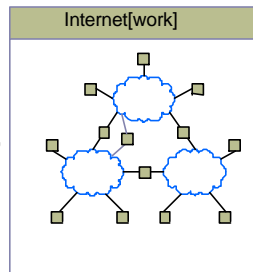
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## Second Step: Internet[work]

- A collection of interconnected networks
- Host: network endpoints (computer, PDA, light switch, ...)
- Router: node that connects networks
- Internet vs. internet



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## Challenge

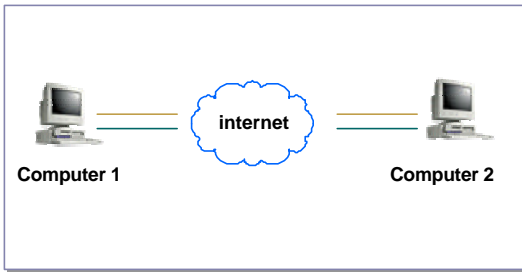
- Many differences between networks
  - Address formats
  - Performance – bandwidth/latency
  - Packet size
  - Loss rate/pattern/handling
  - Routing
- How to translate between various network technologies

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### Third Step: How To Find Nodes?

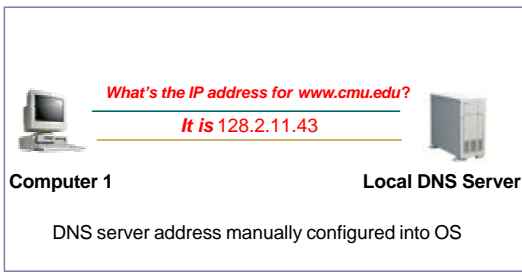


### Naming



- Humans use readable host names
  - E.g. [www.cmu.edu](http://www.cmu.edu)
  - Globally unique (can correspond to multiple hosts)
- Naming system translates to physical address
  - E.g. DNS translates name to IP Address (e.g. 128.2.11.43)
  - Address reflects location in network

### Domain Name System

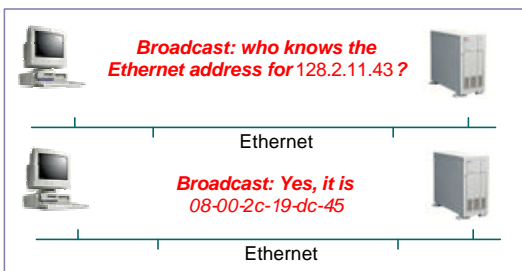


### Packet Routing/Delivery

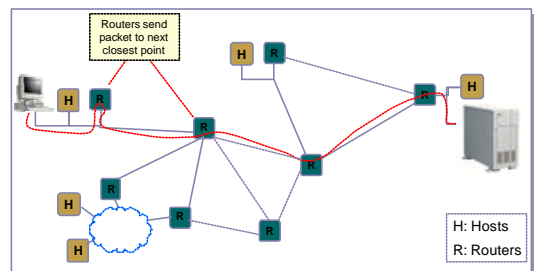


- Each network technology has different local delivery methods
- Address resolution provides delivery information within network
  - E.g., ARP maps IP addresses to Ethernet addresses
  - Local, works only on a particular network
- Routing protocol provides path through an internetwork

### Network: Address Resolution Protocol



### Internetwork: Datagram Routing



## Routing



- Forwarding tables at each router populated by routing protocols.
- Original Internet: manually updated
- Routing protocols update tables based on "cost"
  - Exchange tables with neighbors or everyone
  - Use neighbor leading to shortest path

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## Fourth Step: Application Demands



- Reliability
  - Corruption
  - Lost packets
- Flow and congestion control
- Fragmentation
- In-order delivery
- Etc...

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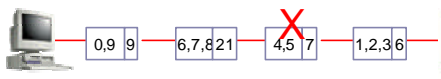
## What if the Data gets Corrupted?



Problem: Data Corruption



Solution: Add a *checksum*



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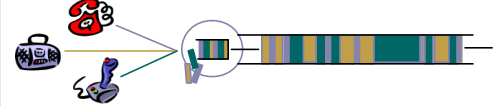
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## What if Network is Overloaded?



Problem: Network Overload



Solution: Buffering and Congestion Control

- Short bursts: buffer
- What if buffer overflows?
  - Packets dropped
  - Sender adjusts rate until load = resources
- Called "congestion control"

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## What if the Data gets Lost?



Problem: Lost Data



Solution: Timeout and Retransmit



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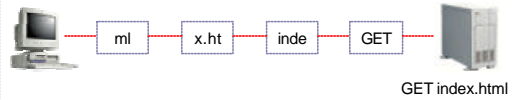
## What if the Data Doesn't Fit?



Problem: Packet size

- On Ethernet, max IP packet is 1.5kbytes
- Typical web page is 10kbytes

Solution: Fragment data across packets



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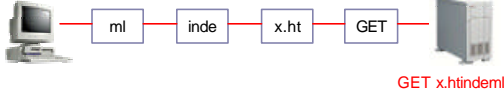
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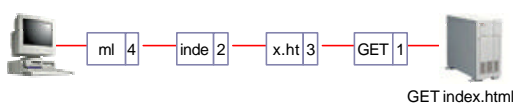
## What if the Data is Out of Order?



Problem: Out of Order



Solution: Add Sequence Numbers



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## Network Functionality Summary



- Link
- Multiplexing
- Routing
- Addressing/naming (locating peers)
- Reliability
- Flow control
- Fragmentation
- Etc....

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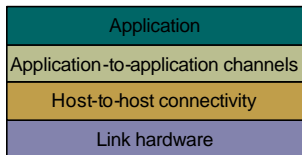
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## What is Layering?



- Modular approach to network functionality
- Example:



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## Protocols



- Module in layered structure
- Set of rules governing communication between network elements (applications, hosts, routers)
- Protocols define:
  - Interface to higher layers (API)
  - Interface to peer
    - Format and order of messages
    - Actions taken on receipt of a message

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## Layering Characteristics



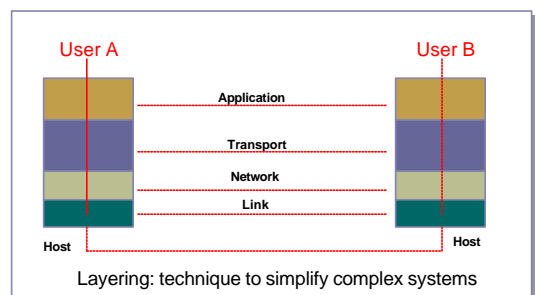
- Each layer relies on services from layer below and exports services to layer above
- Interface defines interaction
- Hides implementation - layers can change without disturbing other layers (black box)

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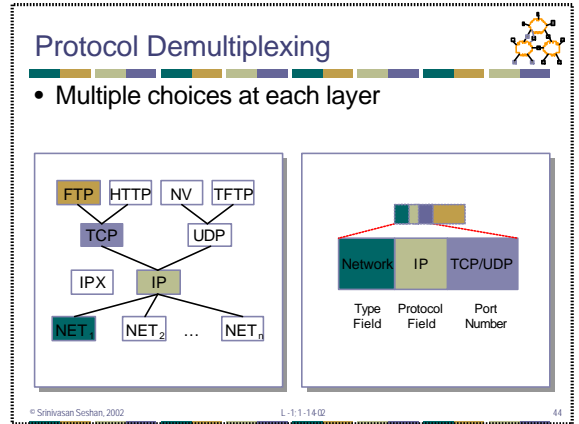
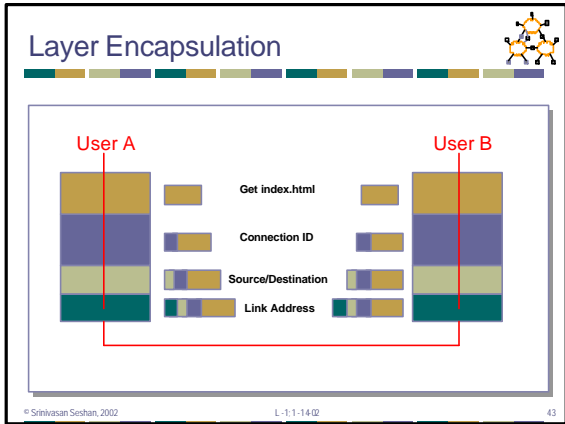
## Layering



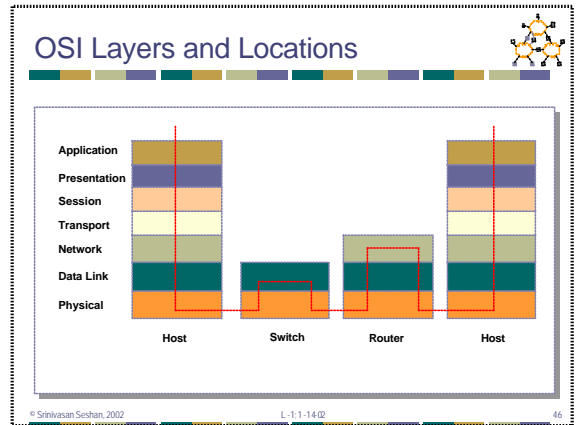
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- ### E.g.: OSI Model: 7 Protocol Layers
- Physical: how to transmit bits
  - Data link: how to transmit frames
  - Network: how to route packets
  - Transport: how to send packets end2end
  - Session: how to tie flows together
  - Presentation: byte ordering, security
  - Application: everything else
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- ### Example: Transport Layer
- First end-to-end layer
  - End-to-end state
  - May provide reliability, flow and congestion control
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- ### Example: Network Layer
- Point-to-point communication
  - Network and host addressing
  - Routing
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## Is Layering Harmful?



- Sometimes..
  - Layer N may duplicate lower level functionality (e.g., error recovery)
  - Layers may need same info (timestamp, MTU)
  - Strict adherence to layering may hurt performance

## Class Coverage



- No coverage of physical and data link layer
  - Students expected to know this
- Focus on network to application layer
- We will deal with:
  - Protocol rules and algorithms
  - Investigate protocol trade-offs
  - Why this way and not another?

## Lecture Topics



### Traditional

- Layering
- Internet architecture
- Routing (IP)
- Transport (TCP)
- Queue management (FQ, RED)
- Naming (DNS)

### Recent Topics

- Multicast
- Mobility
- Active networks
- QOS
- Security
- Network measurement
- Overlay networks
- P2P applications

## Next Lecture: Design Considerations



- How to determine split of functionality
  - Across protocol layers
  - Across network nodes
- Assigned Reading
  - [SRC84] End-to-end Arguments in System Design
  - [Cla88] Design Philosophy of the DARPA Internet Protocols