Practical Foundations for Programming Languages (Second Edition)

Errors and Corrections

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Corrections

The following errors in the second edition have been called to my attention.

- Back cover: Andrew Pitts’s affiliation is printed as “University of Cambridge University Press” rather than “University of Cambridge.” (This is the only mistake that is not under my control!)

- Page 4: In the description of abt’s, operators with no children should also be deemed “leaves.”

- Page 23, after displayed formula (3.6), “even from rules (2.2)” should read “even from rules (2.8)”.


- Page 82, Section 10.3: the definition of the recursor should be

  $$\text{iter}\{\langle z,e_0\rangle;x'.\langle s(x' \cdot 1),[x' \cdot 1,x' \cdot r/x,y]e_1\rangle\}(e).$$

- Page 75, Exercise 9.6: “then $e$ is also hereditarily terminating . . . ” should be “then $e'$ is also hereditarily terminating . . . ”.

- Page 82, last line: the type of the bound variable $n$ is given as $\text{nat} \times \text{nat}$, but should be $\text{nat}$.

- Page 88, first paragraph: replace “induction” by “inductive”.

- Page 90, “which” in the abstract syntax should be “ifnull”, consistently with the concrete syntax.

- Page 107, $u_1, \ldots, u_n$ should be $u_1, \ldots, u_m$. 
- Page 128 et seq.: Rules (15.2d), (15.4f), (15.6d), (15.9d), and (15.9h) lack premises to require that the principal argument be a value as a condition on taking a step.
- Page 128, Rule (15.2a): add an optional premise $e$ val, and an optional rule to evaluate $e$ if the introduction form is to be eager.
- Page 129, Rule (15.4a): add an optional premise $e$ val, and an optional rule to evaluate $e$ if the generator is to be eager.
- Page 129, Rule (15.6a): add an optional premise $e_2$ val and a corresponding rule to evaluate $e_2$ if the introduction form is to be eager.
- Page 140, Rule (16.4a): the type label on the $\lambda$ should be $\tau_1$, not $\tau$.
- Page 147, line 13: “identification convention”, not “identification covnetion”.
- Page 154, and Exercise 1: the type of $\text{emp}$ should be $\forall t :: T.q[t]$.
- Page 152, line -8: $\langle bs, fs' \rangle$ should be $\text{just}(\langle f, \langle bs, fs' \rangle \rangle)$.
- Page 162, line 3: “In other words $e$ is defined to be . . .” should be “In other words $f$ is defined to be . . .”
- Page 171: “fucntion” should be “function.”
- Page 191, dynamics of application: either a by-name or by-value interpretation is possible.
- Page 191, Rule (22.4h): 0 should be $z$.
- Page 191, Rule (22.4i): $n + 1$ should be $s(n)$.
- Page 219, paragraph 2: should read: “serve as behavioral specifications”.
- Page 228, Section 25.3: “$\Phi \Gamma \vdash a \in \tau \phi$” should be “$\Phi \Gamma \vdash e \in \tau \phi$.”
- Page 261, Rule (29.3a): Conclusion should be $e \triangleright e$ initial.
- Page 265, Exercises 29.4 and 29.5 are orphaned: the modal formulation of exceptions is no longer presented in the text.
- Page 267, Section 30.1: $\tau$ cont cont should be $\tau'$ cont cont.
- Page 268, rule 30.2: the premise should read “$k \div \tau.$”
- Page 283, Exercise 31.4: There is no reason to specify the associated type of a symbol to be $\text{sexpr}$. For “pure” symbols the associated type would be $\text{unit}$, indicating no associated information. Alternatively, the associated type could be a product of attributes, including, for example, the “print name” (a string) associated with the symbol.
• Page 288, Section 32.5: “...the name of the fluid...”

• Page 291, Section 33.1: “A dynamic class is a symbol that is generated...”

• Page 292, last paragraph: Revised per online edition to clarify discussion of disequality of names.

• Page 294, line 11: `isin[a](e;x,e1;e2)` should be `isin[a](e;y,e1;e2)` to align the variable name `y` with its definition to follow.

• Page 298, Exercises 33.3 and 33.4 depend on Exercises 29.4 and 29.5, which rely on the modal formulation of exceptions.

• Page 314, first display equation: the printed edition inexplicably has 
  \[ \tau_{cap} \triangleq \tau_{cmd} \times (\text{nat} \rightarrow \text{nat cmd}) \]
  which instead should be
  \[ \tau_{cap} \triangleq \tau_{cmd} \times (\tau \rightarrow \tau_{cmd}). \]

• Page 344, start of paragraph 2: should read: “finite mapping of the task names”.

• Pages 344-5: Rule 37.10(b) should have a premise \( \neg(e_2 \text{ val}) \) to prevent needless repetition. Rule 37.10(c) should send \( a_1 \) to \( e_1(e_2) \) in the result state, not to the senseless substitution instance.

• Page 366, Section 39.4, displayed example at end of section should be
  \[ (\nu v.(V_0 \otimes V_1)) \otimes U \overset{\Sigma}{\longrightarrow} (\nu v.V_1 \otimes V_2) \otimes U \equiv (V_1 \otimes V_2 \otimes U). \]

• Page 380, displayed program has a missing use of `cmd`, should be
  \[
  \text{fix loop: } \tau_{cmd} \text{ is cmd}\{x \leftarrow \text{acc}; \text{match } x \text{ as } a \cdot y \rightarrow \text{ret } y \text{ ow } \rightarrow \text{emit}(x); \text{do loop}\}
  \]

• Page 381, Rules 4.10. There is no execution rule for `sync(\text{never})`, so that the intended, but unstated, progress theorem fails. A solution is to introduce a new action, say `\#`, representing synchronization on the never-occurring event, and to add the execution rule
  \[ \text{sync(never)} \overset{\#}{\longrightarrow} \text{sync(never)}. \]
  The same correction applies to `DA`, which also has a null event.

• Page 387, Rule 41.2(b): Premise \( \Gamma \vdash e_1 : \tau_1 \text{ cmd } w \) should be \( \Gamma \vdash e_1 : \text{cmd}[w](\tau_1) @ w. \)
• Page 389, Rule 41.6a: The process “run( m)” should be “run( at[w]( m))” to record the site of the spawned process.

• Page 390, Theorems 41.2 and 41.3, the transition judgments should be $p \xrightarrow{w} \alpha \Sigma p'$; the “w” should be above the transition arrow, not below it.

• Page 437, last paragraph: strike “obviously reflexive.” Reflexivity is exactly Theorem 46.13.

• Page 438, Section 46.1, definition of a congruence: require only that the relations be partial equivalence relations (symmetric and transitive, not necessarily reflexive).

• Page 442, proof of Lemma 46.16: Add “By Definition 46.8 and Lemma 46.11 . . . .” for clarity.

• Page 447, Section 47.3, “the following rule”.

• Page 448, Section 47.3, should be “. . . it is enough to show that $\gamma(e_1) \sim_\tau \gamma'(e'_1), . . . .”

• Page 456, Section 48.2, statement and proof of Lemma 48.3, $\Delta, t$ should be $\Delta, t$ type.

• Page 459, Section 48.3, proof of Lemma 48.8: second $\rho$ should be $\rho'$:

\[ d'[\rho'] \cong_{\rho'/[t]} \delta'(\tau_2) e'[\rho']. \]

• Page 462, line 8: the second occurrence of $\rho$ should be $\rho'$.

### Improvements

The following changes improve the presentation in the text.

• Various places: ensure uniform use of “statics” and “dynamics” rather than “static semantics” and “dynamic semantics.”

• Change rules (5.10a) and (5.10i) to have typing premises so that it is possible to prove the newly inserted Exercise 5.4 by induction on Rules (5.10).

• Chapter 6, Section 6.2: The statement and proof of the canonical forms lemma could be improved as follows:

**Lemma 1** (Canonical Forms).

1. If $e : \text{num}$ and $e \text{ val}$, then $e = n$ for some number $n$.

2. If $e : \text{str}$ and $e \text{ val}$, then $e = “s”$ for some string $s$. 
\[ \text{fold}(t.\tau\{e\}) \mapsto \text{fold}(t.\tau\{e'\}) \]

\[ e \mapsto e' \]

\[ \text{fold}(t.\tau\{e\}) \mapsto \text{fold}(t.\tau\{e'\}) \]

\[ e_2 \mapsto e'_2 \]

\[ \text{rec}(t.\tau\{x.e_1;e_2\}) \mapsto \text{rec}(t.\tau\{x.e_1;e'_2\}) \]

\[ [e_2] \mapsto e'_2 \]

\[ \text{gen}(t.\tau\{x.e_1;e_2\}) \mapsto \text{gen}(t.\tau\{x.e_1;e'_2\}) \]

\[ e \mapsto e' \]

\[ \text{unfold}(t.\tau\{e\}) \mapsto \text{unfold}(t.\tau\{e'\}) \]

\[ [\text{map}^+(t.\tau\{x.e_1;e_2\}(e_2)/x|e_1)] \]

\[ [\text{gen}^+(t.\tau\{x.e_1;e_2\}) \mapsto \text{gen}^+(t.\tau\{x.e_1;e'_2\})] \]

\[ e \mapsto e' \]

\[ \text{gen}(t.\tau\{x.e_1;e'_2\}) \mapsto \text{gen}(t.\tau\{x.e_1;e'_2\}) \]

\[ \text{unfold}(t.\tau\{\text{gen}(t.\tau\{x.e_1;e'_2\})\}) \mapsto \text{map}^+(t.\tau\{y.\text{gen}(t.\tau\{x.e_1;e_2\})([e_2/x|e_1])\}) \]

Figure 1: Eager-or-Lazy Dynamics for Inductive and Coinductive Types

**Proof.** Each case is proved by induction on typing, making use of the definition of \( e \) val.

- Chapter 11 and elsewhere: replace \( \text{abort}(e) \) by \( \text{case } e \{ \} \) to avoid the implication that "abort" causes a run-time fault, which it does not (and can not!).

- Page 130, the dynamics is specified to be lazy, though it is possible to specify both eager and lazy semantics using optional premises and rules. The rules given in Figure 1 do so.

- Chapter 33, Section 2: It would be preferable, and consistent with later chapters, to rename the operator \( \text{isof} \) to \( \text{isinref} \).

- Chapter 40 Concurrent Algol. The dynamics specifies \( \text{synchronous} \) communication, in which an emit of a message does not terminate until the
message has been received. An *asynchronous* dynamics may be specified by making the alterations given in Figure 2. These rules define an asynchronous send process, and re-define the dynamics of emit to create such a process.

- Section 33.2. It would be more suggestive to write \( \text{inref}(e_1; e_2) \) in place of \( \text{mk}(e_1; e_2) \) to stress its role as the dynamic analogue of \( \text{in}[a](e) \), much as \( \text{getref}(e) \) is the dynamic analogue of \( \text{get}[a] \) for assignables.

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