PFPL Supplement: Types and Parallelism

Robert Harper
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1 Introduction

The treatment of fork-join parallelism distinguishes two cases, the static and dynamic, according to when the “degree” of parallelism is determined. The static case is formulated using the construct \( \text{par} e_1 = e_2 \text{ and } x_1 = x_2 \text{ in } e \), whose dynamics evaluates \( e_1 \) and \( e_2 \) in parallel, then substitutes their values for the corresponding variables in \( e \). The dynamic case is formulated using a tabulator that creates a sequence of values whose length and content are determined at run-time. All elements are evaluated in parallel, then creating a sequence consisting of their values. The dynamic case can be understood as a cascade of binary parallel splits logarithmic in the length of the sequence.

Although the tabulation mechanism is tied to a type of sequences, the parallel definition mechanism is not associated with a type. It would be more elegant if the definition construct were associated with a type, rather than simply imposed as an ad hoc means to an end. And it would be nice if the dynamic case were a smooth generalization of the static case. Might there be a more systematic solution that encompasses both forms of parallelism?

There is, but to see it requires taking a step back to think about how parallelism arises in the first place. The key is not to somehow impose parallelism in any otherwise sequential dynamics, but rather to expose the parallelism that is naturally present to the maximal extent possible. Eager functional languages give rise to parallelism without even trying. For example, to evaluate a sum of two expressions it is necessary to obtain the value of both summands, but the only constraint on how these evaluations are to be scheduled is that they both be completed before the addition is performed. This example is prototypical: the essence of parallelism is sequentiality. Get the essential dependencies right and the parallelism will take care of itself.

2 Cost Semantics for PCFv

First, let us consider how to express sequential dependencies. The best way to do this is based on the language PCFv (Harper, 2019). The statics is reproduced in Figure 1 for convenience. For the sake of clarity it uses the meta-variable \( v \) for values and \( e \) for computations. The computation \( \text{ret}(v) \) simply returns the given value.

The cost dynamics of PCFv is given in Figure 2, defining the judgment \( e \Downarrow^c v \) is defined for closed computations \( e \), closed values \( v \), and cost graphs \( c \).

3 Parallelism, Revisited

The modal framework allows us to manage dependencies, but it is limited to one dependency at a time, precluding parallelism. What is missing is a way to express the simultaneous dependency of one computation on several, perhaps unboundedly many, prior computations whose relative evaluation order is unconstrained.
\[\Gamma, x : \tau \vdash x : \tau\] (1a)

\[\Gamma \vdash z : \text{nat}\] (1b)

\[\Gamma \vdash v : \text{nat}\]

\[\Gamma \vdash s(v) : \text{nat}\] (1c)

\[\Gamma, x : \tau_1 \vdash \tau_2, y : \tau_1 \vdash e : \tau_2\]

\[\Gamma \vdash \text{fun}\{\tau_1 : \tau_2\}(x \cdot y \cdot e) : \tau_1 \rightarrow \tau_2\] (1d)

\[\Gamma \vdash e \sim \cdot \cdot \tau_2\]

\[\Gamma \vdash \text{fun}\{\tau_1 ; \tau_2\}(x \cdot y \cdot e) : \tau_1 \rightarrow \tau_2\] (1e)

\[\Gamma \vdash v : \tau\]

\[\Gamma \vdash \text{ret}(v) \sim \tau\] (1f)

\[\Gamma \vdash v : \tau\]

\[\Gamma \vdash e_0 \sim \tau\]

\[\Gamma, x : \text{nat} \vdash e_1 \sim \tau\]

\[\Gamma \vdash \text{ifz}\{\tau\}(v; e_0; x \cdot e_1) \sim \tau\] (1g)

\[\Gamma \vdash v_1 : \tau_2 \rightarrow \tau\]

\[\Gamma \vdash v_2 : \tau_2\]

\[\Gamma \vdash \text{ap}(v_1; v_2) \sim \tau\] (1h)

\[\Gamma \vdash v_1 : \text{comp}(\tau_1)\]

\[\Gamma, x : \tau_1 \vdash e_2 \sim \tau_2\]

\[\Gamma \vdash \text{bnd}(v_1; x \cdot e_2) \sim \tau_2\] (1i)

Figure 1: Statics of PCFv

\[\text{ret}(v) \Downarrow^1 v\] (2a)

\[e_0 \Downarrow^c v\]

\[\text{ifz}\{\tau\}(z; e_0; x \cdot e_1) \Downarrow^{\text{inc}} v\] (2b)

\[e/x\ e_1 \Downarrow^c v\]

\[\text{ifz}\{\tau\}(s(e); e_0; x \cdot e_1) \Downarrow^{\text{inc}} v\] (2c)

\[\text{fun}\{\tau_2 : \tau\}(f \cdot x \cdot e), v_2/f, x \cdot e \Downarrow^c v\]

\[\text{ap}\{\tau_2 : \tau\}(f \cdot x \cdot e); v_2) \Downarrow^{\text{inc}} v\] (2d)

\[e \Downarrow^c v\]

\[e \Downarrow^c [v/x] e_2 \Downarrow^{\text{inc}} v_2\]

\[\text{bnd}(\text{comp}(e); x \cdot e_2) \Downarrow^{\text{inc} \cdot \text{inc}} v_2\] (2e)

Figure 2: Cost Dynamics of PCFv
\[\Gamma \vdash v_1 : \tau_1, \quad \Gamma \vdash v_2 : \tau_2\]
\[\Gamma \vdash \otimes(v_1; v_2) : \tau_1 \otimes \tau_2\]
\[\Gamma \vdash v : \tau_1 \otimes \tau_2, \quad \Gamma, x_1 : \tau_1, x_2 : \tau_2 \vdash e \downarrow \tau\]
\[\Gamma \vdash \text{split}(v; x_1, x_2, e) \downarrow \tau\]  \hspace{1cm} (3a)

\[\Gamma \vdash e_1 \not\downarrow \tau_1, \quad \Gamma \vdash e_2 \not\downarrow \tau_2\]
\[\Gamma \vdash \& (e_1; e_2) : \tau_1 \& \tau_2\]
\[\Gamma \vdash v : \tau_1 \& \tau_2, \quad \Gamma, x : \tau_1 \otimes \tau_2 \vdash e \downarrow \tau\]
\[\Gamma \vdash \text{parbnd}(v; x, e) \downarrow \tau\]  \hspace{1cm} (3b)

\[\Gamma \vdash \& (e_1; e_2) \text{ val}\]
\[e_1 \downarrow^{c_1} v_1, \quad e_2 \downarrow^{c_2} v_2, \quad [\otimes(v_1; v_2)/x]e \downarrow^{c} v\]
\[\text{parbnd}(\& (e_1; e_2); x, e) \Downarrow^{1\oplus (c_1 \oplus c_2) \oplus c} v\]  \hspace{1cm} (4)

\[\Gamma \vdash e_1 \text{ val}, \quad \Gamma \vdash e_2 \text{ val}\]
\[\otimes(e_1; e_2) \text{ val}\]
\[\quad \Downarrow^{1\oplus(c_1 \oplus c_2) \oplus c} v\]  \hspace{1cm} (5)

Figure 3: Statics of Eager and Lazy Products

\[\Gamma \vdash v_1, v_2/x_1, x_2 \text{ val}\]
\[\Downarrow^{1\oplus(c_1 \oplus c_2)} v\]  \hspace{1cm} (6a)

\[\text{parbnd}(\& (e_1; e_2); x, e) \Downarrow^{1\oplus (c_1 \oplus c_2) \oplus c} v\]  \hspace{1cm} (6b)

\[\Gamma \vdash \otimes(v_1; v_2) \text{ val}\]
\[\Downarrow^{1\oplus(c_1 \oplus c_2)} v\]  \hspace{1cm} (6c)

\[\text{split}(\otimes(v_1; v_2); x_1, x_2, e) \Downarrow^{1\oplus(c_1 \oplus c_2)} v\]  \hspace{1cm} (6d)

Figure 4: Cost Dynamics of Eager and Lazy Products
\[
\Gamma \vdash v : \text{nat} \quad \Gamma, x : \text{nat} \vdash e \leadsto \tau \\
\Gamma \vdash \text{seqgen}(v \cdot x \cdot e) : \text{seqgen}(\tau) \\
\Gamma \vdash v_1 : \text{seqgen}(\tau) \quad \Gamma, x : \text{seq}(\tau) \vdash e_2 \leadsto \tau_2 \\
\Gamma \vdash \text{seqbnd}(v_1 \cdot x \cdot e_2) \leadsto \tau_2
\] (7a)

(7b)

\[
\text{seqgen}(\pi \cdot x \cdot e_2) \text{ val}
\]

\[
\vdash_0 x e \Downarrow^v_0 v_0 \ldots \left[\vdash_{n-1} x e \Downarrow^v_{n-1} v_{n-1}\right] [\text{seq}\{n\}(v_0, \ldots, v_{n-1})/y] e' \Downarrow^v' v'
\]

\[
\text{seqbnd}\left(\text{seqgen}(\pi \cdot x \cdot e) ; y \cdot e'\right) \Downarrow(v_0 \otimes \cdots \otimes v_{n-1} \oplus v')
\] (7d)

Figure 5: Statics and Cost Dynamics of Sequence Generators

This is provided by the lazy product type, which encapsulates two unevaluated computations. The elimination form, called parallel bind, correspondingly evaluates both computations and binds a variable to the eager pair of their values.\(^{1}\)

The statics for lazy and eager product types is given in Figure 3. The corresponding cost dynamics is given in Figure 4.

There are two types corresponding to the lazy and eager product types, but with dynamically determined sizes. The analogue of a lazy tuple is a sequence generator, which, when evaluated, determines the width and the components of a finite sequence. The analogue of an eager pair is a sequence whose length, and the value of each element, are determined dynamically. The elimination form for the generator type creates a new sequence for use within the specified scope. The standard sequence operations given in Harper (2016) are then used to compute with the sequence.

The combined statics and dynamics for generators is given in Figure 5.

**Exercise 3.1.** Extend the dynamics to account for exceptions using the generalized \textbf{bnd} construct given in Harper (2018).

**References**


\(^{1}\)The binary case may easily be generalized to any \(n \geq 0\) suspended computations evaluated in parallel.