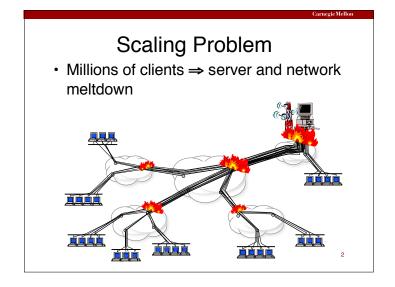
Peer-to-Peer

15-441



P2P System

• Leverage the resources of client machines (peers)

- Computation, storage, bandwidth

Why P2P?

Harness lots of spare capacity

- 1 Big Fast Server: 1Gbit/s, \$10k/month++

- 2,000 cable modems: 1Gbit/s, \$??

- 1M end-hosts: Uh, wow.

Build self-managing systems / Deal with huge scale

 Same techniques attractive for both companies / servers / p2p

• E.g., Akamai's 14,000 nodes

• Google's 100,000+ nodes

Outline

- p2p file sharing techniques
 - Downloading: Whole-file vs. chunks
 - Searching
 - Centralized index (Napster, etc.)
 - Flooding (Gnutella, etc.)
 - Smarter flooding (KaZaA, ...)
 - · Routing (Freenet, etc.)
- Uses of p2p what works well, what doesn't?
 - servers vs. arbitrary nodes
 - Hard state (backups!) vs soft-state (caches)
- Challenges
 - Fairness, freeloading, security, ...

P2P file-sharing

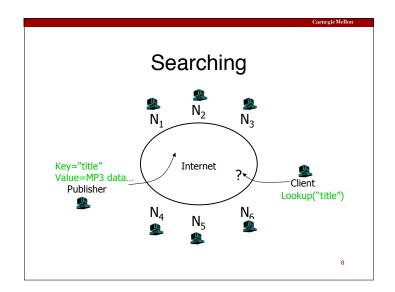
- · Quickly grown in popularity
 - Dozens or hundreds of file sharing applications
 - 35 million American adults use P2P networks --29% of all Internet users in US!
 - Audio/Video transfer now dominates traffic on the Internet

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What's out there?

	Central	Flood	Super- node flood	Route
Whole File	Napster	Gnutella		Freenet
Chunk Based	BitTorrent		KaZaA (bytes, not chunks)	DHTs eDonkey 2000



Searching 2

- · Needles vs. Haystacks
 - Searching for top 40, or an obscure punk track from 1981 that nobody's heard of?
- Search expressiveness
 - Whole word? Regular expressions? File names? Attributes? Whole-text search?
 - (e.g., p2p gnutella or p2p google?)

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Framework

- Common Primitives:
 - Join: how do I begin participating?
 - Publish: how do I advertise my file?
 - Search: how to I find a file?
 - **Fetch**: how to I retrieve a file?

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Next Topic...

- · Centralized Database
 - Napster
- Query Flooding
 - Gnutella
- · Intelligent Query Flooding
 - KaZaA
- Swarming
- BitTorrent
- Unstructured Overlay Routing
 - Freene
- Structured Overlay Routing
 - Distributed Hash Tables

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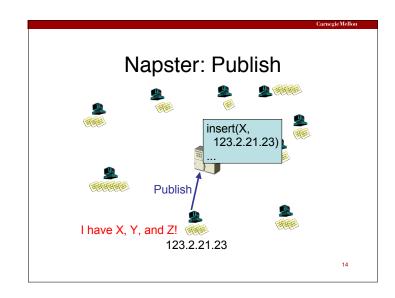
Napster: History

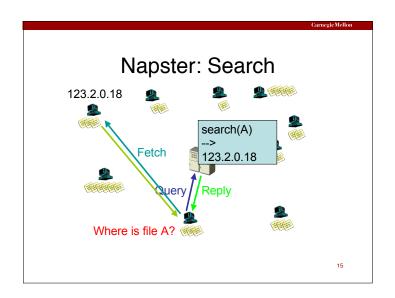
- 1999: Sean Fanning launches Napster
- Peaked at 1.5 million simultaneous users
- Jul 2001: Napster shuts down

Napster: Overiew

- Centralized Database:
 - Join: on startup, client contacts central server
 - Publish: reports list of files to central server
 - Search: query the server => return someone that stores the requested file
 - Fetch: get the file directly from peer

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Napster: Discussion

- Pros:
 - Simple
 - Search scope is O(1)
 - Controllable (pro or con?)
- · Cons:
 - Server maintains O(N) State
 - Server does all processing
 - Single point of failure

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Next Topic...

- Centralized Database
 - Napster
- · Query Flooding
 - Gnutella
- · Intelligent Query Flooding
 - KaZaA
- Swarming
 - BitTorrent
- Unstructured Overlay Routing
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Gnutella: History

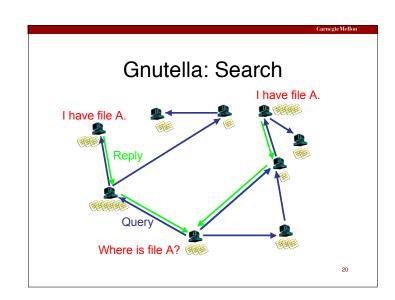
- In 2000, J. Frankel and T. Pepper from Nullsoft released Gnutella
- Soon many other clients: Bearshare, Morpheus, LimeWire, etc.
- In 2001, many protocol enhancements including "ultrapeers"

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Gnutella: Overview

- · Query Flooding:
 - Join: on startup, client contacts a few other nodes; these become its "neighbors"
 - Publish: no need
 - Search: ask neighbors, who ask their neighbors, and so on... when/if found, reply to sender.
 - TTL limits propagation
 - Fetch: get the file directly from peer



Gnutella: Discussion

- · Pros:
 - Fully de-centralized
 - Search cost distributed
 - Processing @ each node permits powerful search semantics
- Cons:
 - Search scope is O(N)
 - Search time is O(???)
 - Nodes leave often, network unstable
- · TTL-limited search works well for haystacks.
 - For scalability, does NOT search every node. May have to re-issue query later

KaZaA: History

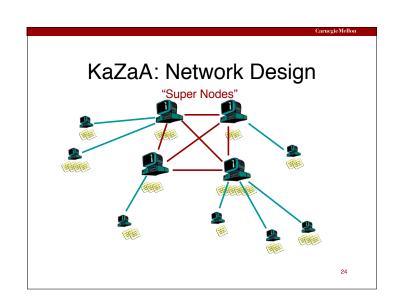
- In 2001, KaZaA created by Dutch company Kazaa BV
- Single network called FastTrack used by other clients as well: Morpheus, giFT, etc.
- Eventually protocol changed so other clients could no longer talk to it
- Most popular file sharing network today with >10 million users (number varies)

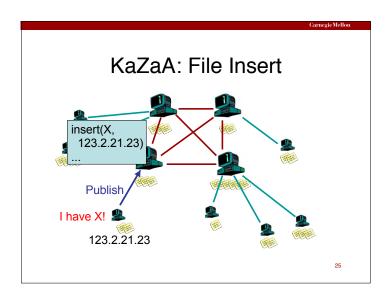
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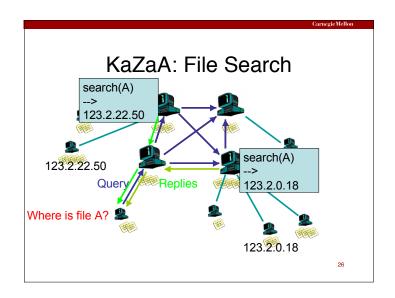
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KaZaA: Overview

- · "Smart" Query Flooding:
 - Join: on startup, client contacts a "supernode" ... may at some point become one itself
 - Publish: send list of files to supernode
 - Search: send query to supernode, supernodes flood query amongst themselves.
 - Fetch: get the file directly from peer(s); can fetch simultaneously from multiple peers







KaZaA: Fetching

- · More than one node may have requested file...
- · How to tell?
 - Must be able to distinguish identical files
 - Not necessarily same filename
 - Same filename not necessarily same file...
- · Use Hash of file
 - KaZaA uses UUHash: fast, but not secure
 - Alternatives: MD5, SHA-1
- · How to fetch?
 - Get bytes [0..1000] from A, [1001...2000] from B
 - Alternative: Erasure Codes

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KaZaA: Discussion

- Pros:
 - Tries to take into account node heterogeneity:
 - Bandwidth
 - · Host Computational Resources
 - Host Availability (?)
 - Rumored to take into account network locality
- Cons:
 - Mechanisms easy to circumvent
 - Still no real guarantees on search scope or search time
- · Similar behavior to gnutella, but better.

Stability and Superpeers

- · Why superpeers?
 - Query consolidation
 - · Many connected nodes may have only a few files
 - Propagating a query to a sub-node would take more b/w than answering it yourself
 - Caching effect
 - · Requires network stability
- Superpeer selection is time-based
 - How long you've been on is a good predictor of how long you'll be around.

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BitTorrent: History

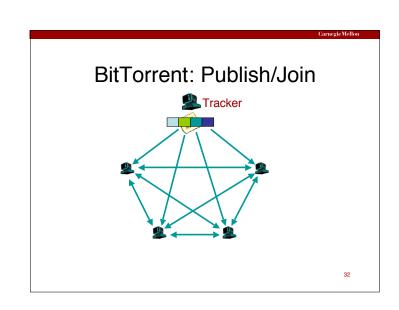
- In 2002, B. Cohen debuted BitTorrent
- · Key Motivation:
 - Popularity exhibits temporal locality (Flash Crowds)
 - E.g., Slashdot effect, CNN on 9/11, new movie/game release
- Focused on Efficient Fetching, not Searching:
 - Distribute the same file to all peers
 - Single publisher, multiple downloaders
- Has some "real" publishers:
 - Blizzard Entertainment using it to distribute the beta of their new game

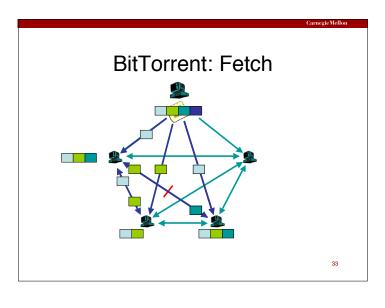
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BitTorrent: Overview

- Swarming:
 - Join: contact centralized "tracker" server, get a list of peers.
 - Publish: Run a tracker server.
 - Search: Out-of-band. E.g., use Google to find a tracker for the file you want.
 - Fetch: Download chunks of the file from your peers. Upload chunks you have to them.
- Big differences from Napster:
 - Chunk based downloading (sound familiar? :)
 - "few large files" focus
 - Anti-freeloading mechanisms





BitTorrent: Sharing Strategy

- · Employ "Tit-for-tat" sharing strategy
 - A is downloading from some other people
 - · A will let the fastest N of those download from him
 - Be optimistic: occasionally let freeloaders download
 - · Otherwise no one would ever start!
 - Also allows you to discover better peers to download from when they reciprocate
- Goal: Pareto Efficiency
 - Game Theory: "No change can make anyone better off without making others worse off"
 - Does it work? (don't know!)

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BitTorrent: Summary

- · Pros:
 - Works reasonably well in practice
 - Gives peers incentive to share resources; avoids freeloaders
- · Cons:
 - Pareto Efficiency relatively weak condition
 - Central tracker server needed to bootstrap swarm
 - (Tracker is a design choice, not a requirement.
 Could easily combine with other approaches.)

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Next Topic...

- · Centralized Database
 - Napster
- Query Flooding
 - Gnutella
- · Intelligent Query Flooding
- KaZaASwarming
- BitTorrent
- Unstructured Overlay Routing
 - Freene
- Structured Overlay Routing
 - Distributed Hash Tables (DHT)

Distributed Hash Tables

- · Academic answer to p2p
- Goals
 - Guaranteed lookup success
 - Provable bounds on search time
 - Provable scalability
- Makes some things harder
 - Fuzzy queries / full-text search / etc.
- · Read-write, not read-only
- Hot Topic in networking since introduction in ~2000/2001

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DHT: Chord Summary

- · Routing table size?
 - -Log N fingers
- · Routing time?
 - -Each hop expects to 1/2 the distance to the desired id => expect O(log N) hops.

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DHT: Discussion

- Pros:
 - Guaranteed Lookup
 - O(log N) per node state and search scope
- Cons:
 - No one uses them? (only one file sharing app)
 - Supporting non-exact match search is hard

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When are p2p / DHTs useful?

- · Caching and "soft-state" data
 - Works well! BitTorrent, KaZaA, etc., all use peers as caches for hot data
- Finding read-only data
 - Limited flooding finds hay
 - DHTs find needles
- BUT

A Peer-to-peer Google?

- Complex intersection gueries ("the" + "who")
 - Billions of hits for each term alone
- · Sophisticated ranking
 - Must compare many results before returning a subset to user
- Very, very hard for a DHT / p2p system
 - Need high inter-node bandwidth
 - (This is exactly what Google does massive clusters)

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P2P: Summary

- Many different styles; remember pros and cons of each
 - centralized, flooding, swarming, unstructured and structured routing
- Lessons learned:
 - Single points of failure are very bad
 - Flooding messages to everyone is bad
 - Underlying network topology is important
 - Not all nodes are equal
 - Need incentives to discourage freeloading
 - Privacy and security are important
 - Structure can provide theoretical bounds and guarantees

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Writable, persistent p2p

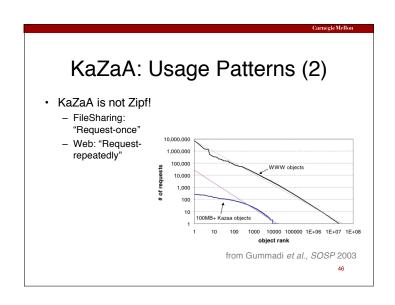
- Do you trust your data to 100,000 monkeys?
- · Node availability hurts
 - Ex: Store 5 copies of data on different nodes
 - When someone goes away, you must replicate the data they held
 - Hard drives are *huge*, but cable modem upload bandwidth is tiny - perhaps 10 Gbytes/day
 - Takes many days to upload contents of 200GB hard drive. Very expensive leave/replication situation!

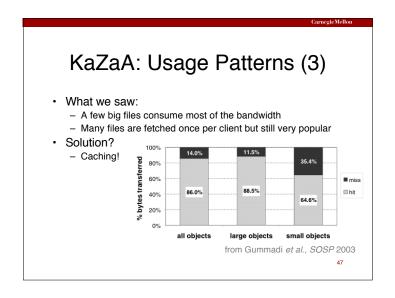
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Extra Slides

KaZaA: Usage Patterns • KaZaA is more than one workload! - Many files < 10MB (e.g., Audio Files) - Many files > 100MB (e.g., Movies) **Body of the company of the compa





Freenet: History

- In 1999, I. Clarke started the Freenet project
- · Basic Idea:
 - Employ Internet-like routing on the overlay network to publish and locate files
- · Addition goals:
 - Provide anonymity and security
 - Make censorship difficult

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Freenet: Overview

- · Routed Queries:
 - Join: on startup, client contacts a few other nodes it knows about; gets a unique node id
 - Publish: route file contents toward the file id. File is stored at node with id closest to file id
 - Search: route query for file id toward the closest node id
 - Fetch: when query reaches a node containing file id. it returns the file to the sender

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Freenet: Routing Tables

- id file identifier (e.g., hash of file)
- next_hop another node that stores the file id
- file file identified by id being stored on the local node
- Forwarding of query for file id
 - If file id stored locally, then stop
 - Forward data back to upstream requestor
 - If not, search for the "closest" id in the table, and forward the message to the corresponding next_hop
 - If data is not found, failure is reported back
 - Requestor then tries next closest match in routing

Id next_hop file

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Freenet: Routing Properties

- "Close" file ids tend to be stored on the same node
 - Why? Publications of similar file ids route toward the same place
- Network tend to be a "small world"
 - Small number of nodes have large number of neighbors (i.e., ~ "six-degrees of separation")
- Consequence:
 - Most queries only traverse a small number of hops to find the file

Freenet: Anonymity & Security

- Anonymity
 - Randomly modify source of packet as it traverses the network
 - Can use "mix-nets" or onion-routing
- · Security & Censorship resistance
 - No constraints on how to choose ids for files => easy to have to files collide, creating "denial of service" (censorship)
 - Solution: have a *id* type that requires a private key signature that is verified when updating the file
 - Cache file on the reverse path of queries/publications => attempt to "replace" file with bogus data will just cause the file to be replicated more!

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Freenet: Discussion

- · Pros:
 - Intelligent routing makes queries relatively short
 - Search scope small (only nodes along search path involved); no flooding
 - Anonymity properties may give you "plausible deniability"
- Cons:
 - Still no provable guarantees!
 - Anonymity features make it hard to measure, debug