

15-441 Computer Networking

Lecture 17 - Physical Layer, Link Layer Basics

Links

- How to make computers talk across a “wire”
- How to share the wire

Application
Transport
Network
Datalink
Physical

10-28-2002 Lecture 17 2

From Signals to Packets

Analog Signal

“Digital” Signal

Bit Stream 0 0 1 0 1 1 1 0 0 0 1

Packets

Packet Transmission

Sender Receiver

10-28-2002 Lecture 17 3

Link Layer: Implementation

- Implemented in “adapter”
 - E.g., PCMCIA card, Ethernet card
 - Typically includes: RAM, DSP chips, host bus interface, and link interface

application transport network link physical

data link protocol

phys. link

adapter card

frame

10-28-2002 Lecture 17 4

Outline

- Physical media is analog
 - Modulation – signals to bits
- Bit stream vs. packets
 - Framing – how to make packets
- Corruption
 - Error detection & recovery
- Sharing
 - Media access

10-28-2002 Lecture 17 5

Why do we care?

- How much bandwidth can I get out of a specific wire (transmission medium)?
- What limits the physical size of the network?
- How can multiple hosts communicate over the same wire at the same time?
- How can I manage bandwidth on a transmission medium?
- How do the properties of copper, fiber, and wireless compare?

10-28-2002 Lecture 17 6

Modulation

- Modulation : method of encoding digital (or analogue) signals onto a waveform
- Sender changes the nature of the signal in a way that the receiver can recognize.
 - Similar to radio: AM or FM
- Digital transmission: encodes the values 0 or 1 in the signal.
 - It is also possible to encode multi-valued symbols
- Amplitude modulation: change the strength of the signal, typically between on and off.
 - Sender and receiver agree on a "rate"
 - On means 1, Off means 0
- Similar: frequency or phase modulation.
- Can also combine modulation types.

10-28-2002 Lecture 17 7

Amplitude and Frequency Modulation

10-28-2002 Lecture 17 8

Modulation

- A (periodic) signal can be viewed as a sum of sine waves of different strengths.
 - Corresponds to energy at a certain frequency
- Every signal has an equivalent representation in the frequency domain.
 - What frequencies are present and what is their strength (energy)
- Similar to radio and TV signals.

10-28-2002 Lecture 17 9

Modulation

- Noise: "random" energy is added to the signal.
- Attenuation: some of the energy in the signal leaks away.
- Dispersion: attenuation and propagation speed are frequency dependent.
 - Changes the shape of the signal
- **Effects limit the data rate that a channel can sustain.**
 - » But affects different technologies in different ways
- **Effects become worse with distance.**
 - » Tradeoff between data rate and distance

10-28-2002 Lecture 17 10

Modulation

- What is the bit pattern?
 - Easy - 101010...
- What about now?
 - 111111000000 or 11110000?
 - Hard to tell where a bit starts and ends
- What about now?
 - 11111111 or 00000000?
 - Hard to tell when what is a 0 and what is a 1
- Any suggestions to solve this?

10-28-2002 Lecture 17 11

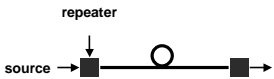
Modulation

- Manchester
 - Used by Ethernet
 - 0=low to high transition, 1=high to low transition
 - Transition for every bit simplifies clock recovery
 - Not very efficient
 - Doubles the number of transitions
 - Circuitry must run twice as fast

10-28-2002 Lecture 17 12

Copper Wire


- Unshielded twisted pair
 - Two copper wires twisted - avoid antenna effect
 - Grouped into cables: multiple pairs with common sheath
- Coax cables.
 - One conductor is placed inside the other conductor
 - Holds the signal in place and keeps out noise
- At end of span, either regenerate electronically or amplify.
- Electronic repeaters are potentially slow, but can eliminate noise.



10-28-2002 Lecture 17 13

Fiber

- Uses total internal reflection
- usually laser diode or LED source
- still subject to chromatic dispersion
- Need repeaters



10-28-2002 Lecture 17 14

Wireless Technologies

- Great technology: no wires to install, convenient mobility, ..
- High attenuation limits distances.
 - Wave propagates out as a sphere
 - Signal strength reduces quickly $(1/\text{distance})^3$
- High noise due to interference from other transmitters.
 - Use MAC and other rules to limit interference
 - Aggressive encoding techniques to make signal less sensitive to noise
- Other effects: fading, security, ..

10-28-2002 Lecture 17 15

Media Comparison

Media	BW	Range	Wiring Cost
Twisted Pair	100Mbps	100m	10c/ft
Coax	1Gbps	1km	10c/ft
Fiber	1Gbps	10km	30c/ft
Wireless	22Mbps	500ft	0

10-28-2002 Lecture 17 16

Outline

- Physical media is analog
 - Modulation – signals to bits
- Bit stream vs. packets
 - Framing – how to make packets
- Corruption
 - Error detection & recovery
- Sharing
 - Media access

10-28-2002 Lecture 17 17

Framing

- A link layer function, defining which bits have which function.
- Minimal functionality: mark the beginning and end of packets (or frames).
- Some techniques:
 - frame delimiter characters with character stuffing
 - frame delimiter codes with bit stuffing
 - out of band delimiters (e.g. FDDI control symbols)
 - synchronous transmission (e.g. SONET)

10-28-2002 Lecture 17 18

Framing

- Length delimited
 - Beginning of frame has length
 - Single corrupt length can cause problems
 - Must have start of frame character to resynchronize
 - Resynchronization can fail if start of frame character is inside packets as well

10-28-2002 Lecture 17 19


Framing

- Byte stuffing
 - Special start of frame byte (e.g. 0xFF)
 - Special escape byte value (e.g. 0xFE)
 - Values actually in text are replaced (e.g. 0xFF by 0xFEFF and 0xFE by 0xFEFE)
 - Worst case – can double the size of frame
- Bit stuffing
 - Special bit sequence (0x01111110)
 - 0 bit stuffed after any 11111 sequence

10-28-2002 Lecture 17 20

Ethernet Framing

- Preamble is 7 bytes of 10101010 followed by one byte of 10101011
 - Does not use byte/bit stuffing for preamble
 - Allows receivers to recognize start of transmission after idle channel



10-28-2002 Lecture 17 21

Consistent Overhead Byte Stuffing

- Run length encoding applied to byte stuffing
- Encoding
 - Add implied 0 to end of frame
 - Each 0 is replaced with (number of bytes to next 0) + 1
 - What if no 0 within 255 bytes? 255 value indicates 254 bytes followed by no zero
 - Worst case – no 0's in packet – 1/254 overhead
 - Possible optimization to encode series of 0's

10-28-2002 Lecture 17 22

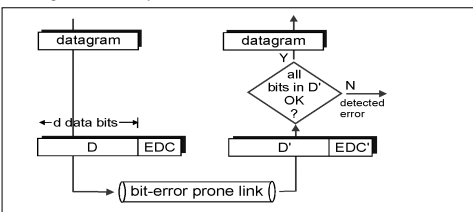
Outline

- Physical media is analog
 - Modulation – signals to bits
- Bit stream vs. packets
 - Framing – how to make packets
- Corruption
 - Error detection & recovery
- Sharing
 - Media access

10-28-2002 Lecture 17 23

Error Detection

- EDC= Error Detection and Correction bits (redundancy)
- D = Data protected by error checking, may include header fields
- Error detection not 100% reliable!
 - Protocol may miss some errors, but rarely
 - Larger EDC field yields better detection and correction



10-28-2002 Lecture 17 24

Parity Checking

Single Bit Parity:
Detect single bit errors

← d data bits → parity bit

0111000110101011 0

10-28-2002 Lecture 17 25

Error Detection - Checksum

- Used by TCP, UDP, IP, etc..
- Ones complement sum of all words/shorts/bytes in packet
- Simple to implement
- Relatively weak detection
 - Easily tricked by common error patterns

10-28-2002 Lecture 17 26

Internet Checksum

- Goal: detect "errors" (e.g., flipped bits) in transmitted segment – Already covered in lecture 7

Sender	Receiver
<ul style="list-style-type: none"> Treat segment contents as sequence of 16-bit integers Checksum: addition (1's complement sum) of segment contents Sender puts checksum value into checksum field in header 	<ul style="list-style-type: none"> Compute checksum of received segment Check if computed checksum equals checksum field value: <ul style="list-style-type: none"> NO - error detected YES - no error detected. But maybe errors nonetheless?

10-28-2002 Lecture 17 27

Error Detection – Cyclic Redundancy Check (CRC)

- Polynomial code
 - Treat packet bits a coefficients of n-bit polynomial
 - Choose r+1 bit generator polynomial (well known – chosen in advance)
 - Add r bits to packet such that message is divisible by generator polynomial
- Better loss detection properties than checksums

10-28-2002 Lecture 17 28

Error Detection – CRC

- View data bits, D, as a binary number
- Choose r+1 bit pattern (generator), G
- Goal: choose r CRC bits, R, such that
 - $\langle D, R \rangle$ exactly divisible by G (modulo 2)
 - Receiver knows G, divides $\langle D, R \rangle$ by G. If non-zero remainder: error detected!
 - Can detect all burst errors less than r+1 bits
- Widely used in practice (ATM, HDCL, Ethernet)

← d bits → ← r bits → bit pattern

D: data bits to be sent | R: CRC bits

$D \cdot 2^r \text{ XOR } R$ mathematical formula

10-28-2002 Lecture 17 29

CRC Example

Want:
 $D \cdot 2^r \text{ XOR } R = nG$
 equivalently:
 $D \cdot 2^r = nG \text{ XOR } R$
 equivalently:
 if we divide $D \cdot 2^r$ by G,
 want remainder R

$R = \text{remainder} \left[\frac{D \cdot 2^r}{G} \right]$

10-28-2002 Lecture 17 30

Error Recovery

- Two forms of error recovery
 - Error Correcting Codes (ECC)
 - Automatic Repeat Request (ARQ)
- ECC
 - Send extra redundant data to help repair losses
- ARQ
 - Receiver sends acknowledgement (ACK) when it receives packet
 - Sender uses ACKs to identify and resend data that was lost
 - Issues already studied in transport layer

10-28-2002 Lecture 17 31

Error Recovery – Error Correcting Codes (ECC)

Two Dimensional Bit Parity:
Detect and correct single bit errors

				row parity
	$d_{1,1}$...	$d_{1,j}$	$d_{1,j+1}$
	$d_{2,1}$...	$d_{2,j}$	$d_{2,j+1}$

	$d_{i,1}$...	$d_{i,j}$	$d_{i,j+1}$
column parity	$d_{i+1,1}$...	$d_{i+1,j}$	$d_{i+1,j+1}$

101011	101011	parity
111100	101100	error
011101	011101	
001010	001010	parity error

no errors
correctable single bit error

10-28-2002 Lecture 17 32

Outline

- Physical media is analog
 - Modulation – signals to bits
- Bit stream vs. packets
 - Framing – how to make packets
- Corruption
 - Error detection & recovery
- Sharing
 - Media access

10-28-2002 Lecture 17 33

MAC Protocols: A Taxonomy

Three broad classes:

- Channel partitioning
 - Divide channel into smaller “pieces” (time slots, frequency)
 - Allocate piece to node for exclusive use
- Random access
 - Allow collisions
 - “Recover” from collisions
- “Taking turns”
 - Tightly coordinate shared access to avoid collisions

Goal: efficient, fair, simple, decentralized

10-28-2002 Lecture 17 34

Supporting Multiple Channels

- Multiple channels can coexist if they transmit
 - at a different frequency
 - at a different time
 - in a different part of the space.
 - Three dimensional space: frequency, space, time
- Space can be limited using wires or using transmit power of wireless transmitters.
- Frequency multiplexing means that different users use a different part of the spectrum.
 - similar to radio: 95.5 versus 102.5 station
- Controlling time is a datalink protocol issue.
 - Media Access Control (MAC): who gets to send when?

10-28-2002 Lecture 17 35

Baseband vs Carrier Modulation

- Baseband modulation: send the “bare” signal.
- Carrier modulation: use the signal to modulate a higher frequency signal (carrier).
 - Can be viewed as the product of the two signals
 - Corresponds to a shift in the frequency domain

10-28-2002 Lecture 17 36

Amplitude Carrier Modulation

10-28-2002 Lecture 17 37

Frequency Division Multiplexing: Multiple Channels

10-28-2002 Lecture 17 38

Wavelength Division Multiplexing

- Send multiple wavelengths through the same fiber.
 - Multiplex and demultiplex the optical signal on the fiber
- Each wavelength represents an optical carrier that can carry a separate signal.
 - E.g., 16 colors of 2.4 Gbit/second
- Like radio, but optical and much faster

10-28-2002 Lecture 17 39

Multiple Access Protocols

- Single shared communication channel
- Two or more simultaneous transmissions → interference
 - Only one node can send successfully at a time
- Multiple access protocol:*
 - Distributed algorithm that determines how stations share channel, i.e., determine when station can transmit
 - Communication about channel sharing must use channel itself!
- What to look for in multiple access protocols:
 - Synchronous or asynchronous
 - Information needed about other stations
 - Robustness (e.g., to channel errors)
 - Performance

10-28-2002 Lecture 17 40

Summary

- Bandwidth and distance of networks is limited by physical properties of media.
 - Attenuation, noise, ...
- Network properties are determined by transmission medium and transmit/receive hardware.
 - Can do much better with better encoding
- Framing helps provide control functionality
- Error detection and recovery mechanisms depend on error rate of the physical medium
- Multiple users can be supported using space, time, or frequency division multiplexing.
- Next lecture: How to do media access control and much more ☺

10-28-2002 Lecture 17 41